

# Shallows Explanation

## Explanation

The Shallows task can be briefly concluded as *finding the shortest path*. Though there are few differences between Shallows and normal shortest path problem, the basic idea of solution is almost the same.

Firstly, find the *starting point*, in this case, will be the origin port. Step1

Then, we *add lines* which depart port is the origin port. Step2

Next, we add more lines which connect the lines added in step2. Step3

At the end, we have create *all possible routines* that starts with origin port and end with every port. Find out the "shortest path" Step4

While connecting ports, update the information in results array.

```
int i;  
int [] results = new int[ports];    //initlialize the results array  
for(i=0;i<ports;i++){  
    results[i]=0;  
}  
int j;  
int temp = 0;
```

At the beginning of our solution, we create an array results[ ] to store the result. And it will be initialized with 0.

```
for(i=0;i<lanes.length;i++){  
    if(lanes[i].depart==origin){  
        results[lanes[i].arrive]=lanes[i].depth;  
        count++;  
    }  
}
```

Then, we use a for loop to find every line which depart with origin port. And we store the depth in the corresponding position of results[ ]. And count the number of lines we find in this step.

```

    int k = 0;
    while(k<ports-count){
        for(i=0;i<ports;i++){
            if(i == origin){
                results[i] = Integer.MAX_VALUE;
                continue;
            }
            if(results[i]==0)
                continue;
            for(j=0;j<lanes.length;j++){
                if(lanes[j].depart==i){
                    if(results[i]<lanes[j].depth && results[i]!=0)
                        temp = results[i];
                    else
                        temp = lanes[j].depth;
                    if(temp>results[lanes[j].arrive])
                        results[lanes[j].arrive]=temp;
                }
            }
        }
        k++;
    }
}

```

For the example given,

Consider maximumDraughts(5, lanes, 0) with lanes of the form {depart, arrive, depth} as follow:

{0, 1, 9} : Lane from port 0 to port 1 with shallow point 9 units deep  
 {0, 2, 2} : Lane from port 0 to port 2 with shallow point 2 units deep  
 {0, 3, 1} : Lane from port 0 to port 3 with shallow point 1 units deep  
 {1, 2, 7} : Lane from port 1 to port 2 with shallow point 7 units deep  
 {1, 3, 2} : Lane from port 1 to port 3 with shallow point 2 units deep  
 {2, 3, 8} : Lane from port 2 to port 3 with shallow point 8 units deep  
 {4, 2, 9} : Lane from port 4 to port 2 with shallow point 9 units deep

Now we have results[ ] which contain {0,9,2,1,0}

Meaning that:

0-0: 0

0-1: 9

0-2: 2

0-3: 1

0-4: 0

And we define 0-0 as interger.maxsize

Then, we ignore all *non-existed routines* (for now) , in this case we use

0-1: 9

0-2: 2

0-3: 1

Next, for port =1, we find every line that connect to port1:

{1, 2, 7} : Lane from port 1 to port 2 with shallow point 7 units deep

{1, 3, 2} : Lane from port 1 to port 3 with shallow point 2 units deep

We now have

0-1-2 = 0-2: 7

0-1-3 = 0-3: 2

The results[ ] will be updated according to this:

0-1: 9

0-2: 7

0-3: 2

Also, we repeat this process for port = 2, 3.

That's what *nested for loops* do in the code above.

If there is a routine which make port 0 connect to a new port (port = 4), we also update that in result array.

And we need a *big while loop* to include all of this, which will circle for the number of “non-existed routines”. (Considering *the worst case*) Cause after each time we finish the nested loops, the origin ports might now be able to *reach to a new port*, meaning that those “non-existed routines” perhaps are now existed.

At the end we return the results[ ].

## Time complexity

$O(P^2 L)$