# EECS151 : Introduction to Digital Design and ICs

## Lecture 14 – Gate Delays

### Bora Nikolić

**Moore's Law Could Ride EUV for 10 More Years**

September 30, 2021, EETimes - ASML plans to introduce new extreme ultraviolet (EUV) lithography equipment that will extend the longevity of Moore's Law for at least ten years, according to executives at the world's only supplier of the tools, which are crucial for the world's most advanced silicon.

Starting in the first half of 2023, the company plans to offer customers equipment that takes EUV numerical aperture (NA) higher to 0.55 NA from the existing 0.33 NA. The company believes that the new equipment will help chip makers reach process nodes well beyond the current threshold (2nm) for at least another 10 years, according to ASML vice president Teun van Gogh, in an interview with EE Times.

EETimes

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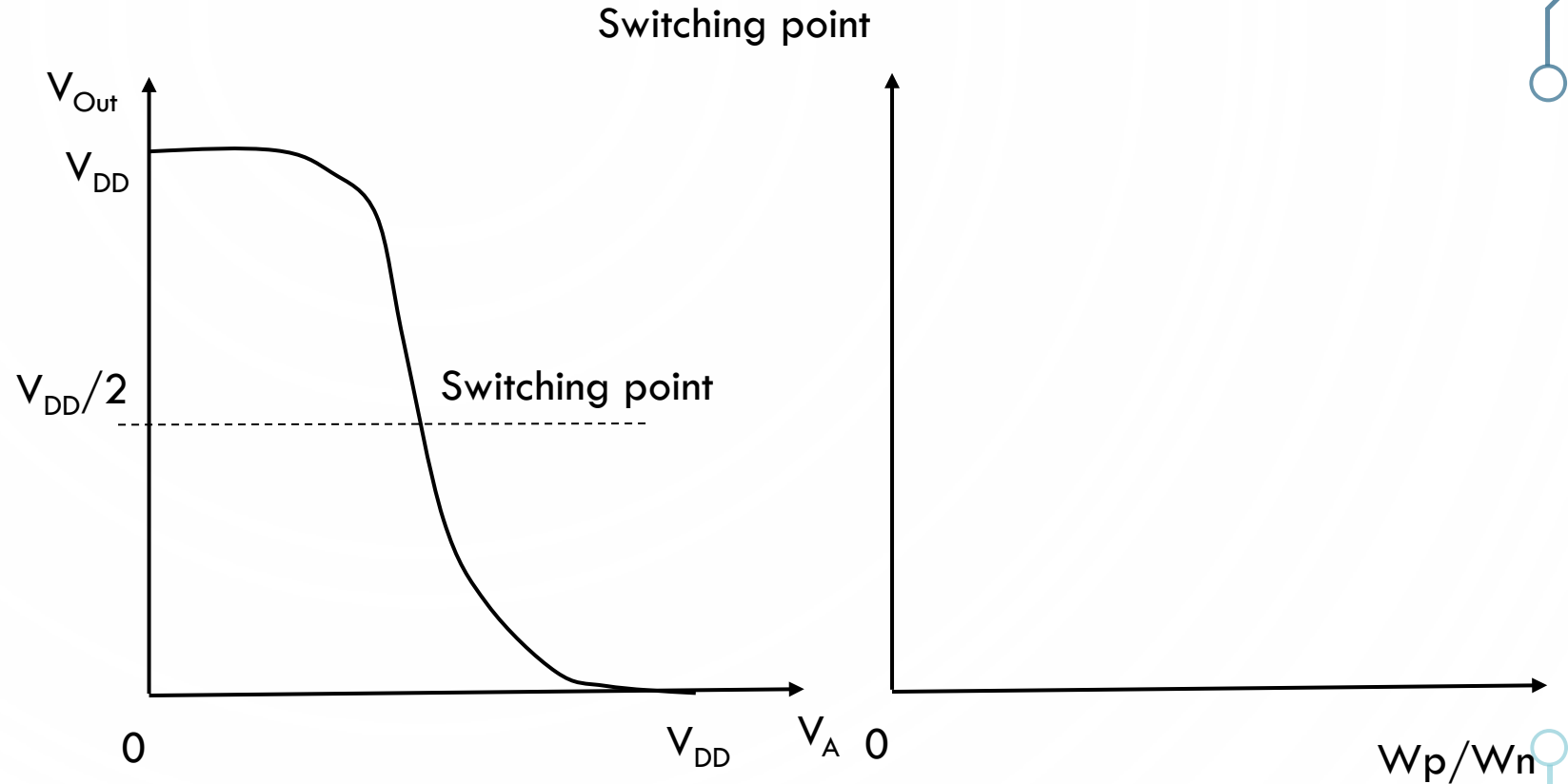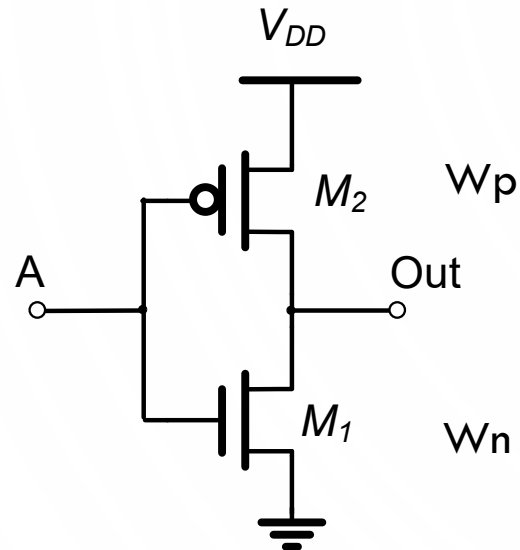UNIVERSITY OF CALIFORNIA

# Review

- CMOS allows for convenient switch level abstraction

- CMOS pull-up and pull-down networks are complementary

  - Graph models for CMOS gates

- Transistor sizing affects gate performance

# CMOS Sizing
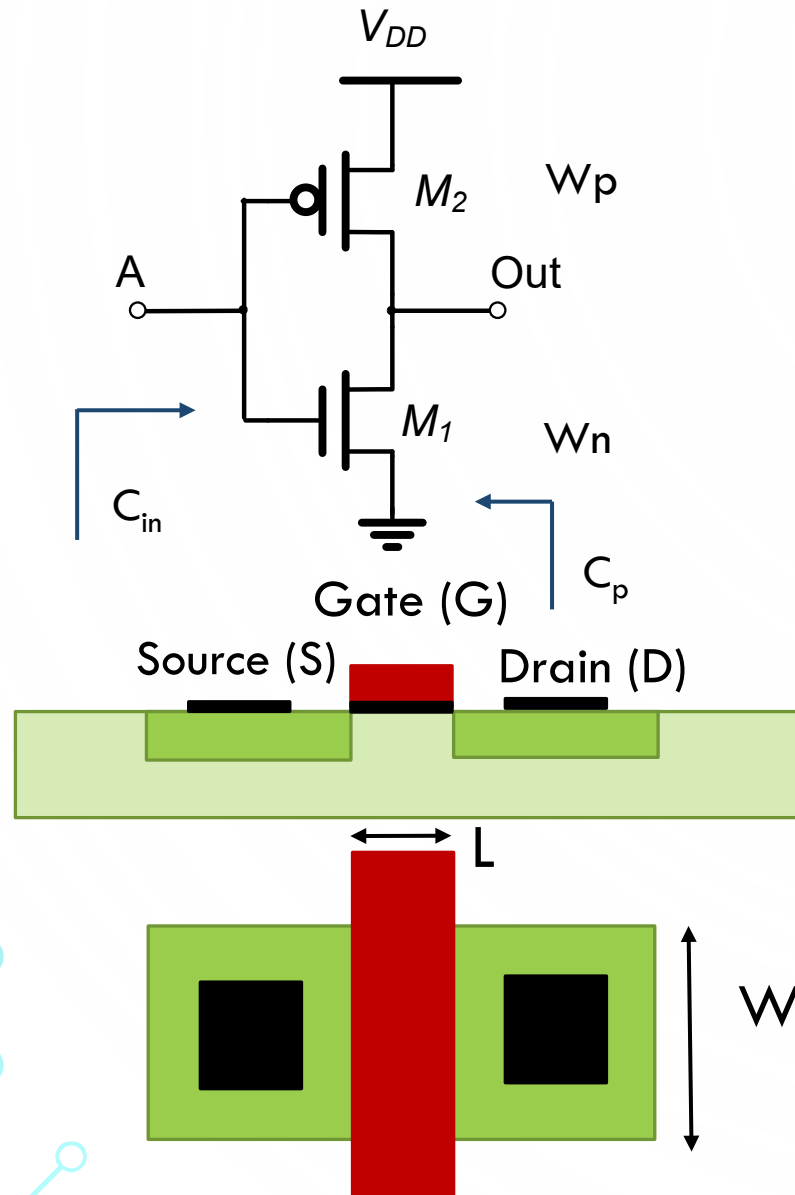
# Transistor Sizing

- Impact of Wp/Wn on VTC



- In the past, Wp > Wn (see Rabaey, 2nd ed)

- In modern processes (finFET), Wp = Wn

- Weak dependence on Wp/Wn

# CMOS Delay

# Capacitances

$V_{DD}$

$M_2$    Wp

A    Out

$M_1$    Wn

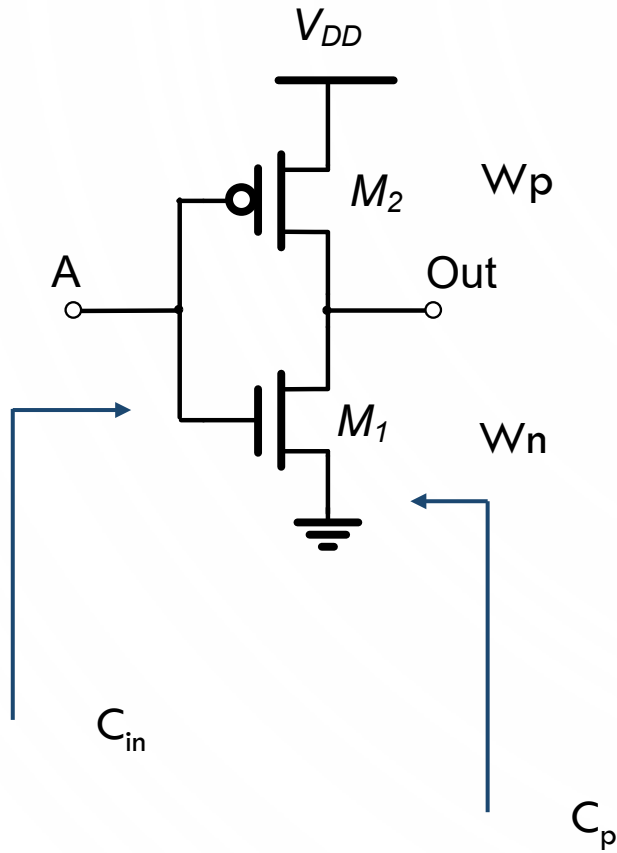$C_{in}$

$C_p$

Gate (G)

Source (S)    Drain (D)

L

W

- $C_{in}$ is largely set by the gate cap
  - ~WL
  - $2 \times W = 2 \times C_{in}$
  - It is non-linear, but we will ignore that

- $C_p$ is largely set by the drain cap
  - ~W (drain area/perimeter)
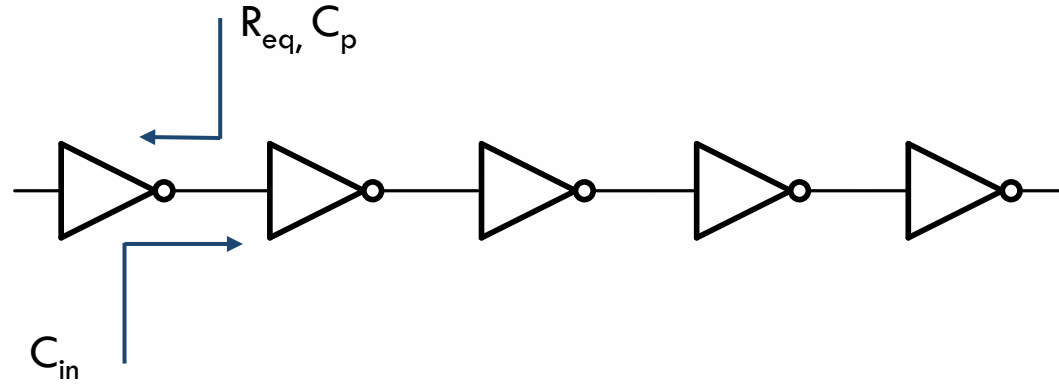  - $2 \times W = 2 \times C_p$

$$C_p = \gamma C_{in}$$

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# Gate Sizing

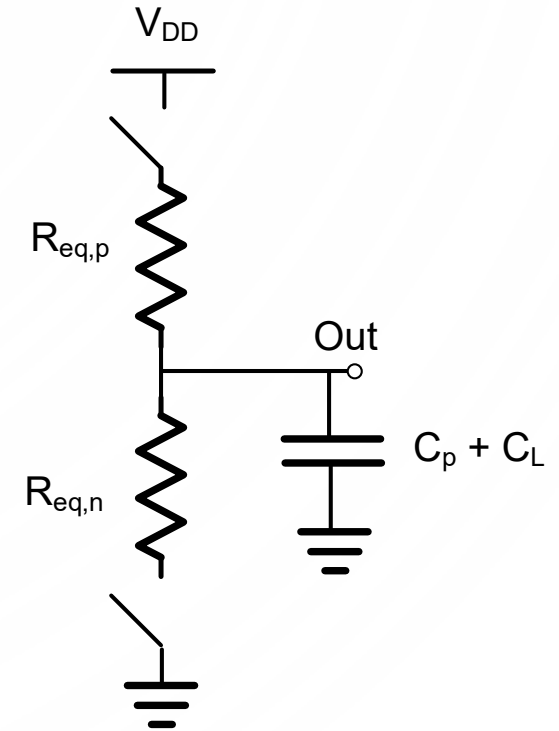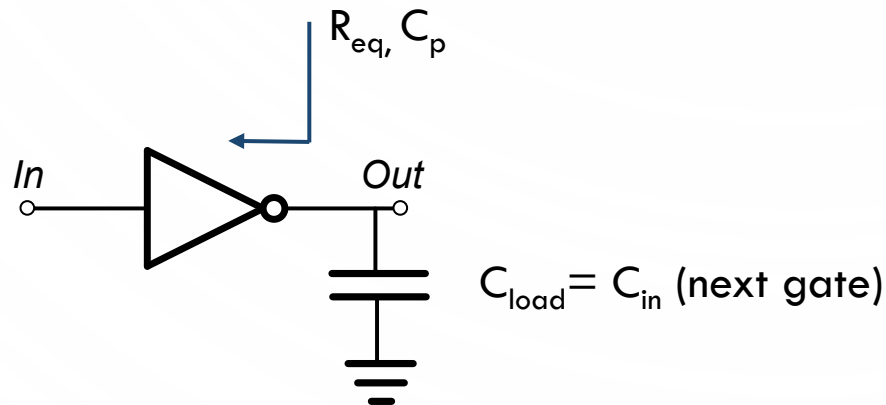- Doubling the gate size (by doubling Ws):



- Doubles $C_{in}$

- Halves equivalent gate resistance
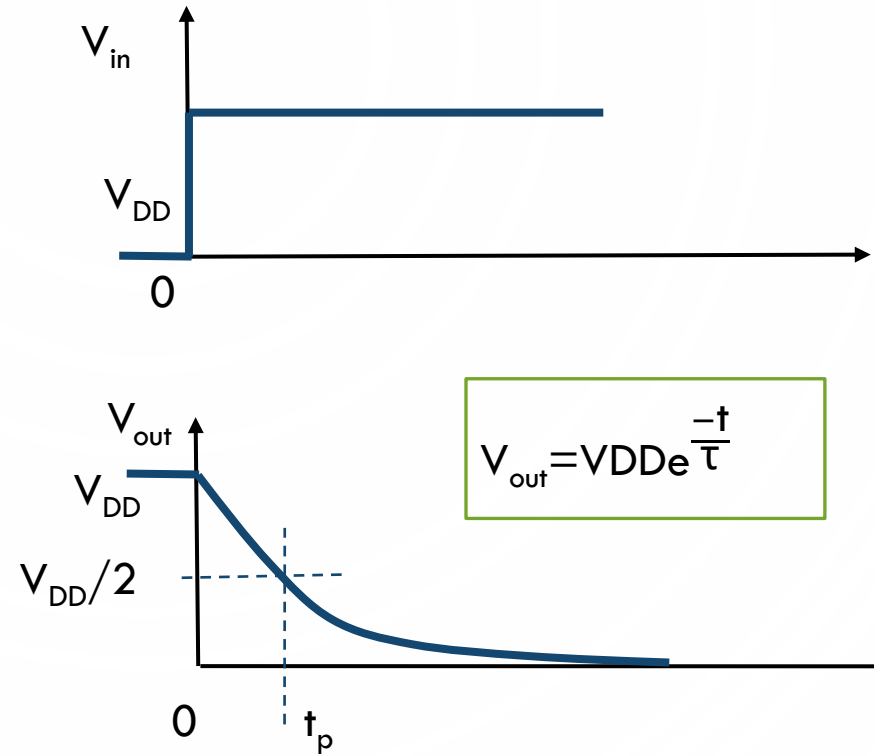
- Doubles $C_p$

# Inverter Delay

- How to time this?



$R_{eq}, C_p$

$C_{in}$

- Each gate has an $R_{eq}$ and drives $C_{in}$ of the next gate



$R_{eq}, C_p$
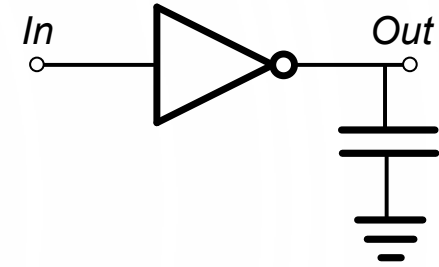
*In*

*Out*

$C_{load} = C_{in}$ (next gate)

$V_{DD}$

$R_{eq,p}$

Out

$C_p + C_L$

$R_{eq,n}$

# Inverter Delay

- High-to-low

$V_{DD}$

$R_{eq,p}$

In

Out

$R_{eq,n}$

$C_p + C_L$

Out

$R_{eq,n}$

$C_p + C_L$

In

$V_{in}$

$V_{DD}$

0

$V_{out}$

$V_{DD}$

$V_{DD}/2$

0    $t_p$

$$V_{out} = VDDe^{\frac{-t}{\tau}}$$
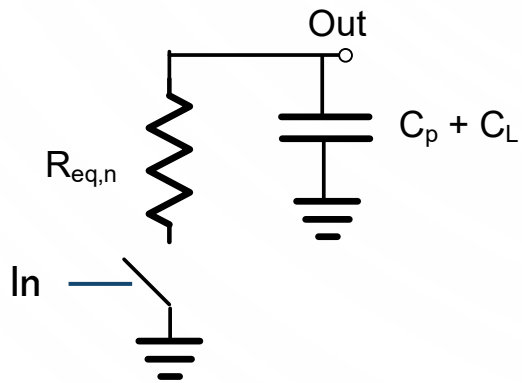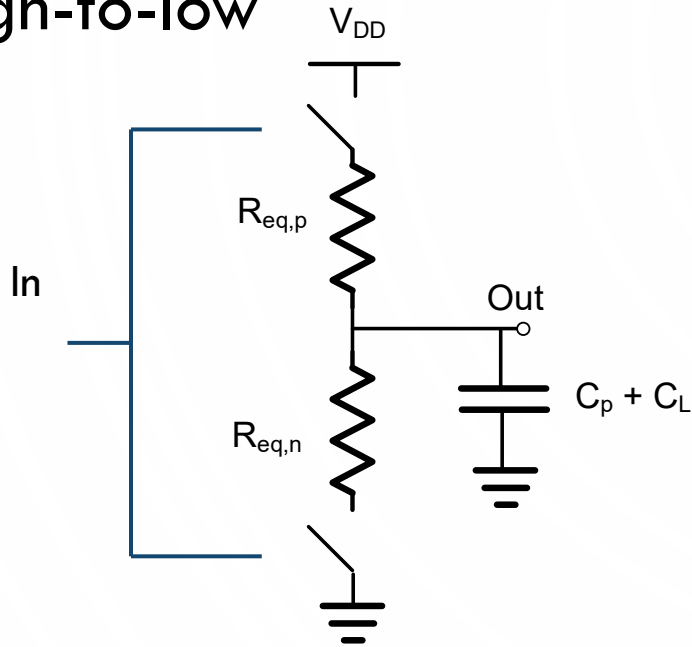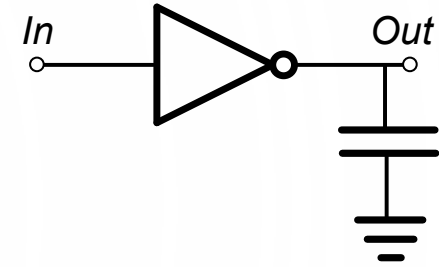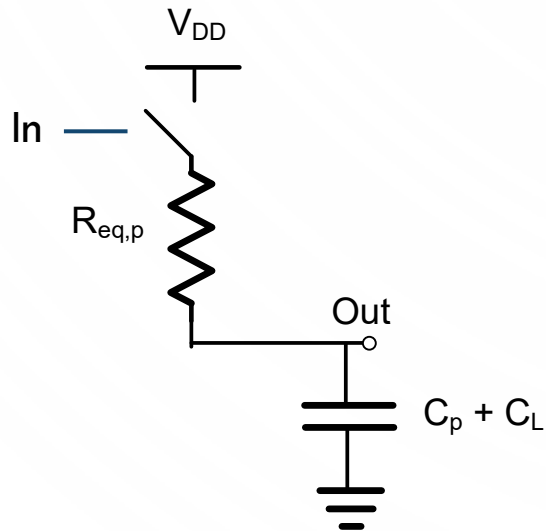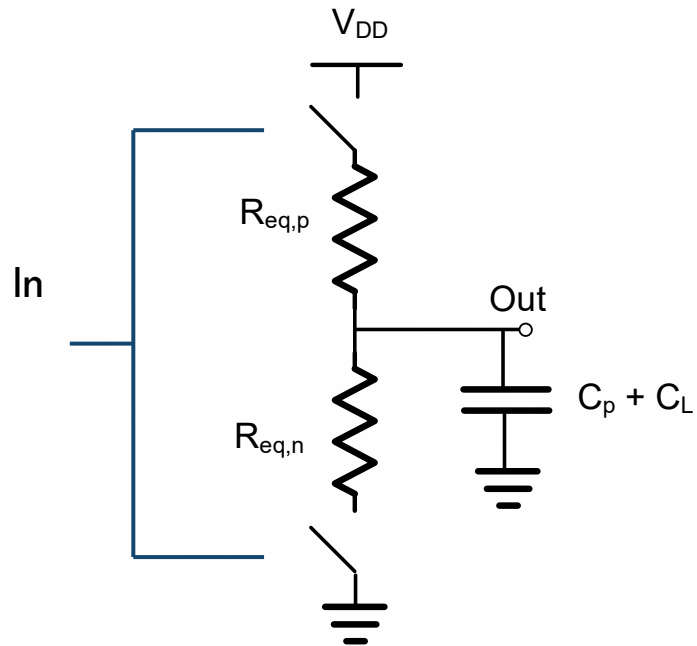
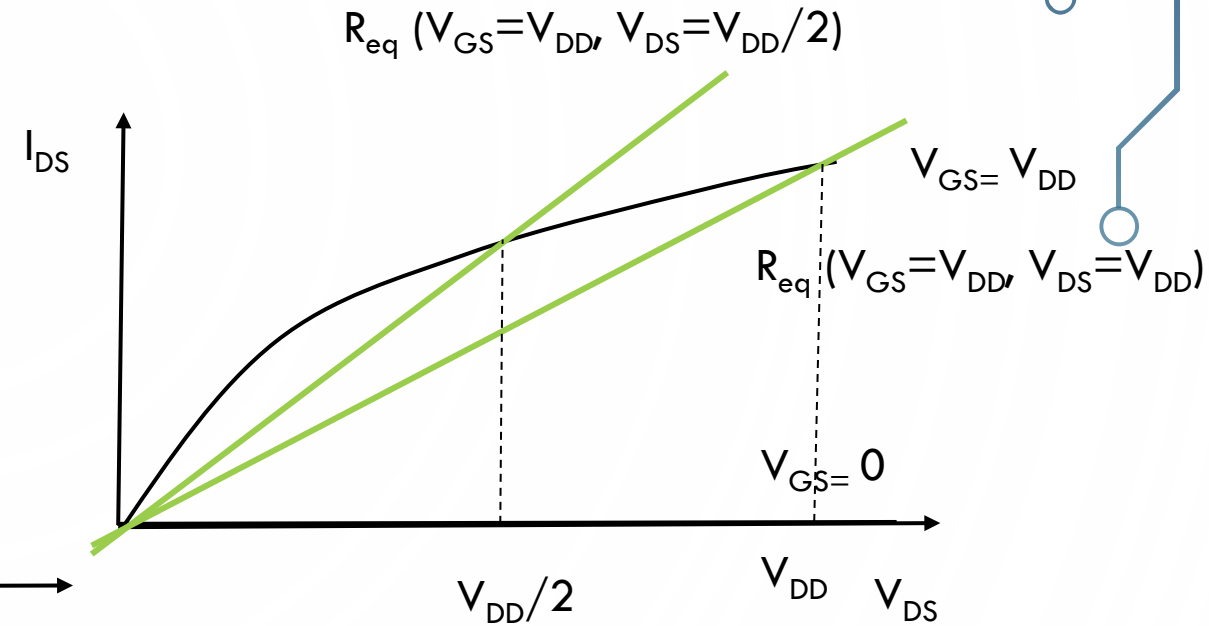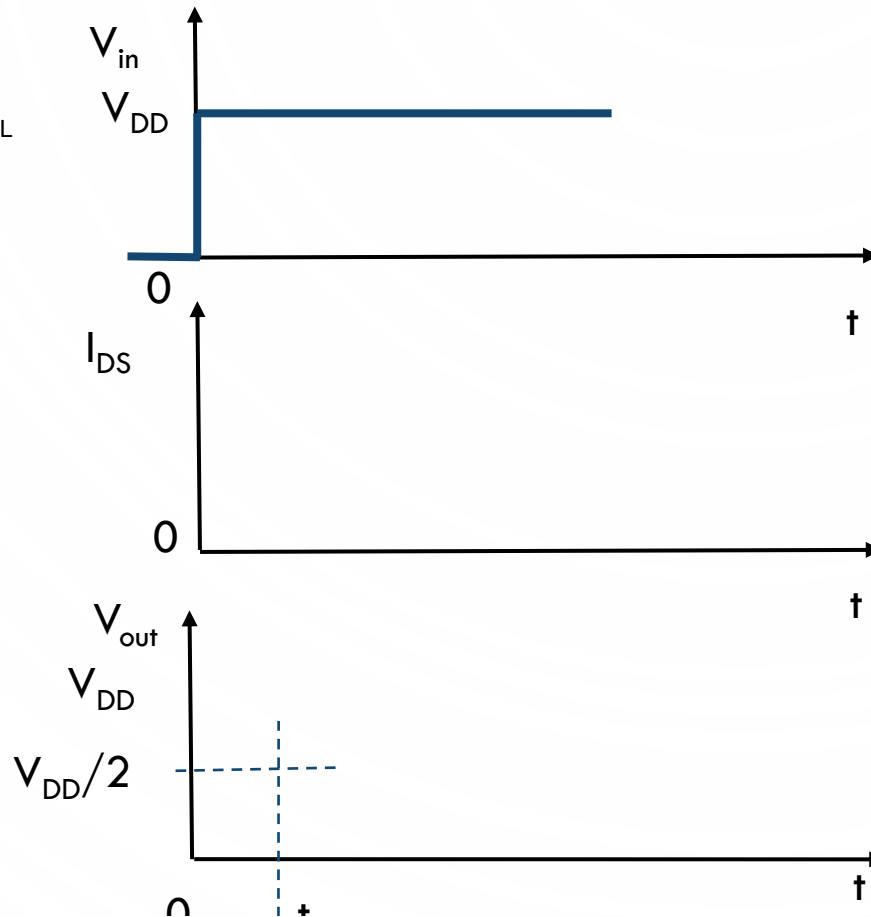$$t_{p,HL} = (\ln 2)\tau = 0.7\ R_{eq,n}(C_p + C_L)$$

$$\tau = R_{eq,n}(C_p + C_L)$$

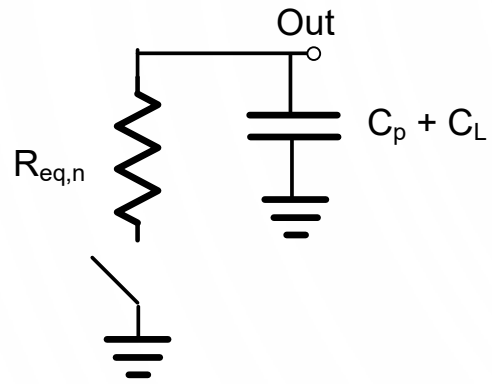In                Out

# Inverter Delay

$V_{DD}$

$R_{eq,p}$

In

Out

$C_p + C_L$

$R_{eq,n}$

$V_{in}$

$V_{DD}$

0

$V_{DD}$    $V_{out}$

$$V_{out}=VDD(1-e^{\frac{-t}{\tau}})$$

$V_{DD}/2$

0    $t_p$

$t_{p,LH}= (ln2)\tau = 0.7\ Req_{,p}(C_p+C_L)$

$V_{DD}$

In

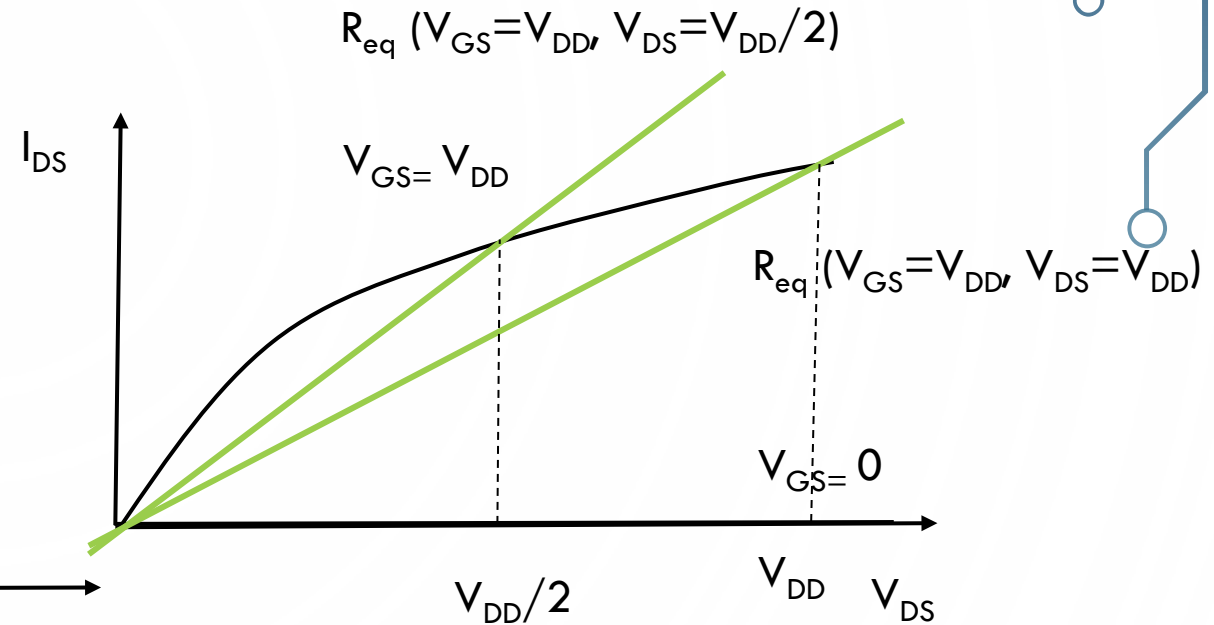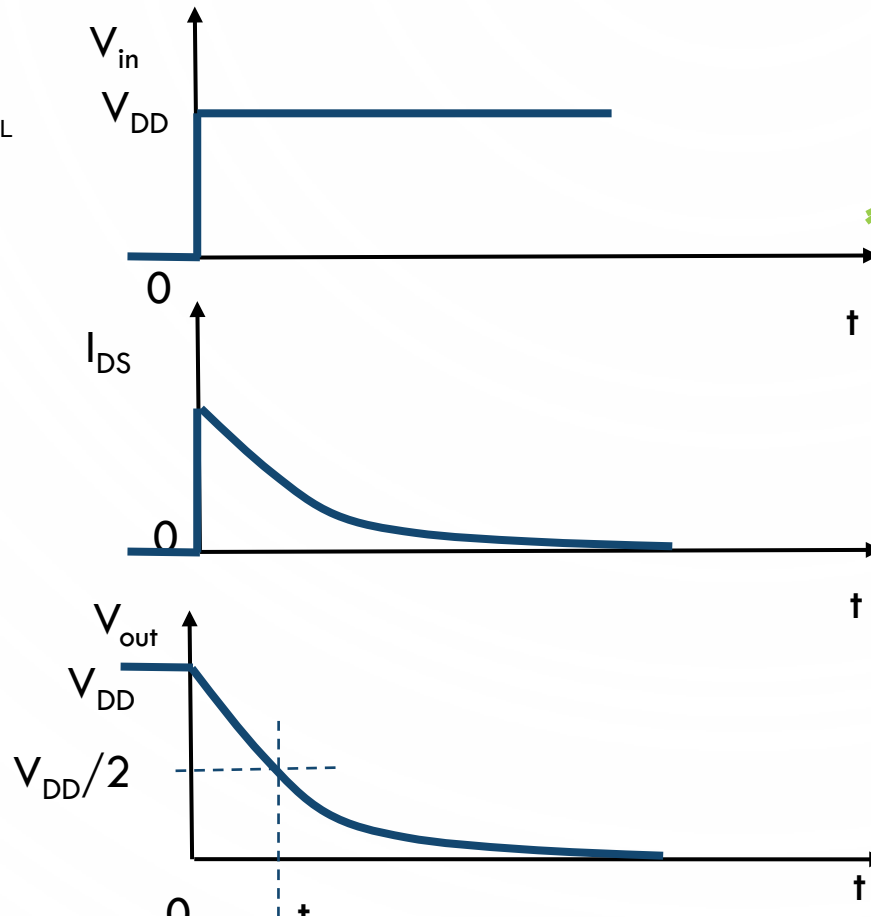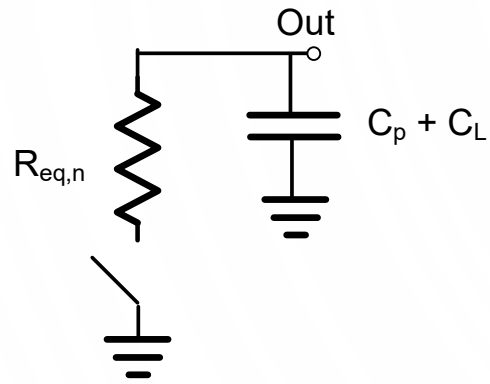$R_{eq,p}$

Out
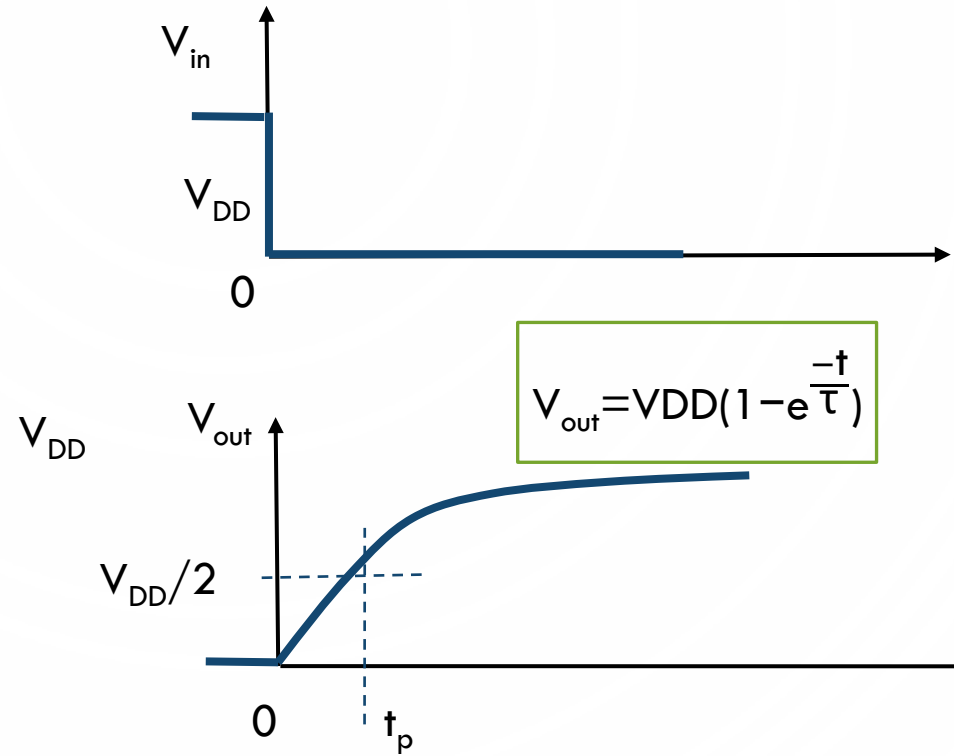
$C_p + C_L$

# Equivalent Resistances

- Transistor $I_{DS}$-$V_{DS}$ trajectory

- Averaging produces $R_{eq}$

# Equivalent Resistances

- Transistor $I_{DS}$-$V_{DS}$ trajectory

- Averaging produces $R_{eq}$

Out

$R_{eq,n}$

$C_p + C_L$

$V_{in}$

$V_{DD}$

0

t

$I_{DS}$

0

t

$V_{out}$

$V_{DD}$

$V_{DD}/2$

0

t

$R_{eq}$ ($V_{GS}$=$V_{DD}$, $V_{DS}$=$V_{DD}/2$)

$I_{DS}$

$V_{GS}$= $V_{DD}$

$R_{eq}$ ($V_{GS}$=$V_{DD}$, $V_{DS}$=$V_{DD}$)

$V_{GS}$= 0

$V_{DD}/2$

$V_{DD}$

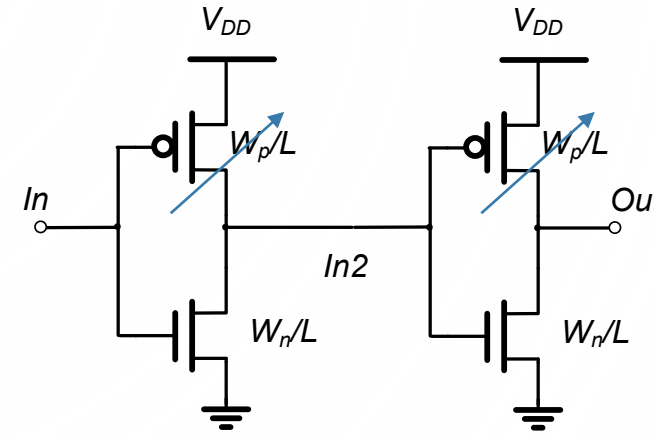$V_{DS}$

$R_{eq} = (R_{eq,start} + R_{eq,mid})/2$

# Optimal P/N Sizing

- Increasing Wp:
  - Reduces $R_p$, increases $C_{in,p}$
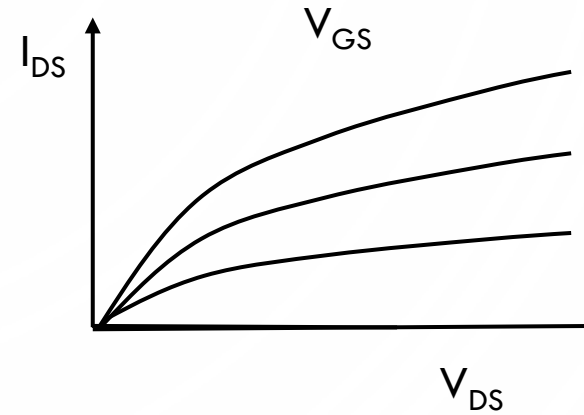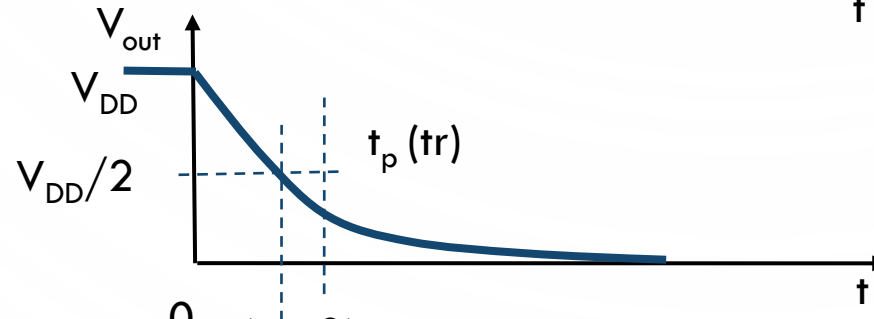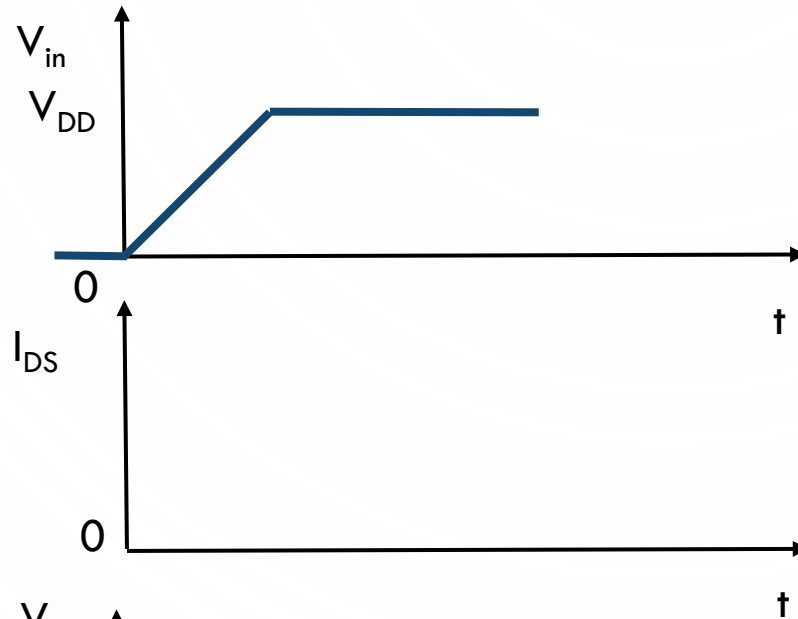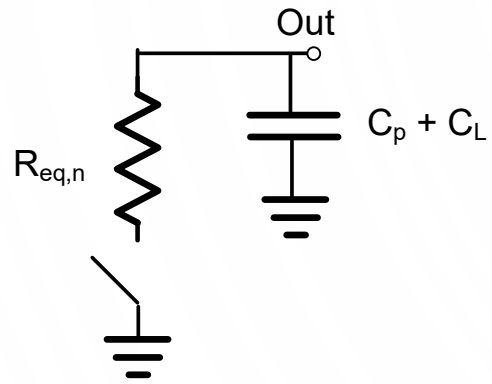  - Reduces $t_{p,LH}$
  - Increases $t_{p,HL}$
- Optimum
  - Wp/Wn = 2 in older technologies, with velocity saturation (like 130nm)
  - Wp/Wn = 1.6 in technologies with strained silicon (e.g. 28nm)
  - Wp/Wn = 1 in finFET technologies
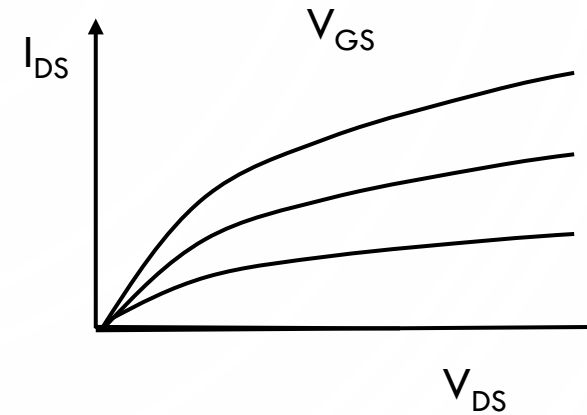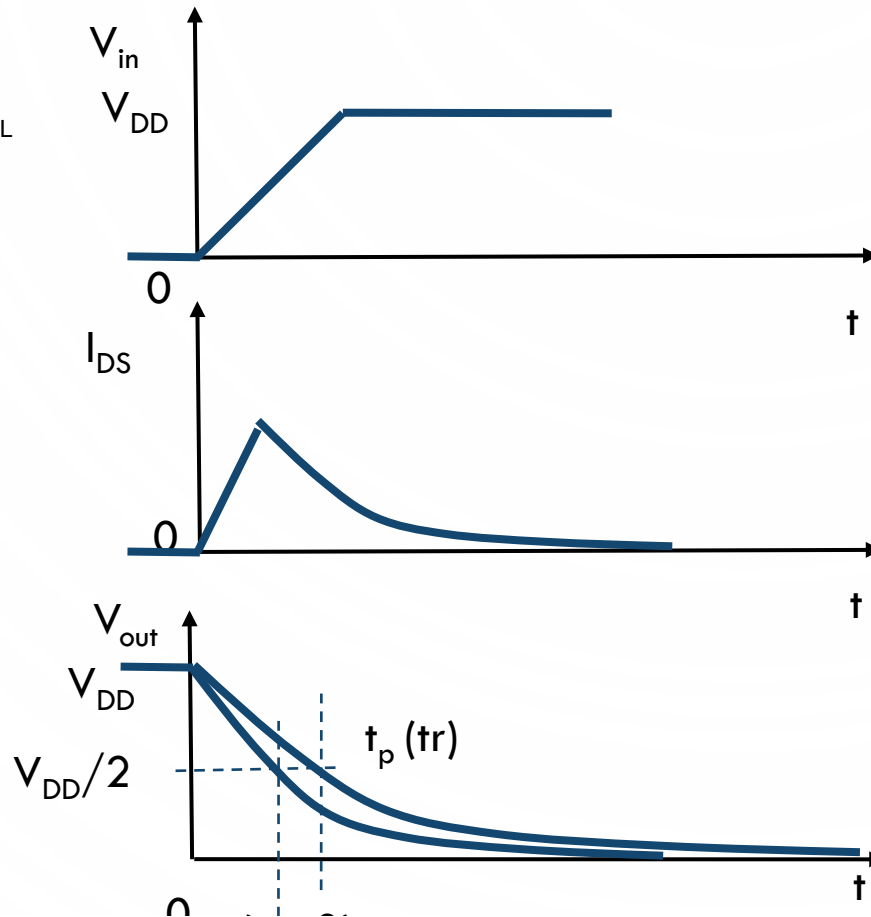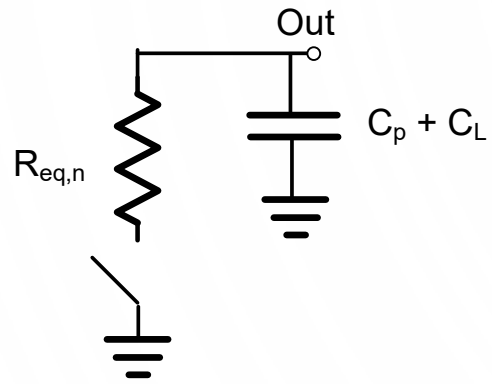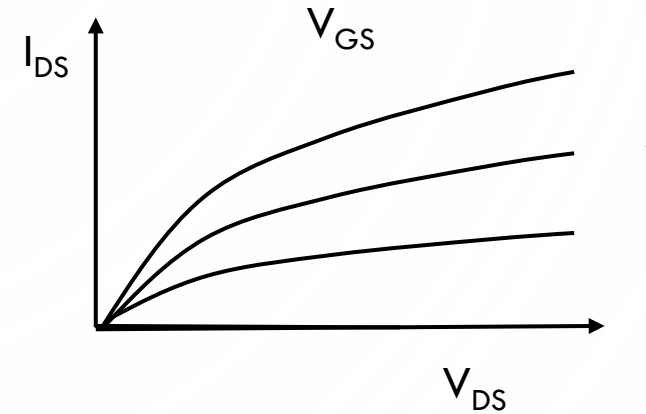
# Impact of Rise/Fall times

- Impacts the $I_{DS}$-$V_{DS}$ trajectory

# Impact of Rise/Fall times

- Impacts the $I_{DS}$-$V_{DS}$ trajectory

# Impact of Supply Voltage

- Lowering VDD, slows down the circuit
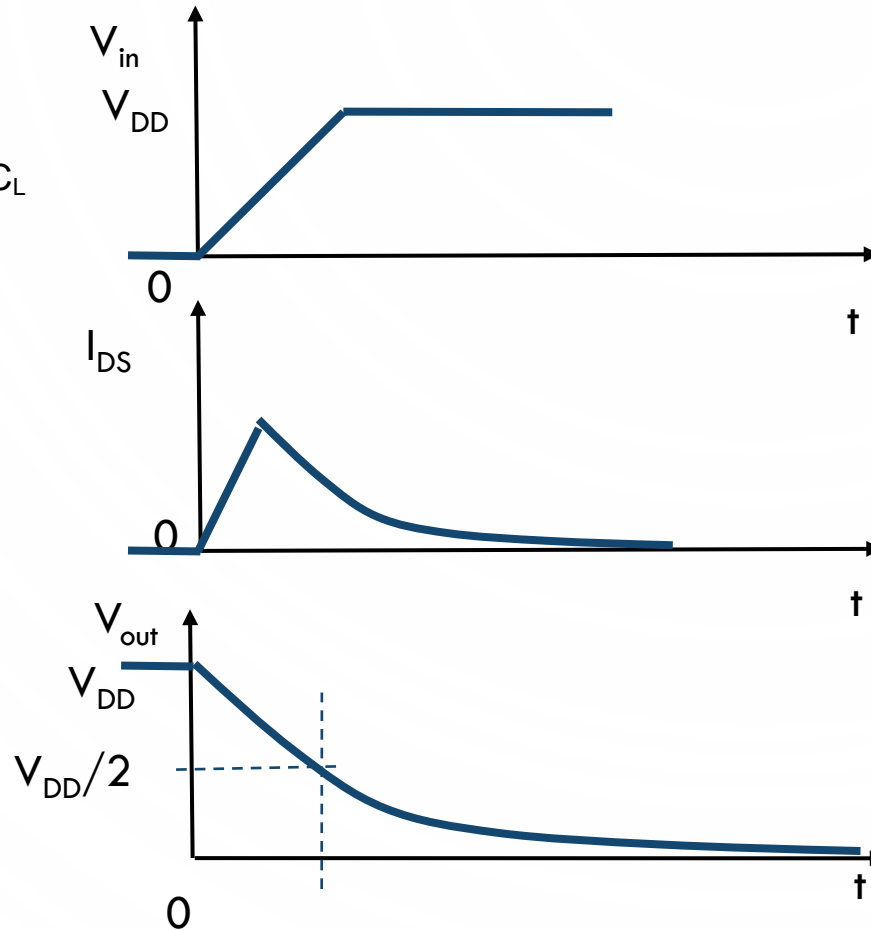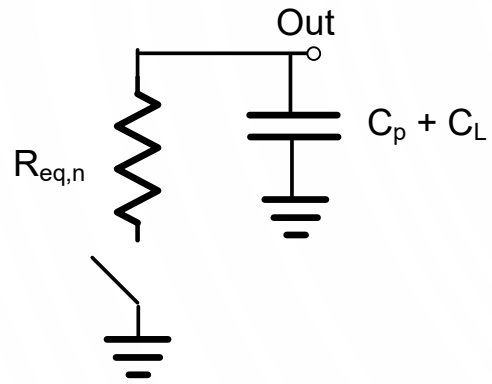
# Quiz: Inverter Delay

- If we double the load capacitance, assuming the default Vout shown in blue, which of the following waveforms shows the new Vout?
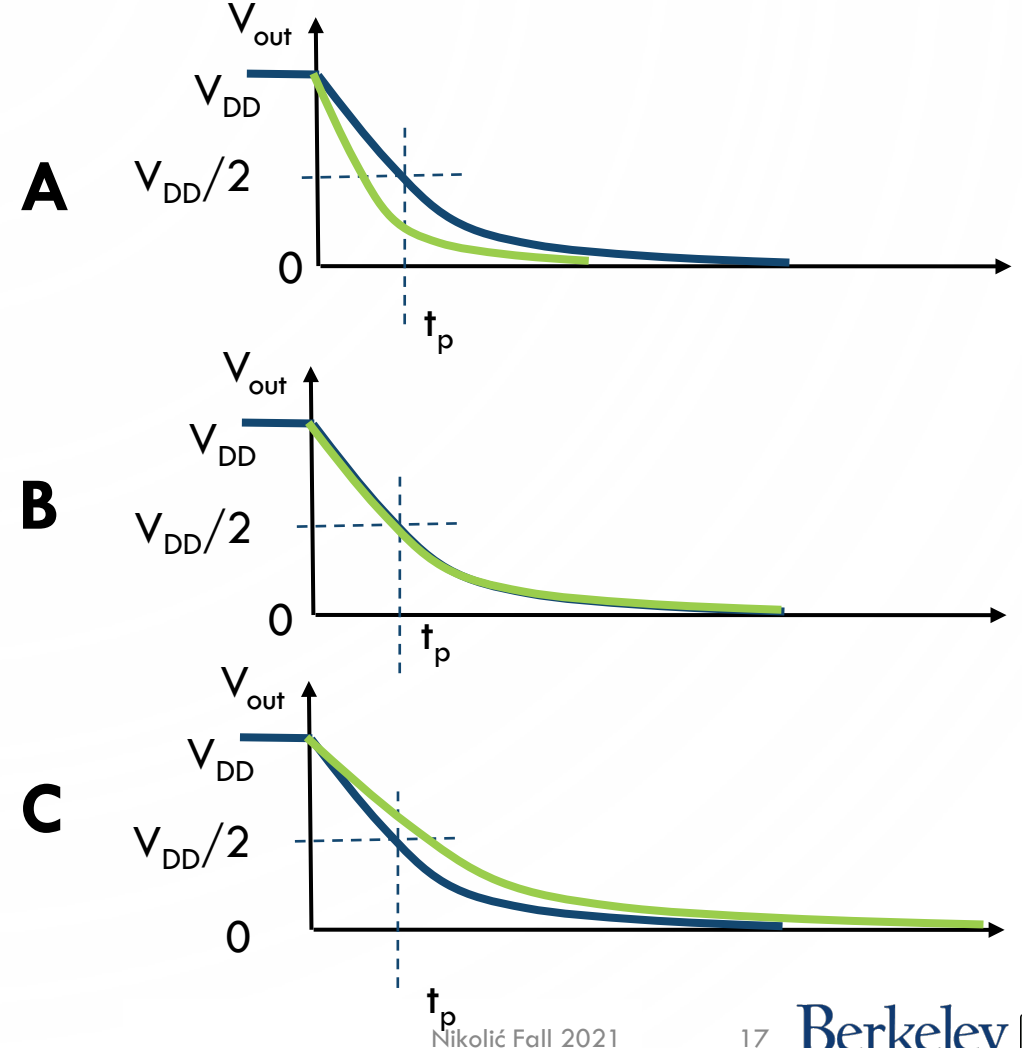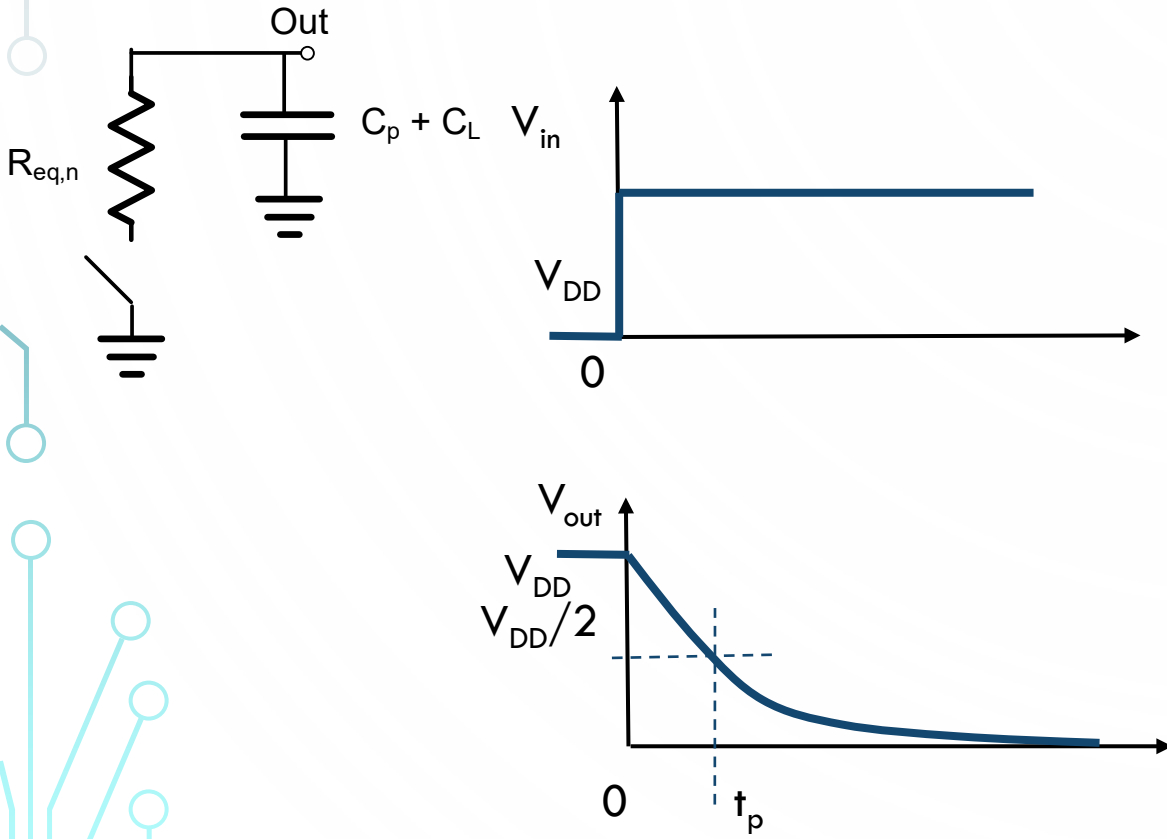
# Sizing CMOS Gates

# Sizing for equal output resistance
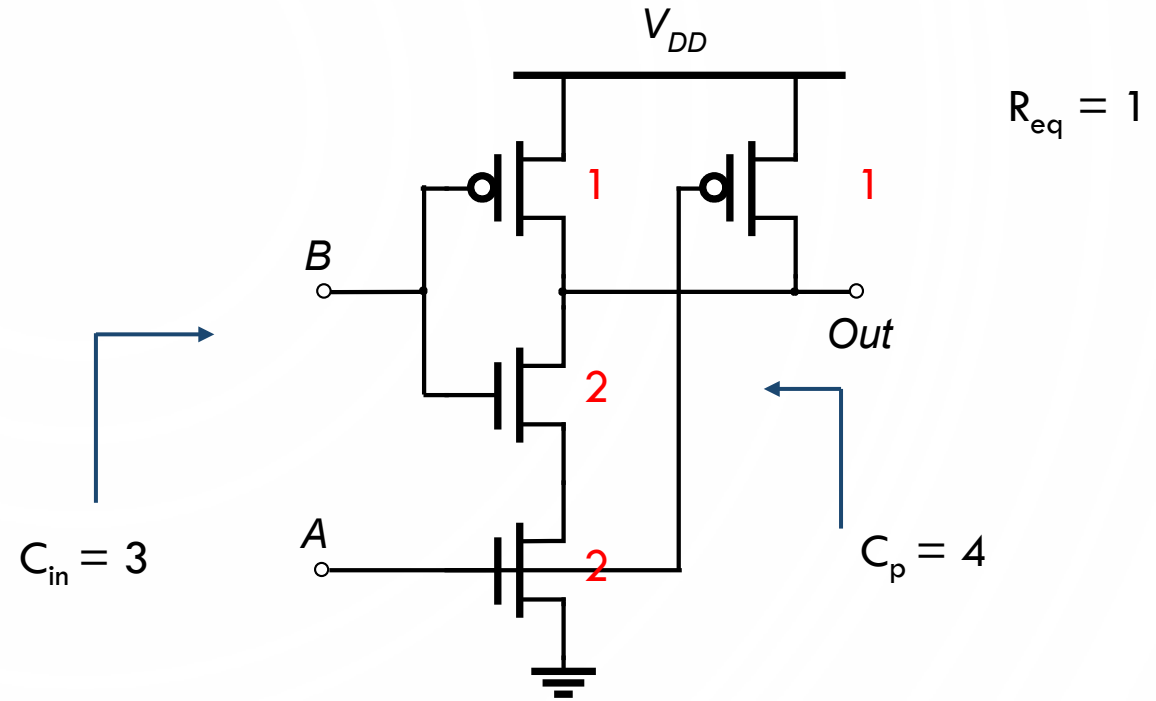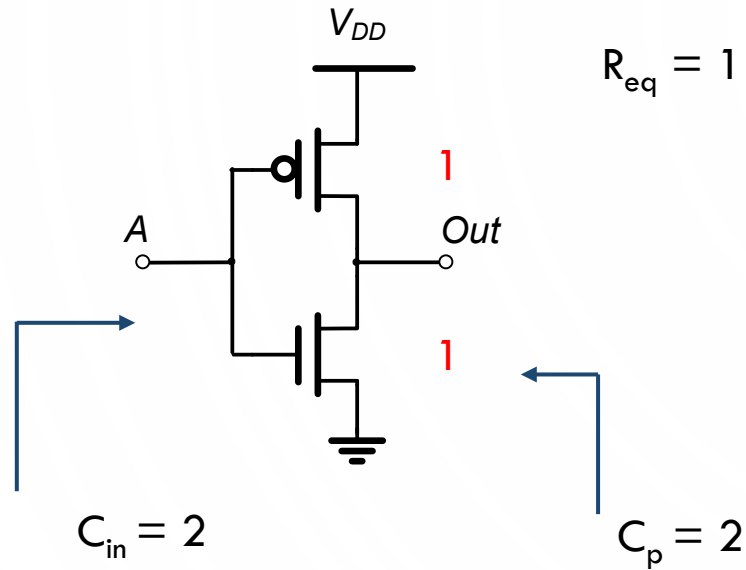
$R_{eq} = 1$

$V_{DD}$

1

$A$ —— $Out$

1

$C_{in} = 2$

$C_p = 2$

$C_{in} = 3$

$V_{DD}$

$R_{eq} = 1$

$B$

1    1

2

$A$

2    $Out$

$C_p = 4$

- In velocity-saturated devices Ion of a stack is 2/3 (not a half) of two devices
  - So the correct upsizing factor is 1.5 (not 2)

- We will use 2, as it makes calculations easier

$V_{DD}$

B

A

$C_{in} = 3$

Out

$C_p = 4$

$V_{DD}$

C

B

A

Out
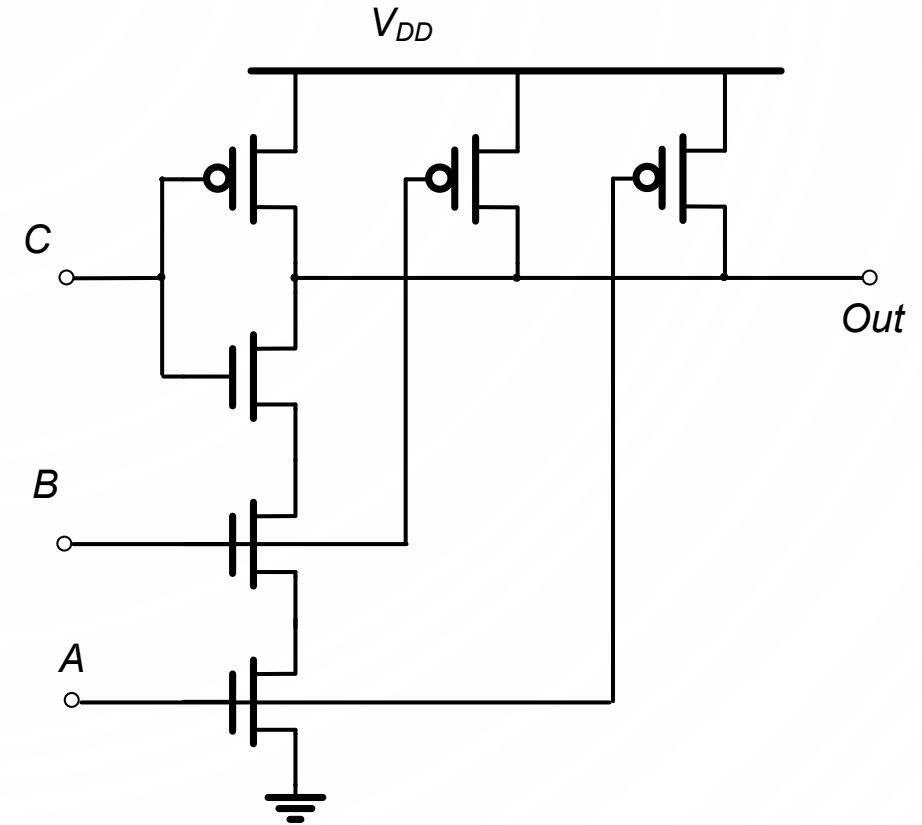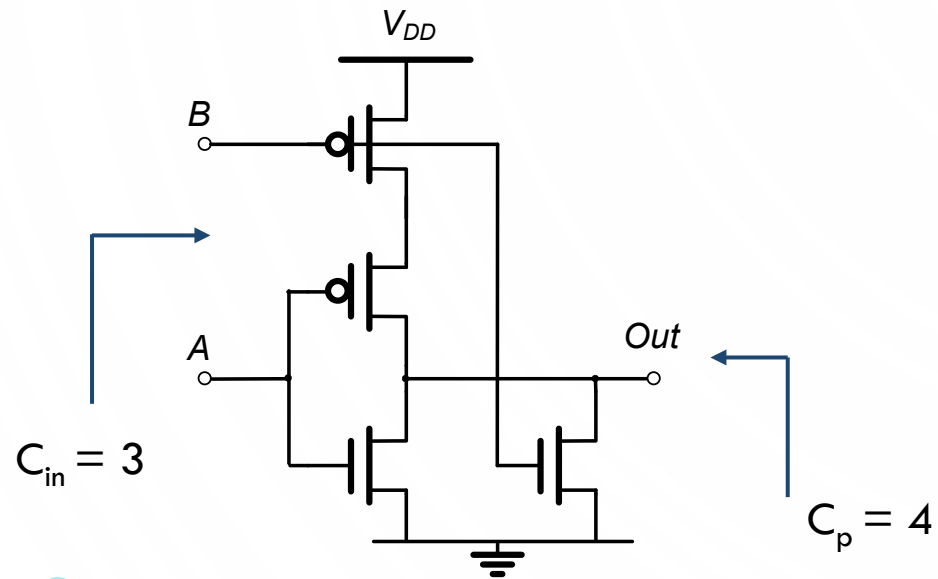
# Stack Ordering
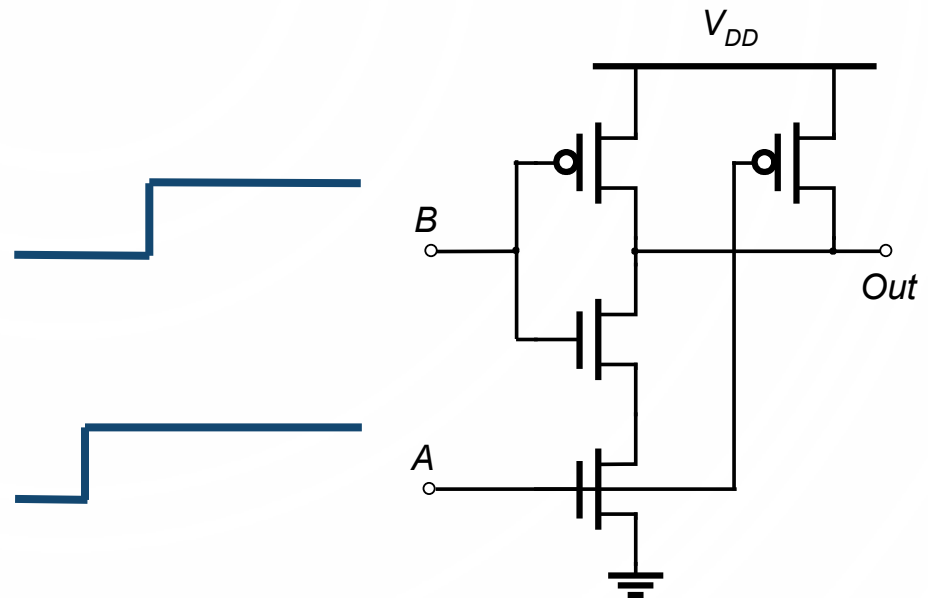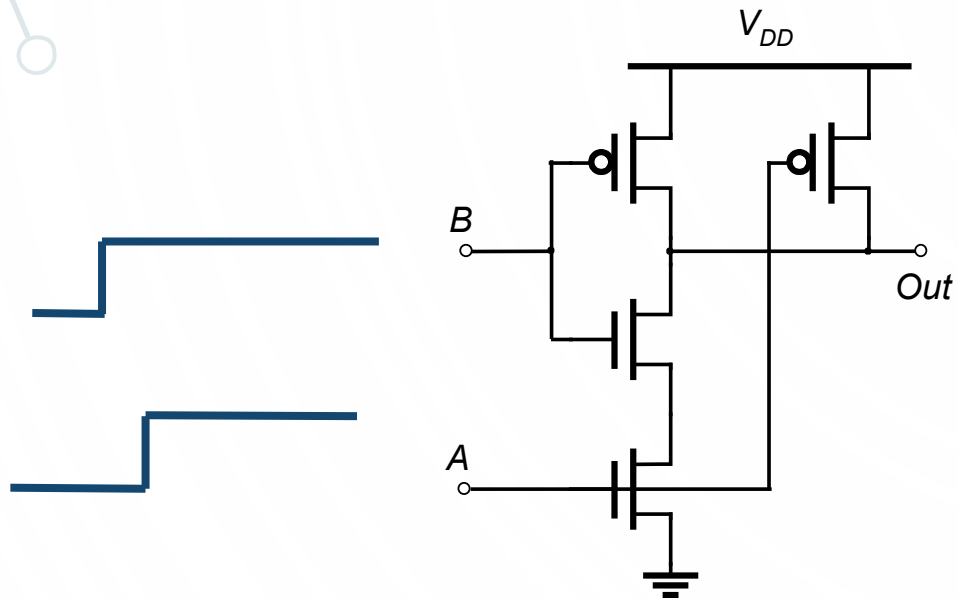


• Critical path goes on top of stack

# Administrivia

- Homework 5 due this week

- Lab 6 (last) this week

- Projects start next week


- There is still time to apply for undergrad scholarships in SoC design!

- Have you thought about doing undergrad research?

# Minimizing Logic Delay

# Inverter RC Delay

$R_{eq}, C_p$

In — Out

$C_L = C_{in}$ (next gate)
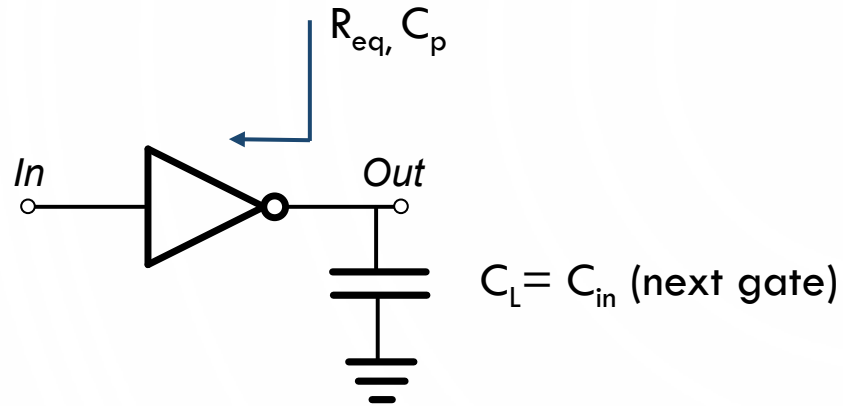
- $t_p = R_{eq}(C_p + C_L) = Req(\gamma\ Cin + C_L)$
  - $\gamma = 1$ (closer to 1.2 in recent processes)

- $t_p = R_{eq}C_{in}(1 + C_L/C_{in}) = \tau_{INV}(1 + f)$
  - Propagation delay is proportional to fanout

- Normalized Delay = 1 + f

Fanout = $f = C_L/C_{in}$

$t_p = \tau_{INV}(1 + f)$

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# Generalizing to Arbitrary Gates

- Delay has two components: d = f + p

- f: *effort delay* = $gh$ (a.k.a. stage effort)
  - Again has two components

- g: *logical effort*
  - Measures relative ability of gate to deliver current
  - $g = 1$ for inverter

- h: *electrical effort* = $C_{out} / C_{in}$
  - Ratio of output to input capacitance
  - Sometimes called fanout

- p: parasitic delay
  - Represents delay of gate driving no load
  - Set by internal parasitic capacitance

# Inverter Delay



$V_{DD}$

$R_{eq} = 1$

1

$A$ —— $Out$

1

$C_{in} = 2$

$C_p = 2$
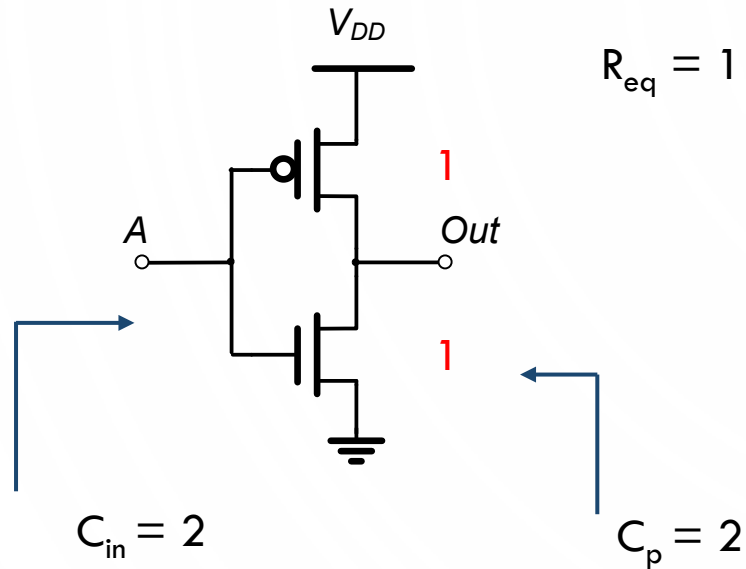
- Parasitic p is the ratio of intrinsic capacitance to an inverter
  - p(inverter) =
- Logical Effort g is the ratio of input capacitance to an inverter
  - g(inverter) =
- Electrical Effort h is the ratio of the load capacitance to the input capacitance
  - h(inverter) =
- Delay = p + f = p + g * h = 1 + f

# NAND2 Gate



$V_{DD}$

$R_{eq} = 1$

A

Out

1

1

$C_{in} = 2$

$V_{DD}$

$R_{eq} = 1$

B

Out

1      1

2

A

2
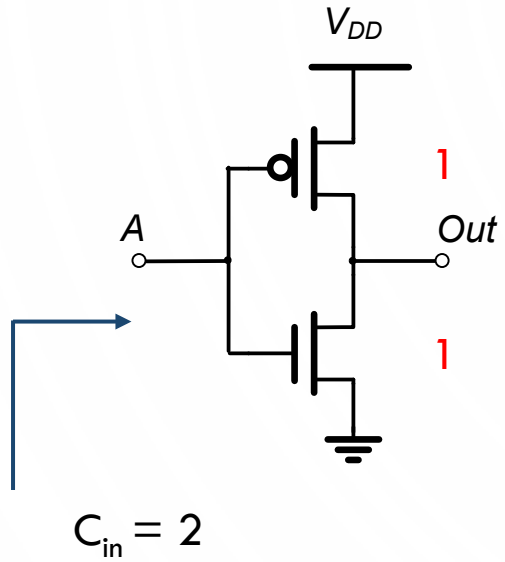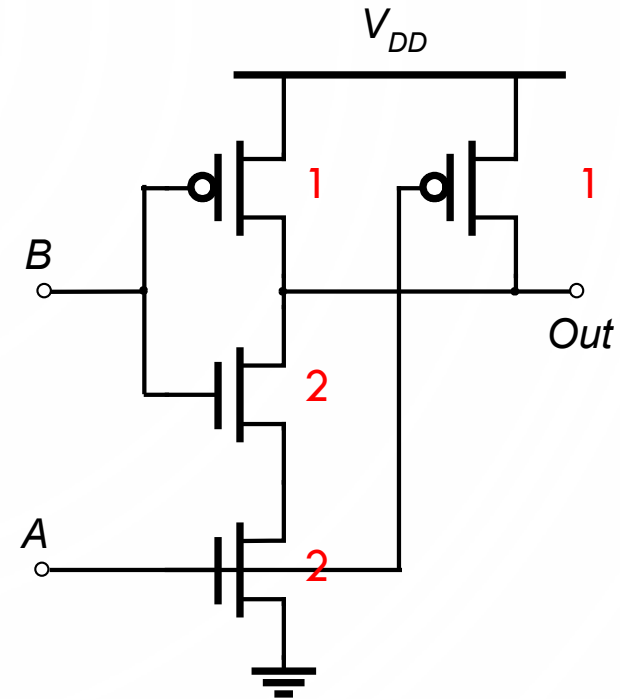
$C_{in} = 3$

# Logical Effort of NAND2 Gate
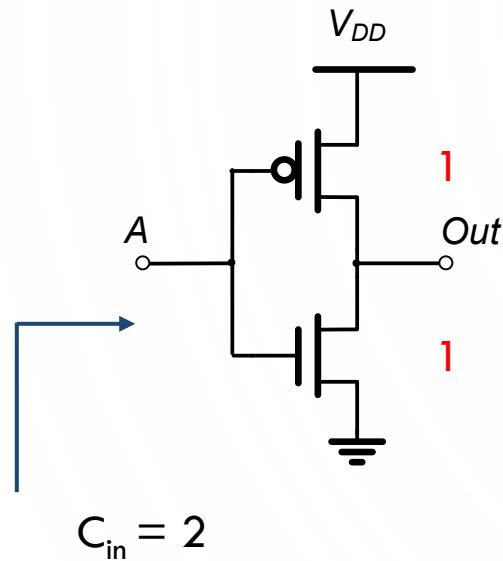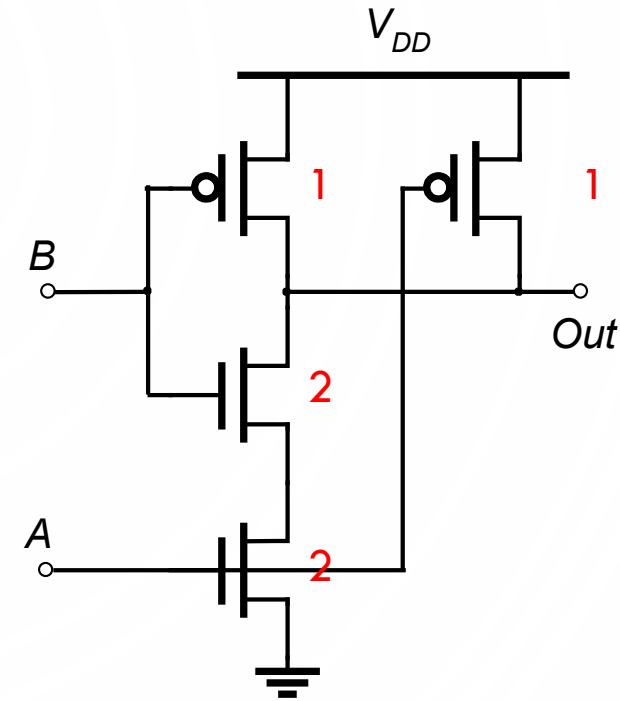
$R_{eq} = 1$

$C_{in} = 2$

$C_{in} = 3$

$R_{eq} = 1$

- In velocity-saturated devices Ion of a stack is 2/3 (not a half) of two devices
  - So the correct upsizing factor is 1.5 (not 2)

- We will use 2, as it makes calculations easier

# NOR2 Gate

$V_{DD}$

$R_{eq} = 1$

A

1

Out

1

$C_{in} = 2$

$V_{DD}$

B

A

Out

$C_{in} = 3$
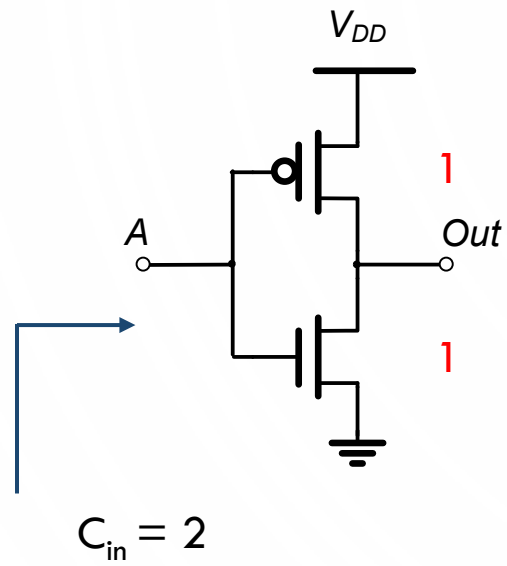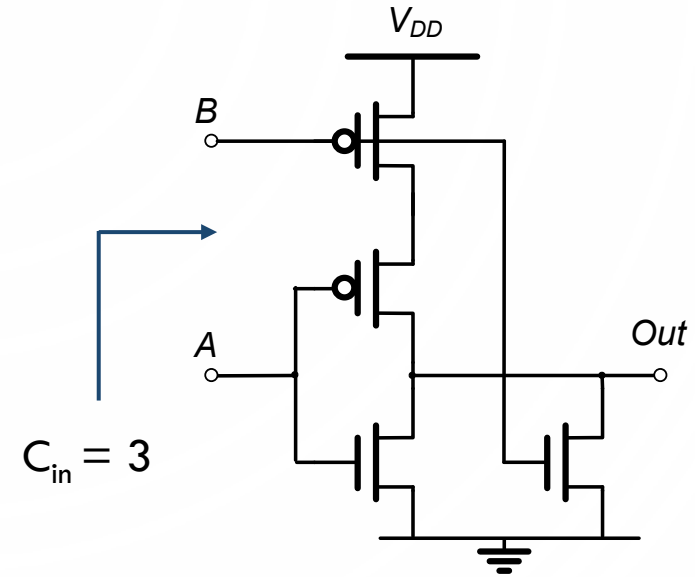
# Example: Inverter Chain
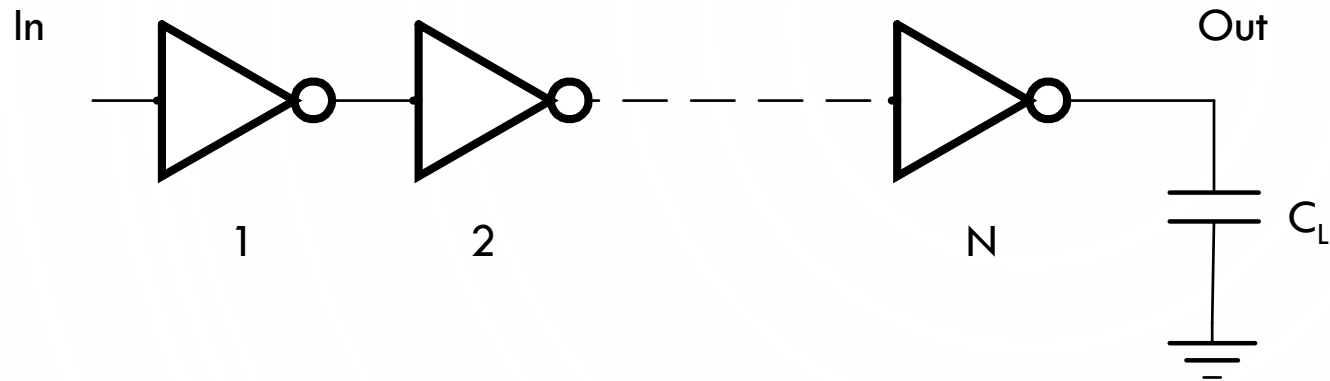


Logical Effort:    g =

Electrical Effort:   h =

Parasitic Delay:   p =

Stage Delay:      d =

Total Delay:      d_total =

# Example: Inverter Chain
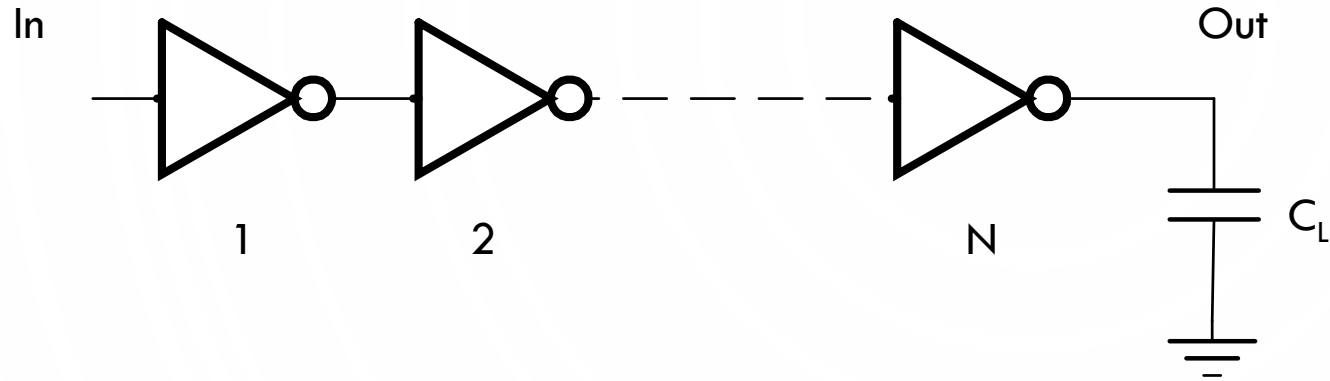


Logical Effort:        $g = 1$

Electrical Effort:   $h = 1$

Parasitic Delay:   $p = 1$

Stage Delay:        $d = 2$

Total Delay:        $d\_total = 2*N$

# Example: FO4 Inverter

- Estimate the delay of a fanout-of-4 (FO4) inverter
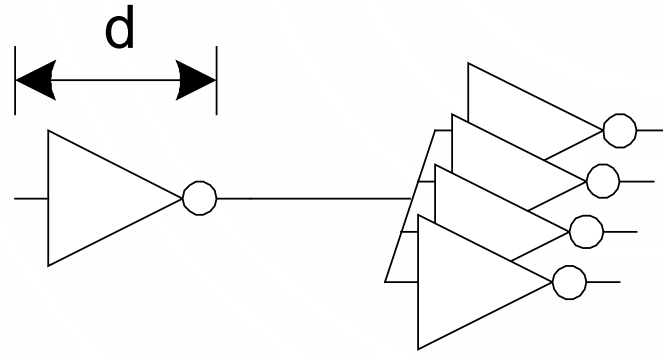


Logical Effort:        g =

Electrical Effort:    h =

Parasitic Delay:    p =

Stage Delay:      d =

# Example: FO4 Inverter

- Estimate the delay of a fanout-of-4 (FO4) inverter
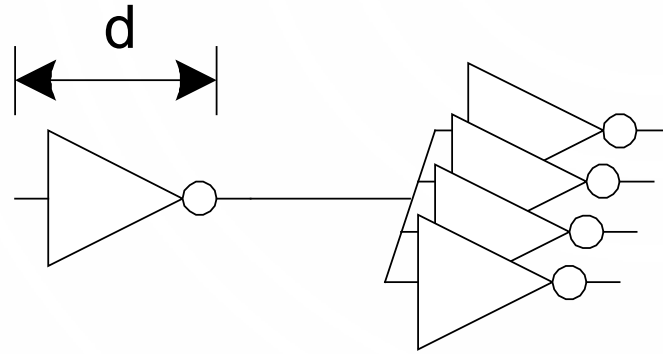


Logical Effort:      $g = 1$

Electrical Effort:    $h = 4$

Parasitic Delay:    $p = 1$

Stage Delay:      $d = 5$

# Multi-stage Logic Networks

- Logical effort generalizes to multistage networks

- *Path Logical Effort*

$$G = \prod g_i$$

- *Path Electrical Effort*

$$H = \frac{C_{\text{out-path}}}{C_{\text{in-path}}}$$

- *Path Effort*

$$F = \prod f_i = \prod g_i h_i$$



$$g_1 = 1 \qquad g_2 = 5/3 \qquad g_3 = 4/3 \qquad g_4 = 1$$
$$h_1 = x/10 \qquad h_2 = y/x \qquad h_3 = z/y \qquad h_4 = 20/z$$
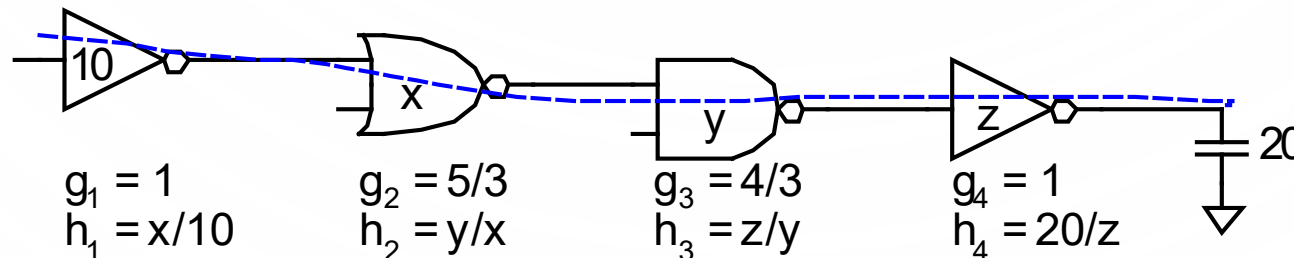
# Branching Effect

$$b = \frac{C_{\text{on path}} + C_{\text{off path}}}{C_{\text{on path}}} \qquad B = \prod b_i$$

G     = 1

H     = 90 / 5 = 18

GH   = 18

$h_1$     = (15 + 15) / 5 = 6

$h_2$     = 90 / 15 = 6

B     = 2

F     = $g_1 g_2 h_1 h_2$ = 36 = BGH
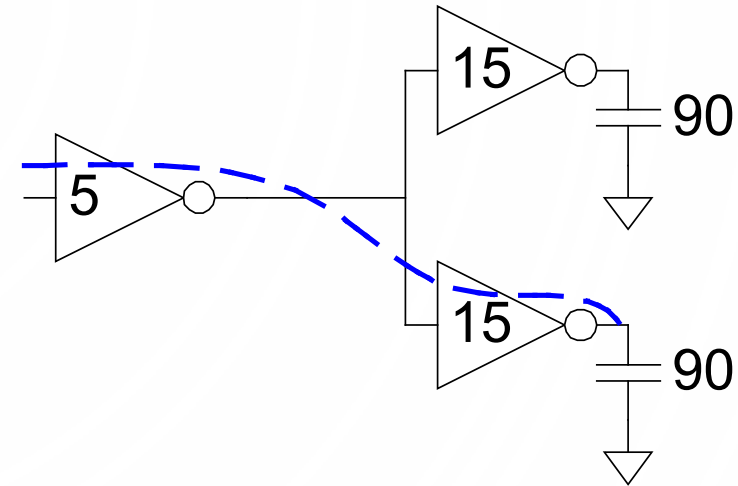
# Designing Fast Circuits

$$D = \sum d_i = D_F + P$$

- Delay is smallest when each stage bears same effort

$$\hat{f} = g_i h_i = F^{\frac{1}{N}}$$

- Thus minimum delay of N stage path is

$$D = NF^{\frac{1}{N}} + P$$

- This is a key result of logical effort
  - Find fastest possible delay
  - Doesn't require calculating gate sizes

# Example: Best Number of Stages

- How many stages should a path use?

  - Minimizing number of stages is not always fastest

- Example: drive 64-bit datapath with unit inverter

$$D = NF^{1/N} + P$$

$$= N(64)^{1/N} + N$$

Initial Driver

Datapath Load

| N: | 1 | 2 | 3 | 4 |
|----|---|---|---|---|
| f: | | | | |
| D: | | | | |

# Example: Best Number of Stages

- How many stages should a path use?
  - Minimizing number of stages is not always fastest

- Example: drive 64-bit datapath with unit inverter

$$D = NF^{1/N} + P$$

$$= N(64)^{1/N} + N$$



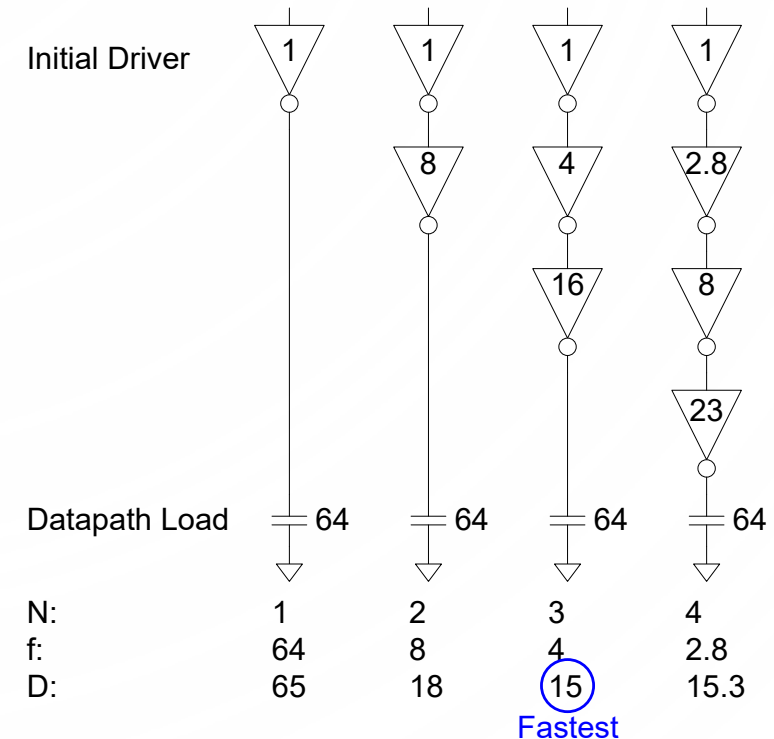| N: | 1 | 2 | 3 | 4 |
|---|---|---|---|---|
| f: | 64 | 8 | 4 | 2.8 |
| D: | 65 | 18 | 15 | 15.3 |
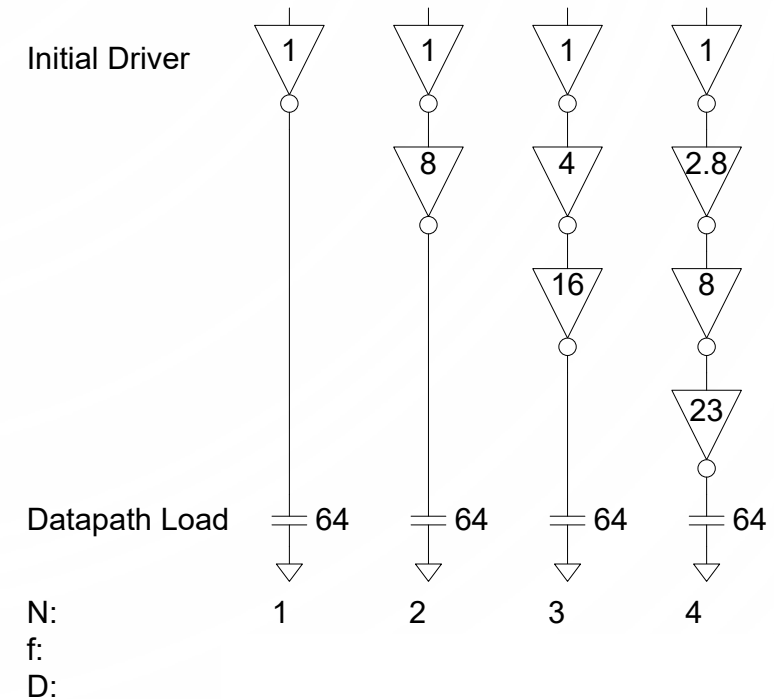
Fastest

# Best Stage Effort

- How many stages should a path use?
  - To drive given capacitance



In ... Out

1    2    N    $C_L$

$$D = NF^{1/N} + Np_{inv}$$
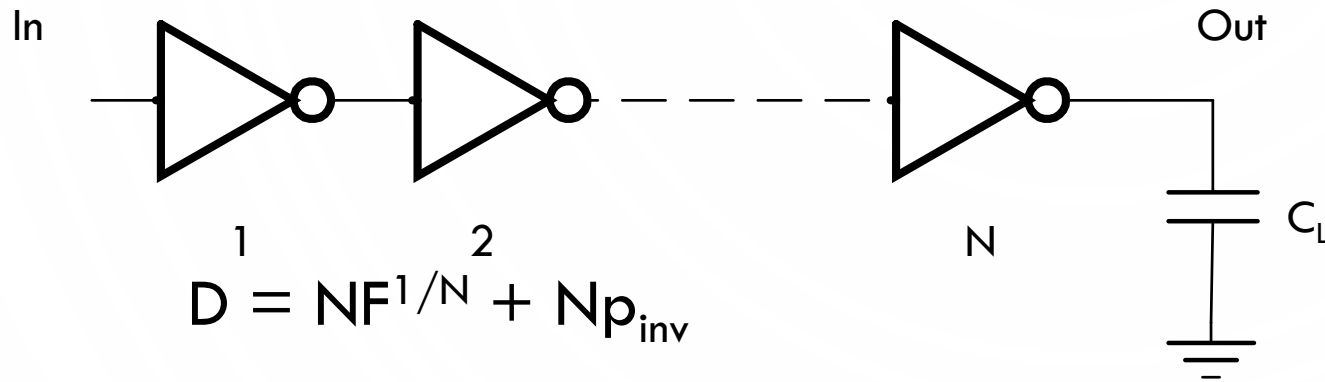
- Define best stage effort

- Neglecting parasitics ($p_{inv} = 0$), we find $\rho = e = 2.718$ ... $\rho = F^{\frac{1}{N}}$

- For $p_{inv} = 1$, solve numerically for $\rho = 3.59$

- Choose 4 – less stages, less energy

# Logical Efforts Method

1) Compute path effort

2) Estimate best number of stages

3) Sketch path with N stages

4) Estimate least delay

5) Determine best stage effort

6) Find gate sizes

$$F = GBH$$

$$N = \log_4 F$$

$$D = NF^{\frac{1}{N}} + P$$

$$\hat{f} = F^{\frac{1}{N}}$$

$$C_{in_i} = \frac{g_i C_{out_i}}{\hat{f}}$$

# Summary

- Delay is a linear function of R and C

- Delay optimization is critical to improve the frequency of the circuit.

- The dimensions of a transistor affect its capacitance and resistance.

- We use RC delay model to describe the delay of a circuit.

- Two delay components:

  - Parasitic delay (p)

  - Effort delay (F)

    - Logical effort (g): intrinsic complexity of the gate

    - Electrical effort (h): load capacitance dependent

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