

EECS 151/251A

Discussion 4

Alisha Menon

9/20/21, 9/21/21, 9/23/21

Administrativa

- Homework 3 is out – due Monday **9/27, 12:00am**
- Homework 4 out this week.

Agenda

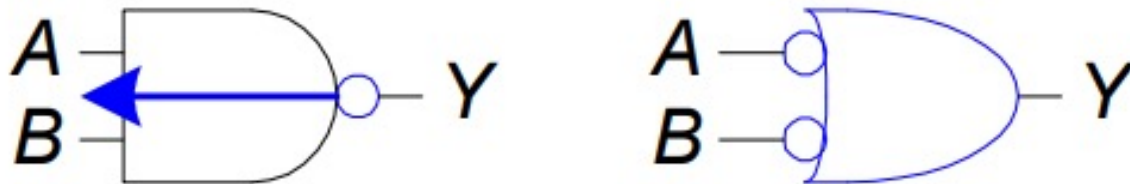
- DeMorgan's Law
 - Bubble pushing
- Karnaugh maps
 - POS
 - SOP
- Finite state machines

DeMorgan's Law: Bubble Pushing

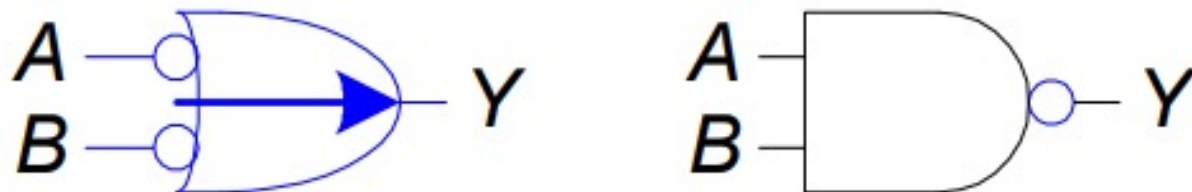
- $(x+y)' = x'y'$
- $(xy)' = x'+y'$
- Bubble = inversion (NOT)

DeMorgan's Law: Bubble Pushing

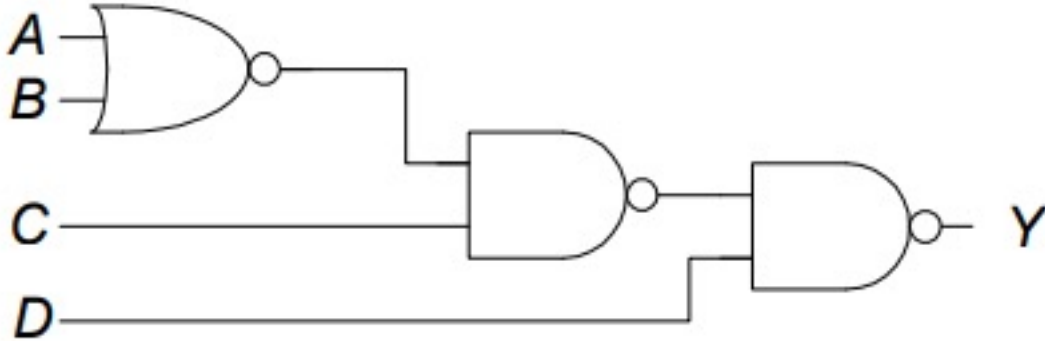
- $(x+y)' = x'y'$
- $(xy)' = x'+y'$
- Bubble = inversion (NOT)
- For a single gate:
 - Swap AND for OR & vice versa
 - Backward pushing: add bubbles to input



- Forward pushing: add bubbles to output



Bubble Pushing Example



SoP & PoS

SoP

PoS

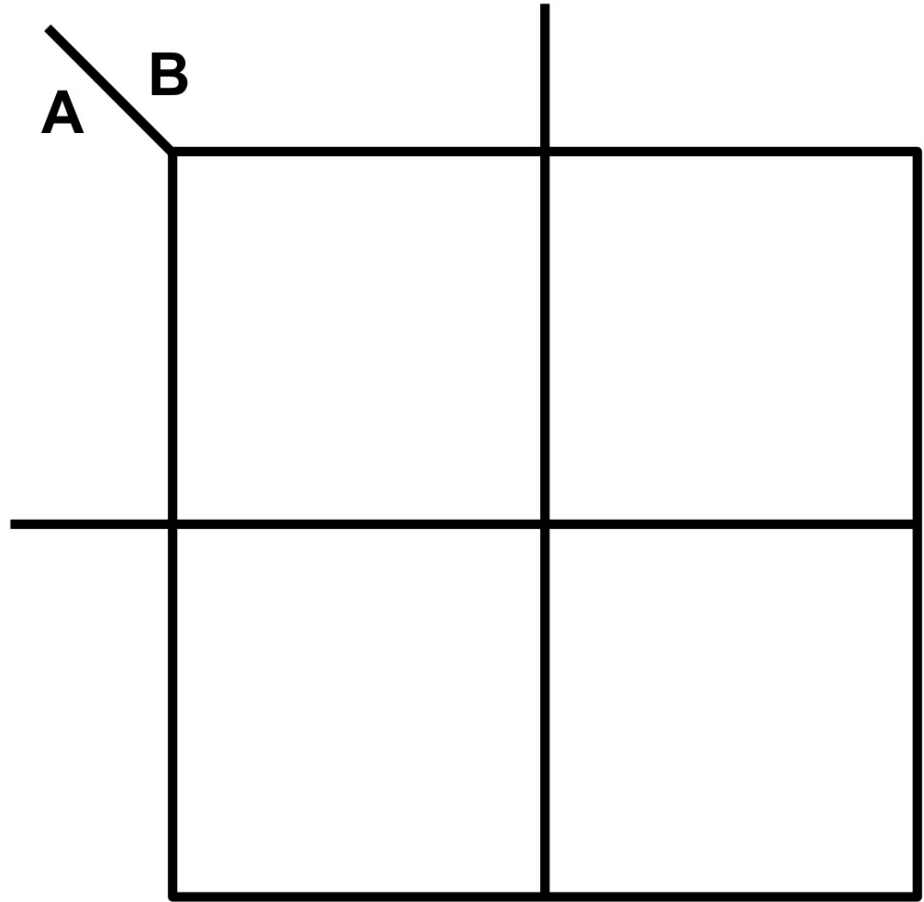
A	B	C	Out
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

K-maps

- K-Maps: visual & systematic Boolean simplification
- 2 important Boolean identities:
 - $(1+A)=1$
 - $(A+\bar{A})=1$
- Leverages **gray coding** to organize neighboring minterms
 - Adjacent minterms only differ by a single bit!
- Key to solving: form groups of 1's by multiples of 2
 - As large & as few as possible
 - Overlapping is OK, wrap boundary where possible
 - Write AND expression for each group
 - Make new SoP expression

K-map example

$$F(A,B) = \bar{A}B + \bar{A}$$



Simplification – Karnaugh maps (SOP)

a	b	c	f
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

	00	01	11	10
0				
1				

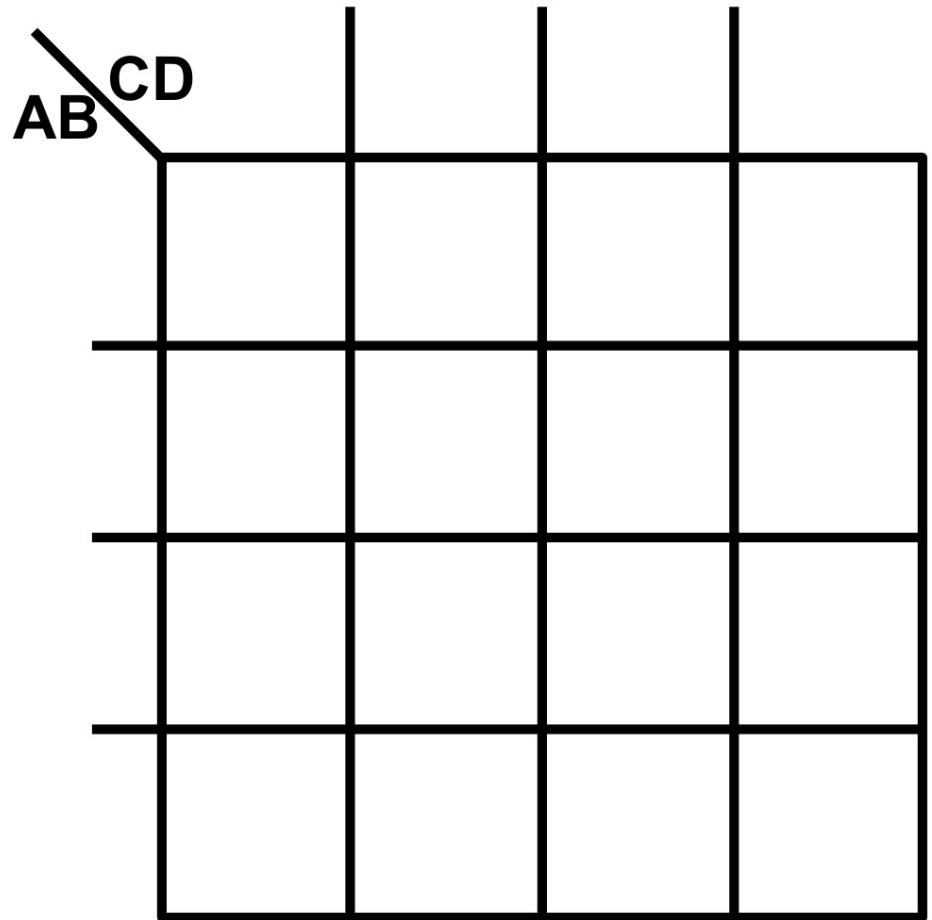
Simplification – Karnaugh maps (POS)

a	b	c	f
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

	00	01	11	10
0				
1				

4-input K-map example

$$\begin{aligned} F(A,B) = & \bar{A}\bar{B}\bar{C}\bar{D} + \bar{A}\bar{B}C\bar{D} + \\ & \bar{A}B\bar{C}\bar{D} + \bar{A}BC\bar{D} + \\ & AB\bar{C}\bar{D} + ABC\bar{D} + \\ & A\bar{B}CD + A\bar{B}C\bar{D} + \\ & ABCD \end{aligned}$$



Finite state machines

FSM review

- Sequential circuit where output depends on **present and past inputs**
- Has finite number of states, and can only be in one state at a time
- Combinational logic used to calculate next state and output
- Represented by state transition diagram



Moore vs. Mealy FSM

	Moore	Mealy
Output function		
# states		
Output synchronous		
Output delay		

Example - Vending Machine FSM

- Dispenses a soda if it receives at least 25 cents
 - Doesn't return change or rollover to next purchase
- Customer can insert three different coins:
 - Quarter – 25¢ – Q
 - Dime – 10¢ – D
 - Nickel – 5¢ – N

Moore Vending Machine

```
module vending_machine(  
    input clk, rst,  
    input Q, D, N,  
    output dispense  
);
```

Mealy Vending Machine

```
module vending_machine(  
    input clk, rst,  
    input Q, D, N,  
    output dispense  
);
```

Verilog Implementation

- Two main sections:
 - State transition (sequential)
 - State/output logic (combinational)

```
module vending_machine()  
  
    // inputs, outputs, clk, rst  
  
    // define state bits  
  
    // define state names as local params  
  
    // state transitions  
  
    // next state and output logic  
  
endmodule
```

Setup and state transitions

```
module vending_machine(  
    input clk, rst,  
    input Q, D, N,  
    output dispense  
);  
  
reg [2:0] NS, CS;  
  
localparam S0 = 3'd0,  
            S5 = 3'd1,  
            S10 = 3'd2,  
            S15 = 3'd3,  
            S20 = 3'd4,  
            S25 = 3'd5;  
  
always @(posedge clk) begin  
    if (rst) CS <= S0;  
    else CS <= NS;  
end  
  
...
```

Moore vs. Mealy combinational logic

```
always @(*) begin
    NS = CS;
    case (CS)
        S0: begin
            if (Q == 1'b1) NS = S25;
            if (D == 1'b1) NS = S10;
            if (N == 1'b1) NS = S5;
        end
        S5: begin
            if (Q == 1'b1) NS = S25;
            if (D == 1'b1) NS = S15;
            if (N == 1'b1) NS = S10;
        end
        ...
        S25: begin
            if (Q == 1'b1) NS = S25;
            if (D == 1'b1) NS = S10;
            if (N == 1'b1) NS = S5;
        end
        default: NS = S0;
    endcase
end

assign dispense = (CS == S25);

endmodule
```

```
reg dispense;
always @(*) begin
    NS = CS;
    dispense = 1'b0;
    case (CS)
        S0: begin
            if (Q == 1'b1) begin
                NS = S0;
                dispense = 1'b1;
            end
            if (D == 1'b1) NS = S10;
            if (N == 1'b1) NS = S5;
        end
        ...
        S15: begin
            if (Q == 1'b1) begin
                NS = S0;
                dispense = 1'b1;
            end
            if (D == 1'b1) begin
                NS = S0;
                dispense = 1'b1;
            end
            if (N == 1'b1) NS = S10;
        end
        ...
        default: NS = S0;
    endcase
end

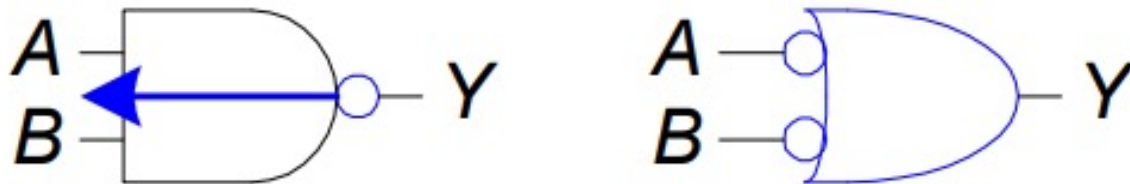
endmodule
```

DeMorgan's Law: Bubble Pushing

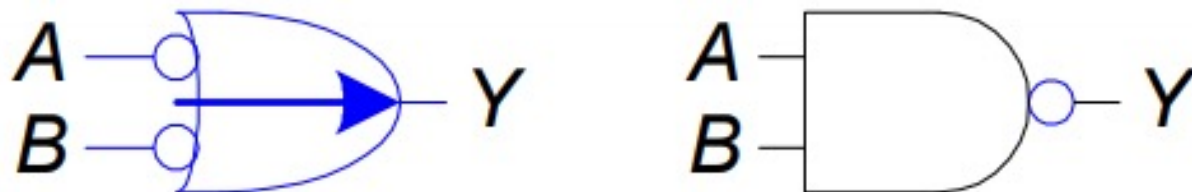
- $(x+y)' = x'y'$
- $(xy)' = x'+y'$
- Bubble = inversion (NOT)

DeMorgan's Law: Bubble Pushing

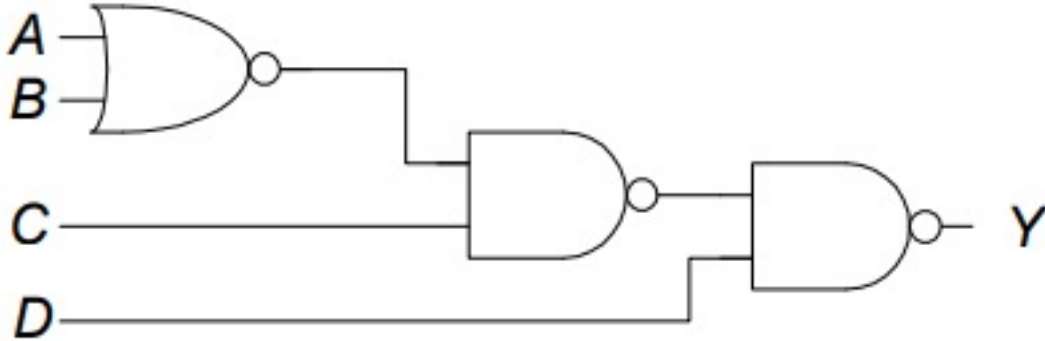
- $(x+y)' = x'y'$
- $(xy)' = x'+y'$
- Bubble = inversion (NOT)
- For a single gate:
 - Swap AND for OR & vice versa
 - Backward pushing: add bubbles to input



- Forward pushing: add bubbles to output



Bubble Pushing Example



SoP & PoS

SoP

PoS

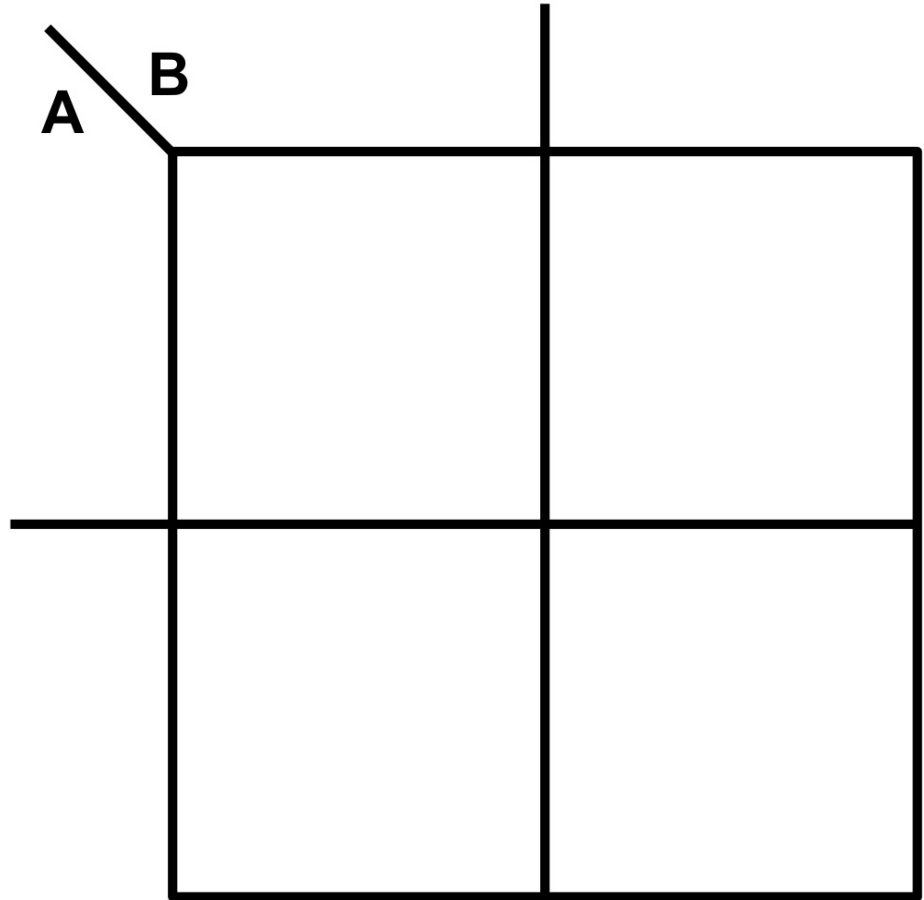
A	B	C	Out
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

K-maps

- K-Maps: visual & systematic Boolean simplification
- 2 important Boolean identities:
 - $(1+A)=1$
 - $(A+\bar{A})=1$
- Leverages **gray coding** to organize neighboring minterms
 - Adjacent minterms only differ by a single bit!
- Key to solving: form groups of 1's by multiples of 2
 - As large & as few as possible
 - Overlapping is OK, wrap boundary where possible
 - Write AND expression for each group
 - Make new SoP expression

K-map example

$$F(A,B) = \bar{A}B + \bar{A}$$



Simplification – Karnaugh maps (SOP)

a	b	c	f
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

	00	01	11	10
0				
1				

Simplification – Karnaugh maps (POS)

a	b	c	f
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	0

	00	01	11	10
0				
1				

4-input K-map example

$$\begin{aligned} F(A,B) = & \bar{A}\bar{B}\bar{C}\bar{D} + \bar{A}\bar{B}C\bar{D} + \\ & \bar{A}B\bar{C}\bar{D} + \bar{A}BC\bar{D} + \\ & AB\bar{C}\bar{D} + ABC\bar{D} + \\ & A\bar{B}CD + A\bar{B}C\bar{D} + \\ & ABCD \end{aligned}$$

