

UNIVERSITY OF HOUSTON

MIDTERM 1 REVIEW

COSC 3320

Algorithms and Data Structures

Note

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1 Exercises

Exercise 1: Big- \mathcal{O}

Consider the following functions:

$$n^{1.5}, n \log n, 2^n, n!, n^n, 1/n$$

Rank the listed functions by order of growth, i.e., give an ordering f_1, f_2, \dots, f_5 such that $f_1 = \mathcal{O}(f_2)$, $f_2 = \mathcal{O}(f_3)$, and so on. Justify your ordering.

Exercise 2: Permutations

A permutation of a sequence $(s_0, s_1, \dots, s_{n-1})$ is a sequence with the same terms, but in a different order. For example, the sequence $(3, 4, 6)$ admits 6 permutations:

$$\begin{array}{cc} (3, 4, 6) & (3, 6, 4) \\ (4, 3, 6) & (4, 6, 3) \\ (6, 3, 4) & (6, 4, 3) \end{array}$$

1. Give a decrease and conquer algorithm to output all permutations of a sequence of n distinct elements.
2. Prove that this algorithm is correct.
3. Give a recurrence for the runtime of this algorithm.
4. Solve this recurrence.

Exercise 3: Hamming Weight

The *Hamming Weight* of a binary number n is the number of bits set to 1 in the binary representation of n . For example, the Hamming Weight of 14 is 3, since

$$14 = 1110_2$$

Give a divide and conquer algorithm to compute the Hamming Weight of a non-negative integer n .

Exercise 4: Recurrences

Consider the following recurrence:

$$T(n) = 3T\left(\frac{n}{2}\right) + n^2$$

1. Show that $T(n) = \mathcal{O}(n^3)$ by induction
2. Solve the recurrence using the DC Recurrence Theorem.

Exercise 5: Sorting

Consider the following sorting algorithm:

```
1: func SELECTION-SORT(arr):
2:   n = LEN(arr)
3:   if n ≤ 1:
4:     return arr
5:   else:
6:     result = SELECTION-SORT(arr[0 : n - 1])  ▷ Sort all but the last element
7:     last-element = arr[n - 1]
```

```
8:      let  $i$  be the first index such that  $\text{arr}[i - 1] \leq \text{last-element} \leq \text{arr}[i]$ 
9:      result.INSERT(last-element,  $i$ )     $\triangleright$  Insert last-element into correct position
10:     return result
```

1. Prove the correctness of this algorithm by induction.
2. Write a recurrence, $T(n)$, for the runtime of this algorithm. Assume finding i in line 8 is done via linear search.
3. Solve this recurrence in terms of asymptotic complexity, i.e., find a function $f(n)$ such that $T(n) = \mathcal{O}(f(n))$.
4. How does this recurrence change if we find i via binary search? How does this affect the asymptotic complexity?