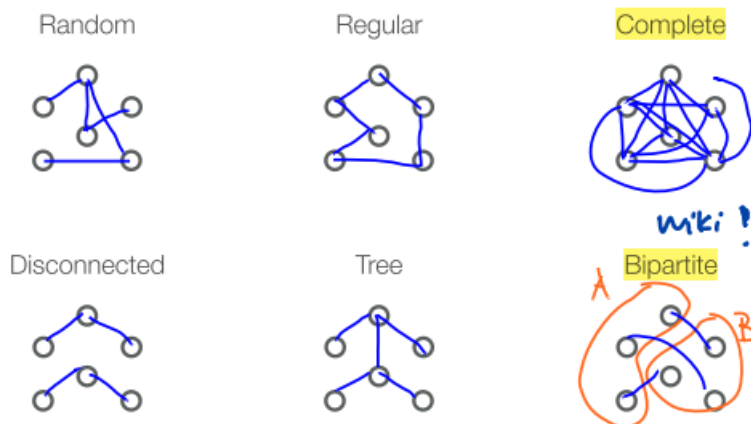


# 02\_Structural Analysis

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## 1. Some Graph Knowledge

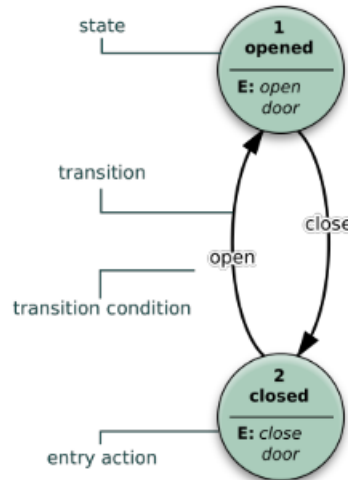


## 2. Finite State Automaton

A state diagram or state-transition graph represents a Finite Automaton

Nodes  $\rightarrow$  States  
 Edges  $\rightarrow$  Transitions  
 $\tau: U \rightarrow \{0, 1\}$   
 $\tau_{12}(u) = \begin{cases} 1 & u = \text{"close"} \\ 0 & u = \text{"open"} \end{cases}$

See for instance: Booth, T.L., 1967. Sequential machines and automata theory. John Wiley & Sons.



where  $\tau$  is some kind of transition activation function

### 3. Bi-Partite Graph and Regular Graph

#### Definition: Bi-Partite Graph

A bipartite graph, also called a bigraph, is a set of graph vertices decomposed into **two disjoint sets** such that **no two graph vertices within the same set are adjacent**.

#### Definition: Regular Graph

A regular graph is a graph where each vertex has the same number of neighbors; i.e. every vertex has the same degree or valency

### Matching in Bi-Partite Graph

#### Definition: Matching

A matching in a Bipartite Graph is a set of the edges chosen in such a way that no two edges share an endpoint.

#### Definition: Maximum Matching

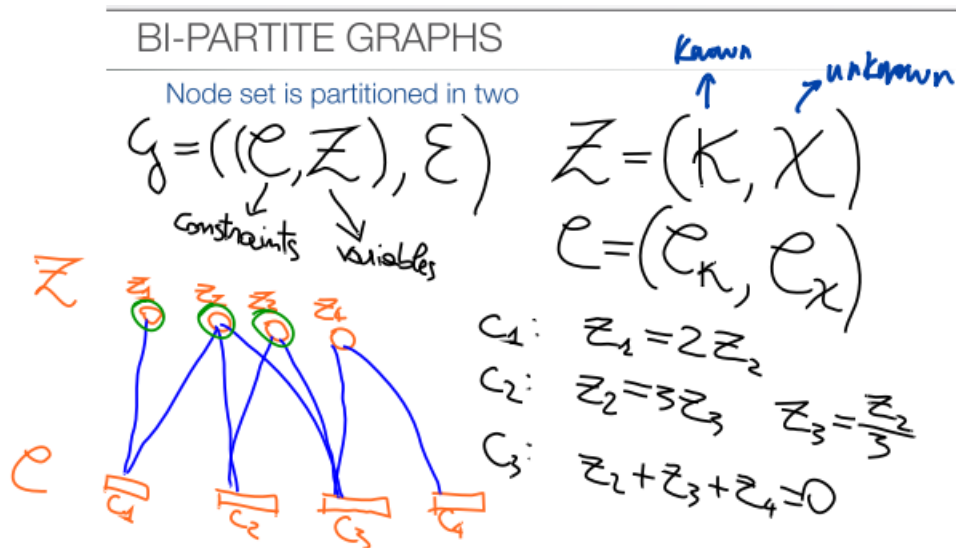
A maximum matching is a matching of maximum size (maximum number of edges)

#### Other Definitions:

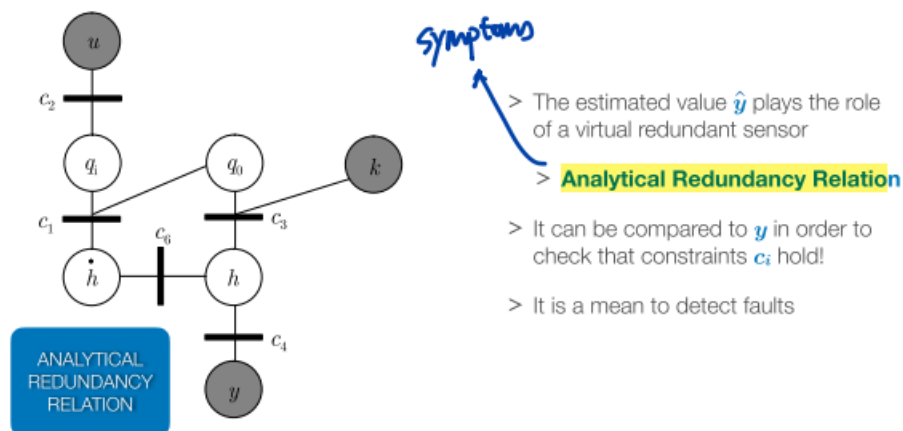
- An edge is said to be weak with respect to  $M$  if it does not belong to  $M$ . A vertex is weak with respect to  $M$  if it is only incident to weak edges
- An  $M$ -augmenting path is an alternating path whose end vertices are both weak with respect to  $M$

A matching  $M$  in a graph  $G$  is maximum **if and only if** there exists no  $M$ -augmenting path in  $G$

### Bi-Partite Usage 1: Define Variable Relations



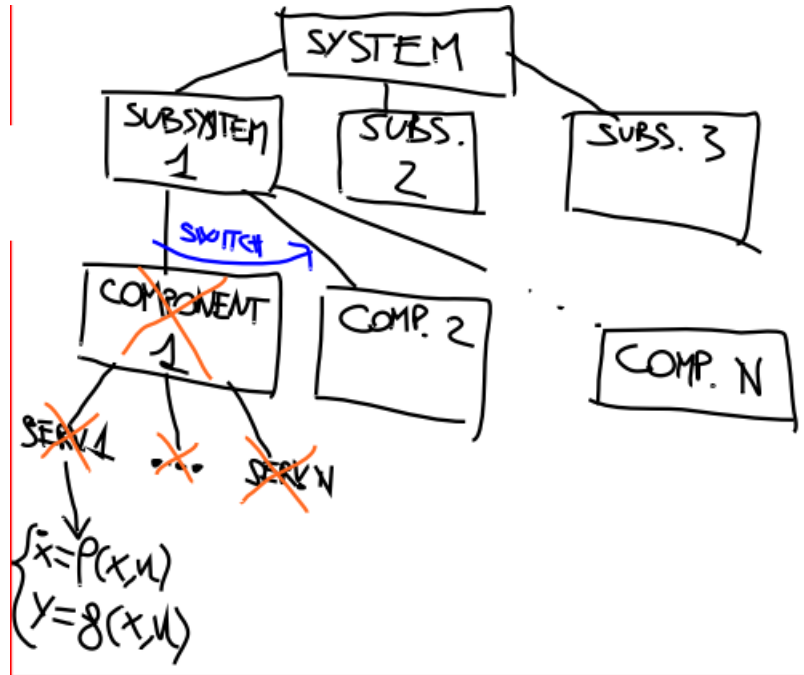
## Bi-Partite Usage 2: Discover Analytical Redundancy Relation



The output  $y$  is measured, but can also be computed from  $u$  and  $k$

### 3. Components and Services

Components and Service model can be used to **organize knowledge** of the system in a **hierarchical way**



## Components

- **Components** provide **services**
- e.g. A tank integrates inflow - outflow to produce a stored mass

## Service

- A service  $s_i$  is described by a 6-tuple

$$\begin{aligned}
 S &= \{cons, prod, proc, r_{gst}, enable, res\} \\
 cons &= \{q_i, q_o\} \\
 prod &= \{h\} \\
 proc &= \{h = q_i - q_o; h = \int h dt + h_o\} \\
 r_{gst} &= \{1\} \\
 enable &= \{1\} \\
 res &= \{vessel, pipes\}
 \end{aligned}$$

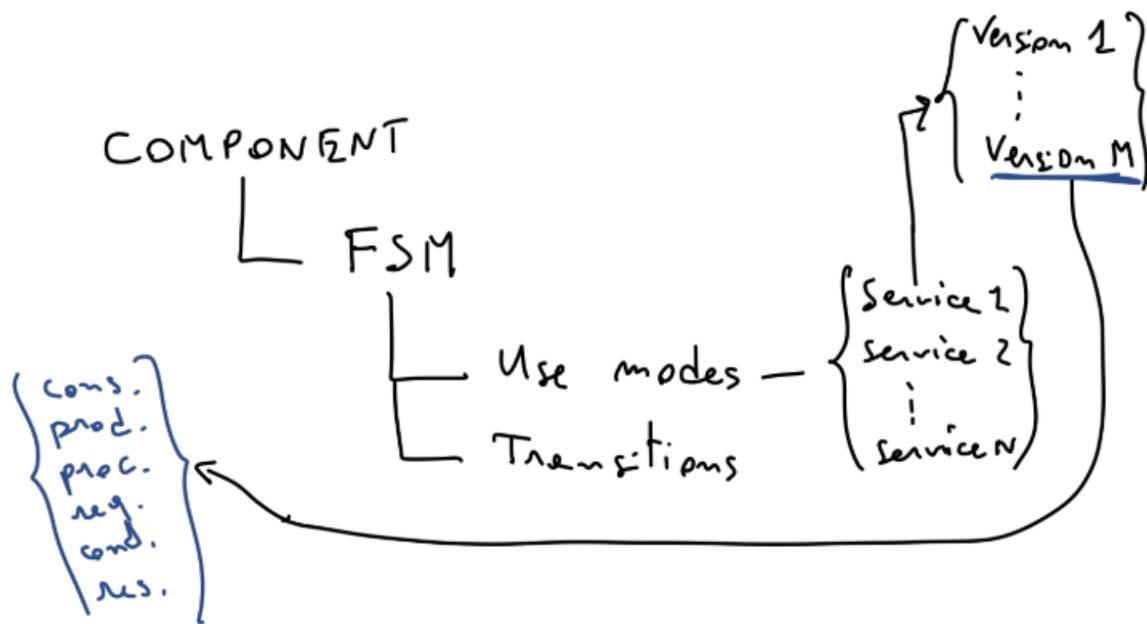
- Consumed Variables
- Produced Variables
- Processes
- Request Signals
- Enable Signals
- Resources
- A services come in different **versions**
- Availability of services depends on component **use mode**

## General Component Model

$\langle \text{component } k \rangle ::= \langle \text{state transition graph } G(M(k), \tau(k), m^0(k)) \rangle$

where

$\langle M(k) \rangle ::= \langle \text{set of use-modes } \{m_i(k), i \in I_m(k)\} \rangle$   
 $\langle \tau(k) \rangle ::= \langle \text{set of transitions } \{\tau_{ij}(k), i, j \in I_m(k)\} \rangle$   
 $\langle m^0(k) \rangle ::= \langle \text{initial use-mode} \rangle$   
 $\langle \text{use-mode } m_i(k) \rangle ::= \langle \text{set of services } S_i(k) \subseteq S(k) \rangle$   
 $\langle \text{service } s_l(k) \rangle ::= \langle \text{pre-ordered versions} \rangle$   
 $\qquad \qquad \qquad \{s_l^j(k), j \in J(s_l(k))\} >$   
 $\langle \text{version } s_l^j(k) \rangle ::= \langle \text{consumed vars } cons_l^j(k),$   
 $\qquad \qquad \qquad \text{produced vars } prod_l^j(k),$   
 $\qquad \qquad \qquad \text{procedures } proc_l^j(k), \text{ request } req_{st_l}^j(k),$   
 $\qquad \qquad \qquad \text{activation cond. } activ_l^j(k),$   
 $\qquad \qquad \qquad \text{hardware and software resources } res_l^j(k) \rangle$   
 $\langle \text{transition } \tau_{ij}(k) \rangle ::= \langle \text{condition } c_{ij}(k), \text{ origin } m_i(k),$   
 $\qquad \qquad \qquad \text{destination } m_j(k) \rangle .$



## Usage

- For **fault diagnosis**
  - As a simple way to do diagnosis of complex systems where components are described by discrete states (e.g. a component providing a critical service is offline instead than online)
- For fault accommodation via switching of hardware redundant components
  - If another component can provide the same service, you know that you can swap it in

## 3. FTA: Fault Tree Analysis

A graphical way to analyze **system dependability properties**

### Assumed Knowledge

- Components and services, how they are interconnected
- A structural, qualitative knowledge is enough

## Fault Tree

### Fault Tree Illustration

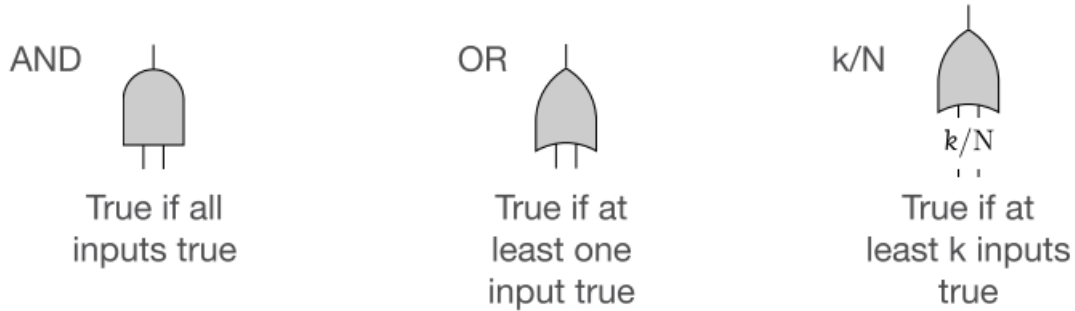
- A **Directed Acyclic Graph**
- **Leaves** represent failure of components (**Basic Events**)

- Components relations are via **Logic Gates**
- **Root** node is the Top Event: System Failure
- fault tree is used to **model fault** not normality!

### Events

- binary variables: **true means failure, false means healthy**

### Gates



### Definition: Fault Tree

It is formally defined as a **4-tuple**

$$F = \langle BE, G, T, I \rangle$$

- **BE** is the set of **basic events** (binary: true means failure, false means healthy)
- **G** is the set of gates
- $E = BE \cup G$  is the set of events
- $T : G \rightarrow \{AND, OR, NOT, K/N\}$  describes gate types
- $I : G \rightarrow \mathcal{P}(E)$  maps gates to their inputs (basically represents the graph's edges set)
  - $\mathcal{P}(X)$  stands for power set of a set  $X$ , i.e. set of all possible subsets of  $X$
  - If  $T(g) = k/N$ , then  $|I(g)| = N$

## Fault Tree Analysis

### Semantic Function of a Fault Tree

To evaluate a  $F$ , a semantic (i.e. logic) function is introduced

$$\pi_F : \mathcal{P}(BE) \times E \mapsto \{0, 1\}$$

- $\pi_F(S, e)$  is true if the event  $e$  is "failure" when all the basic elements in  $S$  are failed

- We use the shorthand notation  $\pi_F(S)$  when  $e$  is the top event  
in this case the function tells us **whether the system as a whole** will fail assuming a given subset of components are failed

### Cut Sets and Minimal Cut Sets

$C \subseteq BE$  is a cut set of  $FTF$  if  $\pi_F(C) = 1$ . A minimal cut set (MCS) is a cut set of which no subset is a cut set, i.e. formally  $C \subseteq BE$  is an MCS if  $\pi_F(C) = 1 \wedge \forall C' \subset C : \pi_F(C') = 0$

### Structure Function.

Let us introduce the  $FT$  structure function  $f$

$$f : \{0, 1\}^{|BE|} \mapsto \{0, 1\}$$

$f(e_1, \dots, e_N)$  is **true** when the failure values  $e_1, \dots, e_N$  for the  $N$  Basic Elements will result in a failure at the top event

### Disjoint Normal Form

A DNF is a disjunction of several conjunction, e.g.:

$$f(A, B, C, D) = (A \wedge B) \vee (C \wedge D)$$

Let us express  $f$  as a **Disjoint Normal Form**, then every conjunction in  $f$  is a MCS

## Other Form of Fault Tree

There are many other form of fault trees, such as

### Probabilistic Fault Trees

- You **propagate probabilities** of failures rather than deterministic failure events
- Good to estimate system MTTF or reliability

## 4. FMEA: Failure Mode and Effects Analysis

FMEA is a method to analyze system failure scenarios

- **How** can the failure of components lead to failure of the whole system
- the main difference with the FTA lies in the “**how**”

### Assumed Knowledge

- **Components and services**, how they are **interconnected**
- **Fault/failure modes** of each components
- **Effects** of each mode at component level
- **How** each effect will **affect connected components**
- A structural, qualitative knowledge is enough

### Faults, effects and propagation

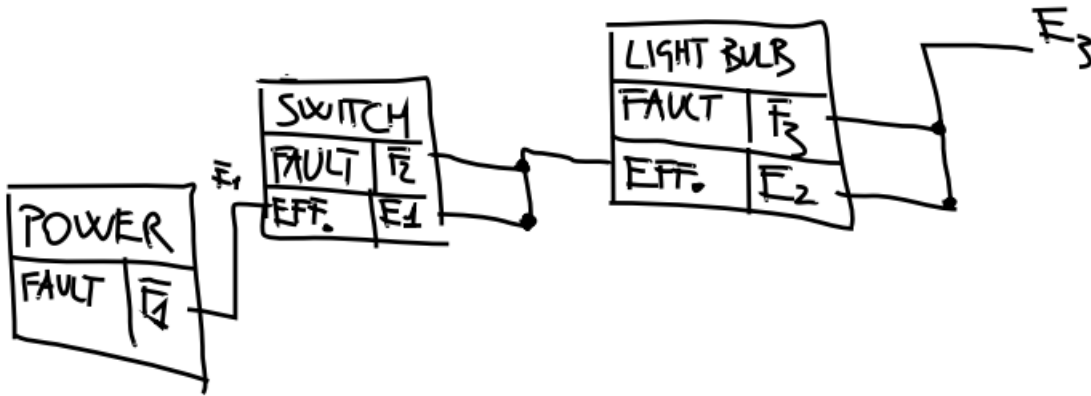


Each components can fail to provide its service either because:

- internal faults (basic event in FT)
- **propagation** of faults from components it depends on

## Graphical Representation

### Graphical Book Representation



## Boolean Matrix Representation

### Matrix Representation

The **fault propagation** is actually a boolean mapping between fault vector  $f$  and effect vector  $e$

- We can formally represent it via a boolean matrix  $M$  and a boolean operation  $\otimes$

$$e = M \otimes f$$

- The operator  $\otimes$  implements a disjoint normal form, as in the FTA

$$e_{(i)} = (M_{(i,1)} \wedge f_{(1)}) \vee (M_{(i,2)} \wedge f_{(2)}) \vee \dots \vee (M_{(i,N)} \wedge f_{(N)})$$

- For hardware redundancy, we need to model it in the  $f$ s

### Hierarchical Composition

$$e_{loc} = M_{loc} \otimes \begin{bmatrix} f_{loc} \\ e_{anc} \end{bmatrix}$$

$$e_{loc} = \left( M_{loc} \otimes \begin{bmatrix} I & 0 \\ 0 & M_{anc} \end{bmatrix} \right) \otimes \begin{bmatrix} f_{loc} \\ f_{anc} \end{bmatrix}$$

## Inverse Inference

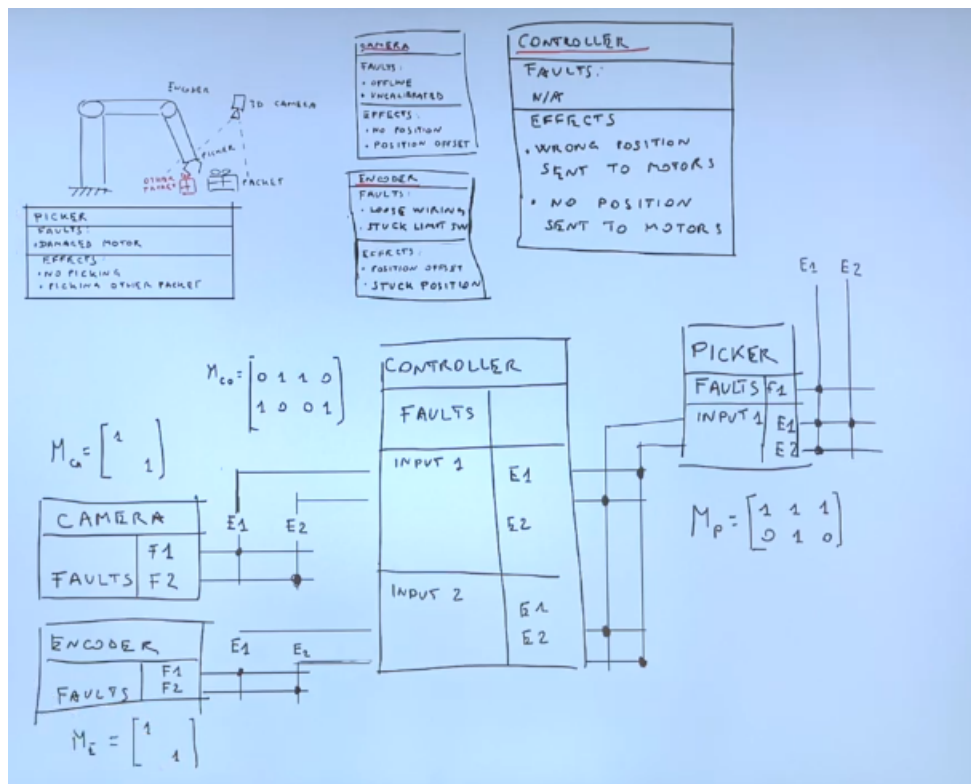
from effects to possible causes

$$f = M^T \odot e$$

The  $\odot$  is defined as:

$$f_{(i)} = \left( M_{(i,1)}^T == e_{(1)} \right) \wedge \left( M_{(i,2)}^T == e_{(2)} \right) \wedge \dots \wedge \left( M_{(i,N)}^T == e_{(N)} \right)$$

## **Example:**



$$\begin{aligned}
e_P &= M_P \otimes \begin{bmatrix} f_P \\ e_{co} \end{bmatrix} & e_{co} &= M_{co} \otimes \begin{bmatrix} e_{ca} \\ e_E \end{bmatrix} \\
e_P &= M_P \otimes \left[ \begin{array}{c|c} I & \\ \hline & M_{co} \end{array} \right] \otimes \begin{bmatrix} f_P \\ e_{ca} \\ e_E \end{bmatrix} & e_{ca} &= M_{ca} \otimes f_{ca} \\
& & e_E &= M_E \otimes f_E \\
\text{Also } e_{co} &= M_{co} \otimes \left[ \begin{array}{c|c} M_{ca} & \\ \hline & M_E \end{array} \right] \otimes \begin{bmatrix} f_{ca} \\ f_E \end{bmatrix} \\
e_P &= M_P \otimes \left[ \begin{array}{c|c} I & \\ \hline & M_{co} \otimes \left[ \begin{array}{c|c} M_{ca} & \\ \hline & M_E \end{array} \right] \end{array} \right] \otimes \begin{bmatrix} f_P \\ f_{ca} \\ f_E \end{bmatrix} \\
&= M_P \otimes \left[ \begin{array}{c|c} I & \\ \hline & M_{co} \end{array} \right] \otimes \left[ \begin{array}{c|c} I & \\ \hline & M_{ca} \quad M_E \end{array} \right] \otimes \begin{bmatrix} f_P \\ f_{ca} \\ f_E \end{bmatrix} \\
&= \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 0 \end{bmatrix} \otimes \begin{bmatrix} 1 & 0 & 1 & 1 & 0 \\ & 1 & 0 & 0 & 1 \end{bmatrix} \otimes \begin{bmatrix} 1 & & & & \\ & 1 & & & \\ & & 1 & & \\ & & & 1 & \\ & & & & 1 \end{bmatrix} \otimes \begin{bmatrix} f_P \\ f_{ca} \\ f_E \end{bmatrix} \}
\end{aligned}$$

$$\begin{aligned}
&\begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 0 \end{bmatrix} \otimes \begin{bmatrix} 1 & 0 & 1 & 1 & 0 \\ & 1 & 0 & 0 & 1 \end{bmatrix} \otimes \begin{bmatrix} 1 & & & & \\ & 1 & & & \\ & & 1 & & \\ & & & 1 & \\ & & & & 1 \end{bmatrix} \\
&= \begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 & 0 \end{bmatrix} \\
e_{P1} &: f_P \vee f_{ca1} \vee f_{ca2} \vee f_{E1} \vee f_{E2} \\
e_{P2} &: f_{ca2} \vee f_{E1}
\end{aligned}$$

## Summary

- Some Graph Knowledge
  - **Bipartite Graph**
- Finite State Automaton
- Bipartite Graph
  - Define Variable Relations
  - **Discover Analytical Redundancy Relations**
- Components and Service

- **General Component Model**
  - Usage: Fault Diagnosis
- FTA: Fault Tree Analysis
  - Fault Tree
  - Cut set and minimal cut set
  - **Fault Tree Analysis: Semantic Model + Structure Function + DNF**
- FMEA: Failure Mode and Effects Analysis
  - **Graphical Representation**
  - **Boolean Matrix Representation + Inverse Inference**