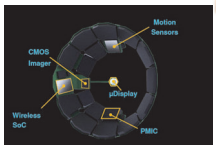
EECS151 : Introduction to Digital Design and ICs

Lecture 5 – Verilog III

Bora Nikolić

Mojo Lens - AR Contact Lenses for Real People
Michael Wiemer and Renaldi Winoto, Mojo Vision



HotChips 33

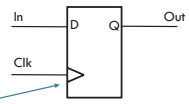
Review

- Verilog is the most-commonly used HDL
- We have seen combinatorial constructs
  - Assign statement
  - Always blocks
- Practice is the best way to learn a new language…

Sequential Logic, Take 2

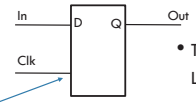
Latches and Flip-Flops

- Flip-flop is edge-triggered, latch is level-sensitive
- D Flip-flop
  - Rising edge



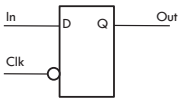
Signifies 'edge triggered'

- D Latch
  - Transparent HIGH



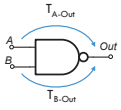
Level sensitive if there is no 'edge'

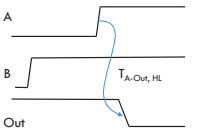
  - Transparent LOW



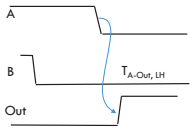
Timing

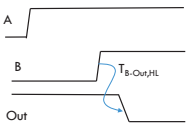
- Combinational logic timing





A is arriving late
(is in the critical path)



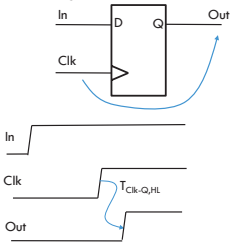


B is arriving late
(is in the critical path)

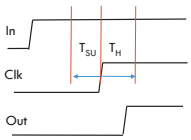
HL and LH transition differ
$t_{A\text{-}Out}$ and $t_{B\text{-}Out}$ differ
In CMOS, propagation delay depends on:
- Gate type, size (output resistance)
- Capacitive loading
- Input slope

Timing

- Flip-flop timing
  (latch timing will be covered later)

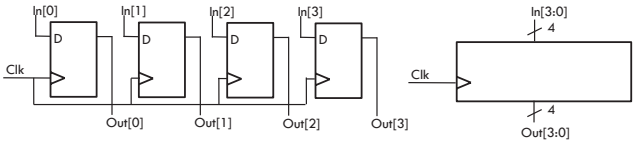


- Setup and hold times
  - Data cannot change in the interval of setup time **before** the clock edge to hold time **after** the clock edge

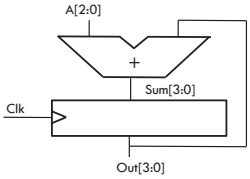


Register

- 4-bit register



- Accumulator



Administrivia

- Homework 2 is due this Friday
  - Homework 3 wil be posted this week
- Lab 3 this week

Sequential Logic in Verilog

State Elements in Verilog

Always blocks are the only way to specify the "behavior" of state elements. Synthesis tools will turn state element behaviors into state element instances.

D-flip-flop with synchronous set and reset example:

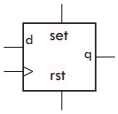
```
module dff(q, d, clk, set, rst);
  input d, clk, set, rst;
  output q;
  reg q;

  always @(posedge clk)
    if (rst)
      q <= 1'b0;
    else if (set)
      q <= 1'b1;
    else
      q <= d;
endmodule
```

keyword

"always @ (posedge clk)" is key to flip-flop inference.

This gives priority to reset over set and set over d.

On FPGAs, maps to native flip-flop.

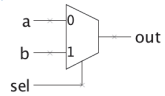


Unlike logic gates, there are no primitive flip-flops in Verilog. Although, it is possible to instantiate FPGA or standard-cell specific flip-flops.

The Sequential always Block

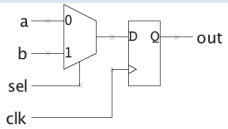
Combinational

```
module comb(input a, b, sel,
            output reg out);
  always @(*) begin
    if (sel) out = b;
    else out = a;
  end
endmodule
```



Sequential

```
module seq(input a, b, sel, clk,
           output reg out);
  always @(posedge clk) begin
    if (sel) out <= b;
    else out <= a;
  end
endmodule
```



Latches vs. Flip-Flops

Flip-Flop

```
module flipflop
(
  input clk,
  input d,
  output reg q
);

  always @(posedge clk)
  begin
    q <= d;
  end
endmodule
```

Latch

```
module latch
(
  input clk,
  input d,
  output reg q
);

  always @(clk or d)
  begin
    if ( clk )
      q <= d;
  end
endmodule
```

Importance of the Sensitivity List

- The use of posedge and negedge makes an always block sequential (edge-triggered)

D-Register with synchronous clear

```
module dff_sync_clear(
  input d, clearb, clock,
  output reg q);

  always @(posedge clock)
    begin
      if (!clearb) q <= 1'b0;
      else q <= d;
    end
endmodule
```

always block entered only at each positive clock edge

D-Register with asynchronous clear

```
module dff_async_clear(
  input d, clearb, clock,
  output reg q);

  always @(negedge clearb or posedge clock)
    begin
      if (!clearb) q <= 1'b0;
      else q <= d;
    end
endmodule
```

always block entered immediately when (active-low) clearb is asserted

Note: The following is incorrect syntax: *always @(clear or negedge clock)*
If one signal in the sensitivity list uses *posedge/negedge*, then all signals must.

- *Assign any signal or variable from only one always block.*

Be wary of race conditions: always blocks with same trigger execute concurrently…

Blocking vs. Nonblocking Assignments

- Verilog supports two types of assignments within always blocks, with subtly different behaviors.
  - Blocking assignment (=): evaluation and assignment are immediate

```
always @(*) begin
  x = a | b;      // 1. evaluate a|b, assign result to x
  y = a ^ b ^ c;  // 2. evaluate a^b^c, assign result to y
  z = b & ~c;     // 3. evaluate b&(~c), assign result to z
end
```

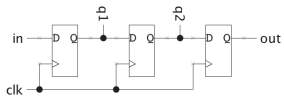
  - Nonblocking assignment (<=): all assignments deferred to end of simulation time step after all right-hand sides have been evaluated (even those in other active always blocks)

```
always @(*) begin
  x <= a | b;     // 1. evaluate a|b, but defer assignment to x
  y <= a ^ b ^ c; // 2. evaluate a^b^c, but defer assignment to y
  z <= b & ~c;    // 3. evaluate b&(~c), but defer assignment to z
  // 4. end of time step: assign new values to x, y and z
end
```

Sometimes, as above, both produce the same result. Sometimes, not!

Assignment Styles for Sequential Logic

What we want:
Register-based digital delay line
(a.k.a. shift-register)



Will non-blocking and blocking assignments both produce the desired result?

```
module nonblocking(
  input in, clk,
  output reg out
);
  reg q1, q2;
  always @(posedge clk) begin
    q1 <= in;
    q2 <= q1;
    out <= q2;
  end

endmodule
```

```
module blocking(
  input in, clk,
  output reg out
);
  reg q1, q2;
  always @(posedge clk) begin
    q1 = in;
    q2 = q1;
    out = q2;
  end

endmodule
```

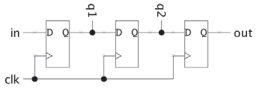

Use Nonblocking for Sequential Logic

```
always @(posedge clk) begin
    q1 <= in;
    q2 <= q1;   // uses old q1
    out <= q2;  // uses old q2
  end
```

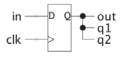
```
always @(posedge clk) begin
    q1 = in;
    q2 = q1;   // uses new q1
    out = q2;  // uses new q2
  end
```

("old" means value before clock edge, "new" means the value after most recent assignment)

"At each rising clock edge, q1, q2, and out simultaneously receive the old values of in, q1, and q2."

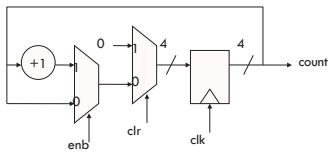


"At each rising clock edge, q1 = in. After that, q2 = q1. After that, out = q2. Therefore out = in."



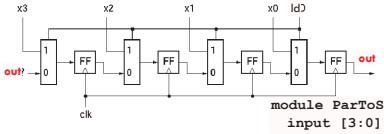
- Blocking assignments **do not** reflect the intrinsic behavior of multi-stage sequential logic
- Guideline: use **nonblocking** assignments for sequential `always` blocks

Example: A Simple Counter




```
// 4-bit counter with enable and synchronous clear
module counter(input clk,enb,clr,
               output reg [3:0] count);
  always @(posedge clk) begin
    count <= clr ? 4'b0 : (enb ? count+1 : count);
  end
endmodule
```


Example - Parallel to Serial Converter




```
module ParToSer(ld, X, out, clk);
  input [3:0] X;
  input ld, clk;
  output out;

  reg [3:0] Q;
  wire [3:0] NS;

  assign NS =
    (ld) ? X : {Q[0], Q[3:1]};

  always @ (posedge clk)
     Q <= NS;

  assign out = Q[0];
endmodule
```

Specifies the muxing with "rotation"

forces Q register (flip-flops) to be rewritten every cycle

connect output

Simplified Verilog Guidelines

- Combinational logic:
  - Continuous Assignment:

    assign a = b & c;
  - Always block with @(*)

    always @(*) **begin**

        a = b & c; // blocking statement

    end
- Sequential logic:
  - Always block with @(posedge clk)

    always @(posedge clk) **begin**

        a <= b & c; // nonblocking statement

    end

assign statement → **wire**

always statement → **reg**

Verilog in EECS 151/251A

- We use behavioral modeling at the bottom of the hierarchy
- Use instantiation to 1) build hierarchy and,
  2) map to FPGA and ASIC resources not supported by synthesis.
- Favor continuous assign and avoid always blocks unless:
  - No other alternative: ex: state elements, case
  - Helps readability and clarity of code: ex: large nested if else
- Use named ports.
- Verilog is a big language. This is only an introduction.
  - Harris & Harris book chapter 4 is a good source.
  - ***Be careful of what you read on the web.*** Many bad examples out there.
  - We will be introducing more useful constructs throughout the semester. Stay tuned!

Verilog vs. SystemVerilog

- **`always`** statements in Verilog can be used to infer flip-flops, latches or logic
  - Depends on the sensitivity list and the statement
  - Easy to create confusion
- System Verilog adds disambiguation:
  - **`always_ff`** for flip-flops
  - **`always_latch`** for latches
  - **`always_comb`** for combinational logic

Verilog Testbenches

Simulating the Circuit

- Once you have a circuit in Verilog (device under test, or DUT), you would like to test it
- Instantiate the DUT and supply its inputs via a testbench
  - Simple
  - Comprehensive
  - Random
- **`initial`** statement supplies the stimuli

Testbench basics

- Example clock

```
reg clk;

initial clk = 0;
always #(`CLOCK_PERIOD/2) clk <= ~clk;
```

- Example inputs

```
initial begin
in <= 4'h0;
  @(negedge clk) in<= 4'h1;
    ...
```

(sets up inputs on the negedge, so they are ready at the posedge)

Only small DUTs can be tested exhaustively

SystemVerilog

- SystemVerilog adds many more verification features
  - We will touch on assertions and covers

  (relates to CS70)

Final Thoughts on Verilog Examples

Verilog looks like C, but it describes hardware:

Entirely different semantics: multiple physical elements with parallel activities and temporal relationships.

A large part of digital design is knowing how to write Verilog that gets you the desired circuit. First understand the circuit you want then figure out how to code it in Verilog. If you try to write Verilog without a clear idea of the desired circuit, you will struggle.

As you get more practice, you will know how to best write Verilog for a desired result.

Be suspicious of the synthesis tools! Check the output of the tools to make sure you get what you want.

Clicker Question

- How many stimuli to exhaustively test a 32-b adder?
  - A) 32
  - B) 64
  - C) 65,536
  - D) 4,294,967,296
  - E) 18,446,744,073,709,551,616

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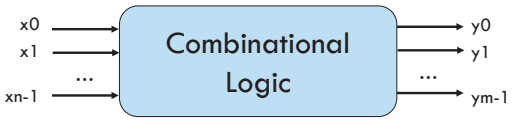
Combinational Logic

Combinational Logic

- The outputs depend *only* on the current values of the inputs.
  - Memoryless: compute the output values using the current inputs.



Combinational Logic Example

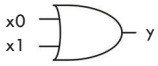
Boolean Equations:

```
y = x0 OR x1
  = x0 + x1
```

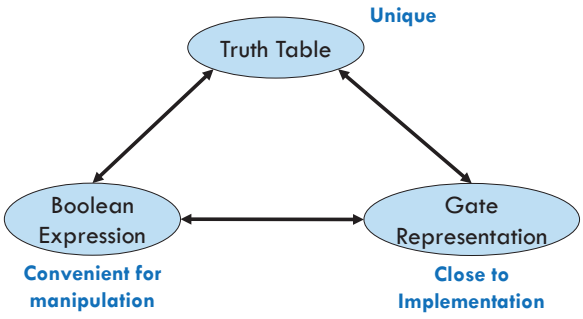
Truth Table Description:

x0	x1	y
0	0	0
0	1	1
1	0	1
1	1	1

Gate Representations:



Relationship Among Representations



Boolean Algebra

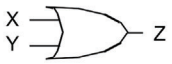
Boolean Algebra Background

- Logic: The study of the principles of reasoning.
- The 19th Century Mathematician, George Boole, developed a math. system (algebra) involving logic, Boolean Algebra.
  - His variables took on TRUE, FALSE.
- Later Claude Shannon (father of information theory) showed (in his Master's thesis!) how to map Boolean Algebra to digital circuits.

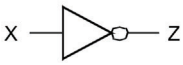
Boolean Algebra Fundamentals

- Two elements {0, 1}
- Two binary operators: AND (·) OR (+)
- One unary operator: NOT (¯ , ´)

X	Y	Z
0	0	0
0	1	0
1	0	0
1	1	1



X	Y	Z
0	0	0
0	1	1
1	0	1
1	1	1



X	Z
0	1
1	0

Boolean Operations

- Given two variables (x, y), 16 logic functions

X	Y	F_0	F_1	F_2	F_3	F_4	F_5	F_6	F_7	F_8	F_9	F_A	F_B	F_C	F_D	F_E	F_F
0	0	0	0	0	0	0	0	0	0	1	1	1	1	1	1	1	1
0	1	0	0	0	0	1	1	1	1	0	0	0	0	1	1	1	1
1	0	0	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1
1	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1

Laws of Boolean Algebra

- Identities, null elements:
  - X+0=X, X•1=X
  - X+1=1, X•0=0
- Idempotency:
  - X+X=X, X•X=X
- Complements:
  - X+X´=1, X•X´=0
- Commutativity:
  - X+Y=Y+X, X•Y=Y•X
- Associativity:
  - (X + Y) + Z= X + (Y + Z)= X + Y + Z
  - (X • Y) • Z = X • (Y • Z) = X • Y • Z
- Distributivity:
  - X • (Y+Z) = (X•Y) + (X•Z)
  - X + (Y•Z) = (X+Y) • (X+Z)
- Duality:
  - AND → OR and vice versa
  - 0 → 1 and vice versa
  - Leave literals unchanged

$$\{F(x_1, x_2, ..., x_n, 0, 1, +, \bullet)\}^D = \{F(x_1, x_2, ..., x_n, 1, 0, \bullet, +)\}$$

Literals are variables or their complements

Proving Distributive Law

- X • (Y+Z) = (X•Y) + (X•Z)

X	Y	Z	(Y+Z)	X • (Y+Z)	(X•Y)	(X•Z)	(X•Y) + (X•Z)
0	0	0					
0	0	1					
0	1	0					
0	1	1					
1	0	0					
1	0	1					
1	1	0					
1	1	1					

Proving Distributive Law

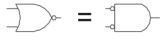
- X • (Y+Z) = (X•Y) + (X•Z)

X	Y	Z	(Y+Z)	X • (Y+Z)	(X•Y)	(X•Z)	(X•Y) + (X•Z)
0	0	0	0	0	0	0	0
0	0	1	1	0	0	0	0
0	1	0	1	0	0	0	0
0	1	1	1	0	0	0	0
1	0	0	0	0	0	0	0
1	0	1	1	1	0	1	1
1	1	0	1	1	1	0	1
1	1	1	1	1	1	1	1

DeMorgan's Law

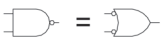
- Theorem for complementing a complex function.

(x + y)' = x' y'



x	y	x'	y'	(x + y)'	x' y'
0	0				
0	1				
1	0				
1	1				

(x y)' = x' + y'

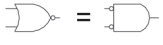


x	y	x'	y'	(x y)'	x' + y'
0	0				
0	1				
1	0				
1	1				

DeMorgan's Law

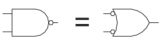
- Procedure for complementing a complex function.

(x + y)' = x' y'



x	y	x'	y'	(x + y)'	x' y'
0	0	1	1	1	1
0	1	1	0	0	0
1	0	0	1	0	0
1	1	0	0	0	0

(x y)' = x' + y'



x	y	x'	y'	(x y)'	x' + y'
0	0	1	1	1	1
0	1	1	0	1	1
1	0	0	1	1	1
1	1	0	0	0	0

Summary

- Sequential logic uses flip-flops and (sometimes) latches
- Flip-flops and latches are inferred in Verilog
  - Always blocks
- Practice is the best way to learn a new language…
- Blocking and non-blocking assignments
- Combinational logic block outputs depend only on its inputs
- Boolean algebra can be used for manipulation and simplification of Boolean equations