ECAP5-DPROC

RISC-V processor

Architecture Document

Revision: DRAFT-1.0.0

Issue Date: June 15, 2023



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1 Introduction

1.1 Purpose

This documents aims at defining the requirements for ECAP5-DPROC as well as describing its architecture. Both user and product requirements will be covered.

1.2 Intended Audience and Use

This document targets hardware engineers who shall implement ECAP5-DPROC by refering to the described architecture. It is also intended for system engineers working on the integration of ECAP5-DPROC in ECAP5. Finally, this document shall be used as a technical reference by software engineers configuring ECAP5-DPROC through hardware-software interfaces.

1.3 Product Scope

ECAP5-DPROC is an implementation of the RISC-V instruction set architecture targetting *Educational Computer Architecture Platform 5* (ECAP5). It will provide the main means of software execution in ECAP5.

1.4 Definitions and Acronyms

hardware-configurable

software-configurable

The bit indexing shall be described somewhere.

Byte size as well.

1.5 References

Date	Version	Title
December 13, 2019	20191213	The RISC-V Instruction Set Manual Volume I: User-Level ISA
March 22, 2019	0.13.2	RISC-V External Debug Support
September 7, 2002	B.3	WISHBONE System-on-Chip (SoC) Interconnection Architecture for Portable IP Core

2 Overall Description

2.1 User needs

ECAP5 is the primary user for ECAP5-DPROC. ECAP5-DPROC could however be used as a standalone RISC-V processor. The following requirements define the user needs.

ID	U_INSTRUCTION_SET_01
Description	ECAP5-DPROC shall implement the RV32I instruction set.

In order to improve the usability of ECAP5-DPROC, it shall have a *von Neumann* architecture as it only requires one memory interface.

ID	U_MEMORY_INTERFACE_01
Description	ECAP5-DPROC shall access both instructions and data through a unique memory interface.

ID	U_MEMORY_INTERFACE_02
Description	ECAP5-DPROC's unique memory interface shall be compliant with the Wishbone specification.

ID	U_MEMORY_INTERFACE_03
Description	ECAP5-DPROC's unique memory interface shall be designed such that memory protocols can be interchanged at compile time.

ID	U_RESET_01
Description	ECAP5-DPROC shall provide a signal which shall hold ECAP5-DPROC in a reset state while asserted.

The polarity of the reset signal mentionned in U_RESET_01 is not specified by the user.

ID	U_BOOT_ADDRESS_01
Description	The address at which ECAP5-DPROC jumps after the reset signal is deasserted shall be hardware-configurable.

The address mentionned in $U_BOOT_ADDRESS_01$ can be either configured through hardware signals or can be selected at compile time.

ID	U_HARDWARE_INTERRUPT_01
Description	ECAP5-DPROC shall provide an signal which shall interrupt ECAP5-DPROC's execution flow while asserted.

ID	U_HARDWARE_INTERRUPT_02
Description	ECAP5-DPROC shall jump to a software-configurable address when it is interrupted.

The memory address at which ECAP5-DPROC shall jump to when interrupted is not specified by the user.

ID	U_DEBUG_01
Description	ECAP5-DPROC shall be compliant with the RISC-V External Debug Support specification.

There is no performance goal required by ECAP5 for ECAP5-DPROC as ECAP5 is an educational platform.

2.2 Assumptions and Dependencies

Describe what the assumptions for the product are: Targeting the ecp5 family, based around opensource toolchains.

3 Requirements

3.1 External Interface Requirements

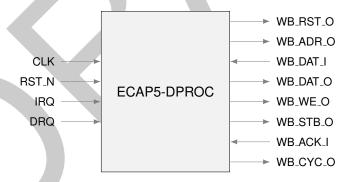


Figure 1: Schematic view of the external interface of ECAP5-DPROC

NAME	TYPE	WIDTH	DESCRIPTION
CLK	I	1	Clock input.
RST₋N	I	1	Hardware reset. Active low.
IRQ	I	1	External interrupt request.
DRQ	I	1	Debug request.

Table 1: ECAP5-DPROC control signals

NAME	TYPE	WIDTH	DESCRIPTION
WB_RST_O	0	1	TBD
WB_ADR_O	0	32	TBD
WB_DAT_I	I	32	TBD
WB_DAT_O	0	32	TBD
WB_WE_O	0	1	TBD
WB_STB_O	0		TBD
WB_ACK_I	I	1	TBD
WB_CYC_O	0		TBD

Table 2: ECAP5-DPROC memory interface signals

ID	I_CLK_01
Description	All inputs and outputs of ECAP5-DPROC shall belong to CLK's clock domain.
Refers to	

ID	I_RESET_01
Description	The RST_N signal shall hold ECAP5-DPROC in a reset state while asserted.
Refers to	U_RESET_01

ID	I_RESET_02
Description	RST_N polarity shall be active low.
Refers to	

ID	I_IRQ_01
Description	ECAP5-DPROC shall jump to a software-configurable address when input IRQ is asserted.
Refers to	U_HARDWARE_INTERRUPT_01, U_HARDWARE_INTERRUPT_02

ID	I_DIRQ_01
Description	TBD
Refers to	

ID	I_MEMORY_INTERFACE_01
Description	Signals from table 2 shall be compliant with the Wishbone specification.
Refers to	U_MEMORY_INTERFACE_02

Behavioral specification for symbols in table 2 is outlined in the functional requirements

section, subsection 3.2.5.

3.2 Functional Requirements

3.2.1 Register file

ID	F_REGISTERS_01
Description	ECAP5-DPROC shall implement 31 user-accessible general purpose registers ranging from $x0$ to $x31$.
Refers to	U_INSTRUCTION_SET_01

ID	F_REGISTERS_02
Description	Register x0 shall be hardwired to the constant zero.
Refers to	U_INSTRUCTION_SET_01

ID	F_REGISTERS_03
Description	ECAP5-DPROC shall implement a pc user-accessible register storing the address of the current instruction.
Refers to	U_INSTRUCTION_SET_01

3.2.2 Instruction decoding

Figure 2 outlines the different instruction encodings for the RV32I instruction set. The opcode parameter is a unique identifier for each instruction. The instruction encoding is infered from the opcode as there can only be one encoding per opcode.

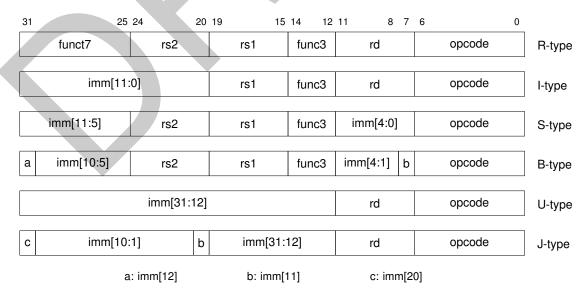


Figure 2: Instruction encodings of the RV32I instruction set

Immediate encoding

Only one immediate value can be encoded in one instruction. The value can be reconstructed from fragments of the following format : imm[x] representing the x^{th} bit or imm[x:y] representing bits from the x^{th} to the y^{th} both included.

ID	F_INSTR_IMMEDIATE_01
Description	Immediate values shall be sign-extended.
Refers to	U_INSTRUCTION_SET_01

ID	F_INSTR_IMMEDIATE_02
Description	The value of an instruction immediate shall be the concatenation of immediate fragments from the instruction encoding.
Refers to	U_INSTRUCTION_SET_01

ID	F_INSTR_IMMEDIATE_03					
Description	Missing immediate fragments shall be replaced by zeros.					
Refers to	U_INSTRUCTION_SET_01					

RV32I is called a Load/Store ISA, meaning that instructions inputs and outputs are passed through registers or through an instruction immediate. There are specific instructions for loading and storing data into memory.

Instruction inputs

ID	F_INSTR_FIRST_INPUT_01					
Description	Instructions encoded using the R-type, I-type, S-type and B-type shall take as their first input the value stored in the register designated by the rs1 parameter.					
Refers to	U_INSTRUCTION_SET_01					

ID	F_INSTR_FIRST_INPUT_02					
Description	Instructions encoded using the U-type and J-type shall take as their first input the immediate value encoded in the instruction.					
Refers to	U_INSTRUCTION_SET_01					

ID	F_INSTR_SECOND_INPUT_01					
Description	nstructions encoded using the R-type, S-type and B-type shall take as their econd input the value stored in the register designated by the rs2 parameter.					
Refers to	U_INSTRUCTION_SET_01					

ID	F_INSTR_SECOND_INPUT_02					
Description	Instructions encoded using the I-type shall take as its second input the immediate value encoded in the instruction.					
Refers to	U_INSTRUCTION_SET_01					

ID	F_INSTR_THIRD_INPUT_01					
Description	Instructions encoded using the S-type and B-type shall take as their third input the immediate value encoded in the instruction.					
Refers to	U_INSTRUCTION_SET_01					

Instruction outputs

ID	F_INSTR_OUTPUT_01						
Description	Instructions encoded using the R-type, I-type, U-type and J-type shall store their result in the register designated by the rd parameter.						
Refers to	U_INSTRUCTION_SET_01						

ID	F_INSTR_OUTPUT_02						
Description	Instructions encoded using the S-type and B-type do not produce any result.						
Refers to	U_INSTRUCTION_SET_01						

Instruction variants

ID	F_INSTR_VARIANT_01
Description	Instructions encoded using the R-type, I-type, S-type and B-type shall use the func3 parameter as a behavior variant selector.
Refers to	U_INSTRUCTION_SET_01

ID	F_INSTR_VARIANT_02					
Description	Instructions encoded using the R-type shall use the func7 parameter as a secondary behavior variant selector.					
Refers to	U_INSTRUCTION_SET_01					

Opcodes

Table 3 outlines the different opcodes values of the RV32I instruction set. Cells marked as *noimp* are for opcodes that are not implemented in ECAP5-DPROC.

opcode[1:0]	11							
opcode[4:2]	000	001	010	011	100	101	110	111
opcode[6:5]	000	001	010	011	100	101	110	111
00	LOAD	noimp	noimp	MISC-MEM	OP-IMM	AUIPC	noimp	noimp
01	STORE	noimp	noimp	noimp	OP	LUI	noimp	noimp
10	noimp	noimp	noimp	noimp	noimp	noimp	noimp	noimp
11	BRANCH	JALR	noimp	JAL	SYSTEM	noimp	noimp	noimp

Table 3: Opcode values for the RV32I instruction set.

ID	F_OPCODE_ENCODING_01
Description	Instructions which use the LUI opcode shall be decoded as an U-type instruction.
Refers to	U_INSTRUCTION_SET_01

ID	F_OPCODE_ENCODING_02
Description	Instructions which use the AUIPC opcode shall be decoded as an U-type instruction.
Refers to	U_INSTRUCTION_SET_01

ID	F_OPCODE_ENCODING_03
Description	Instructions which use the JAL opcode shall be decoded as a J-type instruction.
Refers to	U_INSTRUCTION_SET_01

ID	F_OPCODE_ENCODING_04
Description	Instructions which use the JALR opcode shall be decoded as an I-type instruction.
Refers to	U_INSTRUCTION_SET_01

ID	F_OPCODE_ENCODING_05
Description	Instructions which use the BRANCH opcode shall be decoded as a B-type instruction.
Refers to	U_INSTRUCTION_SET_01

ID	F_OPCODE_ENCODING_06
Description	Instructions which use the LOAD opcode shall be decoded as an I-type instruction.
Refers to	U_INSTRUCTION_SET_01

ID	F_OPCODE_ENCODING_07
Description	Instructions which use the STORE opcode shall be decoded as a S-type instruction.
Refers to	U_INSTRUCTION_SET_01

ID	F_OPCODE_ENCODING_08
Description	Instructions which use the OP-IMM opcode shall be decoded as an I-type instruction.
Refers to	U_INSTRUCTION_SET_01

ID	F_OPCODE_ENCODING_09
Description	Instructions which use the OP opcode shall be decoded as a R-type instruction.
Refers to	U_INSTRUCTION_SET_01

ID	F_OPCODE_ENCODING_10
Description	Instructions which use the MISC-MEM opcode shall be decoded as an I-type instruction.
Refers to	U_INSTRUCTION_SET_01

ID	F_OPCODE_ENCODING_11
Description	Instructions which use the SYSTEM opcode shall be decoded as an I-type instruction.
Refers to	U_INSTRUCTION_SET_01

3.2.3 Instructions behaviors

LUI

ID	F_LUI_01
Description	The LUI behavior shall be applied when the opcode is LUI.
Refers to	U_INSTRUCTION_SET_01

ID	F_ADDI_02
Description	The output of LUI shall be the value of its first input.
Rationale	The LUI instruction shall load the 20 upper bits of the instruction immediate into the destination register and fill the remaining bits with zeros. This is the default behavior for instruction immediates as stated in F_INSTR_IMMEDIATE_02 and F_INSTR_IMMEDIATE_03.
Refers to	U_INSTRUCTION_SET_01

AUIPC

ID	F_AUIPC_01
Description	The AUIPC behavior shall be applied when the opcode is AUIPC.
Refers to	U_INSTRUCTION_SET_01

ID	F_AUIPC_02
Description	The output of AUIPC shall be the sum of its first input and the address of the AUIPC instruction.
Refers to	U_INSTRUCTION_SET_01

JAL

ID	F_JAL_01
Description	The JAL behavior shall be applied when the opcode is JAL.
Refers to	U_INSTRUCTION_SET_01

ID	F_JAL_02
Description	The pc register shall be updated with the sum of the address of the JAL instruction with the first instruction input.
Refers to	U_INSTRUCTION_SET_01

ID	F_JAL_03
Description	The output of JAL shall be the address of the JAL instruction incremented by 4.
Rationale	The JAL instruction shall output the address to the following instruction for it to be used as a <i>return address</i> in the case of a function call.
Refers to	U_INSTRUCTION_SET_01

JALR

ID	F_JALR_01
Description	The JALR behavior shall be applied when the opcode is JALR and func3 is 0x0.
Refers to	U_INSTRUCTION_SET_01

ID	F_JALR_02
Description	The pc register shall be updated with the sum of the first and second inputs of the JALR instruction.
Refers to	U_INSTRUCTION_SET_01

ID	F_JALR_03
Description	The output of JALR shall be the address of the JALR instruction incremented by 4.
Rationale	The JALR instruction shall output the address to the following instruction for it to be used as a <i>return address</i> in the case of a function call.
Refers to	U_INSTRUCTION_SET_01

BEQ

ID	F_BEQ_01
Description	The BEQ behavior shall be applied when the opcode is BRANCH and func3 is 0x0.
Refers to	U_INSTRUCTION_SET_01

ID	F_BEQ_02
Description	When the first and second instruction inputs are equal, the $_{\mbox{\scriptsize pc}}$ register shall be updated with the signed sum of the address of the BEQ instruction with the third input.
Refers to	U_INSTRUCTION_SET_01

BNE

ID	F_BNE_01
Description	The BNE behavior shall be applied when the opcode is BRANCH and func3 is 0x1.
Refers to	U_INSTRUCTION_SET_01

ID	F_BNE_02
Description	When the first and second inputs are not equal, the pc register shall be updated with the signed sum of the address of the BNE instruction with the third input.
Refers to	U_INSTRUCTION_SET_01

BLT

ID	F_BLT_01
Description	The BLT behavior shall be applied when the opcode is BRANCH and func3 is 0x4.
Refers to	U_INSTRUCTION_SET_01

ID	F_BLT_02
Description	When the first input is lower than the second input using a signed comparison, the pc register shall be updated with the signed sum of the address of the BLT instruction with the third input.
Refers to	U_INSTRUCTION_SET_01

BGE

ID	F_BGE_01
Description	The BGE behavior shall be applied when the opcode is BRANCH and func3 is 0x5.
Refers to	U_INSTRUCTION_SET_01

ID	F_BGE_02
Description	When the first input is greater or equal to the second input using a signed comparison, the $_{\mbox{\scriptsize pc}}$ register shall be updated with the signed sum of the address of the BGE instruction with the third input.
Refers to	U_INSTRUCTION_SET_01

BLTU

ID	F_BLTU_01
Description	The BLTU behavior shall be applied when the opcode is BRANCH and func3 is 0x6.
Refers to	U_INSTRUCTION_SET_01

ID	F_BLTU_02
Description	When the first input is lower than the second input using an unsigned comparison, the pc register shall be updated with the signed sum of the address of the BLTU instruction with the third input.
Refers to	U_INSTRUCTION_SET_01

BGEU

ID	F_BGEU_01
Description	The BGEU behavior shall be applied when the opcode is BRANCH and func3 is 0x7.
Refers to	U_INSTRUCTION_SET_01

ID	F_BGEU_02
Description	When the first input is greater or equal to the second input using an unsigned comparison, the pc register shall be updated with the signed sum of the address of the BGEU instruction with the third input.
Refers to	U_INSTRUCTION_SET_01

LB

ID	F_LB_01
Description	The LB behavior shall be applied when the opcode is LOAD and func3 is 0x0.
Refers to	U_INSTRUCTION_SET_01

ID	F_LB_02
Description	The output of LB shall be the 8-bit value stored in memory at the address determined by the signed sum of its first and second inputs.
Refers to	U_INSTRUCTION_SET_01

ID	F_LB_03
Description	The remaining bits of the loaded value shall be filled with the value of its 7 th bit.
Refers to	U_INSTRUCTION_SET_01

LH

ID	F_LH_01
Description	The LH behavior shall be applied when the opcode is LOAD and func3 is 0x1.
Refers to	U_INSTRUCTION_SET_01

ID	F_LH_02
Description	The output of LH shall be the 16-bit value stored in memory at the address determined by the signed sum of its first and second inputs.
Refers to	U_INSTRUCTION_SET_01

ID	F_LH_03
Description	The remaining bits of the loaded value shall be filled with the value of its 15 th bit.
Refers to	U_INSTRUCTION_SET_01

LW

ID	F_LW_01
Description	The LW behavior shall be applied when the opcode is LOAD and func3 is 0x2.
Refers to	U_INSTRUCTION_SET_01

ID	F_LW_02
Description	The output of LW shall be the 32-bit value stored in memory at the address determined by the signed sum of its first and second inputs.
Refers to	U_INSTRUCTION_SET_01

LBU

ID	F_LBU_01
Description	The LBU behavior shall be applied when the opcode is LOAD and func3 is 0x4.
Refers to	U_INSTRUCTION_SET_01

ID	F_LBU_02
Description	The output of LBU shall be the 8-bit value stored in memory at the address determined by the signed sum of its first and second inputs.
Refers to	U_INSTRUCTION_SET_01

ID	F_LBU_03
Description	The remaining bits of the loaded value shall be filled with zeros.
Refers to	U_INSTRUCTION_SET_01

LHU

ID	F_LHU_01
Description	The LHU behavior shall be applied when the opcode is LOAD and func3 is 0x5.
Refers to	U_INSTRUCTION_SET_01

ID	F_LHU_02
Description	The output of LHU shall be the 16-bit value stored in memory at the address determined by the signed sum of its first and second inputs.
Refers to	U_INSTRUCTION_SET_01

ID	F_LHU_04
Description	The remaining bits of the loaded value shall be filled with zeros.
Refers to	U_INSTRUCTION_SET_01

SB

ID	F_SB_01
Description	The SB behavior shall be applied when the opcode is STORE and func3 is 0x0.
Refers to	U_INSTRUCTION_SET_01

ID	F_SB_02
Description	The lowest byte of the second input of SB shall be stored in memory at the address determined by the signed sum of its first and third inputs.
Refers to	U_INSTRUCTION_SET_01

SH

ID	F_SH_01
Description	The SH behavior shall be applied when the opcode is STORE and func3 is 0x1.
Refers to	U_INSTRUCTION_SET_01

ID	F_SH_02
Description	The two lowest bytes of the second input of SB shall be stored in memory at the address determined by the signed sum of its first and third inputs.
Refers to	U_INSTRUCTION_SET_01

SW

ID	F_SW_01
Description	The SW behavior shall be applied when the opcode is STORE and func3 is 0x2.
Refers to	U_INSTRUCTION_SET_01

ID	F_SH_02
Description	The value of the second input of SB shall be stored in memory at the address determined by the signed sum of its first and third inputs.
Refers to	U_INSTRUCTION_SET_01

ADDI

ID	F_ADDI_01
Description	The ADDI behavior shall be applied when the opcode is OP-IMM and when func3 is 0x0.
Refers to	U_INSTRUCTION_SET_01

ID	F_ADDI_02
Description	The output of ADDI shall be the signed integer sum of its two inputs.
Refers to	U_INSTRUCTION_SET_01

ID	F_ADDI_03
Description	The output of ADDI shall be truncated to 32-bits.
Refers to	U_INSTRUCTION_SET_01

SLTI

ID	F_SLTI_01
Description	The SLTI behavior shall be applied when the opcode is OP-IMM and when func3 is 0x2.
Refers to	U_INSTRUCTION_SET_01

ID	F_SLTI_02
Description	The output of SLTI shall be 1 when the signed value of its first input is lower that the signed value of its second input. It shall be 0 otherwise.
Refers to	U_INSTRUCTION_SET_01

SLTIU

ID	F_SLTIU_01
Description	The SLTIU behavior shall be applied when the opcode is OP-IMM and when func3 is 0x3.
Refers to	U_INSTRUCTION_SET_01

ID	F_SLTIU_02
Description	The output of SLTI shall be 1 when the unsigned value of its first input is lower that the unsigned value of its second input. It shall be 0 otherwise.
Refers to	U_INSTRUCTION_SET_01

XORI

ID	F_XORI_01
Description	The XORI behavior shall be applied when the opcode is OP-IMM and when func3 is 0x4.
Refers to	U_INSTRUCTION_SET_01

ID	F_XORI_02
Description	The output of XORI shall be the result of a bitwise xor between its two inputs.
Refers to	U_INSTRUCTION_SET_01

ORI

ID	F_ORI_01
Description	The ORI behavior shall be applied when the opcode is OP-IMM and when func3 is 0x6.
Refers to	U_INSTRUCTION_SET_01

ID	F_ORI_02
Description	The output of ORI shall be the result of a bitwise or between its two inputs.
Refers to	U_INSTRUCTION_SET_01

ANDI

ID	F_ANDI_01
Description	The ANDI behavior shall be applied when the opcode is OP-IMM and when func3 is 0x7.
Refers to	U_INSTRUCTION_SET_01

ID	F_ANDI_02
Description	The output of ANDI shall be the result of a bitwise and between its two inputs.
Refers to	U_INSTRUCTION_SET_01

SLLI

ID	F_SLLI_01
Description	The SLLI behavior shall be applied when the opcode is OP-IMM and func3 is 0x1.
Refers to	U_INSTRUCTION_SET_01

ID	F_SLLI_02
Description	The output of SLLI shall be its first input shifted left by the amount specified by the first 5 bits of its second input.
Refers to	U_INSTRUCTION_SET_01

ID	F_SLLI_03
Description	Zeros shall be inserted in the lower bits when shifting.
Refers to	U_INSTRUCTION_SET_01

SRLI

ID	F_SRLI_01
Description	The SRLI behavior shall be applied when the opcode is OP-IMM, func3 is 0x5 and the 30 th bit of its second input is 0.
Refers to	U_INSTRUCTION_SET_01

ID	F_SRLI_02
Description	The output of SRLI shall be its first input shifted right by the amount specified by the first 5 bits of its second input.
Refers to	U_INSTRUCTION_SET_01

ID	F_SRLI_03
Description	Zeros shall be inserted in the upper bits when shifting.
Refers to	U_INSTRUCTION_SET_01

SRAI

ID	F_SRAI_01
Description	The SRAI behavior shall be applied when the opcode is OP-IMM, func3 is $0x5$ and the 30^{th} bit of its second input is 1.
Refers to	U_INSTRUCTION_SET_01

ID	F_SRAI_02
Description	The output of SRAI shall be its first input shifted right by the amount specified by the first 5 bits of its second input.
Refers to	U_INSTRUCTION_SET_01

ID	F_SRAI_03
Description	The most significant bit of the first input shall be inserted in the upper bits when shifting.
Refers to	U_INSTRUCTION_SET_01

ADD

ID	F_ADD_01
Description	The ADD behavior shall be applied when the opcode is OP, func3 is $0x0$ and the 30^{th} bit of its second input is 0 .
Refers to	U_INSTRUCTION_SET_01

ID	F_ADD_02
Description	The output of ADD shall be the signed integer sum of its two inputs.
Refers to	U_INSTRUCTION_SET_01

ID	F_ADD_03
Description	The output of ADD shall be truncated to 32-bits.
Refers to	U_INSTRUCTION_SET_01

SUB

ID	F_SUB_01
Description	The SUB behavior shall be applied when the opcode is OP, func3 is 0x0 and the 30 th bit of its second input is 1.
Refers to	U_INSTRUCTION_SET_01

ID	F_SUB_02
Description	The output of SUB shall be the signed integer difference of its first input minus its second input.
Refers to	U_INSTRUCTION_SET_01

ID	F_SUB_03
Description	The output of SUB shall be truncated to 32-bits.
Refers to	U_INSTRUCTION_SET_01

SLL

ID	F_SLL_01
Description	The SLL behavior shall be applied when the opcode is OP and func3 is 0x1.
Refers to	U_INSTRUCTION_SET_01

ID	F_SLL_02
Description	The output of SLL shall be its first input shifted left by the amount specified by the first 5 bits of its second input.
Refers to	U_INSTRUCTION_SET_01

ID	F_SLL_03
Description	Zeros shall be inserted in the lower bits when shifting.
Refers to	U_INSTRUCTION_SET_01

SLT

ID	F_SLT_01
Description	The SLT behavior shall be applied when the opcode is OP and func3 is 0x2.
Refers to	U_INSTRUCTION_SET_01

ID	F_SLT_02
Description	The output of SLT shall be 1 when the signed value of its first input is lower that the signed value of its second input. It shall be 0 otherwise.
Refers to	U_INSTRUCTION_SET_01

SLTU

ID	F_SLTU_01
Description	The SLTU behavior shall be applied when the opcode is OP and func3 is 0x3.
Refers to	U_INSTRUCTION_SET_01

ID	F_SLTU_02
Description	The output of SLTU shall be 1 when the unsigned value of its first input is lower that the unsigned value of its second input. It shall be 0 otherwise.
Refers to	U_INSTRUCTION_SET_01

XOR

ID	F_XOR_01
Description	The XOR behavior shall be applied when the opcode is OP and func3 is 0x4.
Refers to	U_INSTRUCTION_SET_01

ID	F_XOR_02
Description	The output of XOR shall be the result of a bitwise xor between its two inputs.
Refers to	U_INSTRUCTION_SET_01

SRL

ID	F_SRL_01
Description	The SRL behavior shall be applied when the opcode is OP, func3 is $0x5$ and the 30^{th} bit of its second input is 0 .
Refers to	U_INSTRUCTION_SET_01

ID	F_SRL_02
Description	The output of SRL shall be its first input shifted right by the amount specified by the first 5 bits of its second input.
Refers to	U_INSTRUCTION_SET_01

ID	F_SRL_03
Description	Zeros shall be inserted in the upper bits when shifting.
Refers to	U_INSTRUCTION_SET_01

SRA

ID	F_SRA_01
Description	The SRA behavior shall be applied when the opcode is OP, func3 is $0x5$ and the 30^{th} bit of its second input is 1.
Refers to	U_INSTRUCTION_SET_01

ID	F_SRA_02
Description	The output of SRA shall be its first input shifted right by the amount specified by the first 5 bits of its second input.
Refers to	U_INSTRUCTION_SET_01

ID	F_SRA_03
Description	The most significant bit of the first input shall be inserted in the upper bits when shifting.
Refers to	U_INSTRUCTION_SET_01

OR

ID	F_OR_01		
Description	The OR behavior shall be applied when the opcode is OP and func3 is 0x6.		
Refers to	Refers to U_INSTRUCTION_SET_01		

ID	F_OR_02	
Description	Description The output of OR shall be the result of a bitwise or between its two inputs.	
Refers to U_INSTRUCTION_SET_01		

AND

ID	F_AND_01
Description	The AND behavior shall be applied when the opcode is OP and func3 is 0x7.
Refers to	U_INSTRUCTION_SET_01

ID	F_AND_02			
Description	The output of AND shall be the result of a bitwise and between its two inputs.			
Refers to	U_INSTRUCTION_SET_01			

FENCE TBD

ECALL TBD

EBREAK TBD

3.2.4 Exceptions

ID	F_INSTR_ADDR_MISALIGNED_01		
Description	An Instruction Address Misaligned exception shall be raised when the target address of a taken branch or an unconditional jump if not four-byte aligned.		
Refers to	U_INSTRUCTION_SET_01		

ID	F_MISALIGNED_MEMORY_ACCESS_01			
Description	A Misaligned Memory Access exception shall be raised when the target address of a load/store instruction is not aligned on the referenced type size.			
Refers to	U_INSTRUCTION_SET_01			

3.2.5 Memory interface

ID	F_ENDIANNESS_01		
Description	on Memory accesses shall use the little-endian format.		
Refers to			

Outline requirements to be compliant with the Wishbone specification.

3.2.6 Debugging

3.3 Nonfunctional Requirements

ID	N_FORMAL_PROOF_01			
Description	Each part of ECAP5-DPROC shall be formally proven when possible, otherwise thouroughly tested			
Refers to				

These can be: performance, safety, security, usability, scalability.

4 Configuration

4.1 Constants

Table 4: Constants used for implementing the RISC-V ISA

NAME TYPE WIDTH DESCRIPTION VALUE	NAME
-----------------------------------	------

4.2 Instanciation parameters

ECAP5-DPROC can be parameterized at build-time through instanciation parameters. Default constant values are defined for such parameters.

Table 5: Instanciation parameters of ECAP5-DPROC

NAME	TYPE	WIDTH	DESCRIPTION	DEFAULT VALUE	

5 Functional Partitioning

ECAP5-DPROC is built around a pipelined architecture with the following stages:

- The instruction fetch stage loads the next instruction from memory.
- The decode stage handles the instruction decoding to provide the next stage with the different instruction input values including reading from internal registers.
- The execute stage implements instruction behaviors. This includes performing integer operations as well as accessing memory.
- The write-back stage which handles storing instructions outputs to internal registers.

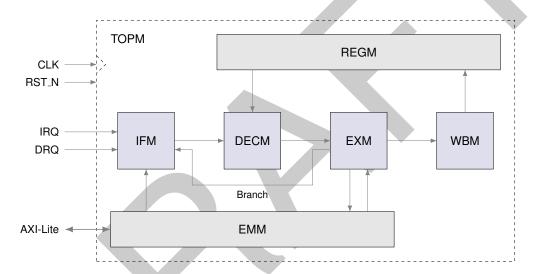


Figure 3: Schematic view of the architecture of ECAP5-DPROC

The design is split into the following functional modules:

- The Top Module (TOPM) which integrates all other modules.
- The External Memory Module (EMM), in charge of accessing memory and peripherals.
- The **Instruction Fetch Module** (IFM), in charge of implementing the instruction fetch stage.
- The **Decode Module** (DECM), in charge of implementing the decode stage.
- The **Register Module** (REGM), implementing the internal registers.
- The **Execute Module** (EXM), in charge of implementing the execute stage.
- The Write-Back Module (WBM), in charge of implementing the write-back stage.

6 Top Module

Handshaking and bubbling



7 External Memory Module



8 Instruction Fetch Module

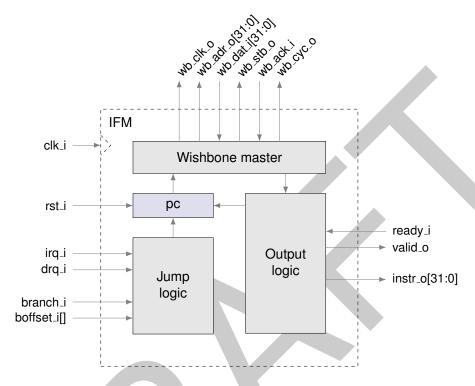


Figure 4: Schematic view of the Instruction Fetch Module

8.1 Interface

The instruction fetch module handles fetching from memory the instructions to be executing. The signals are described in table 6.

Table 6: Instruction Fetch Module interface signals

NAME	TYPE	WIDTH	DESCRIPTION		
clk_i		1	Clock input.		
rst_i	1	1	Reset input.		
JUMP LOC	SIC				
irq_i	I	1	External interrupt request.		
drq₋i	I	1	External debug request.		
branch_i	I	1	Branch request.		
boffset_i	I	TBC	Branch offset. TBC		
WISHBON	WISHBONE MASTER				

wb_clk_o	0	1	Wishbone clock output. This is hardwired to clk₋i.	
wb_adr_o	0	32	Wishbone read address.	
wb_dat_i	I	32	Wishbone read data.	
wb_ack_i	I	1	Acknowledge. Indicates a normal termination of a bus cycle.	
OUTPUT LOGIC				
ready₋i	I	1	Asserted when the output is ready to be received.	
valid_o	0	1	Asserted when the output is ready to be sent.	
instr_o	0	32	Instruction to be executed.	

8.2 Behavior

TBD : Add a description of the timing diagrams. Inputs that are missing from timing diagrams are considered low or undefined.

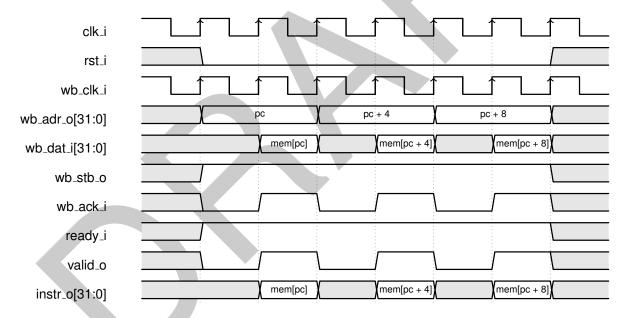


Figure 5: Timing diagram of the behavior of the instruction fetch module

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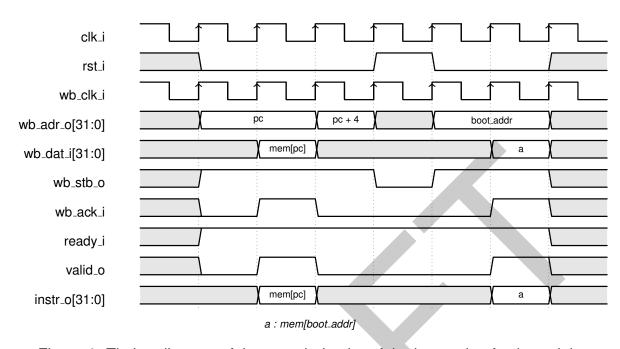


Figure 6: Timing diagram of the reset behavior of the instruction fetch module

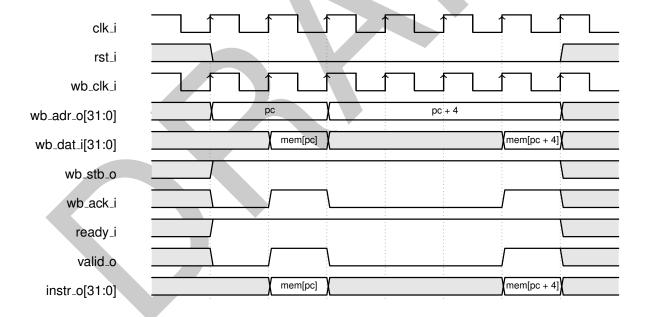


Figure 7: Timing diagram of the memory resource busy behavior of the instruction fetch module

Figure 8: interrupt

Figure 9: branch

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This module doesn't contain any prefetch mechanism as there is no performance requirement for revision 1.0.0. This will lead to a performance bottleneck due to the number of cycles needed for fetching instructions from memory.

Design 8.3

pc stores the value of the next instruction to be loaded from memory. It is connected to the wishbone master which performs the memory read. The read data is transferred to the output logic, in charge of handling the pipeline's handshaking protocol. The value of pc can be either incremented by the output logic, reset by rst_i or loaded with a specific value through the jump logic.

8.3.1 **Jump logic**

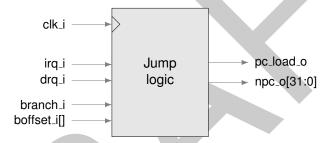


Figure 10: Schematic view of the interface of the Jump Logic

The jump logic loads a value into the pc register based on inputs. Its interface is

described in table 7.

NAME	TYPE	WIDTH	DESCRIPTION
clk_i	I	1	Clock input.
irq_i		1	External interrupt request.
drq₋i	1	1	External debug request.
branch_i	ı	1	Branch request.
boffset_i	I	TBC	Branch offset. TBC
pc_load_o	0	1	Asserted when the pc register shall be updated.
npc_o	0	32	Value used to update the pc register.

Table 7: Jump logic interface signals

Figures 11 and 12 describe the behavior of the jump logic. The values to be loaded in memory are hardcoded and set at compile time.

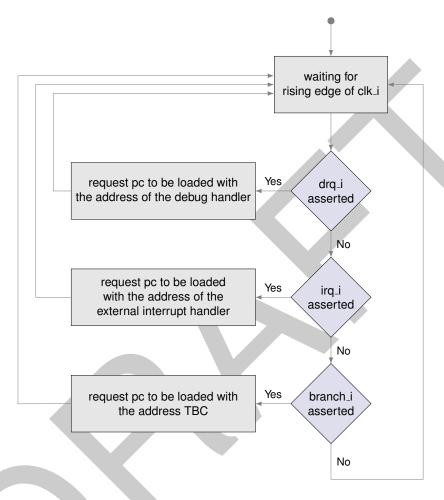


Figure 11: Activity diagram of the jump logic behavior

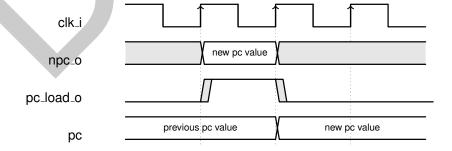


Figure 12: Timing diagram for the output port of the jump logic.

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8.3.2 PC register

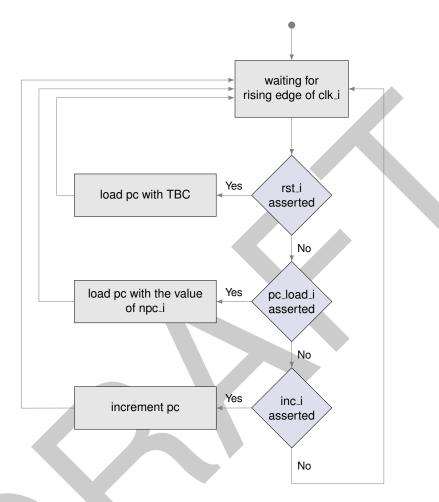


Figure 13: Activity diagram of the pc register

8.3.3 Wishbone master

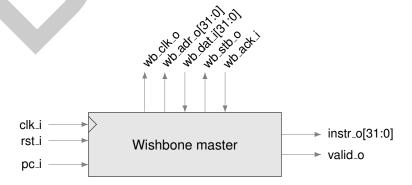


Figure 14: Schematic view of the interface of the Wishbone Master

The wishbone master fetches from memory the instruction to be executed. Its interface is described in table 8. The memory port of the wishbone master is described in figure 15.

Table 8: Instruction Fetch Module Interface signals

NAME	TYPE	WIDTH	DESCRIPTION	
clk₋i	I	1	Clock input.	
rst_i	I	1	Reset input.	
WISHBONE MASTER				
wb_clk_o	0	1	Wishbone clock output. This is hardwired to clk_i.	
wb_adr_o	0	32	Wishbone read address.	
wb_dat_i	I	32	Wishbone read data.	
wb_stb_o	0	1	Strobe output indicates a valid data transfer cycle.	
wb_ack_i	I	1	Acknowledge. Indicates a normal termination of a bus cycle.	
OUTPUT				
instr₋o	0	32	Instruction to be executed.	
valid_o	0	1	Asserted when the instruction output is valid.	

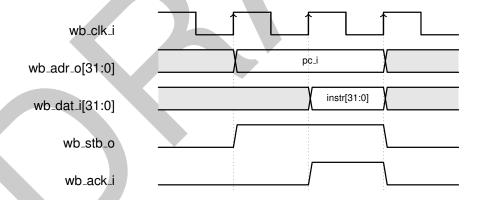


Figure 15: Timing diagram for the memory port of IFM's wishbone master

The output port of the wishbone master has the following properties:

8.3.4 Output logic

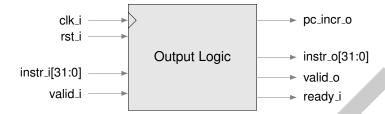


Figure 16: Schematic view of the interface of the Output Logic



9 Decode Module



10 Register Module

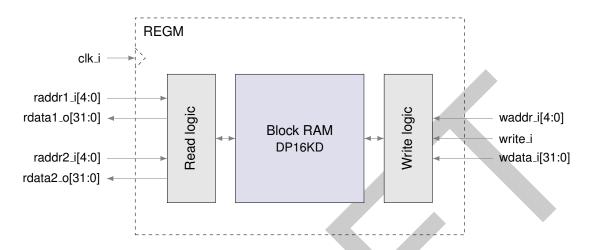


Figure 17: Schematic view of the Register Module

10.1 Interface

The register module implements the 32 internal registers of ECAP5-DPROC. It has two reading port and one writing port. The signals are described in table 9.

NAME	TYPE	WIDTH	DESCRIPTION		
clk_i	I	1	Clock input.		
FIRST READING PORT					
raddr1_i	-	5	Register selector.		
rdata1_o	0	32	Selected register value.		
SECOND READING PORT					
raddr2_i	I	5	Register selector.		
rdata2_o	0	32	Selected register value.		
WRITING PORT					
waddr_i	1	5	Register selector.		
write_i	I	1	Asserted to indicate a write.		
wdata₋i	I	32	Data to be written.		

Table 9: Register Module interface signals

10.2 Behavior

When reading, rdata1_i and rdata2_i output, on the rising edge of clk_i, the value of the register respectively selected by raddr1_i and raddr2_i. A register write happens on the rising edge of clk_i when write_i is asserted, writing the value wdata_i in the register selected by waddr_i.

10.3 Design



11 Execute Module



12 Write-Back Module



13 Debug

