# PRINTABLE VERSION

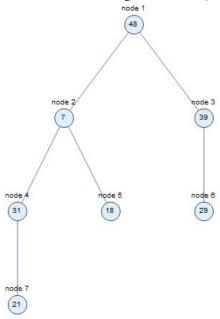
Quiz 7

# You scored 100 out of 100

### **Ouestion 1**

## Your answer is CORRECT.

A graph with nodes and edges is shown below (and we have stored random integers in each node):



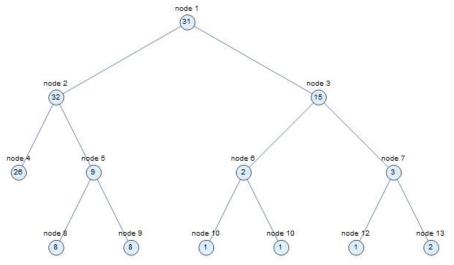
Is this a binary tree?

- a) Yes! This is a binary tree!
- b) No, Node 5 does not have the correct number of children.
- c) Yes! This is a full tree, but it is not binary!
- d) Yes! This is a binary tree, and it is full!

## Question 2

# Your answer is CORRECT.

A full binary tree is said to have the *heap* or *max heap* property if the value stored in a parent node is *always* bigger than or equal to the values stored in *any* of its children. We will use the notation v(p) to denote the value stored at node p, which means that a full binary tree T has property *heap* if, whenever node p is a child of node p, it follows that  $v(p) \ge v(c)$ .



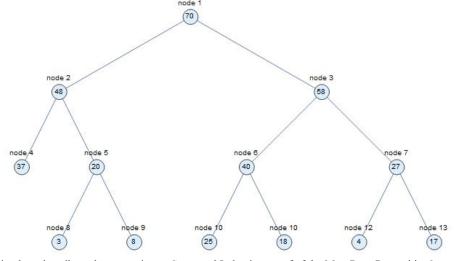
Does the full binary tree above have the heap property?

- a) Yes! This tree has property heap! (Heap, Heap Hurray!)
- b) No! This tree does not have property heap. :(

# **Question 3**

### Your answer is CORRECT.

The main point of this quiz question (and the next three) is to understand how "structural induction" can be used to prove the following proposition: Proposition (The Max-Root Proposition). If a full binary tree, T, has the heap property, then  $v(r) \ge v(n)$  for all nodes n. Here, as usual, r denotes the root of T and v(n) denotes the value (or "key") stored at node n. The full binary tree in the image below satisfies the heap property, and you can check that the Proposition's conclusion is true.



Which, if any, of the following best describes a base case in our Structural Induction proof of the Max-Root Proposition?

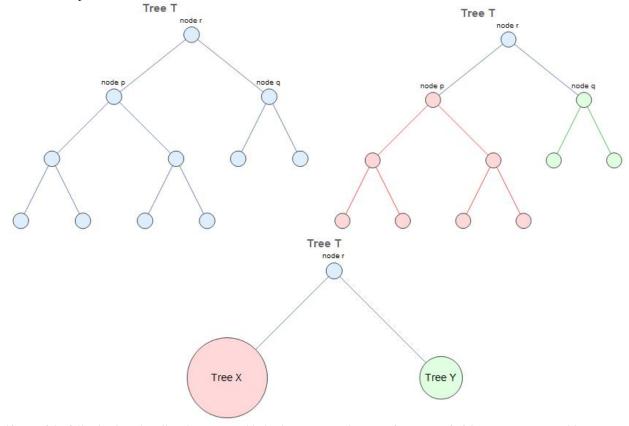
- a)  $\bigcirc$  We would suppose T is a full binary tree with the heap property and then note that  $v(r) \ge v(n)$  for all other nodes n.
- **b)**  $\odot$  We would let T be a tree consisting of one node, r (and note that T is full binary *and* has the heap property for any stored value v(r)). Next we would check that  $v(r) \ge v(n)$  for all nodes n.
- c)  $\bigcirc$  We would suppose T is a full binary tree with the heap property and satisfies  $v(r) \ge v(n)$  for all nodes n.
- **d)**  $\bigcirc$  We would let T be a tree consisting of one node, r (and note that T is full binary *and* has the heap property for any stored value v(r)). Next we would check that  $v(r) \ge v(n)$  for some node n.
- e) We would suppose a full binary heap-property tree T has nodes and then conclude that all the nodes have the same value.

#### **Ouestion 4**

## Your answer is CORRECT.

The main point of this quiz question (and the next two) is to continue our exploration of how "structural induction" can be used to prove the following proposition:

Proposition (The Max-Root Proposition). If a full binary tree, T, has the heap property, then  $v(r) \ge v(n)$  for all nodes n. Here, as usual, r denotes the root of T and v(n) denotes the value (or "key") stored at node n. The image below shows a sample tree, T, that is full binary and is assumed to have the heap property; in fact, three images of T are displayed so as to clarify how to htink about any full binary tree as two subtrees -- X and Y -- joined to the root r.



Which, if any, of the following best describes the "structrual inductive" or "recursive" step for our proof of the Max-Root Proposition?

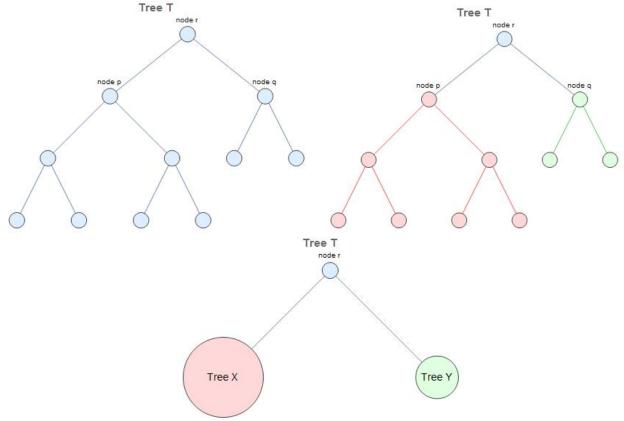
- a)  $\bigcirc$  We would let T be a full binary, heap property tree that consists of a root r joining two subtrees X and Y (each of which have the heap-property). We would then assume the Max-Root Proposition is true for both X and Y and argue that  $v(r) \le v(n)$  for all nodes n.
- **b)**  $\odot$  We would let T be a full binary, heap property tree that consists of a root r joining two subtrees X and Y (each of which have the heap-property). We would then assume the Max-Root Proposition is true for both X and Y and argue that  $v(r) \ge v(n)$  for all nodes n.
- c) We would suppose a full binary heap-property tree T has nodes and then conclude that all the node's have the same value.
- d)  $\bigcirc$  We would suppose the Max-Root Proposition is true for some  $k \in N$  and then show it is true for k+1.
- e) We would suppose the Max-Root Proposition is true for full binary trees with the heap property. This would then imply it is true for all full binary trees with the heap property.

#### Question 5

#### Your answer is CORRECT.

The main point of this quiz question (and the next one) is to continue our exploration of how "structural induction" can be used to prove the Max-Root Proposition:

Proposition (The Max-Root Proposition). If a full binary tree, T, has the heap property, then  $v(r) \ge v(n)$  for all nodes n. Here, as usual, r denotes the root of T and v(n) denotes the value (or "key") stored at node n.



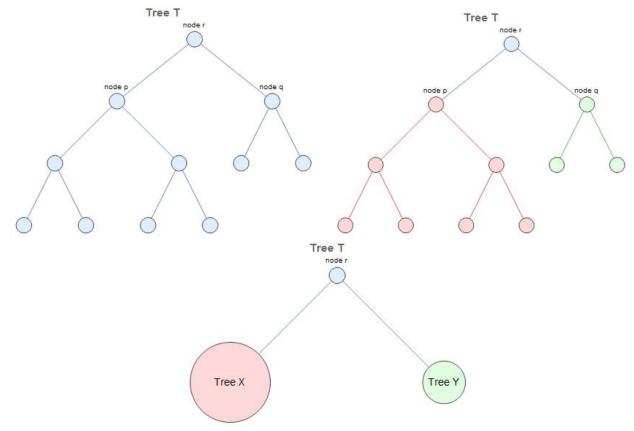
This question continues our work from the previous problem where, as part of our recursive step, we have assumed that a full binary heap-tree T consists of a root r joining two other full binary heap-trees, X and Y, each with respective roots p and q (see the images above for one example), and where the Max-Root Proposition is true for each. We now want to prove that  $v(r) \ge v(n)$  for all nodes n. As a step in this direction, we need to understand why it is only necessary to prove that  $v(r) \ge v(p)$  and  $v(r) \ge v(q)$ . Why do we need to only prove these two inequalities?

- a)  $\bigcirc$  If we prove that  $v(r) \ge v(p)$  and  $v(r) \ge v(q)$  it will follow that the Max-Root Proposition is true for X. This is due to our inductive hypothesis which assumes that  $v(n) \ge v(p)$  for all nodes n in X and that  $v(n) \ge v(q)$  for all nodes n in Y.
- **b)** If we prove that  $v(r) \ge v(p)$  and  $v(r) \ge v(q)$  it will follow that the Max-Root Proposition is true for X. This is due to our inductive hypothesis which assumes that  $v(p) \ge v(n)$  for all nodes n in X and that  $v(q) \ge v(n)$  for all nodes n in Y.
- c) We only need to prove these two inequalities because X is a full binary tree with more than three nodes.
- d) We only need to prove these inequalities because the picture does not show any values stored in nodes that are children of p and q.

# Question 6

# Your answer is CORRECT.

The main point of this quiz question is to continue our exploration of how "structural induction" can be used to prove the Max-Root Proposition: Proposition (The Max-Root Proposition). If a full binary tree, T, has the heap property, then  $v(r) \ge v(n)$  for all nodes n. Here, as usual, r denotes the root of T and v(n) denotes the value (or "key") stored at node n.



This question continues our work from the previous two problems where, as part of our recursive step, we have assumed that a full binary heap-tree T consists of a root r joining two other full binary heap-trees, X and Y, each with respective roots p and q (see the images above for one example), and where the Max-Root Proposition is true for each.. To complete the proof we now need to answer one question: why is it true that  $v(r) \ge v(p)$  and  $v(r) \ge v(q)$ ?

- a)  $\bigcirc$  This is true because, by assumption, v(r) = 0 and all other nodes are assigned negative values.
- b) These inequalities are not true (and so the Max-Root Proposition may be false)!
- c)  $\bigcirc$  This is true because, by assumption,  $v(r) = \infty$  and all other nodes have finite value.
- d) These inequalities are true because if they were false then T would not be a binary tree.
- e)  $\odot$  This is true because, by assumption, T has the heap property.

# Question 7

# Your answer is CORRECT.

Consider the set  $S \subseteq N \times N$  recursively defined by the following rules:

$$(0,0) \in S$$

$$(a, b) \in S \Rightarrow (a+3, b+2) \in S.$$

Which of the following is a true statement about S?

- a)  $\bigcirc$  (3, 2)  $\notin$  S.
- **b)** (0,0) is a "base element" of S.
- c)  $\bigcirc$  (9, 10) is a "base element" of S.
- **d)** |S| = 2
- e)  $\circ$  S = {(0,0)}

#### **Question 8**

# Your answer is CORRECT.

Consider the set  $S \subseteq N \times N$  recursively defined by the following rules:

$$(0,0) \in S$$

$$(a, b) \in S \Rightarrow (a + 29, b + 4) \in S.$$

Suppose we wish to prove a fact about the elements of S using Structural Induction (as we will happen in the next problem). From the options below, which could serve as the "base case" part of our proof?

- a)  $\bigcirc$  Verify that the proposition is true for the single element (29, 4).
- b) Find a contradiction.
- c)  $\bigcirc$  Verify that the proposition is true for the single element (0,0).
- d)  $\bigcirc$  Demonstrate the fact is true on the "base elements" of S -- in this case there are 2 such base elements, (29, 4) and (30, 8).
- e) There is no base case part of a Structural Induction proof.

#### **Question 9**

# Your answer is CORRECT.

Consider the set  $S \subseteq N \times N$  recursively defined by the following rules:

$$(0,0) \in S$$

$$(a, b) \in S \Rightarrow (a + 3, b + 2) \in S.$$

Suppose we wish to use Structural Induction to prove the following proposition about the elements of S:

Proposition. If  $(x, y) \in S$ , then x + y is a multiple of 5.

Which of the following most accurately describes the "inductive" step of our proof?

- a)  $\bigcirc$  Prove That if a + b is a multiple of 5, then (a + 1) + (b + 1) is also a multiple of 5.
- **b)** O Prove that  $((a, b) \in S \land (\exists c \in Z, 5 = c \cdot (a + b))) \Rightarrow (\exists d \in Z, d \cdot (a + 3 + b + 2) = 5)$
- c)  $\bigcirc$  Suppose that the proposition is true for  $(a,b) \in S$  and then prove it is true for the next element of S.
- d) Prove that 0 + 0 = 0 is a multiple of 5.
- e) © Prove that  $((a, b) \in S \land (\exists c \in Z, a+b=c \cdot 5)) \Rightarrow (\exists d \in Z, a+3+b+2=d \cdot 5)$

## Question 10

## Your answer is CORRECT.

Consider the set  $S \subseteq N \times N$  recursively defined by the following rules:

$$(0,0) \in S$$

$$(a, b) \in S \implies (a + 23, b + 6) \in S$$
.

Is the following proposition true?

Proposition. If  $(x, y) \in S$ , then x + y is a multiple of 29.

- a)  $\odot$  Yes! It is true for the base element (0,0) and if we assume  $(a,b) \in S$  satisfies  $a+b=c\cdot 29$  for some c, then  $a+23+b+6=(a+b)+(23+6)=c\cdot 29+29=c\cdot 29+1\cdot 29+(c+1)\cdot 29$ .
- **b)**  $\bigcirc$  No, it is not true because even though the proposition is true for the Base Element(s), the inductive step fails. It does *not* follow that if  $(a, b) \in S$  satisfies 29|(a+b)| then 29|(a+23, b+6)|.