# EECS 151/251A Homework 4

Due Friday, Oct 4<sup>th</sup>, 2019

### Reading

In addition to reviewing the RISC-V ISA and datapath lectures, skim through the RISC-V ISA spec. In particular, focus on the Introduction, Chapter 2 (RV32I Base Integer Instruction Set), and the table on page 130.

## Problem 1: RISC-V Manual Assembly

Manually construct the binary instruction for the following assembly instructions. Provide the solution as a 32-bit binary number.

- (a) add x1, x2, x3
- (b) addi x1, x2, 100
- (c)  $1b \times 1, 4(\times 2)$
- (d) beq x6, x8, 1024

# Problem 2: RISC-V Assembly Programs

Write down the values of the specified registers after the following programs have run:

```
(a)
             li x1, 100
             li x2, 200
             add x3, x1, x2
             sub x3, x3, x1
    x1 = _{_{_{_{_{_{_{_{_{_{_{1}}}}}}}}}}, x2 = _{_{_{_{_{_{_{_{_{_{_{_{1}}}}}}}}}}}, x3 = _{_{_{_{_{_{_{_{_{_{_{1}}}}}}}}}}
(b)
             li x1, 100
             li x0, 200
             add x0, x1, x0
    x0 = ___, x1 = ___
(c)
             li x1, Oxdead
             li x2, Oxbeef
             li x3, 0x1024
             sh x1, 0(x3)
             sh x2, 2(x3)
             1w x4, 0(x3)
```

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(e) Assume the following instructions start at address 0x0.

```
li x1, 0
    jalr x2, x0, 16
    addi x1, x1, 100
    addi x1, x1, 200
    jal x0, end
    jalr x3, x2, 0
    end: nop

x1 = ____, x2 = ___, x3 = ____
```

#### Problem 3: RISC-V Psuedo-instructions

Several psuedo-instructions are defined by the RISC-V assembler which translate to sequences of one or more RISC-V base instructions. For each psuedo-instruction, write the RISC-V assembly instructions that implement it. Refer to page 130 of the RISC-V spec for a list of the base RISC-V instructions.

- (a) nop. Do nothing, don't change the architectural state.
- (b) mv rd, rs. Move the value in register rs to register rd.
- (c) li rd, imm. Load a 32-bit immediate into register rd.
- (d) begz rs, imm. Branch if rs is greater than or equal to zero.
- (e) j imm. Jump to PC += imm and don't link any register.
- (f) bgt rs1, rs2, imm. Branch if rs1 is greater than rs2.

#### Problem 4: RISC-V Instruction Decoder

Consider the complete RV32I datapath drawn in the lecture slides. Write down logical expressions for the following control signals in terms of an instruction's opcode, funct3, and funct7 bits.

You can simplify the expressions by comparing the opcode field against the constants in Table 25.1 in the RISC-V ISA manual. For example: assign sig = (opcode == OP-32) || (opcode == LOAD);

- (a) WBSel
- (b) MemRW, assume 0 = read and 1 = write
- (c) PCSel
- (d) BSel

endmodule

## Problem 5: RISC-V Memory Decoder

The RV32I ISA defines 5 memory load instructions: 1b, 1bu, 1h, 1hu, 1w. Complete the implementation of a Verilog module which converts the raw output of the data memory to the value to be written to the regfile according to the type of load instruction and the memory address.

```
module load_decoder(
  input [31:0] addr, // the byte-address for the load instruction
  input [31:0] raw_data, // the raw data from the DMEM
  input lb, lbu, lh, lhu, lw, // type of load instruction, only 1 is high at a time
  output [31:0] wb_data // writeback data (to the regfile)
);

// Your implementation
```

# Problem 6: RV64I ALU

Refer to Chapter 5 (RV64I Base Integer Instruction Set) in the RISC-V spec. Implement an ALU that supports the add, sll, sub 64-bit integer instructions as well as their 'W' suffix variants.

```
`define ALU_ADDW 1
`define ALU_SUB 2
`define ALU_SUBW 3
`define ALU_SLL 4
`define ALU_SLLW 5
`define ALU_SRA 6
`define ALU_SRAW 7
module rv64_alu(
  input [63:0] a,
  input [63:0] b,
  input [2:0] op, // op can be any of values `define'd above output [63:0] c,
);

// Your implementation
```

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 ${\tt endmodule}$ 

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