Tutorial 9 – Solutions

Exercise 1*:

Consider the following four alternative definitions of TCCS agent M:

- $M_1 \stackrel{\text{def}}{=} \epsilon(3).(\epsilon(2).a.M_1 + b.M_1)$
- $M_2 \stackrel{\text{def}}{=} \epsilon(5).a.M_2 + \epsilon(3).b.M_2$
- $M_3 \stackrel{\text{def}}{=} \epsilon(3).(\epsilon(2).a.M_3 + \tau.M_3)$
- $M_4 \stackrel{\text{def}}{=} \epsilon(5).a.M_4 + \epsilon(3).\tau.M_4$

For which of the above four definitions do we have $M_i \stackrel{\epsilon(4)}{\longrightarrow}$.

- 1. For which of the four definitions do we have $M_i \stackrel{4}{\to}$. In the affirmative cases use the SOS rules for TCCS to prove the delay-transition as well as identify the target process P_i such that $M_i \stackrel{\epsilon(4)}{\longleftrightarrow} P_i$.
 - M_1 has a transition $M_1 \stackrel{4}{\rightarrow} \epsilon(1).a.M_1 + b.M_1$ which we prove below:

$$\begin{array}{c|c} \hline \epsilon(2).a.M_1 \xrightarrow{1} \epsilon(1).a.M1 & b.M_1 \xrightarrow{1} b.M1 \\ \hline \epsilon(2).a.M_1 + b.M_1 \xrightarrow{1} \epsilon(1).a.M_1 + b.M_1 \\ \hline \epsilon(3)(\epsilon(2).a.M_1 + b.M_1) \xrightarrow{4} \epsilon(1).a.M_1 + b.M_1 \\ \hline M_i \xrightarrow{4} \epsilon(1).a.M_1 + b.M_1 \\ \hline \end{array}$$

• M_2 has a transition $M_2 \stackrel{4}{\rightarrow} \epsilon(1).a.M_2 + b.M_2$ which we prove below:

$$\frac{\epsilon(5).a.M_2 \xrightarrow{4} \epsilon(1).a.M_2}{\epsilon(3).b.M_2 \xrightarrow{4} b.M_2}$$

$$\frac{\epsilon(5).a.M_2 + \epsilon(3).b.M_2 \xrightarrow{4} \epsilon(1).a.M_2 + b.M_2}{M_2 \xrightarrow{4} \epsilon(1).a.M_2 + b.M_2}$$

- M₃ cannot delay four time units, since a tau transition i enabled after just three time units.
- M₄ has the same problem as M₃.
- 2. Discuss the general relationship between process terms $\epsilon(d).(P+Q)$ and $\epsilon(d).P+\epsilon(d).Q$. The above processes are equivalent as, unlike action prefixes, delay prefixes are distributive. This is because delays to not resolve non-determinism.

Exercise 2*:

Consider the agent M and the three variants of agent N:

•
$$M \stackrel{\text{def}}{=} \epsilon(3).(\epsilon(2).\bar{a}.M + \bar{b}.M)$$

•
$$N_1 \stackrel{\text{def}}{=} \epsilon(5).b.N_1 + \epsilon(3).a.N_1$$

•
$$N_2 \stackrel{\text{def}}{=} \epsilon(3).(\epsilon(2).a.N_2 + \tau.N_2)$$

•
$$N_3 \stackrel{\text{def}}{=} \epsilon(5).\tau.N_3 + \epsilon(3).b.N_3$$

Indicate the values of i for which a) $M \mid N_i \xrightarrow{3}, b)$ $M \mid N_i \xrightarrow{5}, \text{ and } c)$ $M \mid N_i \xrightarrow{8}.$ In the affirmative cases give proper proofs using the SOS rules for TCCS.

 M | N₁ can delay for both three and five time units but cannot delay for eight time units, since after five time units both b and b are enabled for synchronization and must be engaged. The proofs for the two former are as follows:

$$\frac{\epsilon(3).(\epsilon(2).\bar{a}.M + \bar{b}.M) \xrightarrow{3} \epsilon(2).\bar{a}.M + \bar{b}.M}{M \xrightarrow{3} \epsilon(2).\bar{a}.M + \bar{b}.M} \frac{\epsilon(5).b.N_1 \xrightarrow{3} \epsilon(2).b.N_1}{\epsilon(5).b.N_1 \xrightarrow{3} \epsilon(2).b.N_1} \frac{\epsilon(3).a.N_1 \xrightarrow{3} a.N_1}{\epsilon(5).b.N + \epsilon(3).a.N \xrightarrow{3} \epsilon(2).b.N_1 + a.N_1} \frac{\epsilon(3).a.N_1 \xrightarrow{3} a.N_1}{N_1 \xrightarrow{3} \epsilon(2).b.N_1 + a.N_1}$$

$$M \mid N_1 \xrightarrow{3} \epsilon(2).\bar{a}.M + \bar{b}.M \mid \epsilon(2).b.N_1 + a.N_1$$

$$M \mid N_1 \xrightarrow{3} \epsilon(2).\bar{a}.M + \bar{b}.M \mid \epsilon(2).b.N_1 + a.N_1$$

• $M \mid N_2$ can only delay for no more than three time units as N_2 itself has an enabled tau transition after three time units.

$$\begin{array}{c|c} \hline \epsilon(3).(\epsilon(2).\bar{a}.M + \bar{b}.M) \xrightarrow{3} \epsilon(2).\bar{a}.M + \bar{b}.M \\ \hline M \xrightarrow{3} \epsilon(2).\bar{a}.M + \bar{b}.M \\ \hline M \mid N_2 \xrightarrow{3} \epsilon(2).\bar{a}.M + \bar{b}.M \\ \hline \end{pmatrix} \begin{array}{c|c} \epsilon(3).(\epsilon(2).b.N + \tau.N) \xrightarrow{3} \epsilon(2).b.N_2 + \tau.N_2 \\ \hline \\ N_2 \xrightarrow{3} \epsilon(2).b.N_2 + \tau.N_2 \\ \hline \\ M \mid N_2 \xrightarrow{3} \epsilon(2).\bar{a}.M + \bar{b}.M \mid \epsilon(2).b.N_2 + \tau.N_2 \\ \hline \end{array}$$

 M | N₃ can delay up to three time units after which M and N₃ are enabled to synchronized on the b action and must do so.

$$\frac{\epsilon(3).(\epsilon(2).\bar{a}.M + \bar{b}.M) \xrightarrow{3} \epsilon(2).\bar{a}.M + \bar{b}.M}{M \xrightarrow{3} \epsilon(2).\bar{a}.M + \bar{b}.M} = \frac{\epsilon(5).\tau.N_3 \xrightarrow{3} \epsilon(2).\tau.N_3 \xrightarrow{\epsilon(3).b.N_3} \epsilon(3).b.N_3 \xrightarrow{3} b.N_3}{\epsilon(5).\tau.N + \epsilon(3).b.N \xrightarrow{3} \epsilon(2).\tau.N_3 + b.N_3} = \frac{\epsilon(5).\tau.N_3 \xrightarrow{3} \epsilon(2).\tau.N_3 \xrightarrow{\delta(2).\tau.N_3} b.N_3}{N_3 \xrightarrow{3} \epsilon(2).\tau.N_3 + b.N_3}$$

$$M \mid N_3 \xrightarrow{3} \epsilon(2).\bar{a}.M + \bar{b}.M \mid \epsilon(2).\tau.N_3 + b.N_3$$

Exercise 4:

• TCCS process:

$$T \stackrel{\text{def}}{=} set_5.T_5 + set_{10}.T_{10} + set_{30}.T_{30}$$

$$T_5 \stackrel{\text{def}}{=} \epsilon(5).\overline{to}.T + T$$

$$T_{10} \stackrel{\text{def}}{=} \epsilon(10).\overline{to}.T + T$$

$$T_{30} \stackrel{\text{def}}{=} \epsilon(30).\overline{to}.T + T$$

• Timed Automata: A solution model can be seen in Figure 1.

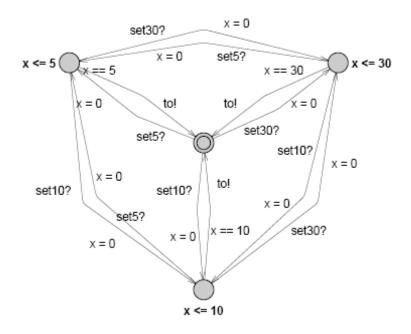


Figure 1

Exercise 5:

- Timed automaton

