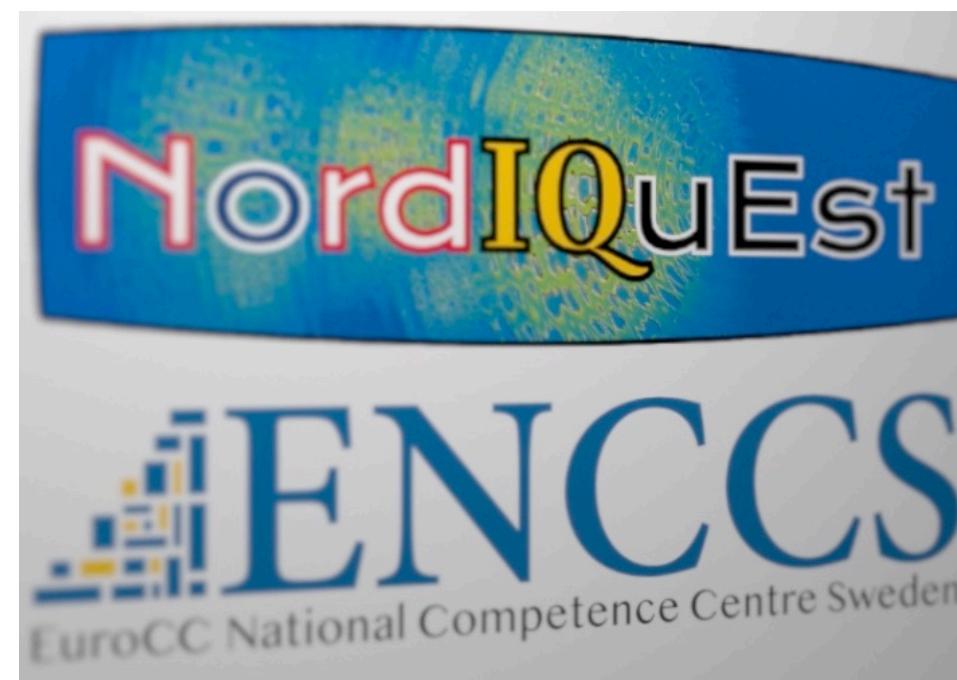
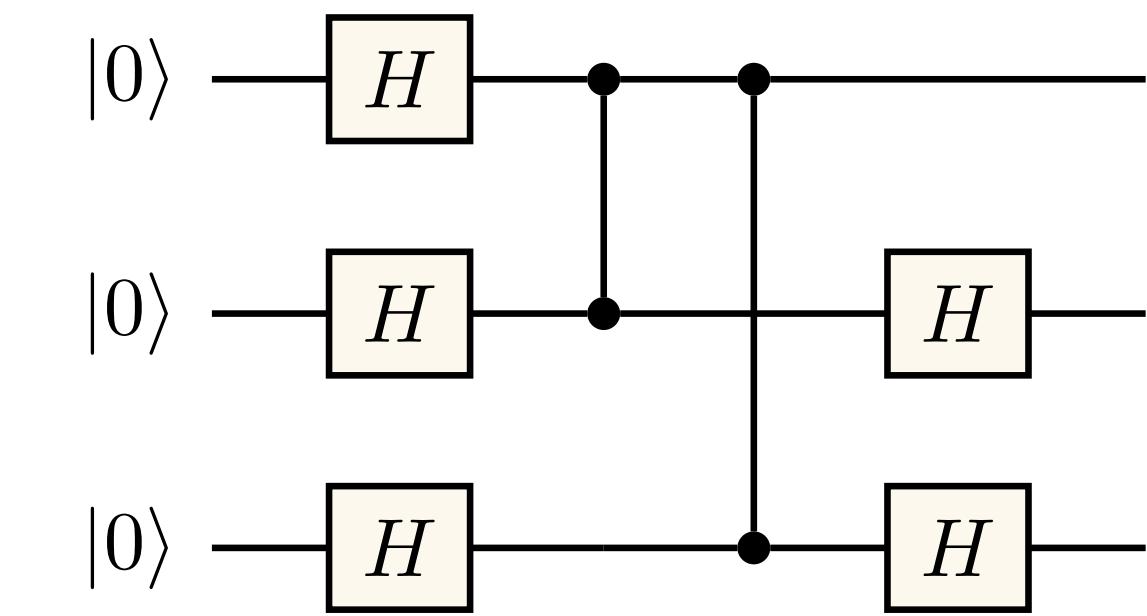


## Quantum states, qubits, logic gates, and algorithms



Anton Frisk Kockum

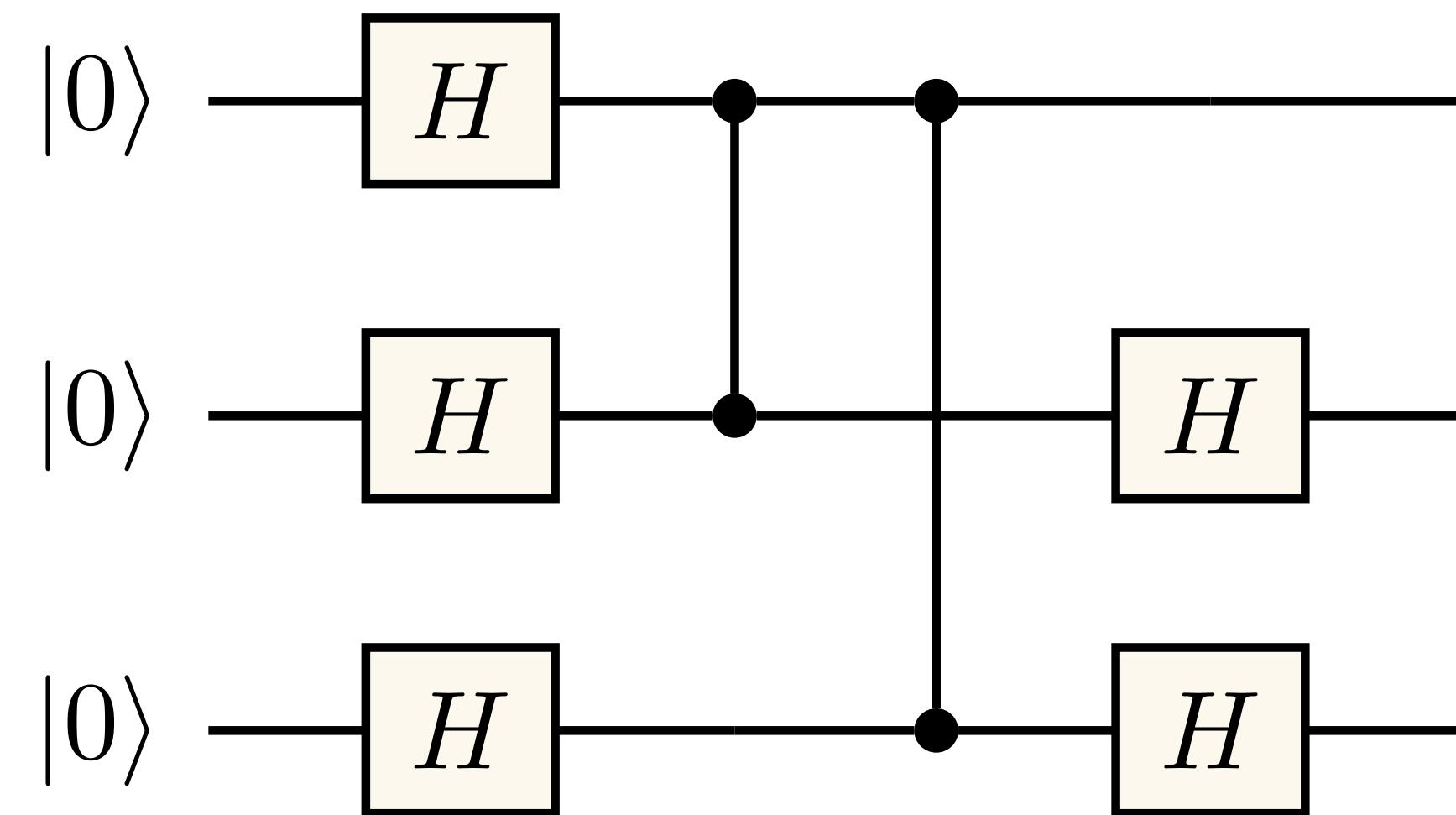
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# Outline

- Components of a quantum circuit
- Quantum bits
- Single-qubit gates
- Multi-qubit gates
- Universal gate sets
- The Solovay-Kitaev theorem
- Quantum algorithms and compilation
- Summary

# Components of a quantum circuit



# Quantum bits

A quantum bit (qubit) can be in a superposition of states

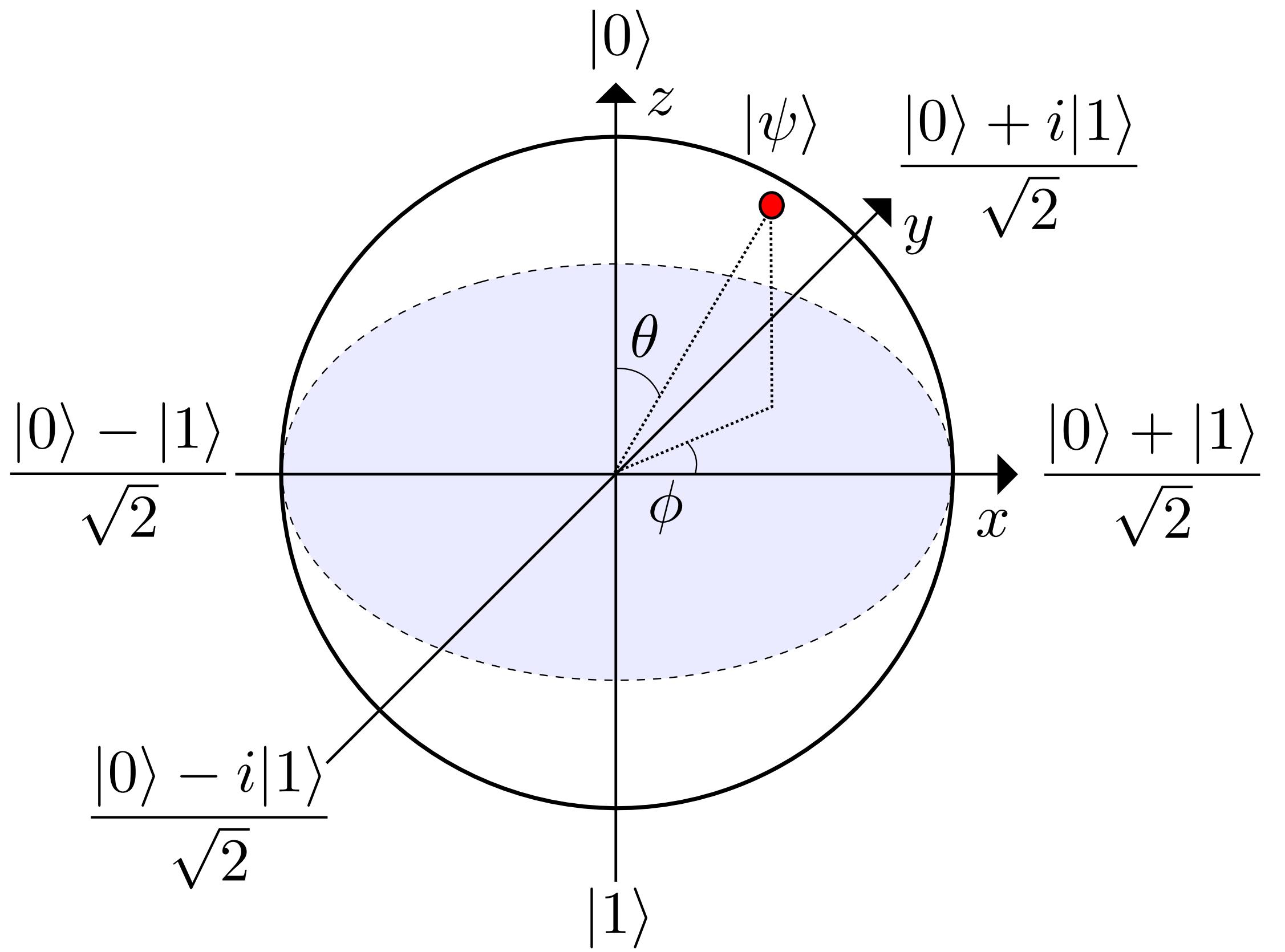
$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle = \alpha \begin{pmatrix} 1 \\ 0 \end{pmatrix} + \beta \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} \alpha \\ \beta \end{pmatrix}$$

$$\alpha, \beta \in \mathbb{C} \quad |\alpha|^2 + |\beta|^2 = 1$$

Can be visualised on the Bloch sphere

$$|\psi\rangle = \cos \frac{\theta}{2}|0\rangle + e^{i\phi} \sin \frac{\theta}{2}|1\rangle$$

Measurements give either 0 or 1 with probabilities  $|\alpha|^2$  and  $|\beta|^2$



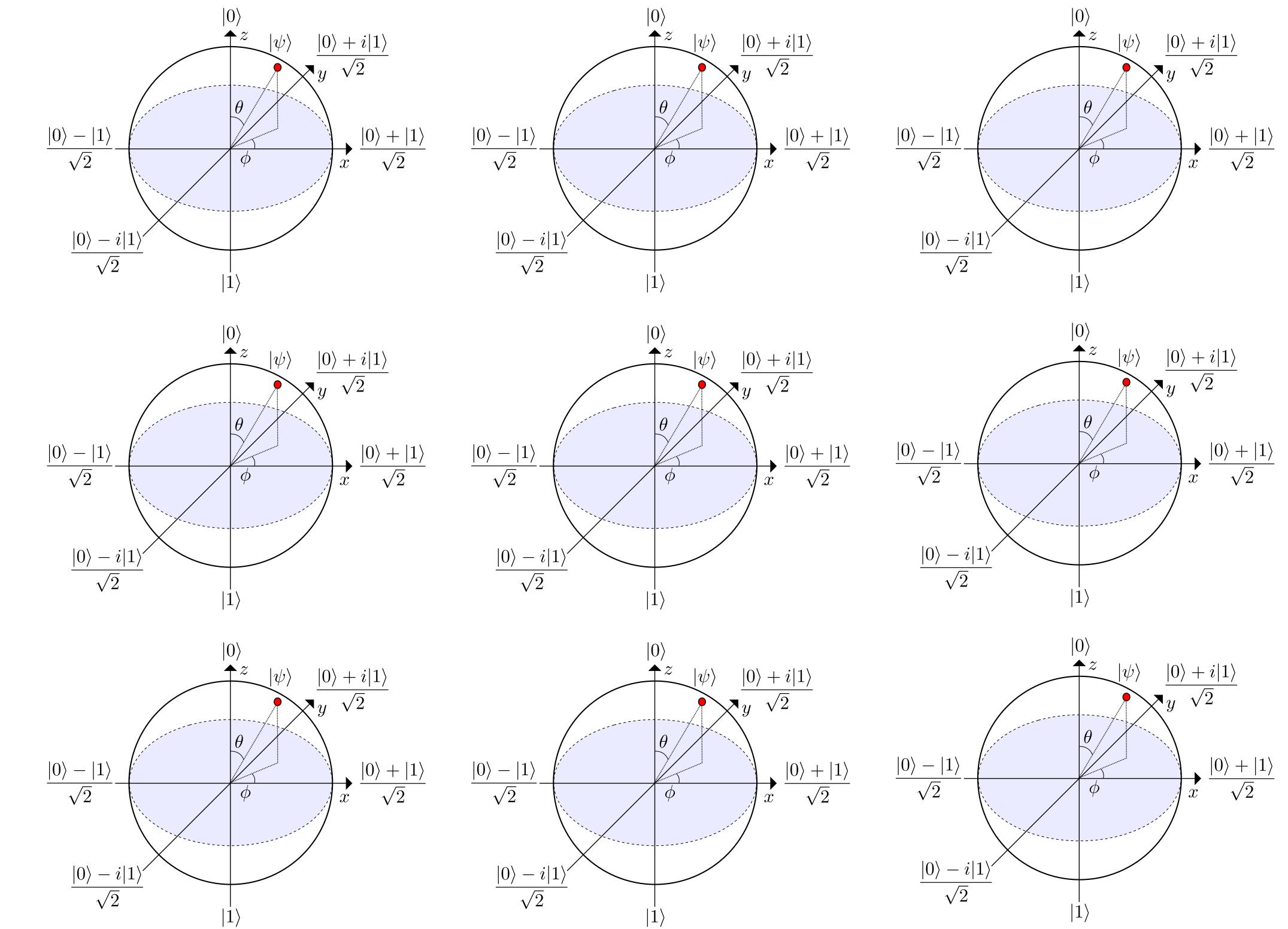
# Multi-qubit states

$N$  qubits can be in a superposition of  $2^N$  states

$|000\dots 00\rangle, |100\dots 00\rangle, |010\dots 00\rangle, \dots, |111\dots 10\rangle, |111\dots 11\rangle$

$N$  classical bits can only be in one state at a time

Storing all the information about a quantum state can require  $\gg N$  classical bits



# Single-qubit gates

Operations changing the state of a qubit must preserve the norm

$$U|\psi\rangle = U(\alpha|0\rangle + \beta|1\rangle) = |\psi'\rangle = \alpha'|0\rangle + \beta'|1\rangle$$

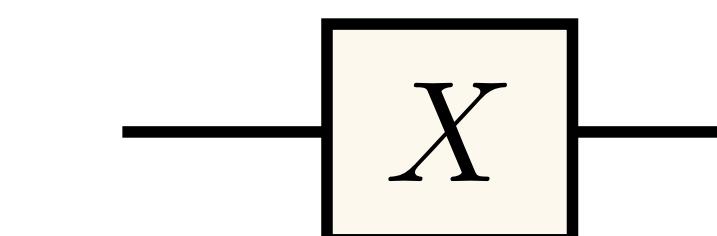
$$|\alpha|^2 + |\beta|^2 = 1 = |\alpha'|^2 + |\beta'|^2$$

They are 2x2 unitary matrices, e.g., the Pauli matrices

$$X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix},$$

$$Y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix},$$

$$Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

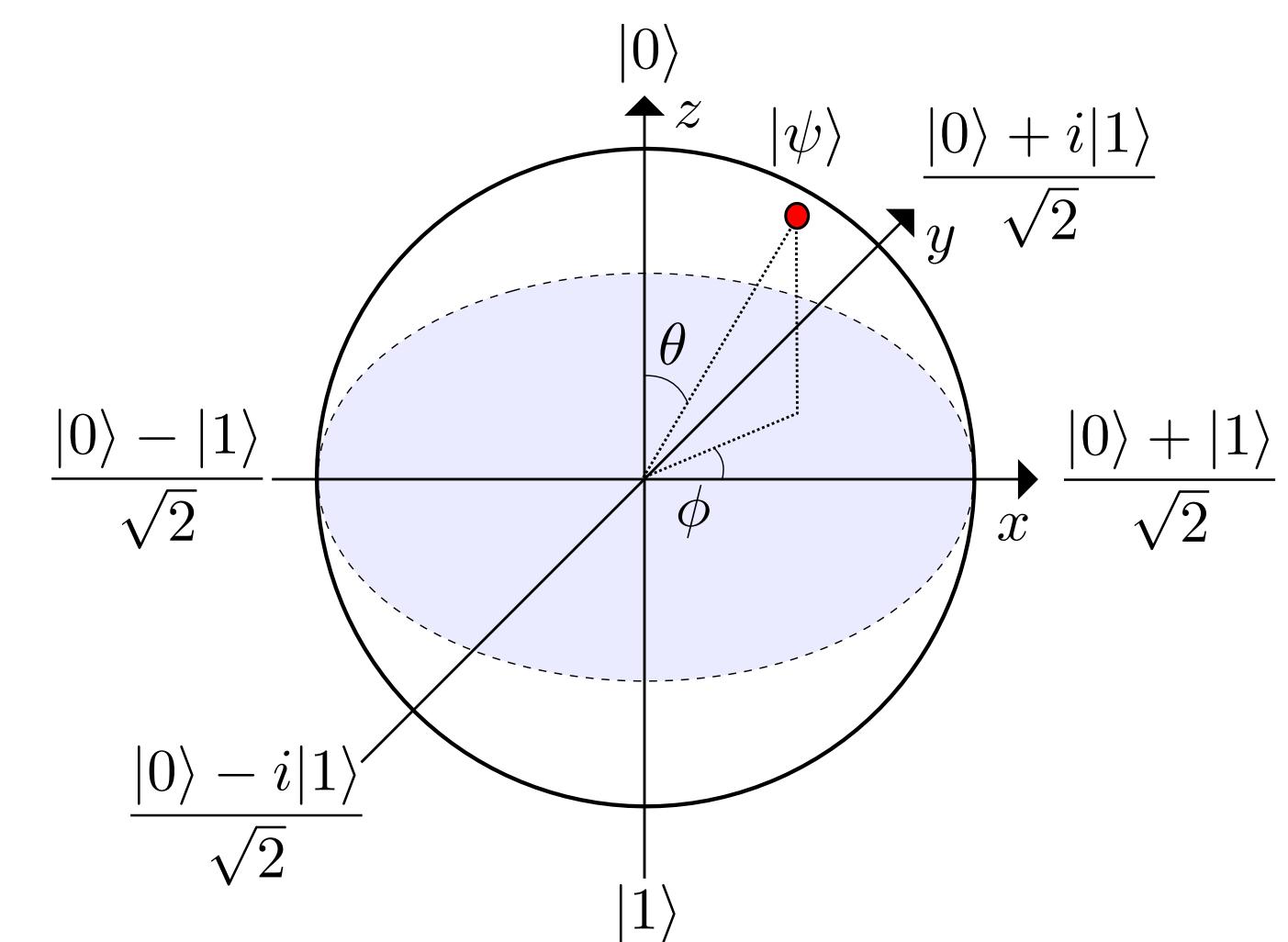


is the quantum NOT gate

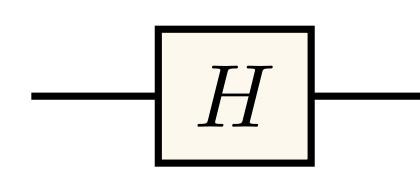
$$X(\alpha|0\rangle + \beta|1\rangle) = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} \beta \\ \alpha \end{pmatrix} = \beta|0\rangle + \alpha|1\rangle$$

Rotations around different axes of the Bloch sphere

$$\begin{aligned} R_x(\theta) &= \exp(-i\theta X/2) \\ &= \cos(\theta/2)I - i \sin(\theta/2)X \\ &= \begin{pmatrix} \cos(\theta/2) & -i \sin(\theta/2) \\ -i \sin(\theta/2) & \cos(\theta/2) \end{pmatrix} \end{aligned}$$

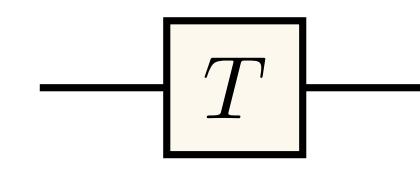


# More single-qubit gates



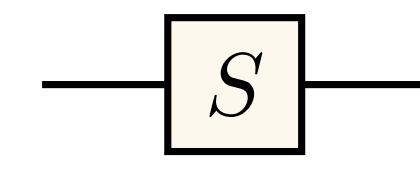
The Hadamard gate

$$H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} = \frac{X + Z}{\sqrt{2}}$$



The T gate ( $\pi/8$  gate)

$$T = \begin{pmatrix} 1 & 0 \\ 0 & \exp(i\pi/4) \end{pmatrix} = \exp(i\pi/8) \begin{pmatrix} \exp(-i\pi/8) & 0 \\ 0 & \exp(i\pi/8) \end{pmatrix}$$



The phase (or S, or P) gate

$$S = \begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix} = T^2$$

# Two-qubit gates

**Controlled-NOT**

$$\text{CNOT} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{array}{c} |00\rangle \quad |01\rangle \quad |10\rangle \quad |11\rangle \\ \hline |00\rangle \\ |01\rangle \\ |10\rangle \\ |11\rangle \end{array} \quad \begin{array}{c} \text{---} \bullet \text{---} \\ \text{---} \oplus \text{---} \end{array}$$

**Controlled-Z**

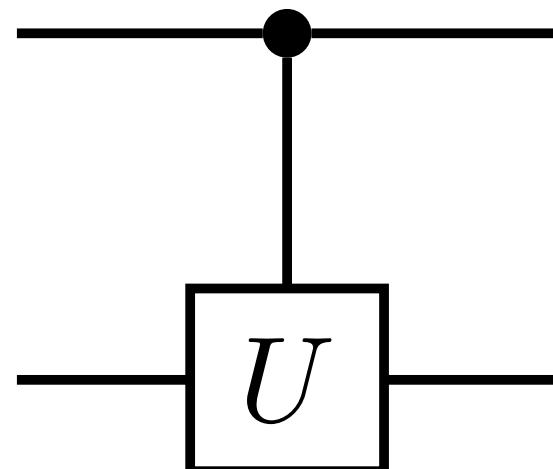
$$\text{CZ} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix} \quad \begin{array}{c} \text{---} \bullet \text{---} \\ \text{---} \bullet \text{---} \\ \text{---} \bullet \text{---} \end{array}$$

**SWAP**

$$\text{SWAP} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \quad \begin{array}{c} \text{---} \times \text{---} \\ \text{---} \times \text{---} \end{array}$$

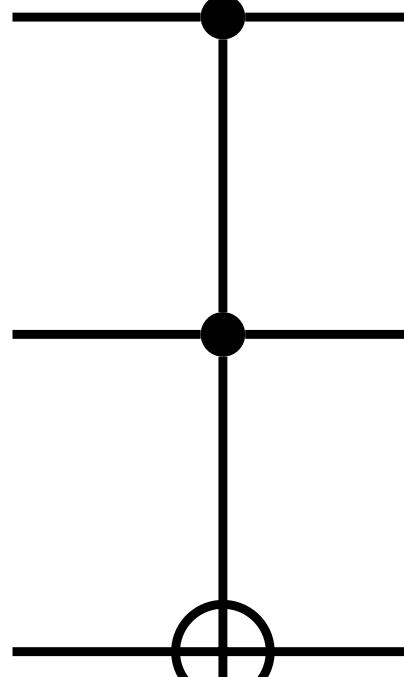
**Controlled unitary**

$$\begin{pmatrix} I_2 & 0_2 \\ 0_2 & U \end{pmatrix}$$

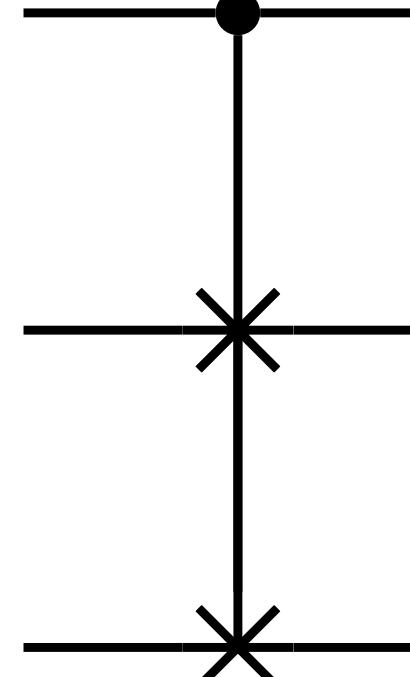


# Three-qubit gates

Controlled-controlled-NOT

$$\text{Toffoli} = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \end{pmatrix}$$


Controlled-SWAP

$$\text{Fredkin} = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \end{pmatrix}$$


Multi-qubit gates can be decomposed into sequences of single- and two-qubit gates

# Universal gate sets

## Classical computing

A gate set is universal if it enables expressing any Boolean function on any number of bits

The NAND gate is universal  
 $\{\text{AND}, \text{OR}\}$  is not a universal gate set

## Quantum computing

A gate set is universal if the gates therein can approximate any unitary transformation on any number of qubits to any precision

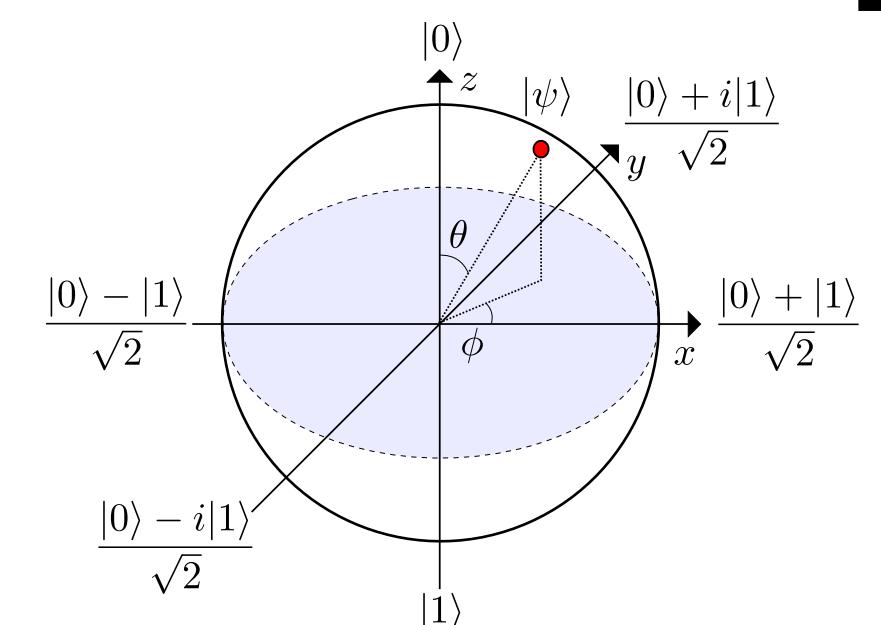
What are requirements for a universal set of quantum gates?

# Universal quantum gate sets

## Failure modes

## Universal gate sets

- Almost anything else than H in {H, CNOT, S}
- Almost any two-qubit gate on its own
- In practice: many single-qubit gates + one or two two-qubit gates



# Quantum versus classical computing

Are there problems that a classical computer  
can't solve but a quantum computer can?



The classical computer, given enough resources, can store all the complex amplitudes specifying the state of the quantum computer and simulate all the gates in a quantum circuit

For the quantum computer to be faster, one thing to worry about is whether the universal gate set can represent the desired algorithm with enough precision without requiring too long circuits

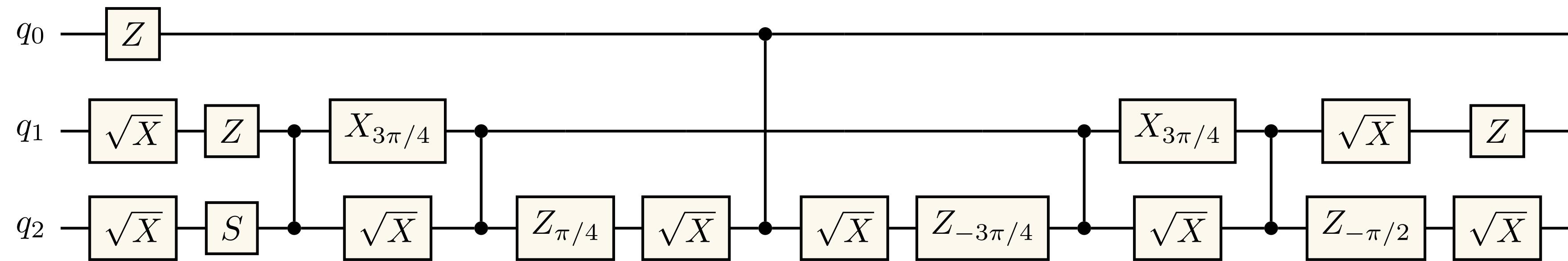
# The Solovay-Kitaev theorem

Let  $G$  be a finite subset of  $SU(2)$  and  $U \in SU(2)$ . If the group generated by  $G$  is dense in  $SU(2)$ , then for any  $\varepsilon > 0$  it is possible to approximate  $U$  to precision  $\varepsilon$  using  $\mathcal{O}\left(\log^4\left[\frac{1}{\varepsilon}\right]\right)$  gates from  $G$ .

For an  $N$ -qubit unitary, at most  $\mathcal{O}\left(4^N \log^4\left[\frac{1}{\varepsilon}\right]\right)$  gates suffice

Precision is thus not a problem in practice for available universal gate sets

# Quantum algorithms and compilation



Quantum algorithms are sequences of gates acting on quantum states

Compilation steps

# Summary

- Qubits can be in superposition states; exponentially many classical bits are required to describe many qubits
- Quantum algorithms are implemented by applying a sequence of single- and two-qubit gates (unitary matrices) to the qubits (states represented as vectors)
- Quantum algorithms need to be compiled to fit on the quantum hardware; the Solovay-Kitaev theorem tells us that universal gate sets can achieve this without prohibitive overhead to ensure precision

