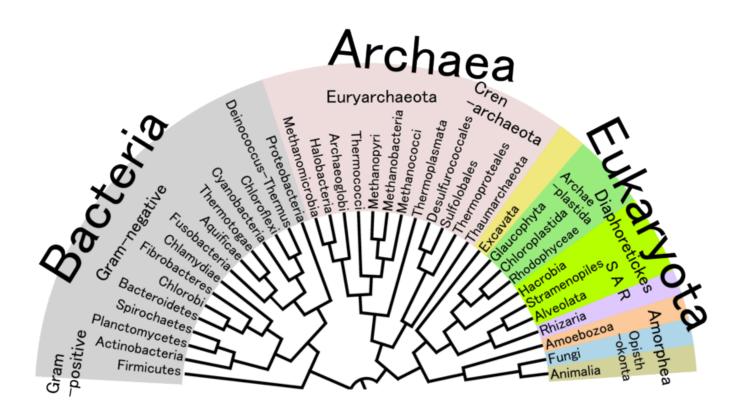
## Phylogenetics

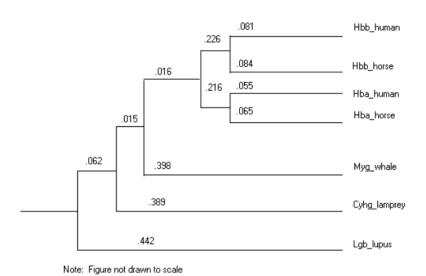


ECES T480/680 Rosen

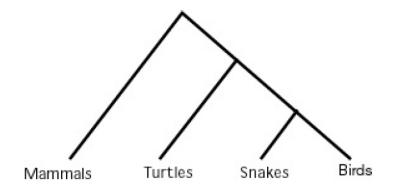
## Constructing Phylogenetic Trees

- Phylogenetic trees illustrate the evolutionary relationships among groups of organisms, or among a family of related nucleic acid or protein sequences
- E.g., how might have this family been derived during evolution

#### **Globin Sequences**



## Hypothetical Tree Relating Organisms



## Phylogenetic Relationships Among Organisms

- Entrez: www.ncbi.nlm.nih.gov/Taxonomy
- Ribosomal database project: rdp.cme.msu.edu/html/
- Tree of Life: phylogeny.arizona.edu/tree/phylogeny.html

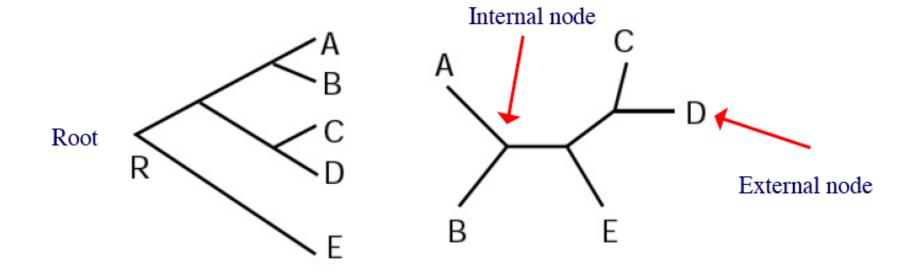
## Phylogeny Applications

- Tree of life: Analyzing changes that have occurred in evolution of different organisms
- Phylogenetic relationships among genes can help predict which ones might have similar functions (e.g., ortholog detection)
- Follow changes occuring in rapidly changing species (e.g., HIV virus)

## **Traditional Methods**

- Traditionally: morphological features (e.g.,number of legs, beak shape, etc.)
- Today: Mostly molecular data (e.g., DNA and protein sequences)

### Rooted vs. Unrooted Trees



Rooted tree

Unrooted tree

Note: Here, each node has three neighboring nodes

## Terminology

- External nodes: things under comparison; operational taxonomic units (OTUs)
- Internal nodes: ancestral units; hypothetical; goal is to group current day units
- Root: common ancestor of all OTUs under study. Path from root to node defines evolutionary path
- Unrooted: specify relationship but not evolutionary path
  - If have an outgroup (external reason to believe certain OTU branched off first), then can root
- Topology: branching pattern of a tree
- Branch length: amount of difference that occurred along a branch

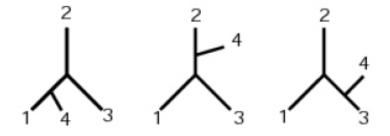
### Tree construction methods

- Distance methods: evolutionary distances are computed for all OTUs and build tree where distance between OTUs "matches" these distances
- Maximum parsimony (MP): choose tree that minimizes number of changes required to explain data
- Maximum likelihood (ML): under a model of sequence evolution, find the tree which gives the highest likelihood of the observed data

## Number of Possible Trees

Given *n* OTUs, there are  $\prod_{i=3}^{n} (2i - 5)$  unrooted trees





OTUs	unrooted trees
3	1
4	3
5	15
10	2,027,025

## Number of possible trees

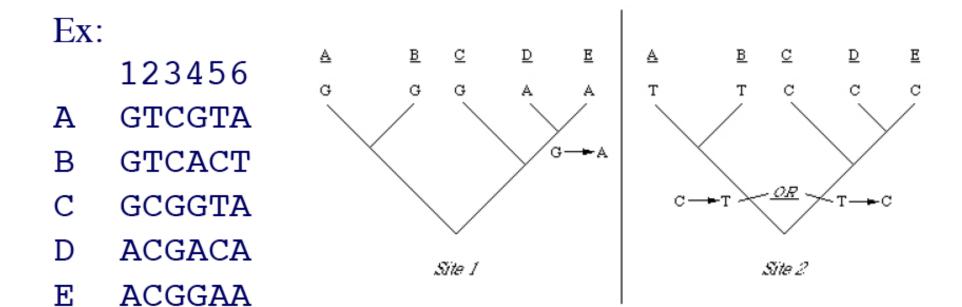
Given *n* OTUs, there are  $\prod_{i=3}^{n} (2i-3)$  rooted trees

Bottom Line: an enumeration strategy over all possible trees to find the best one under some criteria is not feasible!

OTUs	Rooted trees
3	3
4	15
5	105
10	34,459,425

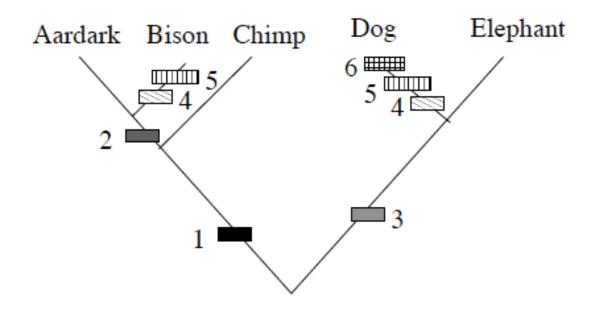
## Parsimony

Find tree which minimizes number of changes needed to explain data



## Example for all sites

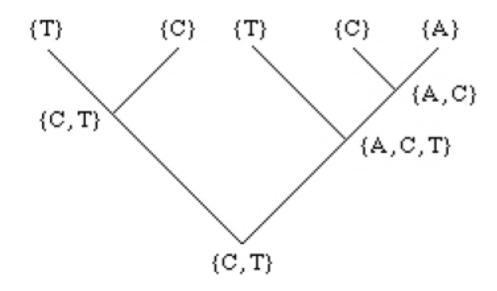
Species	site 1	site 2	site 3	site 4	site 5	site 6
Aardvark	С	A	G	G	T	A
Bison	C	A	G	A	C	A
Chimp	C	G	G	G	T	A
Dog	T	G	C	A	C	T
Chimp Dog Elephant	T	G	C	G	T	A



## Parsimony

- For given example tree and alignment, can do this for all sites, and get away with as few as 8 changes
- Changing the tree (either the topology or labeling of leaves) changes the minimum number of changes need
- Two computational problems
  - (Easy) Given a particular tree, how do you find minimum number of changes need to explain data? (Fitch)
  - (Hard) How do you search through all trees?

## Parsimony: Fitch's Algorithm



Idea: construct set of possible nucleotides for internal nodes, based on possible assignments of children

## Parsimony: Fitch's Algorithm

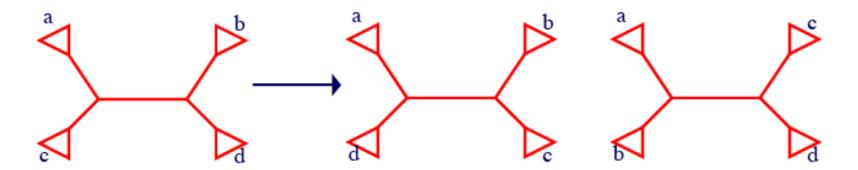
- For each site:
  - Each leaf is labeled with set containing observed nucleotide at that position
  - For each internal node i with children j and k with labels  $S_j$  and  $S_k$

$$S_i = \left\{ \begin{array}{l} S_j \cup S_k \text{ if } S_j \cap S_k \text{ is empty} \\ S_j \cap S_k \text{ otherwise} \end{array} \right.$$

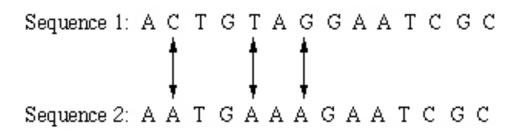
Total # changes necessary for a site is # of union operations

## Parsimony

- How do you search through all trees?
  - Enumerate all trees (too many…)
  - Can use techniques to try to limit the search space (e.g., branch and bound)
  - or use heuristics (many possibilities)
    - E.g., nearest neighbor interchange. Start with a tree and consider neighboring trees. If any neighboring tree has fewer changes, take it as current tree. Stop when no improvements



## Computing Distances between two sequences



Could compute fraction of mismatches between two sequences; however, this is an underestimate of actual distance

## A simple clustering method for building a ROOTED tree

UPGMA (Unweighted Pair Group Method using Arithmetic averages)
Or the Average Linkage Method

Given two disjoint clusters S<sub>i</sub>, S<sub>j</sub> of sequences.

$$d_{ij} = \frac{1}{|S_i| \times |S_j|} d_{pq}$$

Claim that if  $S_k = S_i \cup S_j$ , then distance to another cluster  $S_l$  is:

$$d_{il} |S_{i}| + d_{jl} |S_{j}|$$

$$d_{kl} = \frac{}{|S_{i}| + |S_{j}|}$$

## Algorithm: Average Linkage

#### **Initialization:**

Assign each x<sub>i</sub> into its own cluster S<sub>i</sub>

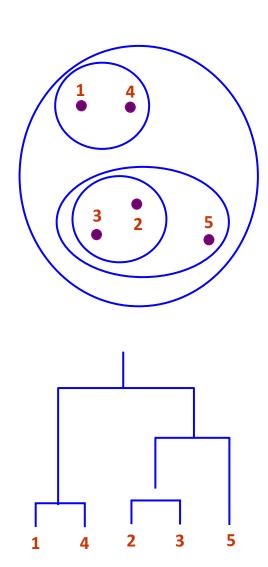
Define one leaf per sequence, height 0

#### **Iteration:**

Find two clusters  $S_i$ ,  $S_j$  s.t.  $d_{ij}$  is min Let  $S_k = S_i \cup S_j$ Define node connecting  $S_i$ ,  $S_j$ , & place it at height  $d_{ij}/2$ Delete  $S_i$ ,  $S_i$ 

#### **Termination:**

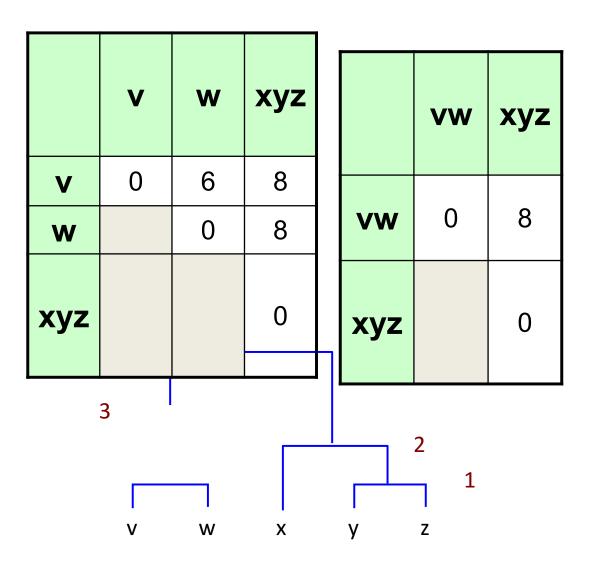
When two clusters i, j remain, place root at height d<sub>ii</sub>/2



## Example

	V	w	X	у	Z
V	0	6	8	8	8
W		0	8	8	8
X			0	4	4
У				0	2
Z					0

W	X	yz
6	8	8
0	8	8
	0	4
		6 8 0 8



### Ultrametric Distances and Molecular Clock

#### **Definition:**

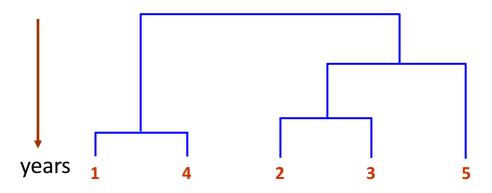
A distance function d(.,.) is ultrametric if for any three distances  $d_{ij} \le d_{ik} \le d_{jk}$ , it is true that

$$d_{ij} \le d_{ik} = d_{jk}$$

#### The Molecular Clock:

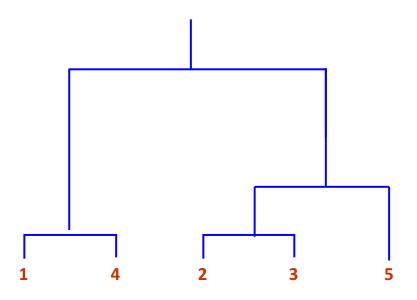
The evolutionary distance between species x and y is  $2\times$  the Earth time to reach the nearest common ancestor

That is, the molecular clock has constant rate in all species



The molecular clock results in ultrametric distances

## Ultrametric Distances & Average Linkage



Average Linkage is guaranteed to reconstruct correctly a binary tree with ultrametric distances

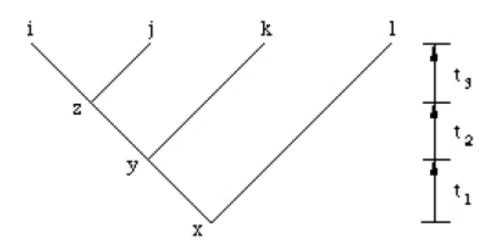
- Given a probabilistic model for nucleotide (or protein) substitution (e.g., Jukes & Cantor), pick the tree that has highest probability of generating observed data
  - I.e., Given data D and model M, find tree T such that Pr(D/T, M) is maximized
- Models gives values  $p_{ij}(t)$ , the probability of going from nucleotide i to j in time t

- Makes 2 independence assumptions
  - Different sites evolve independently
  - Diverged sequences (or species) evolve independently after diverging
- If  $D_i$  is data for *i*th site

$$Pr(D|T, M) = \prod_{i} Pr(D_{i}|T, M)$$

How to calculate  $Pr(D_i/T,M)$ ?

 $p_{xy}(t) \sim \text{prob}$ of going from xto y in time t

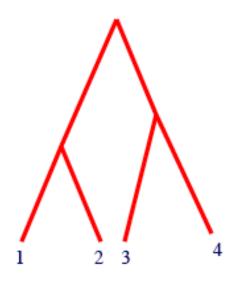


$$Pr(i, j, k, l|T, M) = \sum_{x} \sum_{y} \sum_{z} pr(x) (p_{xl} \cdot (t_1 + t_2 + t_3) \cdot p_{xy}(t_1) \cdot p_{yk}(t_2 + t_3) \cdot p_{yz}(t_2) \cdot p_{zi}(t_3) \cdot p_{zj}(t_3))$$

- Given tree topology and branch lengths, can efficiently calculate Pr(D/T, M) using dynamic programming
  - I.e., don't have to enumerate over all internal states
- Finding best maximum likelihood tree is expensive
  - Must consider all topologies
  - Find best edge lengths for each topology
    - Idea: use some search procedure, e.g., EM, to optimize these lengths

## Assessing Reliability -- The Bootstrap

Say we've inferred the following tree



Would like to get confidence levels that 1 & 2 belong together, and 3&4 belong together

## Assessing the Reliability - The Bootstrap

Say we're given following alignment:

12345678

- 1 GCAGTACT
- 2 GTAGTACT
- 3 ACAATACC
- 4 ACAACACT

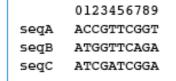
We'll create a pseudosample by choosing sites randomly

until N sites are chosen

(N is length of alignment)

## Bootstrapping





#### Replicate 1

1562314951 seqA CTCCGCTTTC seqB TTCGGTTATT seqC TTCCGTAATT

(a) Step 1

Assemble pseudo-

datasets, repeat

1000 times

#### Replicate 2

5234924418
seqA TCGTTCTTCG
seqB TGGTAGTTTG
seqC TCGAACAATG

#### Replicate 3

5607718907
seqA TCAGGCGTAG
seqB TCAAATGAAA
seqC TCAGGTGAAG

etc

## Bootstrapping

0123456789
seqA ACCGTTCGGT
seqB ATCGATCGGA
seqC ATGGTTCAGA

5071398375

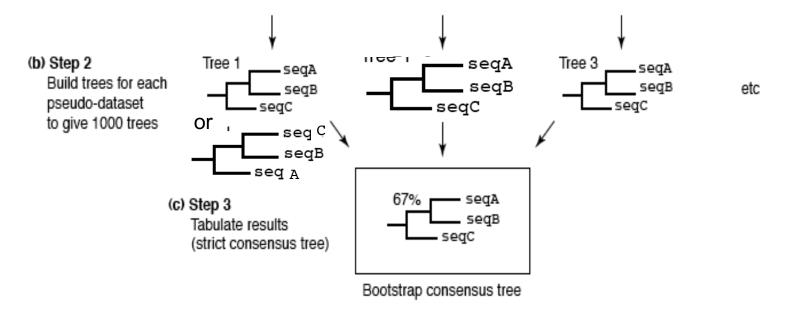
seqA TAGCGTGGGT

seqB TAGTGAGGGT

seqC TAATGAGGAT

4880372653
seqA TGGAGGCCTG
seqB AGGAGGCCTG
seqC TGGAGAGCTG

7748125485
seqA GGTGCCTTGT
seqB GGAGTCTAGT
seqC AATGTGTTGT



## Many to choose from

- Serial Sequence Alignment
  - ClustalW
  - Contralign
  - MUSCLE
  - PROBCONS
  - PROBALIGN
  - Poy
- Serial Tree Inference
  - PAUP
  - Poyt

## Large Data Issues

- Serial
  - Many memory requirements
  - Long time
- Parallel
  - Break into chunks
  - Less time

## Parallel Align& TreeCodes

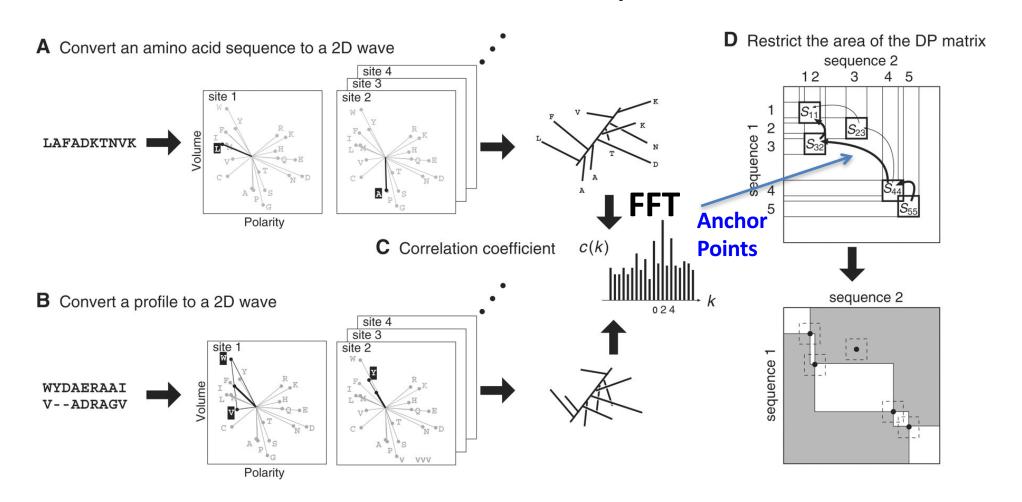
### Alignment

MAFFT

#### **Tree**

- RAxML
- MrBayes
- BEAST(2)
- GARLI
- PhyloBayes
- DPPDIV
- FastTree
- jModelTest2

# MAFFT (Multiple Alignment using Fourier Transform)



## Tree Inference from alignments

ML

- RaxML (several heuristics to reduce search)
- FastTree (nearest neighbor exchanges for ml search)
- Garli (use genetic algorithm for ML search)

Bayesian MCMC (Monte Carlo Markov Chains)

- MrBayes
- PhyloBayes (an infinite mixture model accounting for site-specific amino-acid or nucleotide preferences)
- Beast (relaxed molecular clock and demographic history)

## Tree from different models of Nucleotide Substitution

 DPPDIV (Using Fixed tree topology, change parameters using MCMC)

 jModelTest2 (Likelihood ratio tests, information criterion, and decision theory to get candidate trees)

### Tree Inference – What to use?!

- Almost no systematic comparisons
- ML techniques: RAxML to FastTree (<a href="http://journals.plos.org/plosone/article?id=10.1">http://journals.plos.org/plosone/article?id=10.1</a> 371/journal.pone.0027731)
  - Says FastTree may be faster and just as good as raxml on large datasets
- MCMC: MrBayes is classic
- Ones that don't use alignment seem on the front of the state-of-the-art
  - Suggest UNCOMMON models of molecular evolution