CS101 Algorithms and Data Structures

Fall 2021

Homework 5

Due date: 23:59, October 31, 2021

1. Please write your solutions in English.

2. Submit your solutions to gradescope.com.

3. Set your FULL Name to your Chinese name and your STUDENT ID correctly in Account Settings.

4. If you want to submit a handwritten version, scan it clearly. Camscanner is recommended.

5. When submitting, match your solutions to the according problem numbers correctly.

6. No late submission will be accepted.

7. Violations to any of the above may result in zero grade.

Problem 1: Multiple Choices

Multiple Choices: Each question has one or more correct answer(s). Select all the correct answer(s). For each question, you get 0 point if you select one or more wrong answers, but you get half points if you select a non-empty subset of the correct answers.

Note that you should write you answers of Problem 1 in the table below.

Q(1)	Q(2)	Q(3)
C D	D	A B

(1) Which of the following are true about Binary Trees? (2pts)

A. Every full binary tree is also a complete binary tree.

B. A full binary tree with n nodes has height of $O(\ln n)$.

C. A complete binary tree with n leaves has exactly $2n - 1$ nodes.

D. A full binary tree with n leaves has exactly $2n - 1$ nodes.

E. None of the above.

(2) Which traversals of the left tree and right tree, will produce the same sequence node name? (2pts)
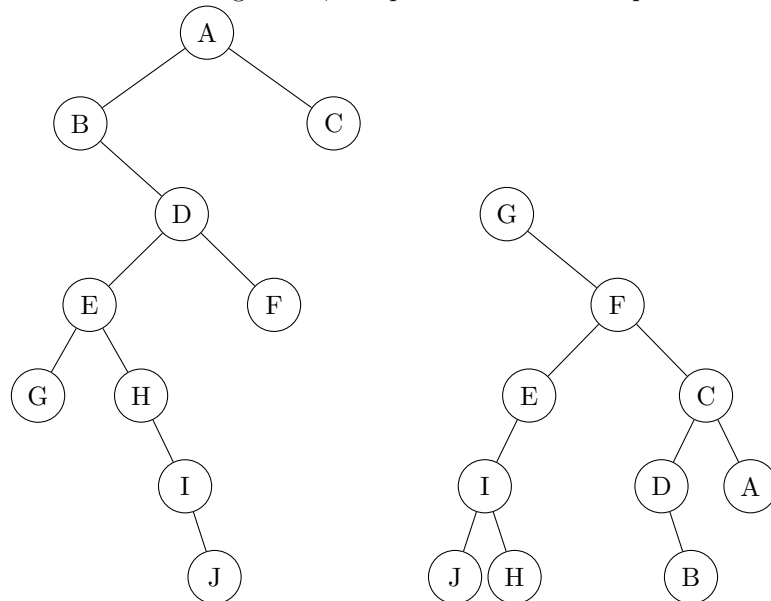
A. left: In-order, right: In-order

B. left: Pre-order, right: Pre-order

C. left: Post-order, right: Post-order

D. left: Post-order, right: In-order

(3) Which of the following statements about the binary heap is **not** true? Note that binary heaps mentioned in this problem are implemented by complete binary trees. (2pts)

A. There exists a heap with seven distinct elements so that the in-order traversal gives the element in sorted order.

B. If item A is an ancestor of item B in a heap (used as a Priority Queue) then it must be the case that the Insert operation for item A occurred before the **Insert** operation for item B.

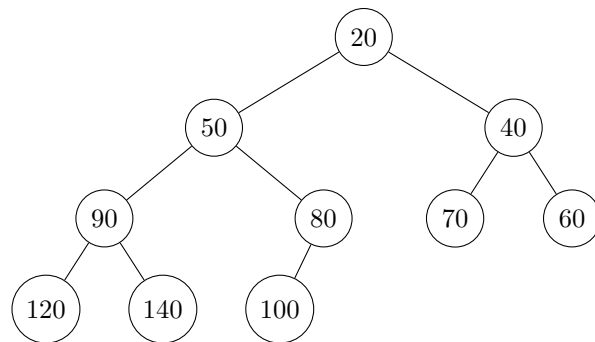C. If array A is sorted from smallest to largest then A (excluding A[0]) corresponds to a min-heap.

D. None of the above.

Problem 2: Construct a Heap

(1) Suppose we construct a min-heap from the following initial heap using Floyd's method. After the construction is completed, what should this heap look like? Please draw it as a binary tree. (5pts)
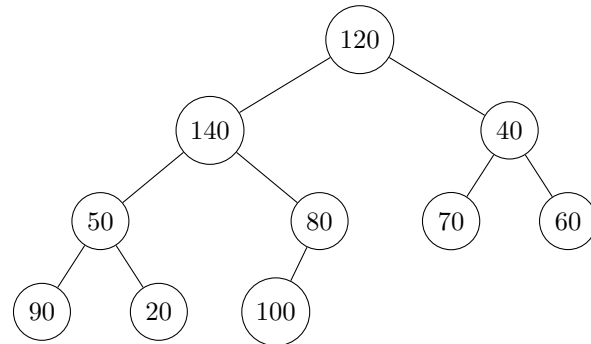
sorted:

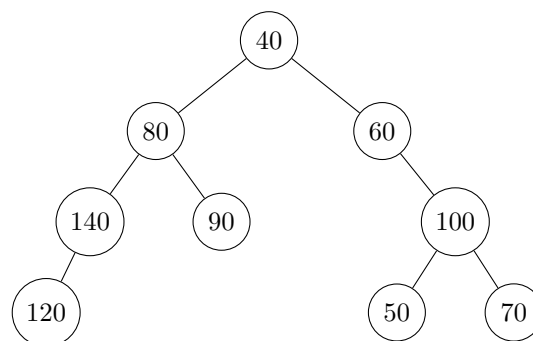(2) Suppose we delete the root from the heap constructed in (1). What will be the post-order traversal of the new heap? (2pts)
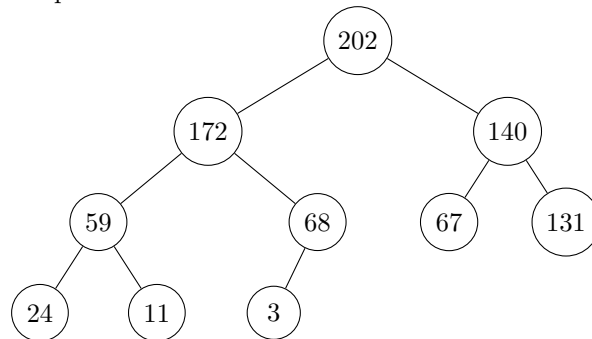**Your answer:** 120, 140, 90, 80, 50, 70, 100, 60, 40.

(3) Suppose there is a binary tree with the same post-order traversal of the heap in (2), and the binary tree has an **in-order** traversal of $120, 140, 80, 90, 40, 60, 50, 100, 70$. Please draw this binary tree below. (5pts)

Problem 3: Heap Sort

You are given such a max heap like this:

Then you need to use array method to show each step of heap sort in increasing order. Fill in the value in the table below. Notice that the value we have put is the step of each value sorted successfully. For each step, you should always make you heap satisfies the requirement of max heap property. (10pts)

index	0	1	2	3	4	5	6	7	8	9	10
value		202	172	140	59	68	67	131	24	11	3

Table 1: The original array to represent max heap.

index	0	1	2	3	4	5	6	7	8	9	10
value		172	68	140	59	3	67	131	24	11	202

Table 2: First value is successfully sorted.

index	0	1	2	3	4	5	6	7	8	9	10
value		140	68	131	59	3	67	11	24	172	202

Table 3: Second value is successfully sorted.

index	0	1	2	3	4	5	6	7	8	9	10
value		131	68	67	59	3	24	11	140	172	202

Table 4: Third value is successfully sorted.

index	0	1	2	3	4	5	6	7	8	9	10
value		68	59	67	11	3	24	131	140	172	202

Table 5: Fourth value is successfully sorted.

index	0	1	2	3	4	5	6	7	8	9	10
value		67	59	24	11	3	68	131	140	172	202

Table 6: Fifth value is successfully sorted.

index	0	1	2	3	4	5	6	7	8	9	10
value		59	11	24	3	67	68	131	140	172	202

Table 7: Sixth value is successfully sorted.

index	0	1	2	3	4	5	6	7	8	9	10
value		24	11	3	59	67	68	131	140	172	202

Table 8: Seventh value is successfully sorted.

index	0	1	2	3	4	5	6	7	8	9	10
value		11	3	24	59	67	68	131	140	172	202

Table 9: Eighth value is successfully sorted.

index	0	1	2	3	4	5	6	7	8	9	10
value		3	11	24	59	67	68	131	140	172	202

Table 10: Last 2 values are successfully sorted.

Problem 4: i-Heap

In class, Prof. Zhang has mentioned the method of the array storage of a binary heap. In order to have a better view of heap, we decide to extend the idea of a binary heap to i-heap. In other words, each node in the heap now has at most i children instead of just two, which is stored in a complete binary tree.

For example, the following heap is a 3-heap.
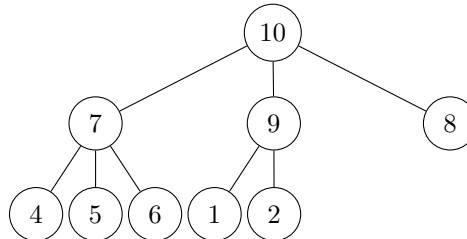
(1) (4pts) If you are given the node with index x, what is the index of its parent and its yth child ($1 \le y \le i$)?

Notice: We assume the root is kept in $A[1]$. For your final answer, please represent it in terms of i, x and y. Please use flooring or ceiling to ensure that your final answer is as tight as possible if you need, i.e. $\lfloor m \rfloor$, $\lceil m \rceil$.

Solution:

Parent:

$$P = \left\lfloor \frac{x + i - 2}{i} \right\rfloor$$

Child:

$$C_y = ix - i + 1 + y$$

(2) (4pts) What is the height of a i-heap of n elements? Please show your steps.

Notice: For your final answer, please represent it in terms of i and n. Please use flooring or ceiling to ensure that your final answer is as tight as possible if you need, i.e. $\lfloor m \rfloor$, $\lceil m \rceil$.

Solution:

$$h = \lfloor \log_i((i-1)n) \rfloor$$

The tth level of the heap has m_t nodes, where:

$$m_t = i^t$$

The root is at 0th level. If each level of the heap is full, the heap should have M_h nodes.(h is the height of the heap.)

$$M_h = \sum_{t=0}^{h} = \frac{i^{h+1} - 1}{i - 1}$$

Then we use mathematical induction to prove the conclusion.

Base case:

When $x = 1$, $h \lfloor \log_i(i-1) \rfloor = 0$, the heap has height 0 and the level is full.

Induction process:

Assume the heap with n nodes has height $h = \lfloor \log_i((i-1)n) \rfloor$

We will show the heap with $n+1$ nodes also satisfies the above equation.

There exists two cases:

1. Each level of the heap with n nodes is full (the tree is complete).
2. The last level of the heap with n nodes is not full.

Case 1:

$$n = \frac{i^{h+1} - 1}{i - 1}$$

Adding one node will cause the height increase.

$$
\begin{aligned}
\lfloor \log_i((i-1)(n+1)) \rfloor &= \lfloor \log_i(i^{h+1} - 1 + i - 1) \rfloor \\
&= \lfloor \log_i\left(i^{h+1} + (i-2)\right) \rfloor \\
&= h + 1
\end{aligned}
$$

Case 2:

Adding the node won't cause the height increase.

$$\frac{i^h - 1}{i - 1} \le n + 1 < \frac{i^{h+1} - 1}{i - 1}$$

$$h < \log_i(i^h - 1) \le \log_i((i-1)(n+1)) < \log_i\left(i^{h+1} - 1\right) < h + 1$$

$$\lfloor \log_i((i-1)(n+1)) \rfloor = h$$

Therefore, by mathematical induction, $h = \lfloor \log_i((i-1)n) \rfloor$ for all $x \ge 1$.

(3) (4pts) Now we want to study which value of i can minimize the comparison complexity of heap-sort. For heap-sort, given a built heap, the worst-case number of comparisons is $\Theta(nhi)$, where $h = \Theta(\log_i n)$ is the height of the heap. Suppose the worst-case number of comparisons is $T(n, i)$. You need to do:

- Explain why $T(n, i) = \Theta(nhi)$.

- Suppose n is fixed, solve for i so that $T(n, i)$ is minimized.

Notice: i is an integer actually. In this problem, we only consider the complexity of comparison, not the accurate number of comparison.

Solution:

In the worst case, all the elements will have to be delivered to the bottom. At each down percolation, the element will be compared i times. The kth level has i^k nodes. Therefore, there are i^k nodes being percolate k times.

$$
\begin{aligned}
T(n) &= \sum_{k=1}^{h} i^k \cdot k \cdot i \\
&= \frac{i^{h+1} \cdot h - 1}{i - 1} \cdot i \\
&= \left(\frac{(i^{h+1} - 1) \cdot h}{i - 1} + \frac{h}{i - 1} \right) \cdot i - \frac{i}{i - 1} \\
&= \left(nhi + \frac{hi}{i - 1} - \frac{i}{i - 1} \right) \\
&= \Theta\left(nhi + h \right) \\
&= \Theta\left(nhi \right)
\end{aligned}
$$

We use derivatives to find minumum.

$$
\begin{aligned}
(nhi)' &= (n \log_i ni)' \\
&= n \left((\log_i n)' \cdot i + \log_i n \cdot i' \right) \\
&= n \cdot \left(-\frac{\ln n}{\ln^2 i \cdot i} \cdot i + \log_i n \right) \\
&= n \left(\log_i n - \frac{\ln n}{\ln^2 i} \right) \\
&= n \left(\frac{\ln n \cdot \ln i - \ln n}{\ln^2 i} \right) \\
\ln i &= 1 \\
i &= e
\end{aligned}
$$

Since i is an integer, we should compare $i_1 = 2$ and $i_2 = 3$.

$$
\begin{aligned}
\frac{n \log_{i_2} n \cdot i_2}{n \log_{i_1} n \cdot i_1} &= \frac{3 \log_3 n \cdot 3}{2 \log_2 n \cdot 2} \\
&= \frac{3}{2} \frac{\frac{\ln n}{\ln 3}}{\frac{\ln n}{\ln 2}} \\
&= \frac{3}{2} \log_3 2 < 1
\end{aligned}
$$

Therefore, when $i = 3$, $T(n, i)$ is minimized.

(4) (5pts) TA Xiaozhang Master has a new idea that if $i = 1$, we only need to do the `BUILD-HEAP` operation to get a sorted array. Since we know from the lecture that `BUILD-HEAP` takes $O(n)$ time, he thinks it can actually sort in $O(n)$ time!

Now you are the student of TA Xiaozhang Master, and you need to judge his statement is true or false.

- If his statement is true, please explain the reason and prove its correctness.

- If his statement is false, please help him find the fallacy in his argument with your own reason. In the heap-sort he proposes, what sorting algorithm is actually performed? What is the worst running time of such a sorting algorithm?

Solution:
The statement is false, for that the complexity of $O(n)$ we examined on the class is only when $i = 2$. When $i = 1$, the build heap operation deteriorate to insertion sort. The time complexity should be $O(n^2)$ for the worst case.

In the worst case:
Only one node doesn't need percolation. For the rest $n - 1$ nodes, the ith nodes need $i - 1$ percolation.

$$\sum_{k=1}^{n-1} k = \frac{n(n-1)}{2}$$

The time complexity in the worst case is $O(n^2)$.
In the best case:
No node need percolation. Time complexity is $O(n)$.
On average, the toatl percolation times equals to the inversion of the array. The time complexity is $O(n+d)$, d is the inversion of the array. When $d = O(n)$, the sorting can be complete in $O(n)$ time.
Therefore, he actually performed insertion sort and the worst-case complexity is $O(n^2)$.