

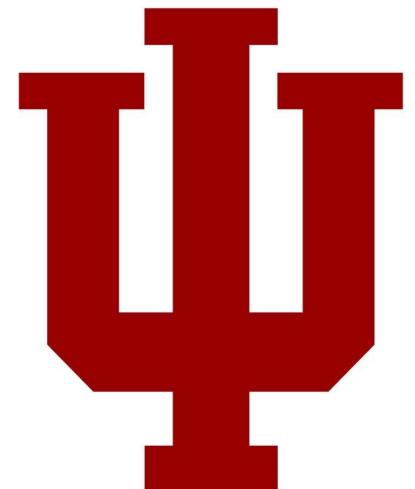
PL: due sooner  
next year

# 06 Physically Unclonable Functions (PUFs)

Engr 399/599: Hardware Security

Andrew Lukefahr

*Indiana University*



Adapted from: Mark Tehranipoor of University of Florida

Course Website

enr599.github.io

Room  
Pass code:  
2-3-5

Write that down!

# Project 1: Hardware Trojan

- Goal: “Corrupt” a working DES implementation
- We give you DES in HW
- You need to:
  - Deploy DES
  - Corrupt DES

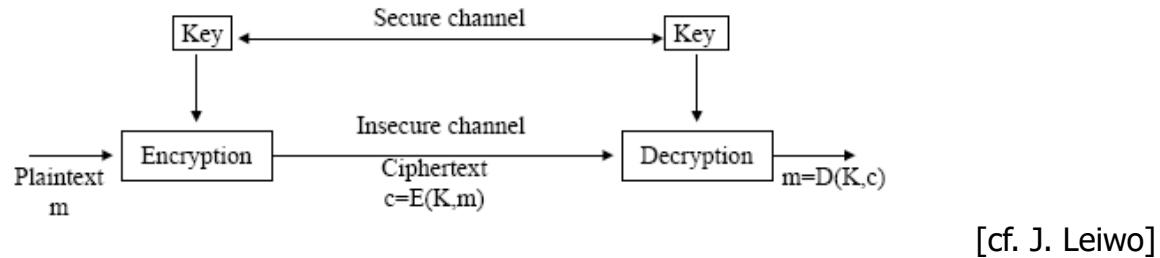
# Group Assignments

- Chris Sozio
- Will Fleming
- Clare Barnes
- Austin Parkes
- Max Harms
- Michael Foster
- Trey Peterson
- Yifan Zhang
- Jack Ruocco

Due next Wednesday (?)

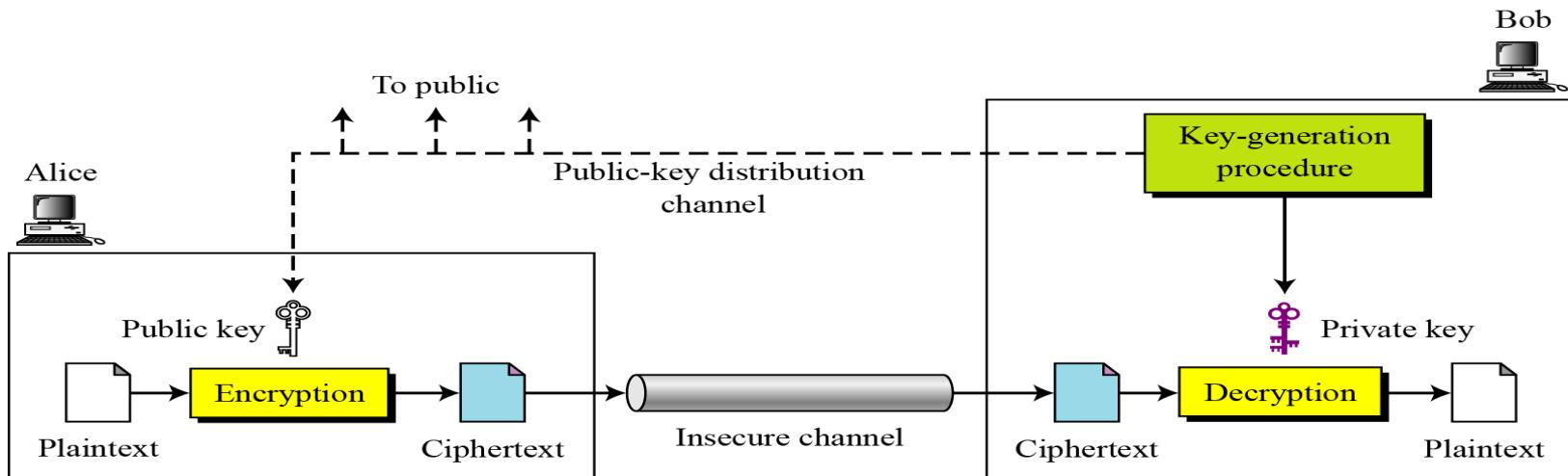
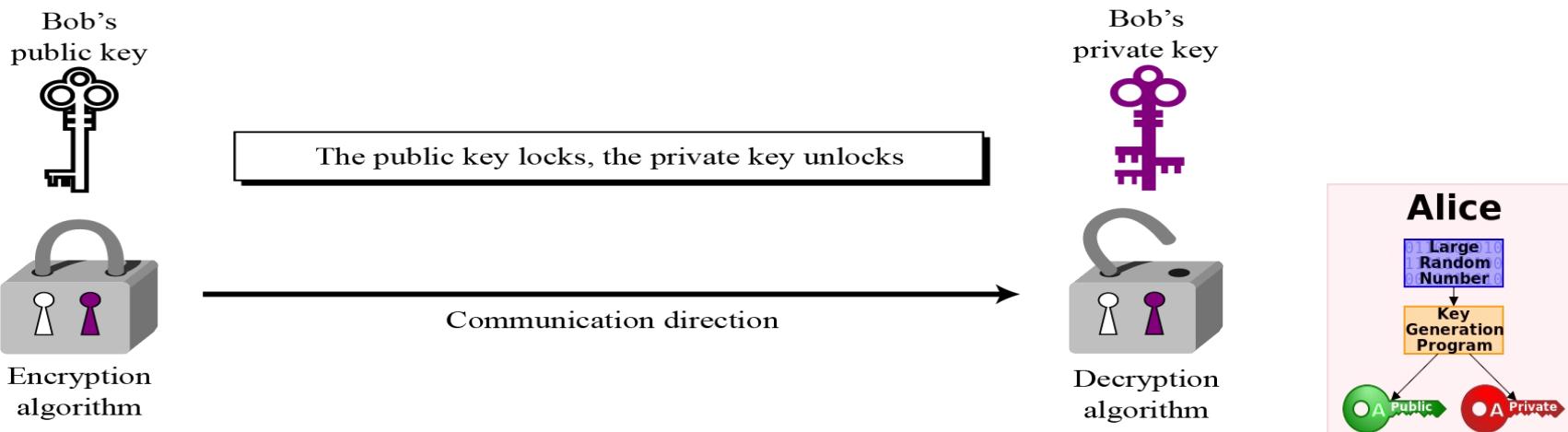
# Symmetric and Asymmetric Cryptosystems (1)

- Symmetric encryption = **secret key encryption**
  - $K_E = K_D$  — called a **secret key** or a **private key**
  - Only sender S and receiver R know the key



- As long as the key remains secret, it also provides **authentication** (= proof of sender's identity)

# Asymmetric Security

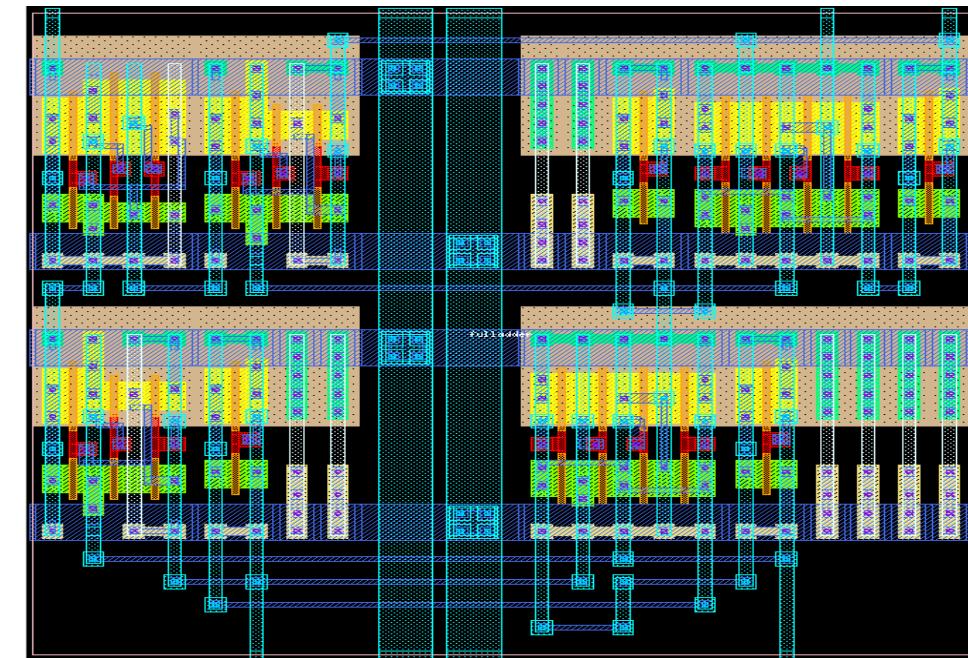
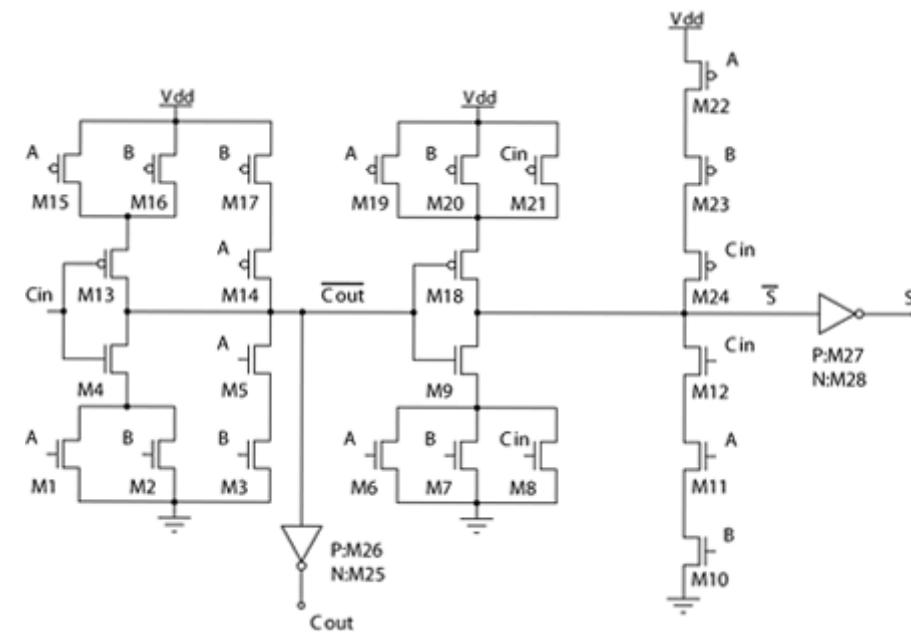
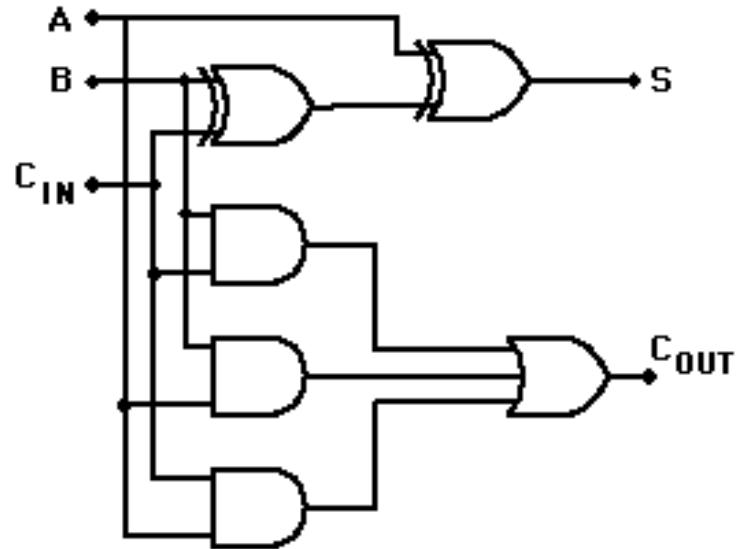


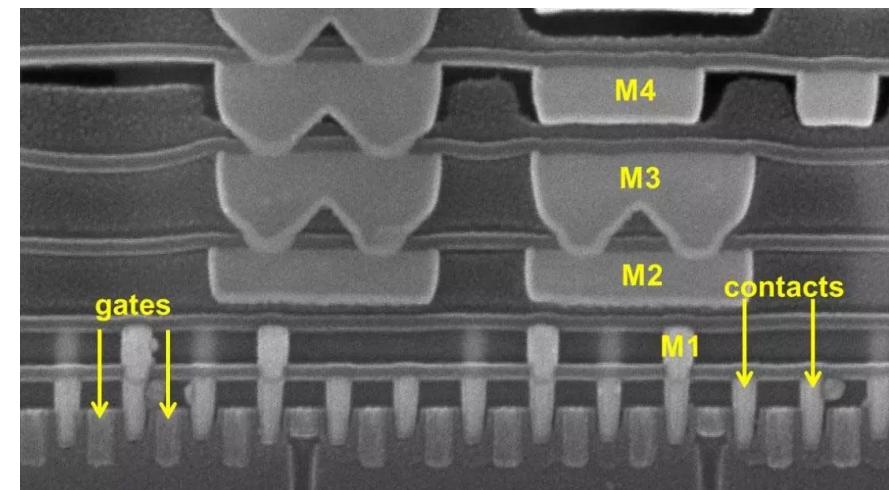
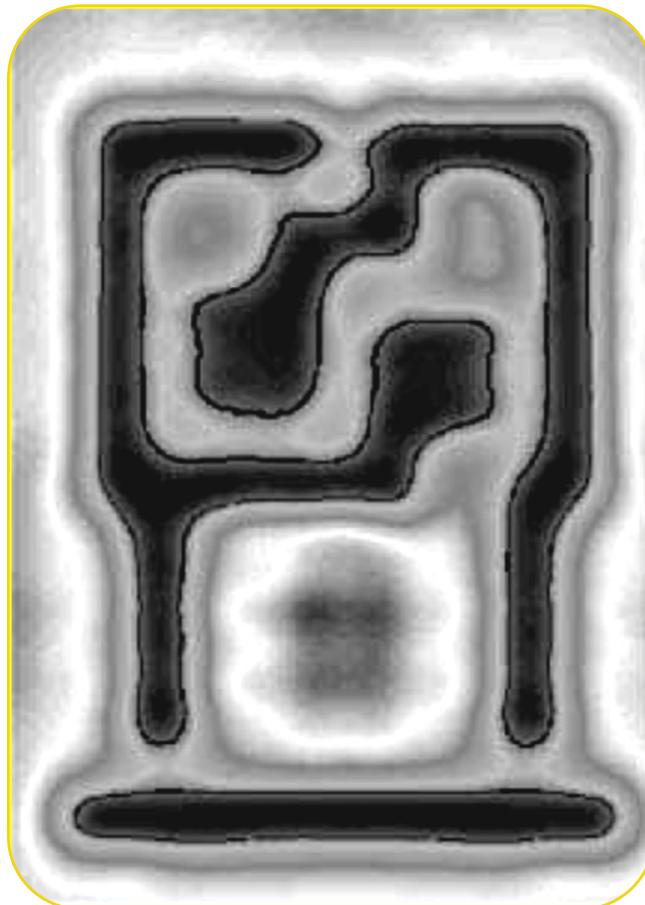
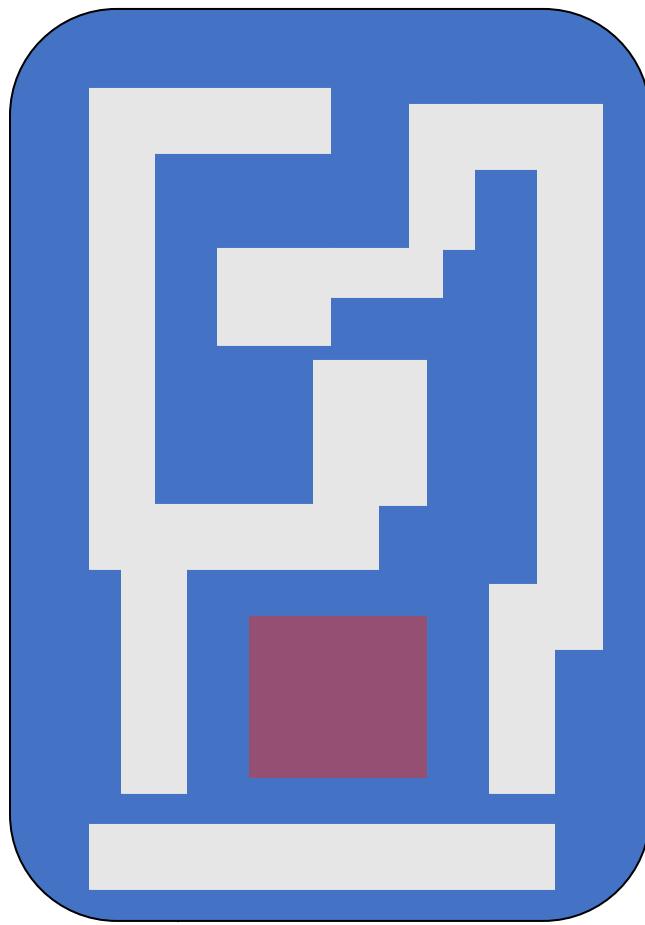
```

module fulladder (a,b,cin,sum,cout);
    input a,b,cin;
    output sum,cout;

    reg sum,cout;
    always @*
    begin
        sum <= a ^ b ^ cin;
        cout <= (a & b) | (a & cin) | (b & cin);
    end
endmodule

```





# Security

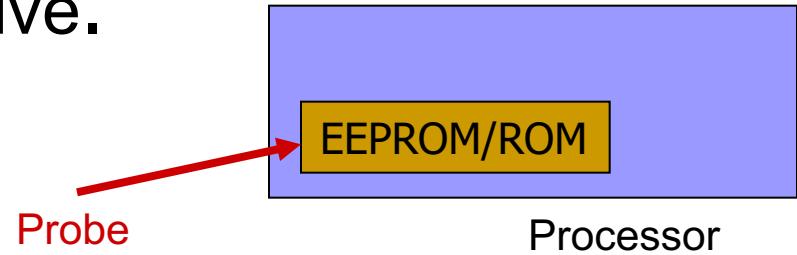
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- Asymmetry b/w the information (secret)
- **One-way functions**
  - Easy to evaluate in one direction but hard to reverse in the other
  - E.g., multiplying large prime number as opposed to factoring them
- **One-way hash functions**
  - Maps a variable length input to a fixed length output
  - **Avalanche property:** changing one bit in the input alters nearly half of the output bits
  - Pre-image resistant, collision resistant
  - Usage: digital signature, secured password storage, file identification, and message authentication code

# Problem

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Storing **digital** information in a device in a way that is resistant to **physical attacks** is difficult and expensive.



- Adversaries can physically extract secret keys from EEPROM while processor is off
- Trusted party must embed and test secret keys in a secure location
- EEPROM adds additional complexity to manufacturing

# IBM 4758

## Problem:

**Storing digital information in a device in a way that is resistant to physical attack is difficult and expensive.**



### **IBM 4758**

Tamper-proof package containing a secure processor which has a secret key and memory

Tens of sensors, resistance, temperature, voltage, etc.

Continually battery-powered

~ \$3000 for a 99 MHz processor and 128MB of memory

# Challenges of algorithmic (mathematical) one-way functions

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## ■ Technological

- ❑ Massive number of parallel devices broke DES
- ❑ Reverse-engineering of secure processors

## ■ Fundamental

- ❑ There is no proof that attacks do not exist
- ❑ E.g., quantum computers could factor two large prime numbers in polynomial time

## ■ Practical

- ❑ Embedded systems applications

# Solution -- POWF

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- Use the chaotic physical structures that are hard to model instead of mathematical one-way functions!
- Physical One Way Functions (POWF)
  - Inexpensive to fabricate
  - Prohibitively difficult to duplicate
  - No compact mathematical representation
  - Intrinsically tamper-resistant

# Feature: Process Variation

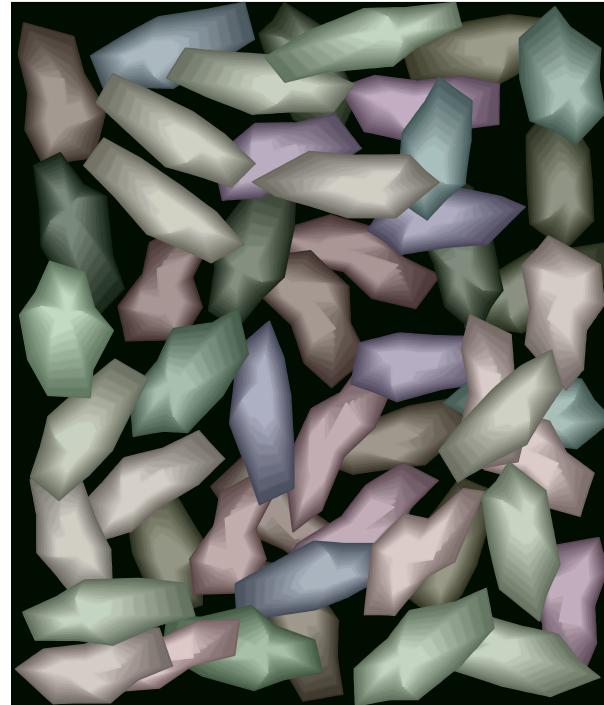
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- Do we expect process variation (length, widths, oxide thickness) in circuit and system?
    - Impact circuit performance
    - Functional failure
    - Major obstacle to the continued scaling of integrated-circuit technology in the sub-45 nm regime
  - Process variations can be turned into a feature rather than a problem?
    - Each IC has unique properties
-

# Solution

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**Extract key information from a complex physical system.**



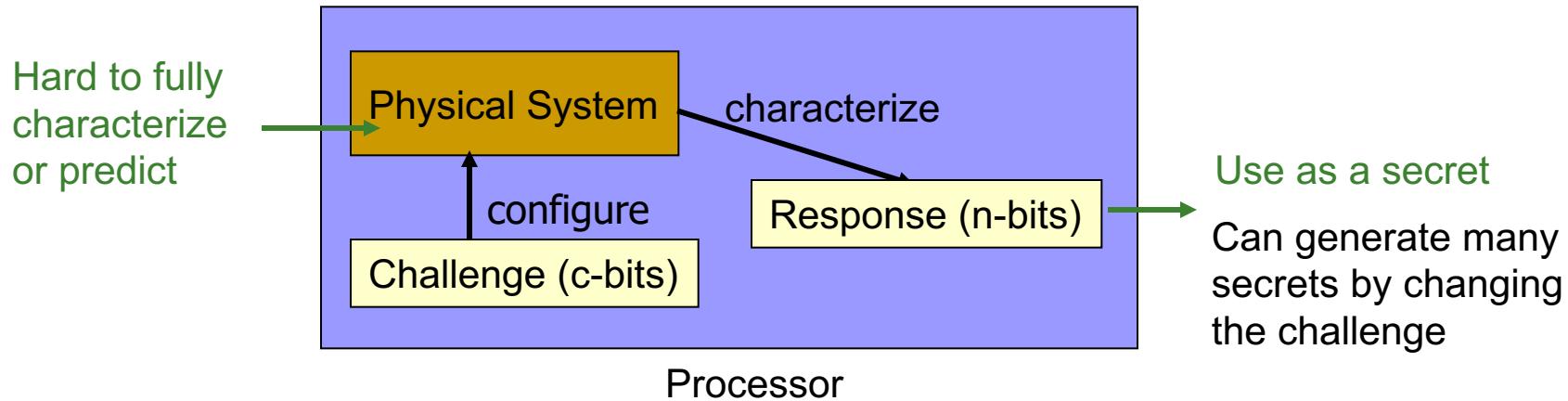
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Devadas, et. al, DAC02

# Physical Unclonable/Random Functions (PUFs)

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- Generate keys from a complex physical system



- Security Advantage
  - Keys are generated on demand → No non-volatile secrets
  - No need to program the secret
  - Can generate multiple master keys
- What can be hard to predict, but easy to measure?

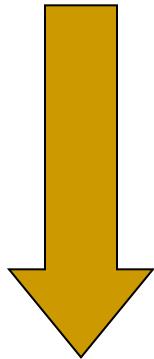
# Definition

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A Physical Random Function or **Physical Unclonable Function (PUF)** is a function that is:

- ❑ Based on a physical system
- ❑ Easy to evaluate (using the physical system)
- ❑ Its output looks like a random function
- ❑ Unpredictable even for an attacker with physical access

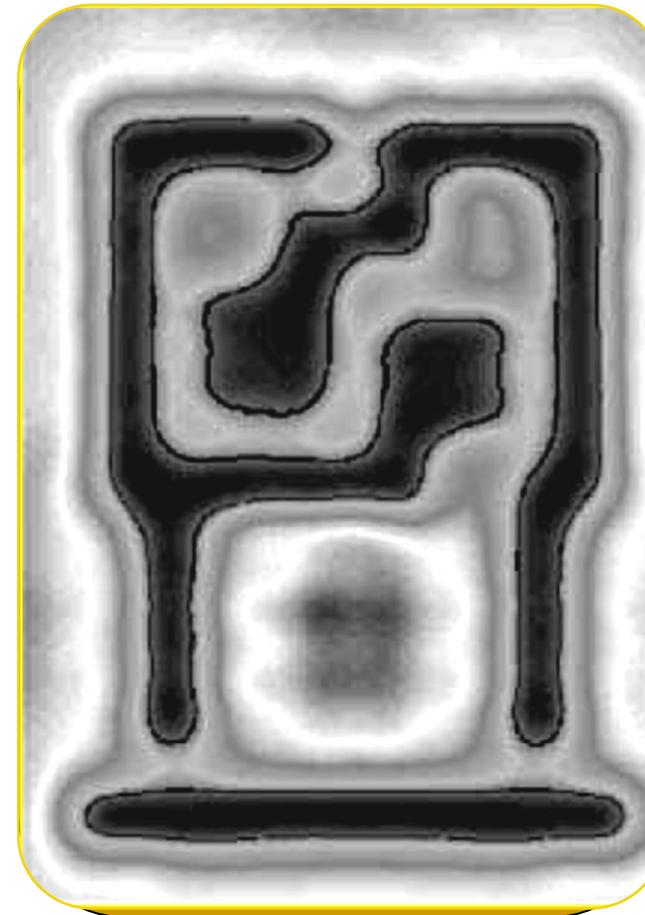
## Sub-Wavelength WYSINWYG



**What You See Is Not What You Get**

Process variations

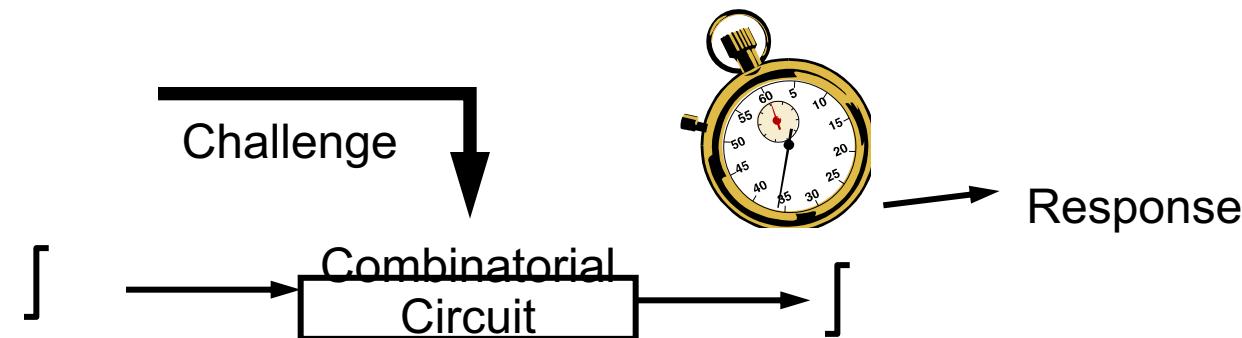
No two transistors have the same parameters



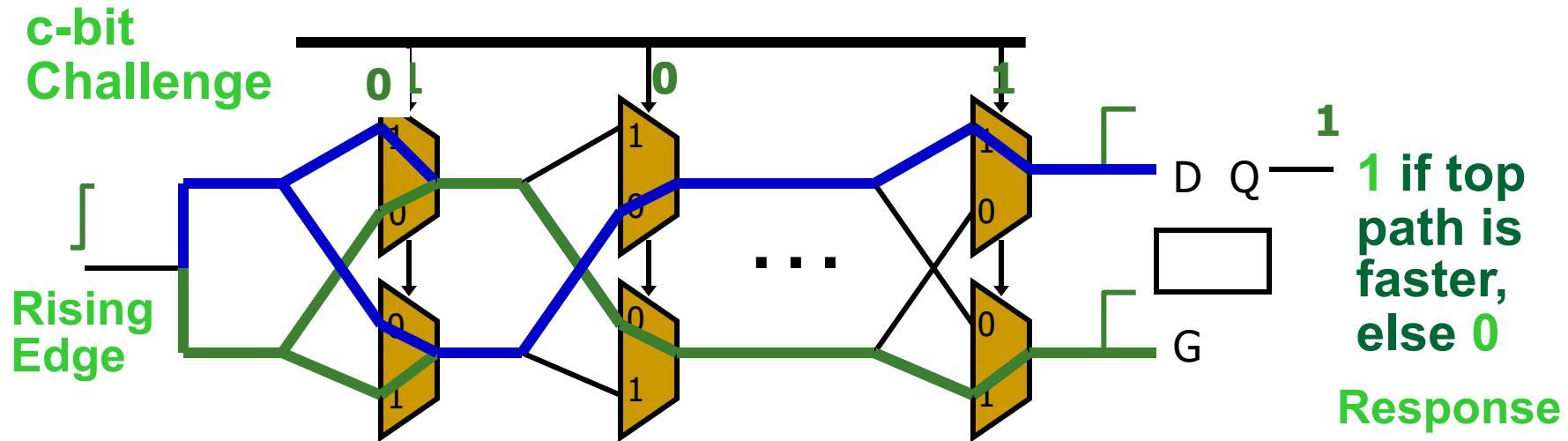
# Silicon PUF – Proof of Concept

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- Because of process variations, no two Integrated Circuits are identical
- Experiments in which *identical circuits with identical layouts* were placed on different FPGAs show that path delays vary enough across ICs to use them for identification.



# A Candidate: Silicon PUF



- Compare two paths with an **identical delay** in design
  - Random process variation determines which path is faster
  - An arbiter outputs 1-bit digital response
- **Path delays in an IC are statistically distributed due to random manufacturing variations**

# Experiments

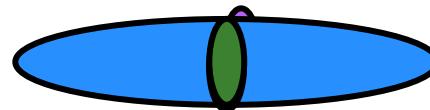
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- Fabricated candidate PUF on multiple ICs, 0.18um TSMC
- Apply 100 random challenges and observe responses

**100 bits of response**

**Distance between Chip X and Y  
responses = 24 bits**

**At 70C measurement  
noise for chip X = 2**



**Measurement noise for Chip X = 0.5**



# Measurement Attacks and Software Attacks

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Can an adversary create a *software clone* of a given PUF chip?

**Distance between Chip X and Y  
responses = 24**

**At 70C measurement  
noise for chip X = 2**

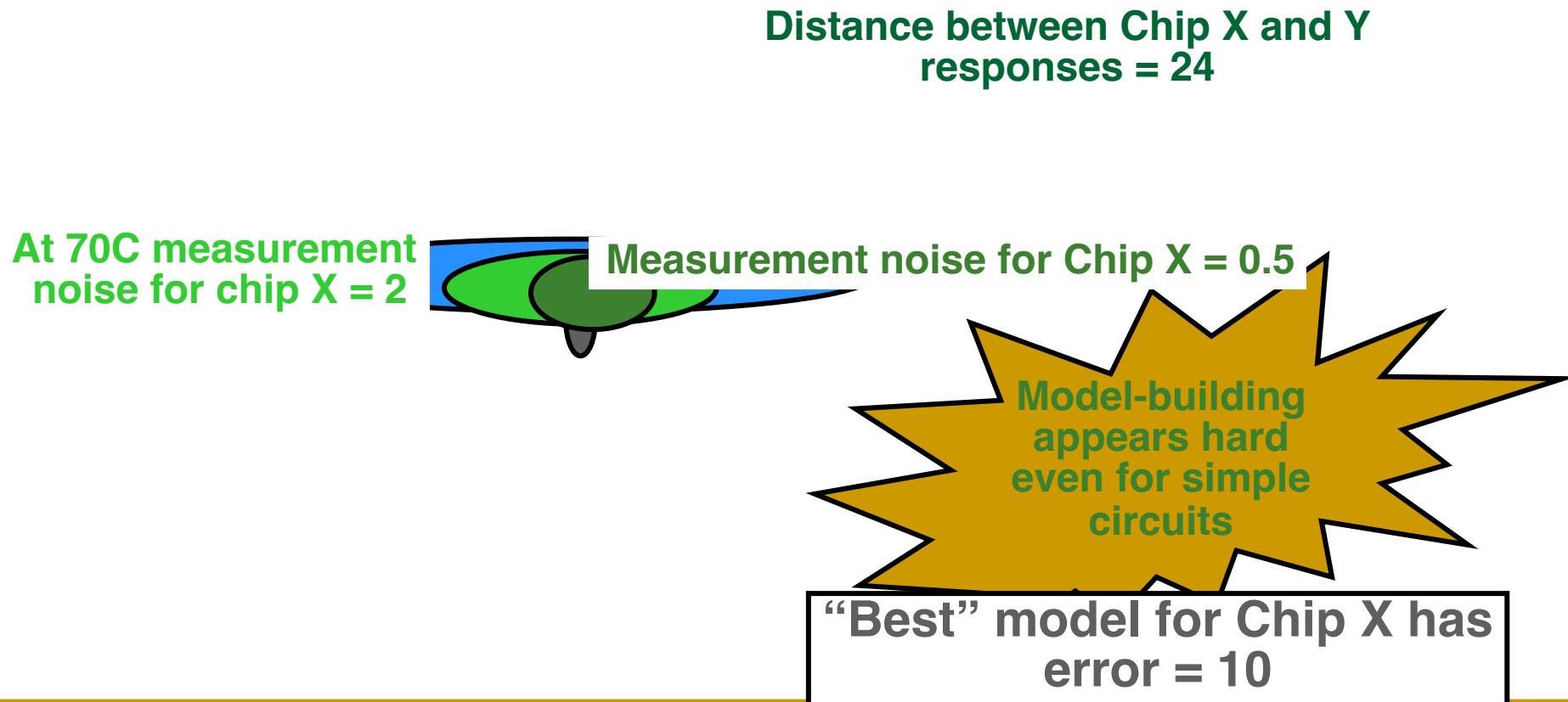


**Measurement noise for Chip X = 0.5**

# Measurement Attacks and Software Attacks

---

Can an adversary create a *software clone* of a given PUF chip?



# Physical Attacks

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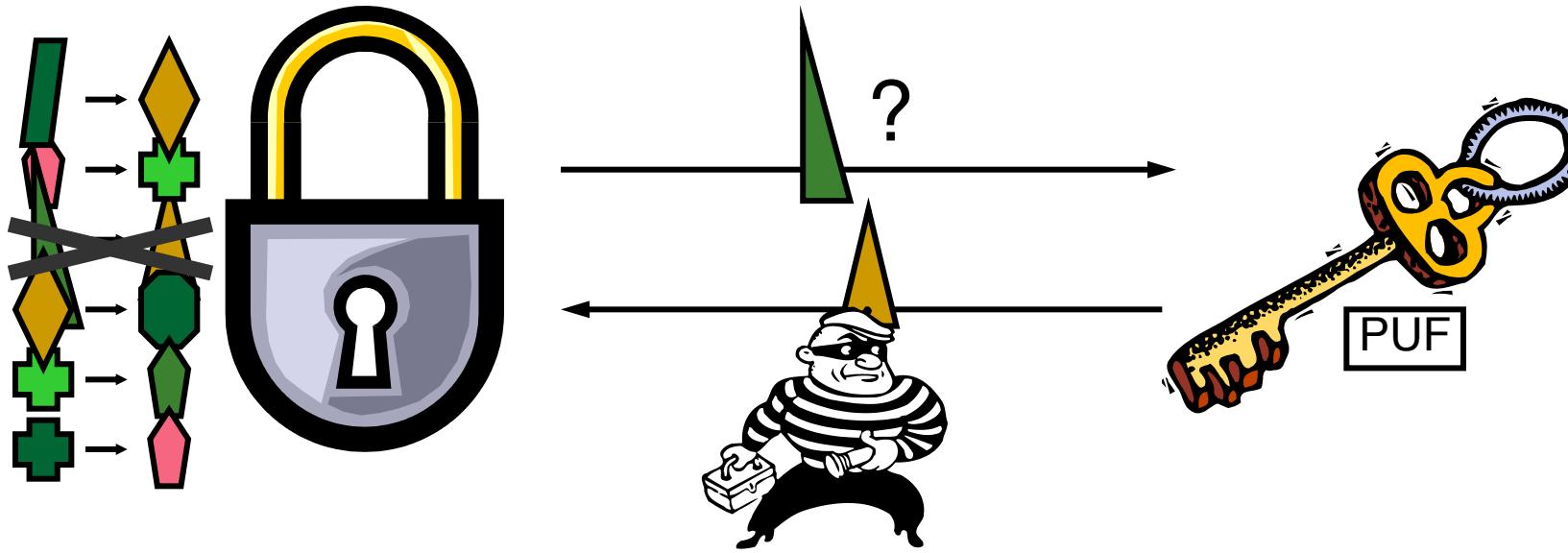
- Make PUF delays depend on overlaid metal layers and package
- Invasive attack (e.g., package removal) changes PUF delays and destroys PUF
- Non-invasive attacks are still possible
  - To find wire delays one needs to find precise relative timing of transient signals as opposed to looking for 0's and 1's
  - Wire delay is not a number but a function of challenge bits and adjacent wire voltages and capacitances

# Using a PUF as an Unclonable Key

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A Silicon PUF can be used as an unclonable key.

- The lock has a database of challenge-response pairs.
- To open the lock, the key has to show that it knows the response to one or more challenges.



# Applications

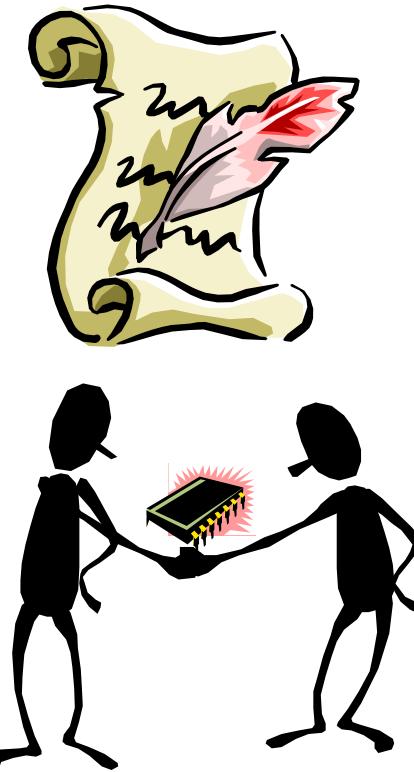
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## ■ Anonymous Computation

Alice wants to run computations on Bob's computer, and wants to make sure that she is getting correct results. A certificate is returned with her results to show that they were correctly executed.

## ■ Software Licensing

Alice wants to sell Bob a program which will only run on Bob's chip (identified by a PUF). The program is copy-protected so it will not run on any other chip.



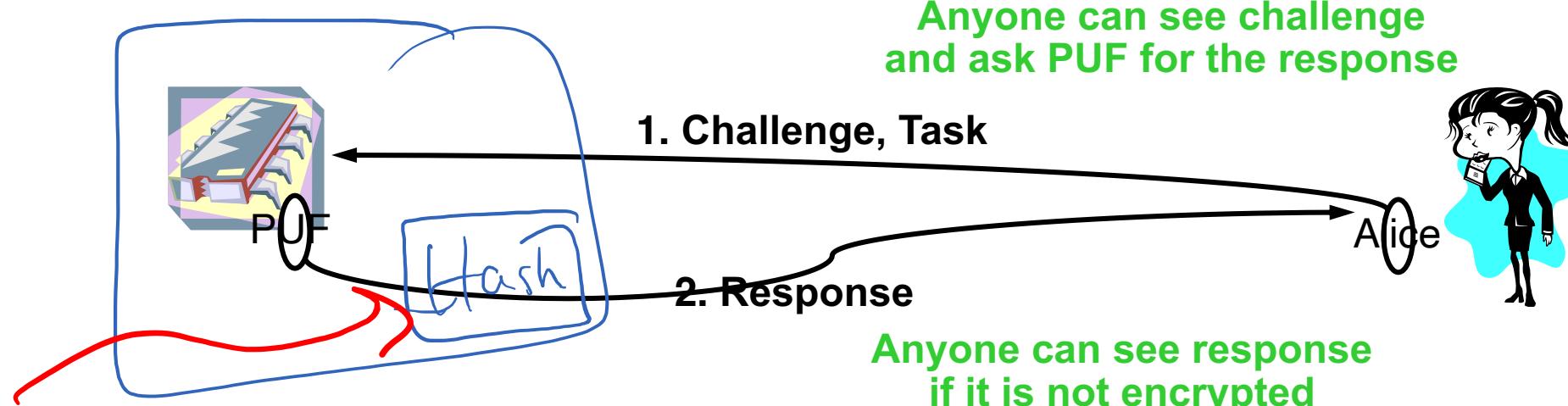
**We can enable the above applications by trusting only a single-chip processor that contains a silicon PUF.**

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# Sharing a Secret with a Silicon PUF

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Suppose Alice wishes to share a secret with the silicon PUF  
She has a challenge response pair that no one else knows,  
which can authenticate the PUF  
She asks the PUF for the response to a challenge



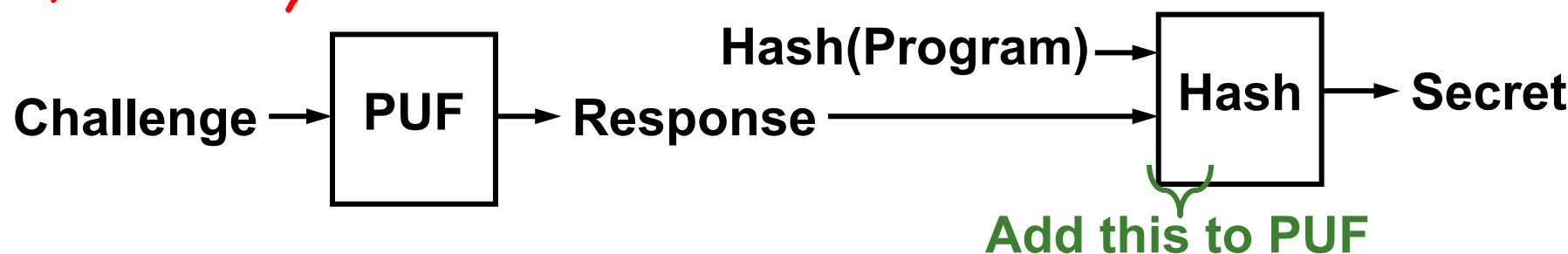
# Restricting Access to the PUF

- To prevent the attack, the man in the middle must be **prevented** from finding out the response.
- **Alice's program** must be able to establish a shared secret with the PUF, **the attacker's program** must not be able to get the secret.

⇒ **Combine response with hash of program.**

- The PUF can only be accessed via the **GetSecret** function:

*Program* →



# Cryptographic Hash Function

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- Crypto hash function  $h(x)$  must provide
  - **Compression** — output length is small
  - **Efficiency** —  $h(x)$  easy to compute for any  $x$
  - **One-way** — given a value  $y$  it is infeasible to find an  $x$  such that  $h(x) = y$
  - **Weak collision resistance** — given  $x$  and  $h(x)$ ,  
infeasible to find  $y \neq x$  such that  $h(y) = h(x)$
  - **Strong collision resistance** — infeasible to find any  $x$  and  $y$ , with  $x \neq y$  such that  $h(x) = h(y)$

← (input  
choose 1)

← choose  
2

# Getting a Challenge-Response Pair

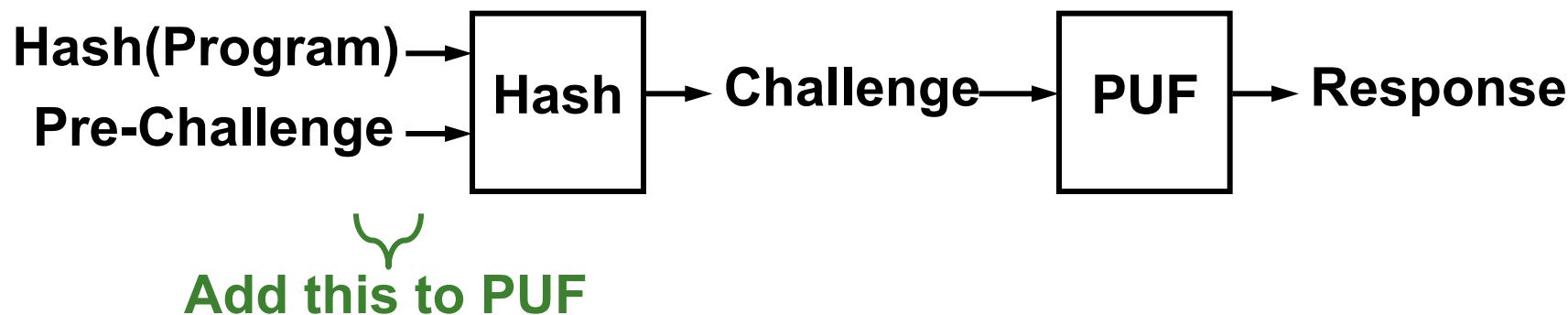
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- Now Alice **can** use a Challenge-Response pair to generate a shared **secret** with the PUF equipped device.
- But Alice **can't** get a Challenge-Response pair in the first place since the PUF **never** releases responses directly.  
⇒ **An extra function that can return responses is needed.**

# Getting a Challenge-Response Pair – 2

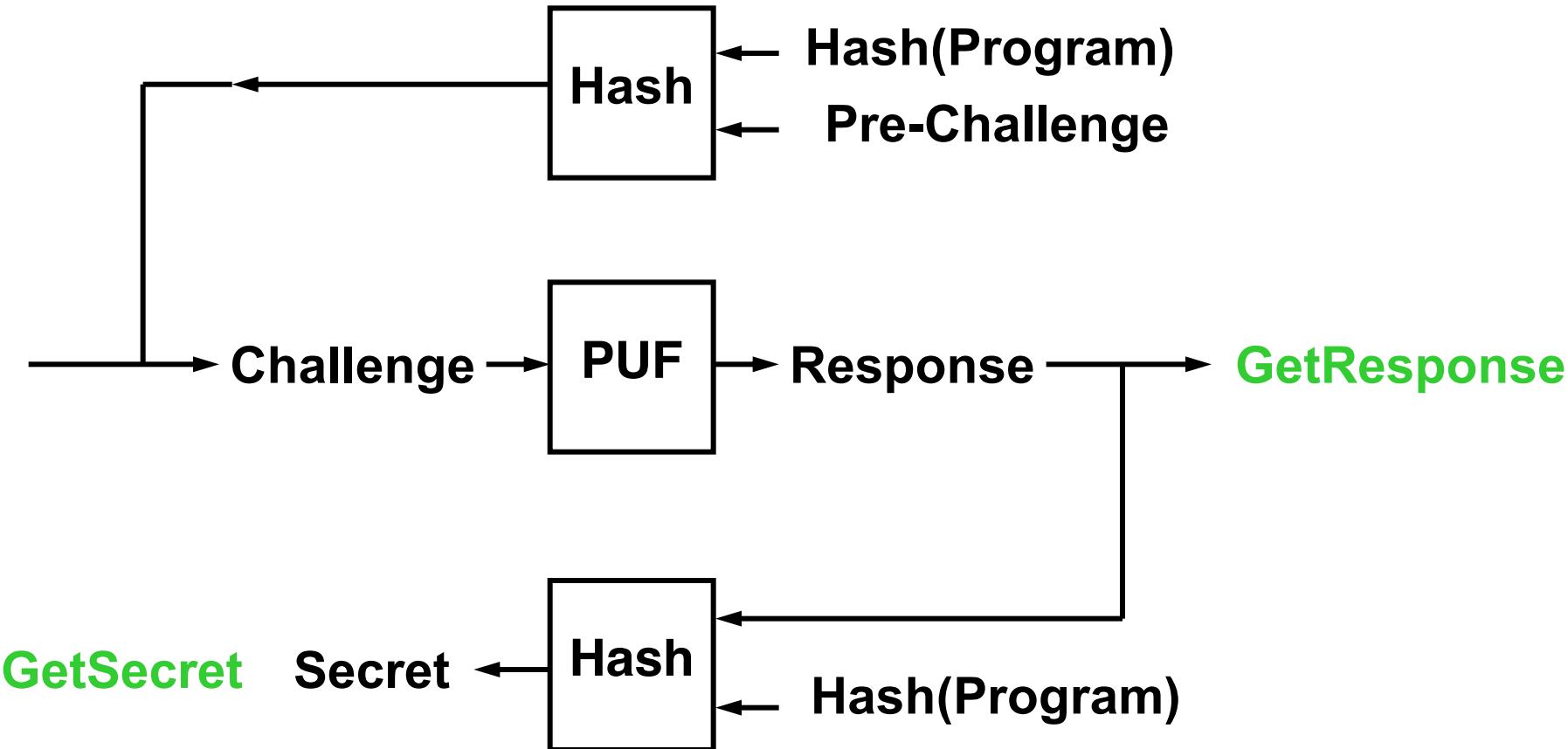
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- Let Alice use a Pre-Challenge.
  - Use program hash to prevent eavesdroppers from using the pre-challenge.
- 
- The PUF has a GetResponse function



# Controlled PUF Implementation

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# Software Licensing

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Program (Ecode, Challenge)

    Secret = GetSecret( Challenge )

    Code = Decrypt( Ecode, Secret )  Hash(Program)

Run Code

Ecode has been encrypted with Secret by Manufacturer

Secret is known to the manufacturer because he knows Response to Challenge and can compute

    Secret = Hash(Hash(Program), Response)

Adversary cannot determine Secret because he does not know Response or Pre-Challenge

If adversary tries a different program, a different secret will be generated because Hash(Program) is different

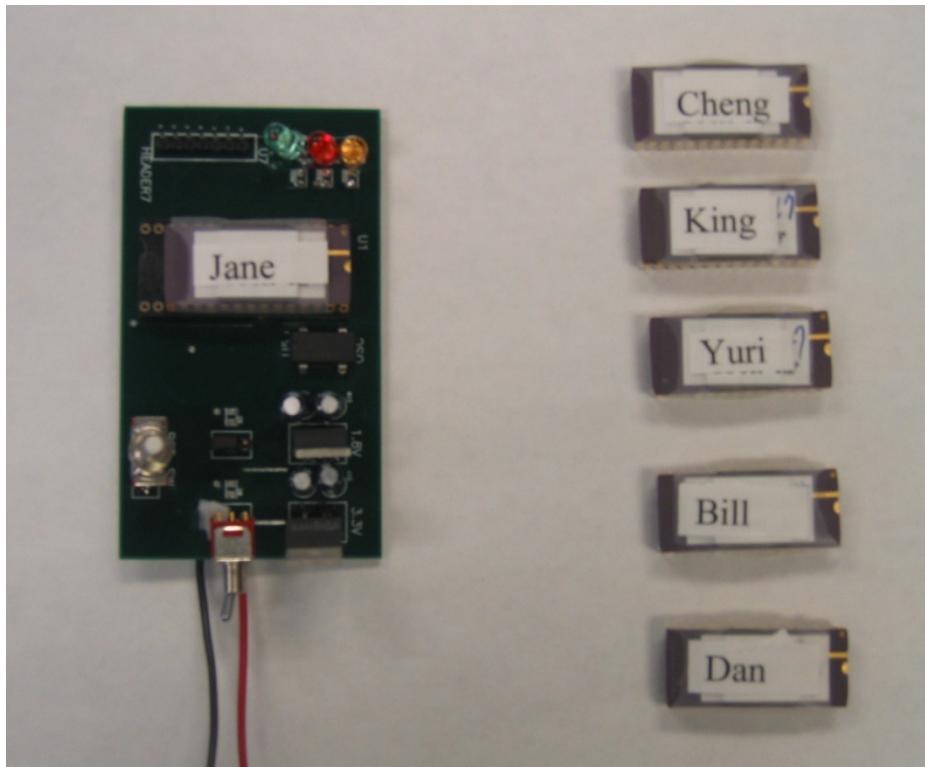
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# More on Physically Unclonable Functions (PUFs)

# PUF Experiments

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- Fabricated 200 “identical” chips with PUFs in TSMC 0.18 $\mu$  on 5 different wafer runs



## Security

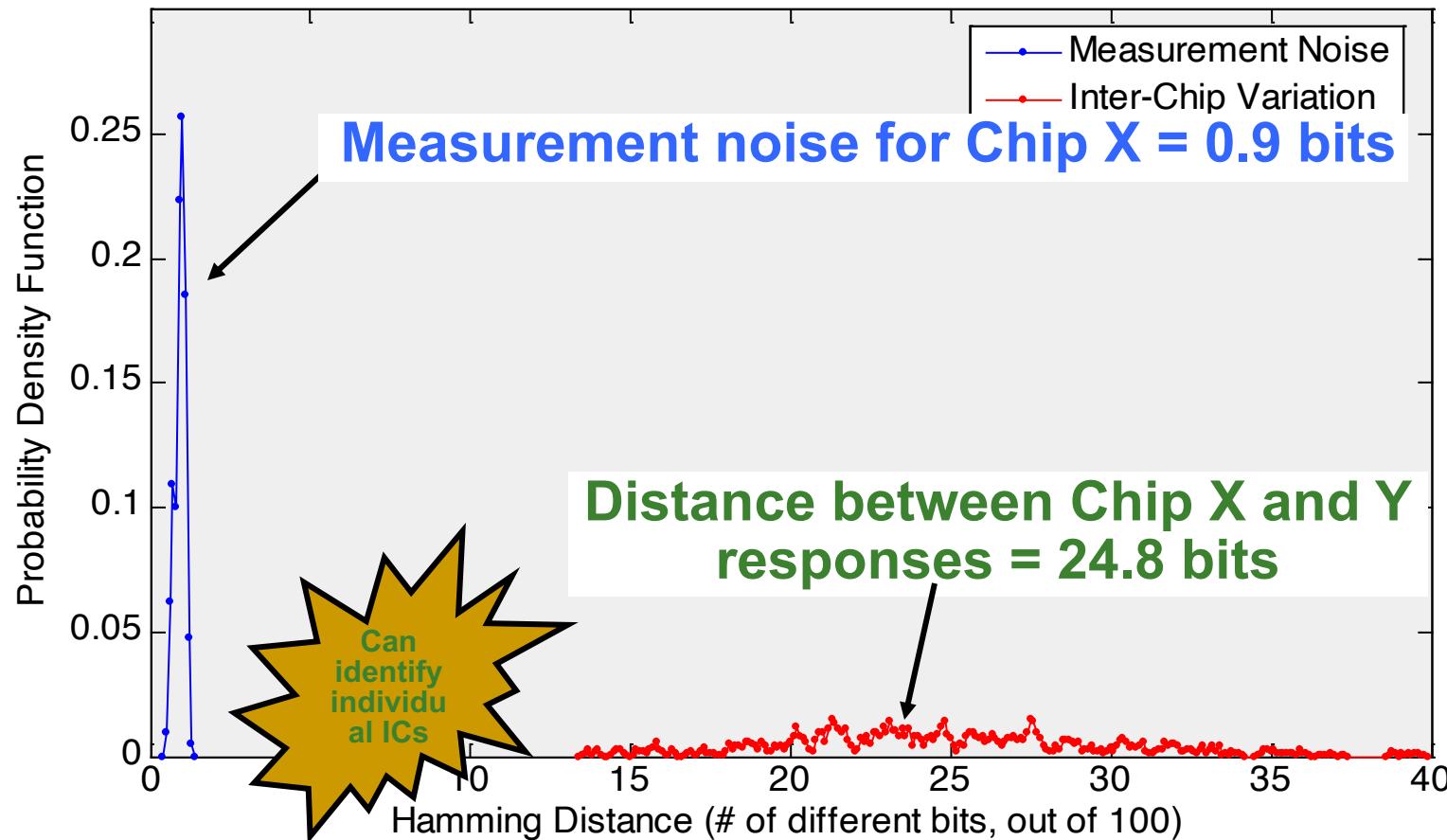
- What is the probability that a challenge produces **different responses** on two different PUFs?

## Reliability

- What is the probability that a PUF output for a challenge **changes with temperature**?
- With voltage variation?

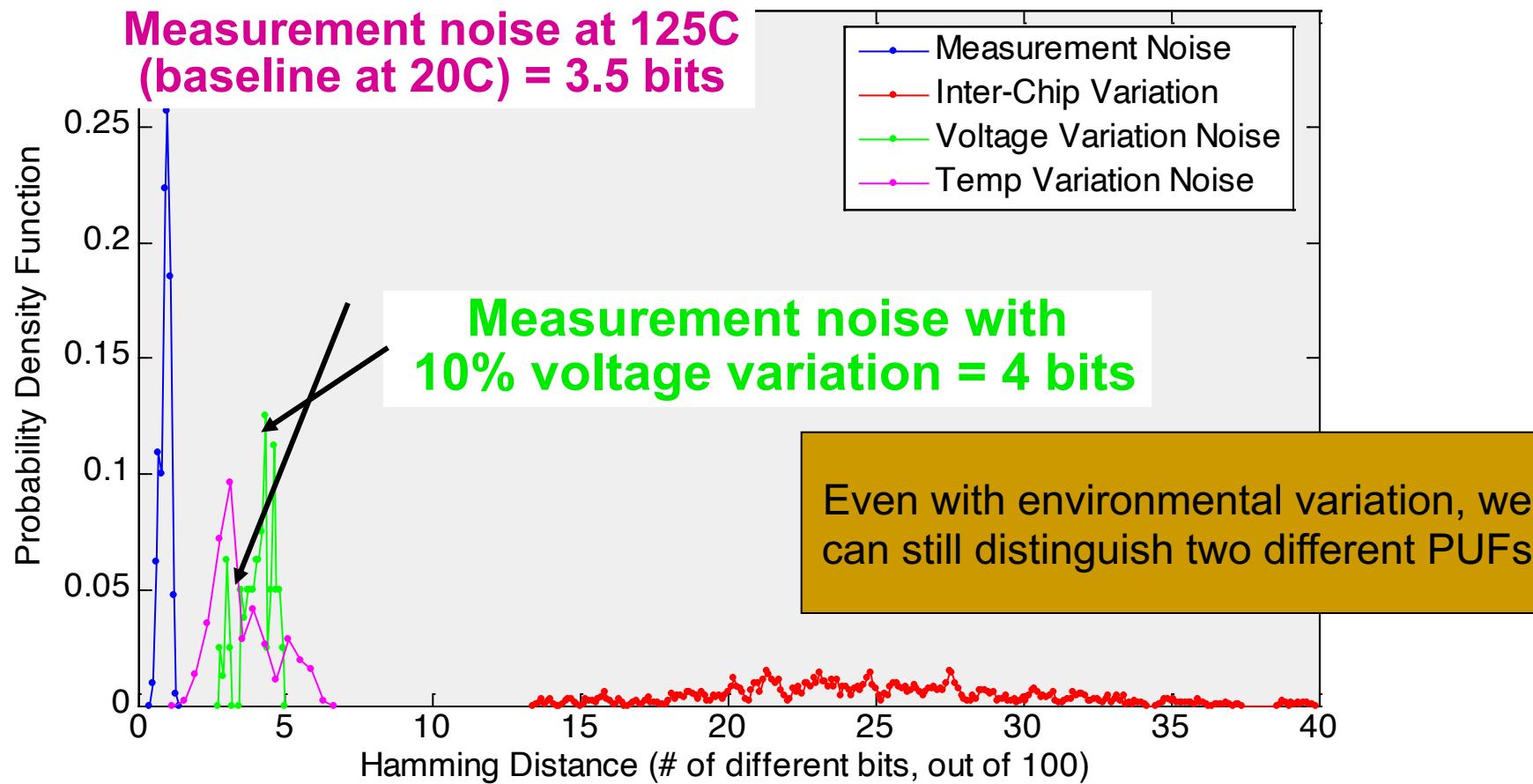
# Inter-Chip Variation

- Apply random challenges and observe 100 response bits



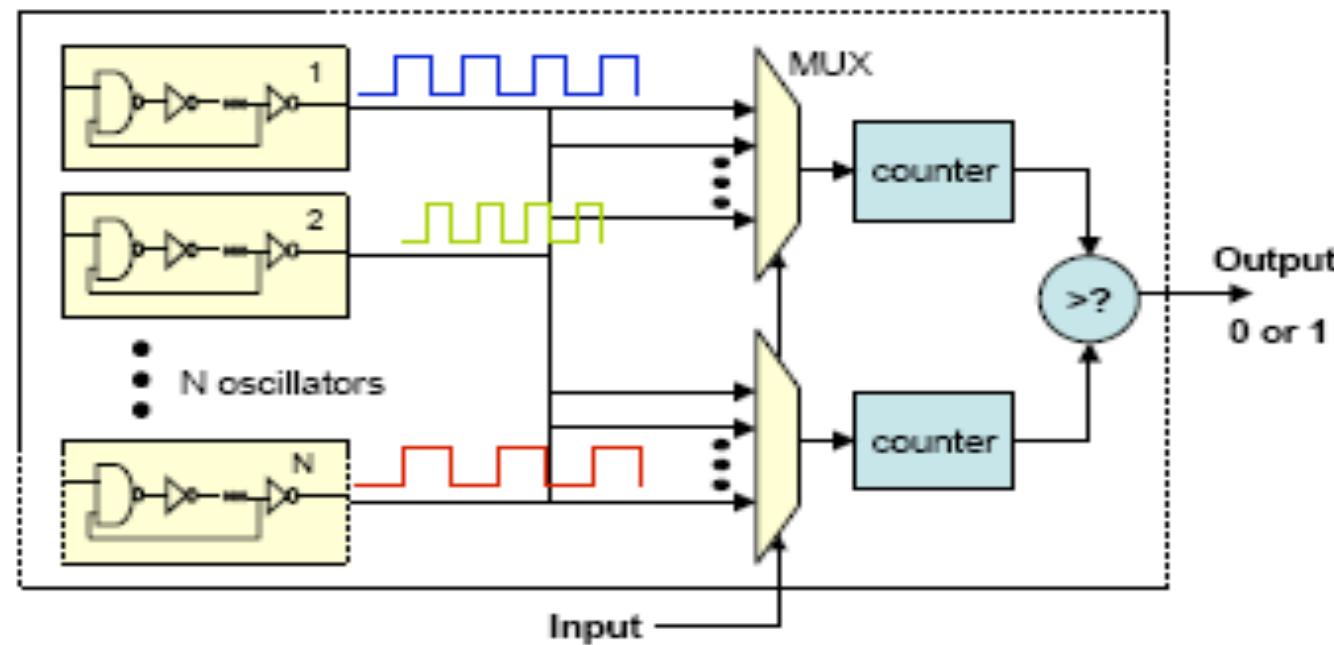
# Environmental Variations

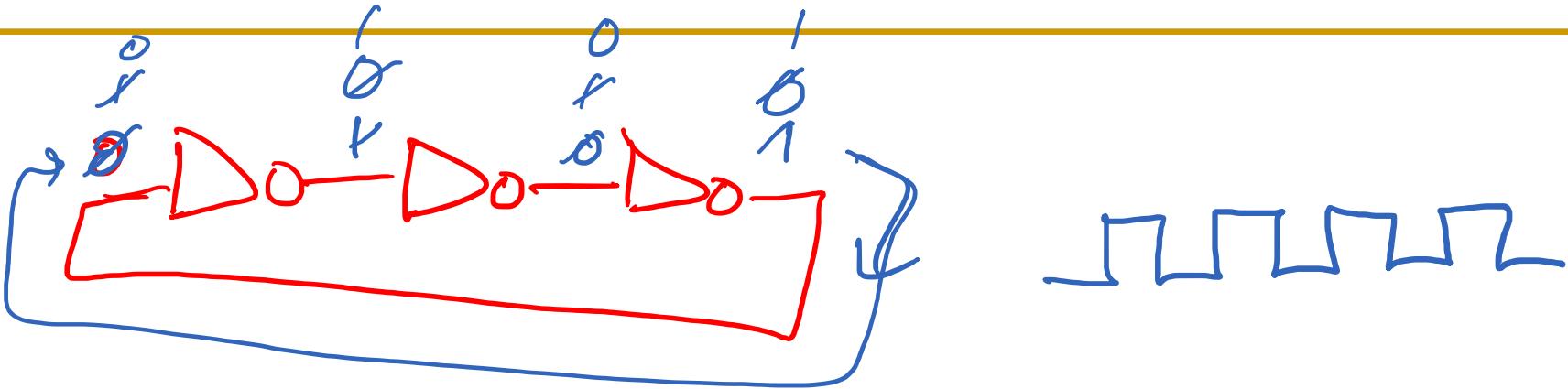
- What happens if we change voltage and temperature?



# Ring-Oscillator (RO) PUF

- The structure relies on delay loops and counters instead of MUX and arbiters
- Better results on FPGA – more stable





Ring Oscillator

# RO PUFs (cont'd)

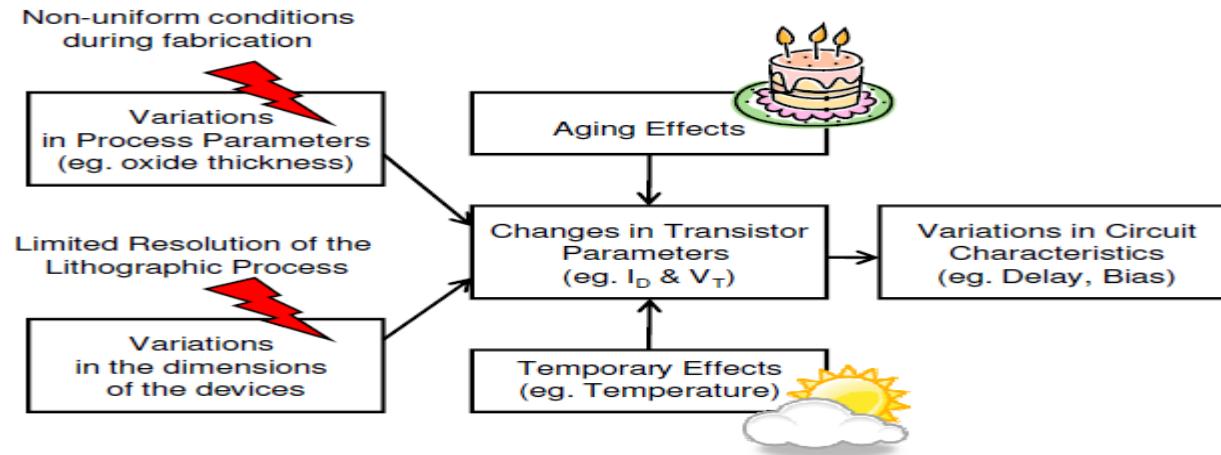
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- Easy to duplicate a ring oscillator and make sure the oscillators are identical
  - Much easier than ensuring the racing paths with equal path segments
- How many bits can we generate from the scheme in the previous page?
  - There are  $N(N-1)/2$  distinct pairs, but the entropy is significantly smaller:  $\log_2(N!)$
  - E.g., 35 ROs can produce 133 bits, 128 ROs can produce 716, and 1024 ROs can produce 8769

Consider the following minimal example, given three ROs:  $RO_A.f < RO_B.f$  and  $RO_B.f < RO_C.f$  implies  $RO_A.f < RO_C.f$ . The total PUF entropy is only  $\log_2(N!)$  bit as there are  $N!$  ways to sort the frequency values.

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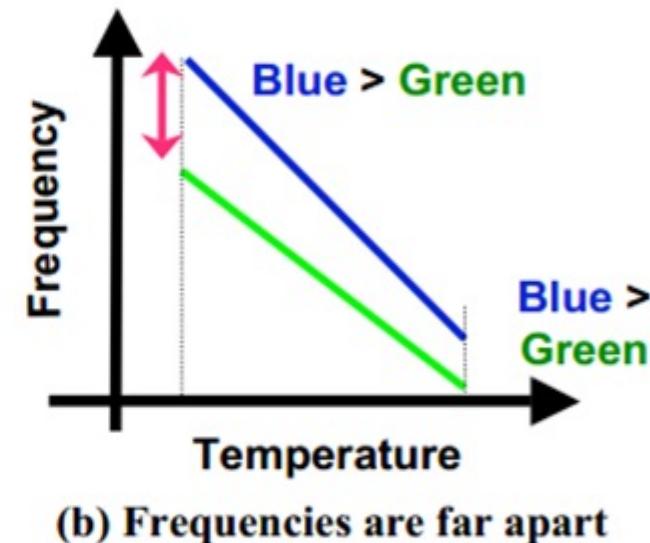
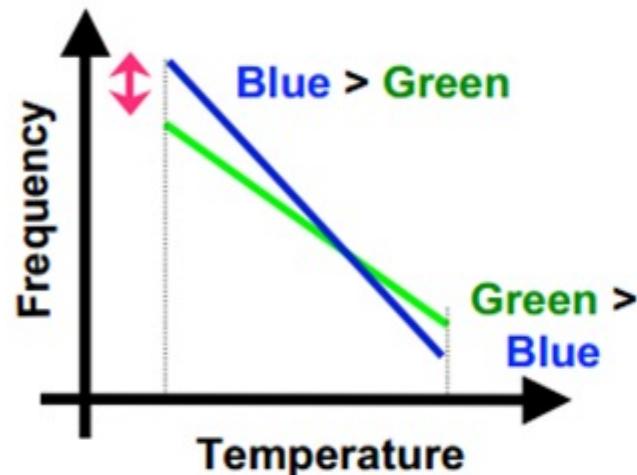
# Reliability of RO PUFs



- Two types of reliability issues:
- Aging:
  - Negative Bias Temperature Instability
  - Hot Carrier Injection (HCI)
  - Temp Dependent Dielectric Breakdown
  - Interconnect Failure
- Temperature
  - Slows down the device

# Reliability Enhancement

- Environmental changes have a large impact on the freq. (and even relative ones)



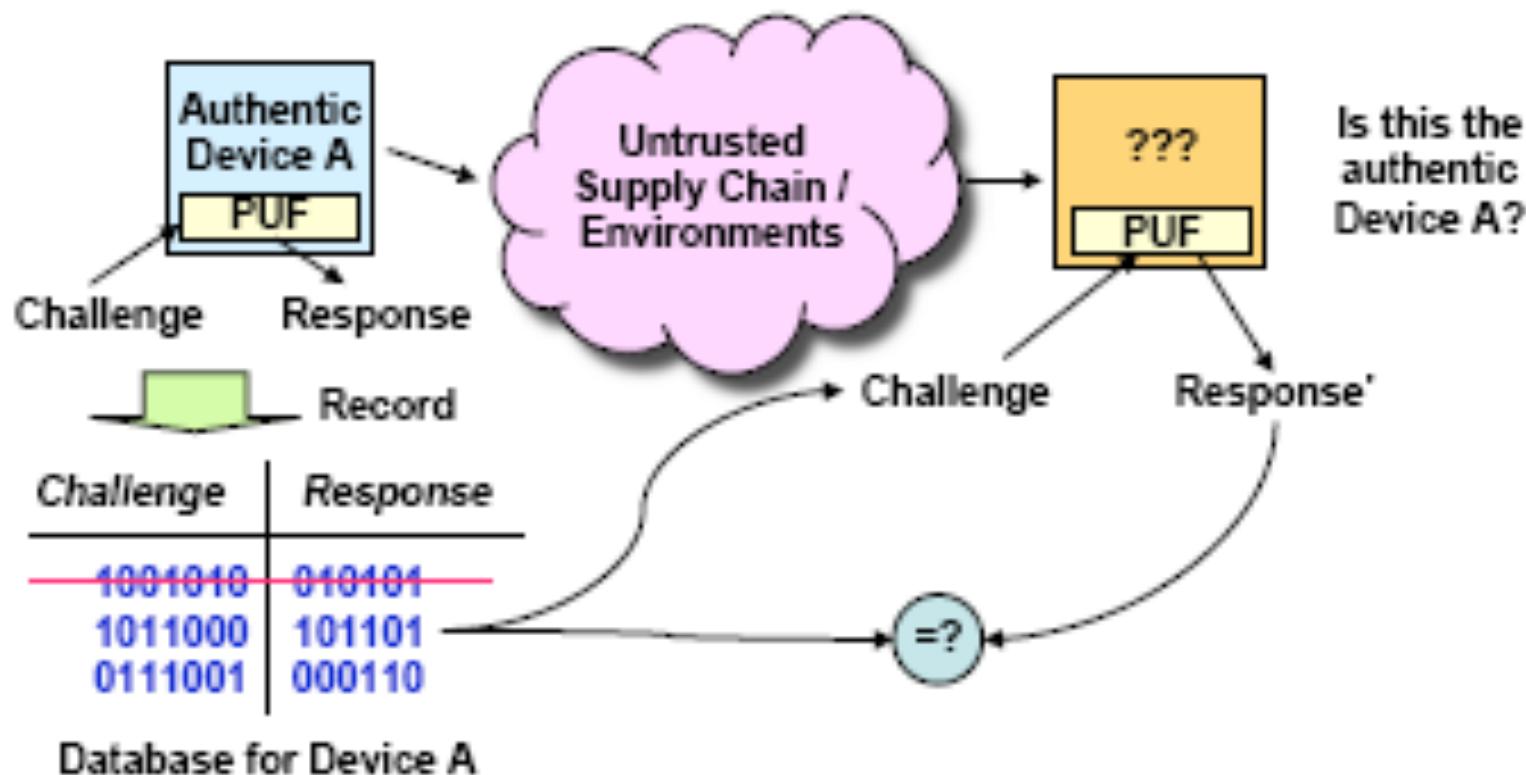
# RO PUFs

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- ROs whose frequencies are far are more stable than the ones with closer frequencies
    - Possible advantage: do not use all pairs, but only the stable ones
    - It is easy to watch the distance in the counter and pick the very different ones.
      - Can be done during enrollment
  - RO PUF allows an easier implementation for both ASICs and FPGAs.
  - The **Arbiter** PUF is appropriate for resource constrained platforms such as RFIDs and the **RO PUF** is better for use in FPGAs and in secure processor design.
-

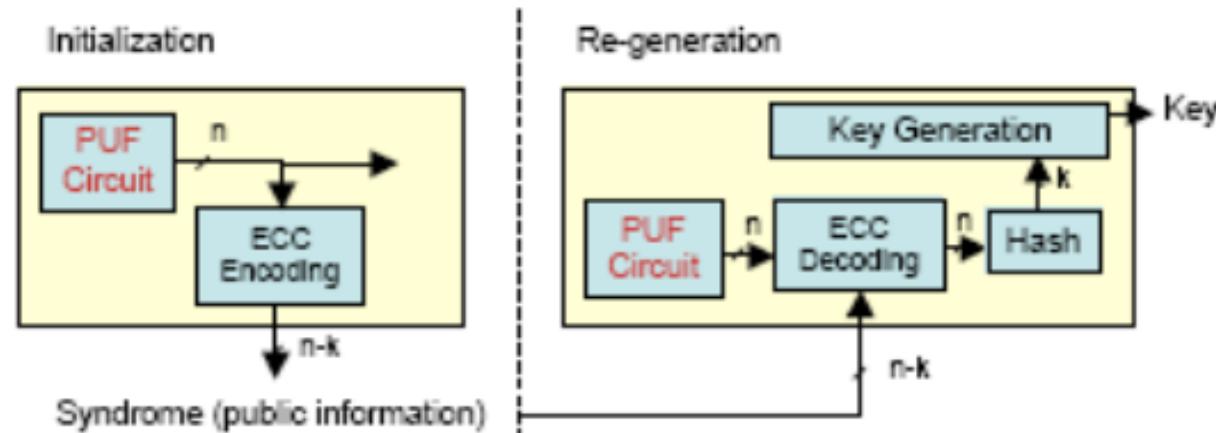
# Applications – Authentication

- Same challenges should not be used to prevent the man-in-the-middle attacks



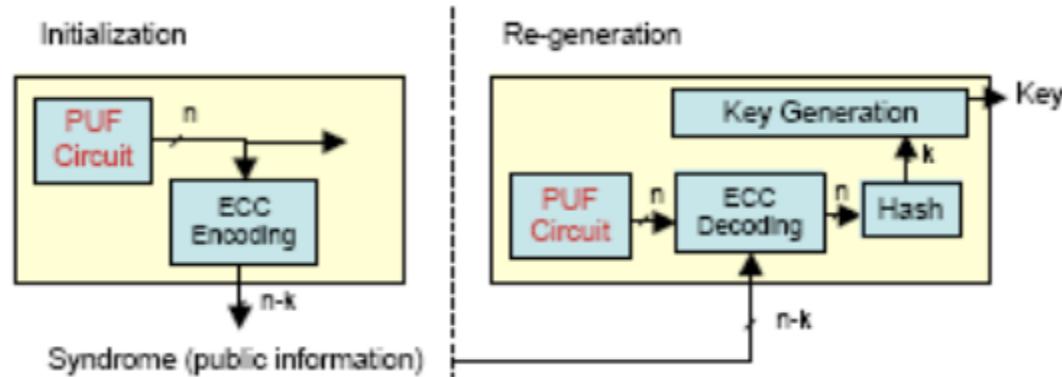
# Application – Cryptographic Key Generation

- The unstability is a problem
- Some crypto protocols (e.g., RSA) require specific mathematical properties that random numbers generated by PUFs do not have
- How can we use PUFs to generate crypto keys?
  - Error correction process: initialization and regeneration
  - There should be a one-way function that can generate the key from the PUF output



# Crypto Key Generation

- Initialization: a PUF output is generated and error correcting code (e.g., BCH) computes the syndrome (public info)
- Regeneration: PUF uses the syndrome from the initial phase to correct changes in the output
- Clearly, the syndrome reveals information about the circuit output and introduces vulnerabilities



# Experiments with RO PUFs

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- Experiments done on 15 Xilinx Virtex4 LX25 FPGA (90nm)
- They placed 1024 ROs in each FPGA as a 16-by-64 array
- Each RO consisted of 5 INVs and 1 AND, implemented using look-up tables
- The goal is to know if the PUF outputs are unique (**for security**) and reproducible (**for reliability and security**)

# Reliability and Security Metrics

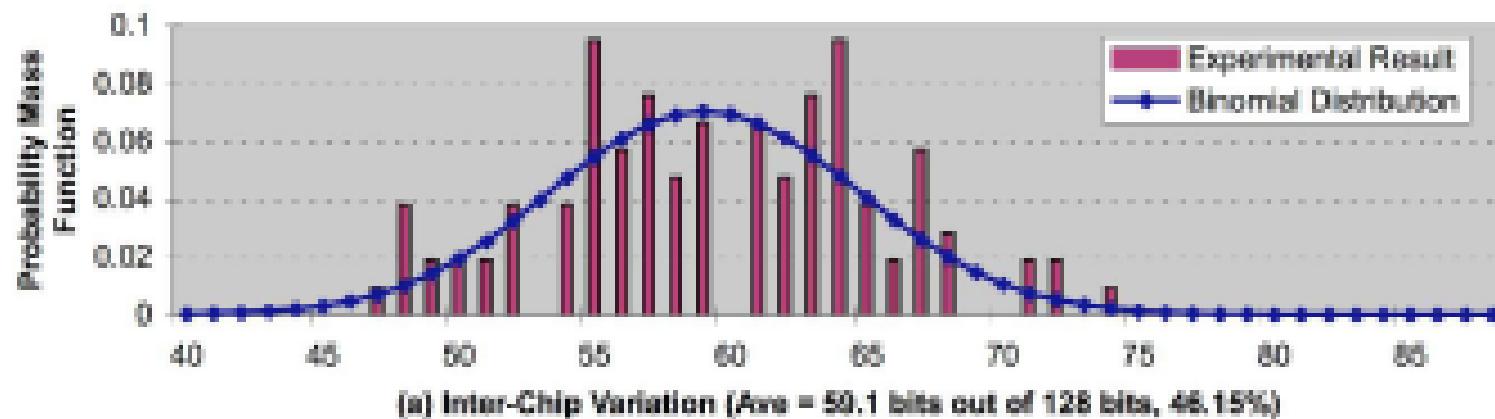
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- *Inter-chip variation:* How many PUF output bits are different between PUF A and PUF B? This is a measure of uniqueness. If the PUF produces uniformly distributed independent random bits, the inter-chip variation should be 50% on average.
- *Intra-chip (environmental) variation:* How many PUF output bits change when re-generated again from a single PUF with or without environmental changes? This indicates the *reproducibility* of the PUF outputs. Ideally, the intra-chip variation should be 0%.

# The Probability Distribution for Inter-chip Variations

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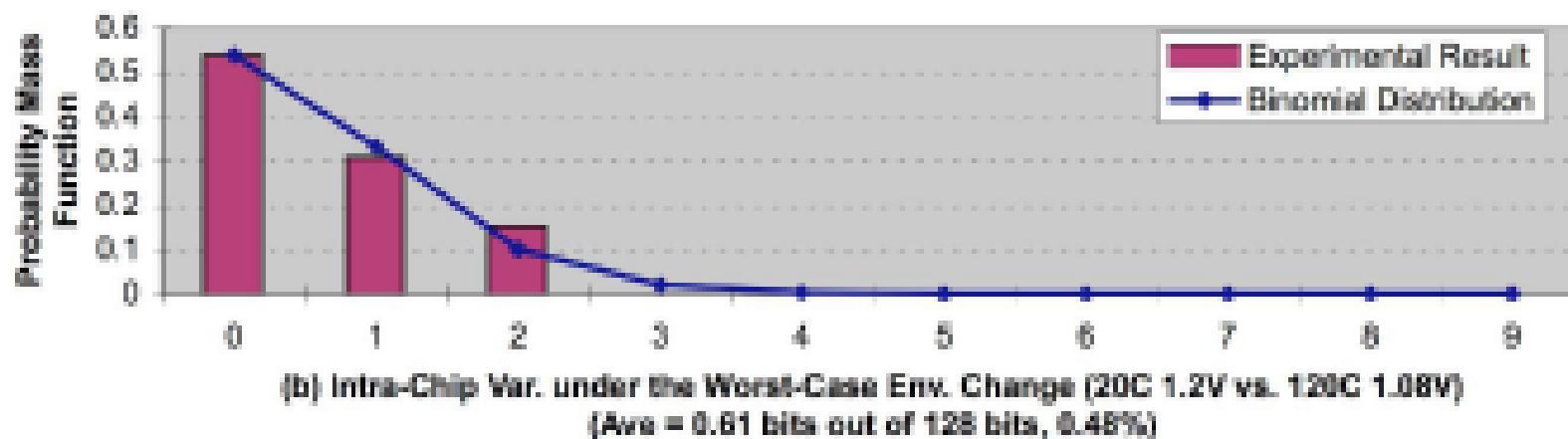
- 128 bits are produced from each PUF
- x-axis: number of PUF o/p bits different b/w two FPGAs; y-axis: probability
- Purple bars show the results from 105 pair-wise comparisons
- Blue lines show a binomial distribution with fitted parameters ( $n=128$ ,  $p = 0.4615$ )
- Average inter-chip variations  $0.4615 \sim 0.5$



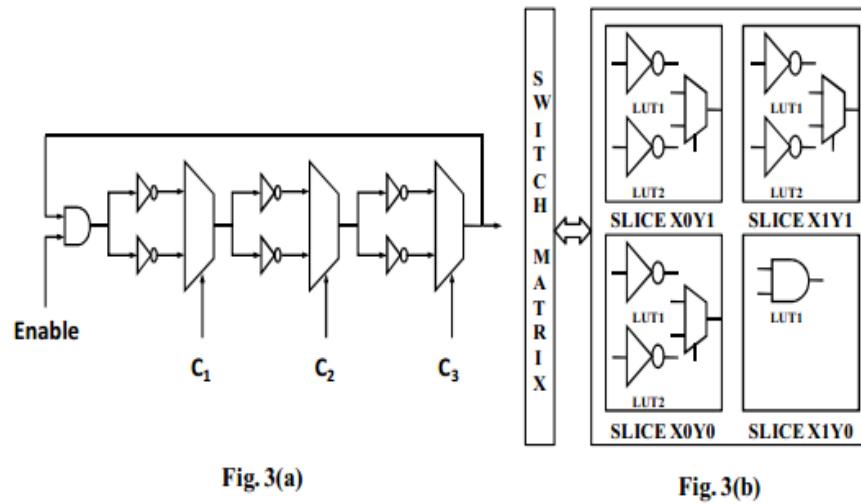
# The Probability Distribution for Intra-chip Variations

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- PUF responses are generated at two different conditions and compared
- Changing the temperature from 20°C to 120°C and the core voltage from 1.2 to 1.08 altered the PUF o/p by ~0.6 bits (0.47%)
- Intra-chip variations is much lower than inter-chip – the PUF o/p did not change from small to moderate



# Configurable Ring Oscillator

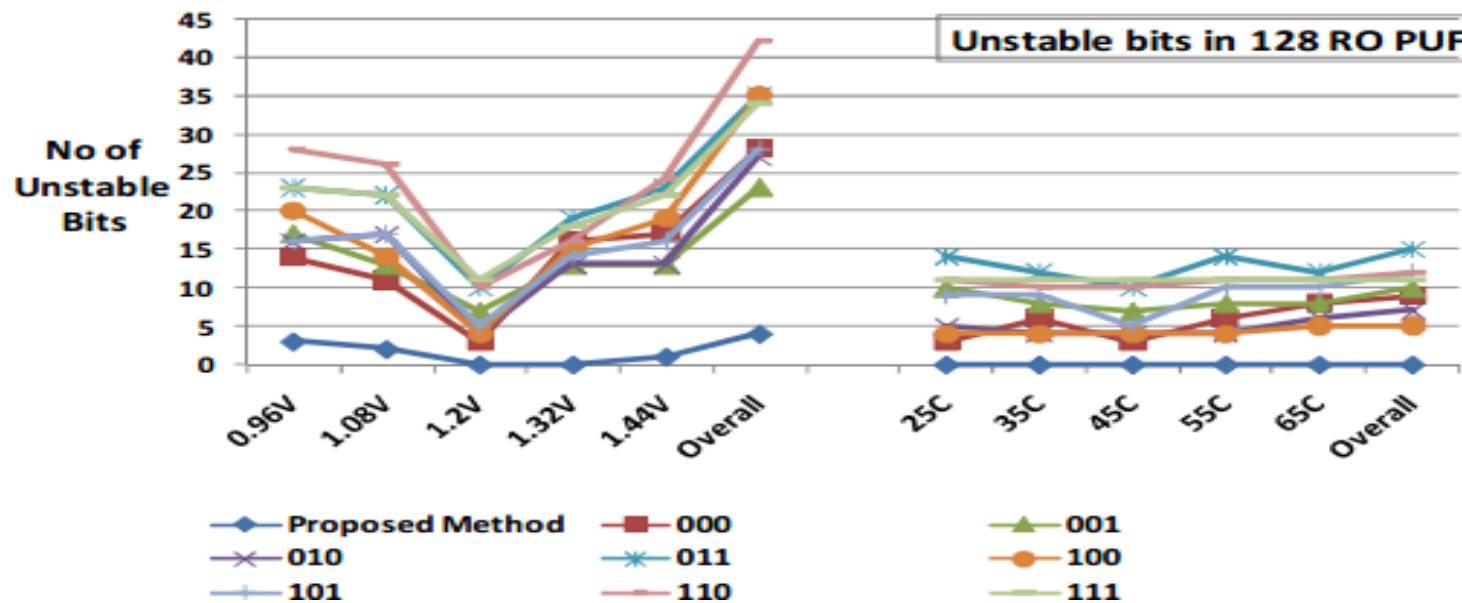


**Table 1.** Frequency differences in a configurable RO pair

$c_1c_2c_3$	Frequency of ROs in CLB i	Frequency of ROs in CLB j	$\Delta f$
000	$f_0$	$f_0$	$ f_0 - f_0 $
001	$f_1$	$f_1$	$ f_1 - f_1 $
010	$f_2$	$f_2$	$ f_2 - f_2 $
011	$f_3$	$f_3$	$ f_3 - f_3 $
100	$f_4$	$f_4$	$ f_4 - f_4 $
101	$f_5$	$f_5$	$ f_5 - f_5 $
110	$f_6$	$f_6$	$ f_6 - f_6 $
111	$f_7$	$f_7$	$ f_7 - f_7 $

- The pair which has the maximum difference in frequency in CRO is selected.

# Configurable Ring Oscillator



- Higher difference in frequency will ensure higher reliability
- Redundancy

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# True Random Number Generator



# Random Numbers in Cryptography

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- The keystream in the one-time pad
  - The secret key in the DES encryption
  - The prime numbers  $p, q$  in the RSA encryption
  - Session keys
  - The private key in digital signature algorithm (DSA)
  - The initialization vectors (IVs) used in ciphers
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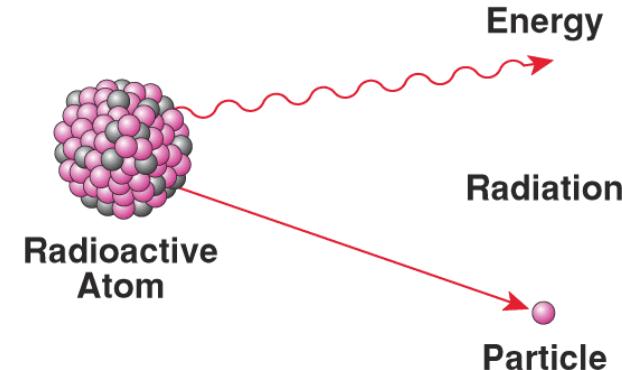
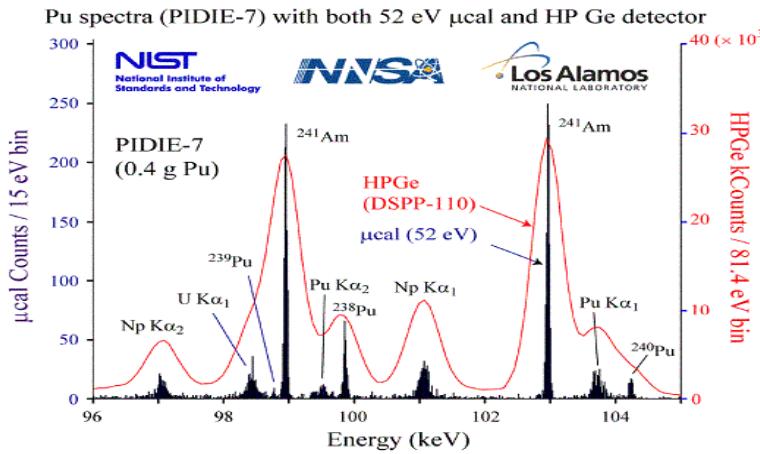
# Pseudo-random Number Generator

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- **Pseudo-random number generator:**
  - A polynomial-time computable function  $f(x)$  that expands a short random string  $x$  into a long string  $f(x)$  that appears random
- **Not truly random in that:**
  - Deterministic algorithm
  - Dependent on initial values (seed)
- **Objectives**
  - Fast
  - Secure

# Sources

- The only truly random number sources are those related to physical phenomena such as **the rate of radioactive decay** of an element or the **thermal noise** of a semiconductor.



- Randomness is bound to natural phenomena. It is impossible to algorithmically generate truly random numbers.

Microcalorimeter (black) and high-purity germanium (red) spectra of a mixture of plutonium isotopes. Minimal thermal noise is achieved at 100 mK. High sensitivity is due to use of a superconducting quantum interference device.

# Good TRNG Design

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## ■ Entropy Source:

- ❑ Randomness present in physical processes such as thermal and shot noise in circuits, brownian motion, or nuclear decay.

## ■ Harvesting Mechanism:

- ❑ The mechanism that does not disturb the physical process but collects as much entropy as possible.

## ■ Post-Processing (optional):

- ❑ Applied to mask imperfections in entropy sources or harvesting mechanism or to provide tolerance in the presence of environmental changes and tampering.

# Set of Requirements

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- ❑ The Design Should be purely digital
  - ❑ The harvesting mechanism should be simple.
    - The unpredictability of the TRNG should not be based on the complexity of the harvesting mechanism, but only on the unpredictability of the entropy source.
  - ❑ No correction circuits are allowed
  - ❑ Compact and efficient design (high throughput per area and energy spent).
  - ❑ The design should be sufficiently simple to allow rigorous analysis.
-

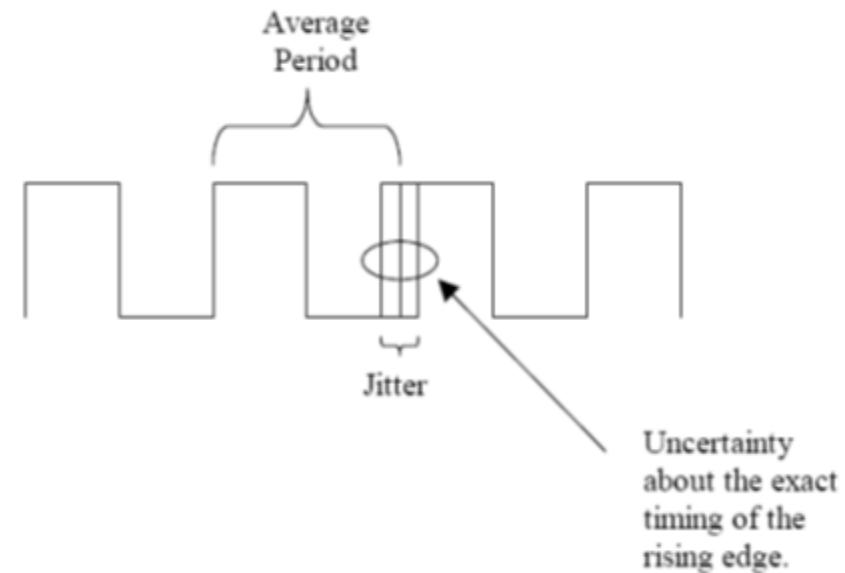
# Method : Clock Jitter

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- Jitter is variations in the significant instants of a clock
- Jitter is nondeterministic (random)

- **Sources of Jitter:**

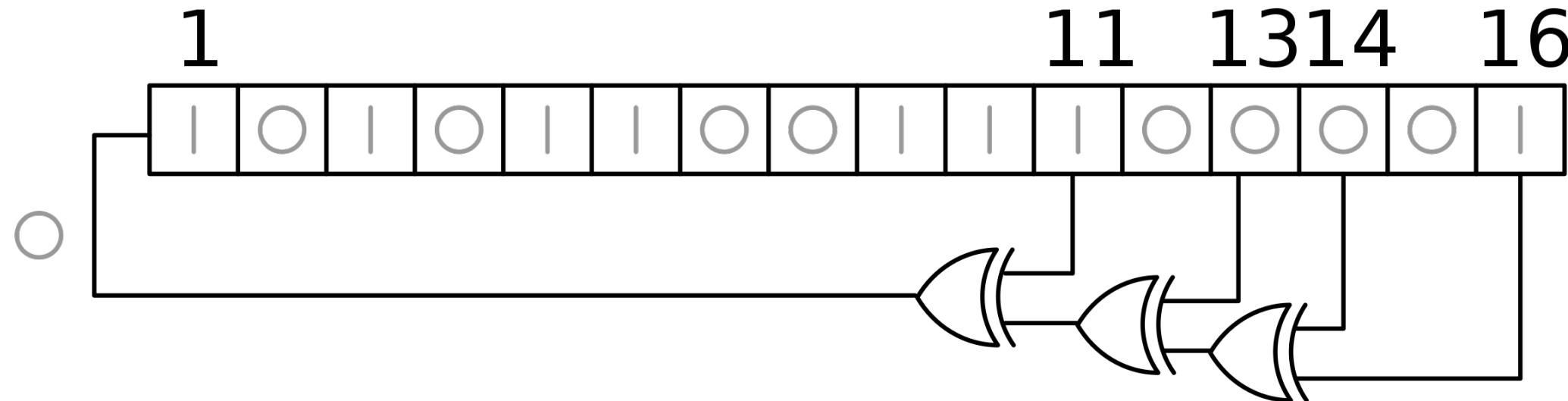
- Semiconductor noise
- Cross-talk
- Power supply variations
- Electromagnetic fields



# LFSR

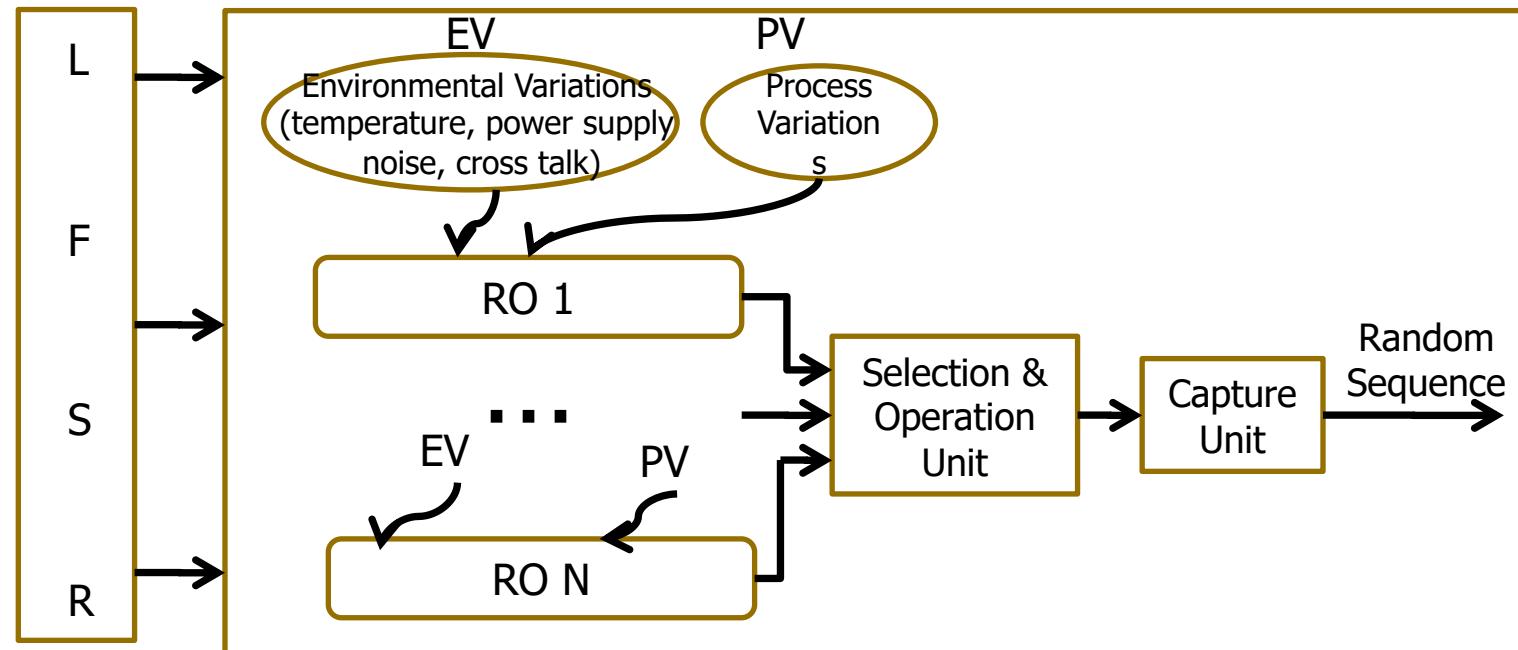
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- Linear Feedback Shift Register - a shift register whose input bit is a linear function of its previous state



# TRNG Structure

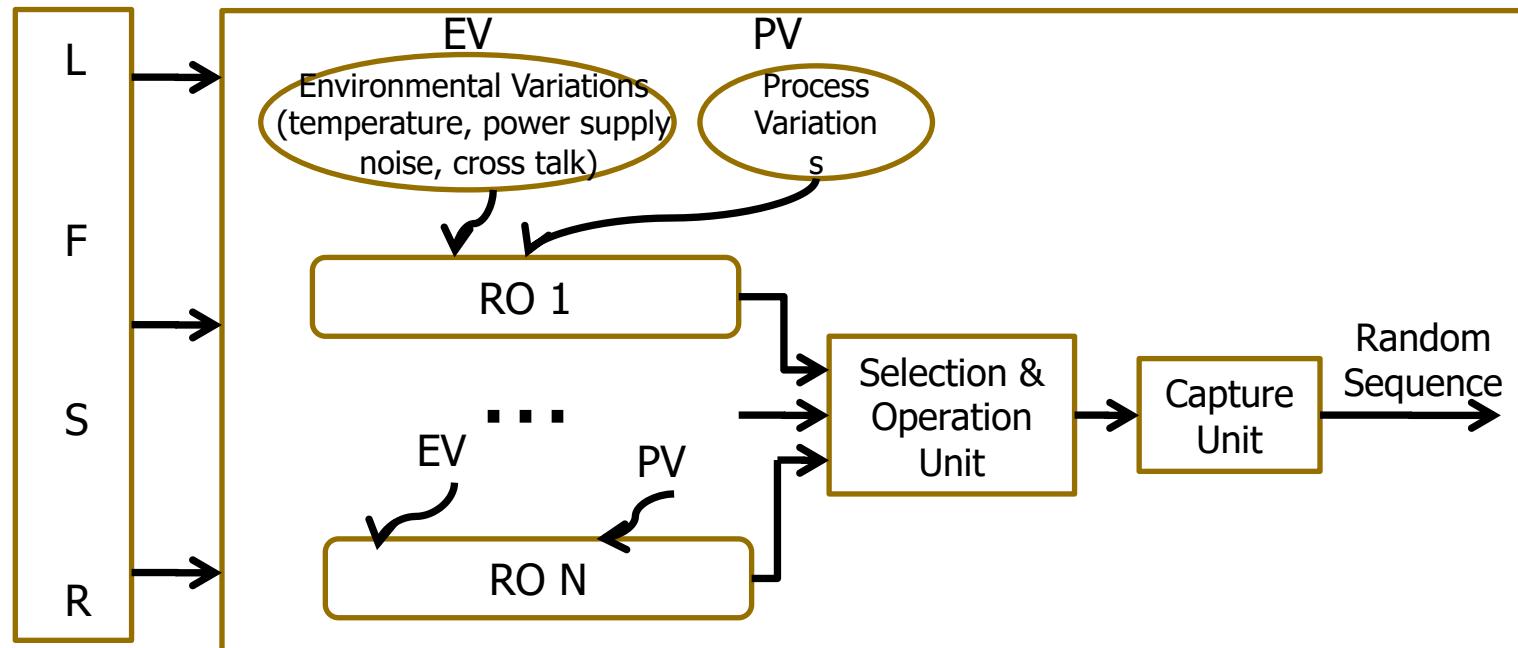
- ❑ **LFSR:** Generate random patterns, causing random switching noise



# TRNG Structure

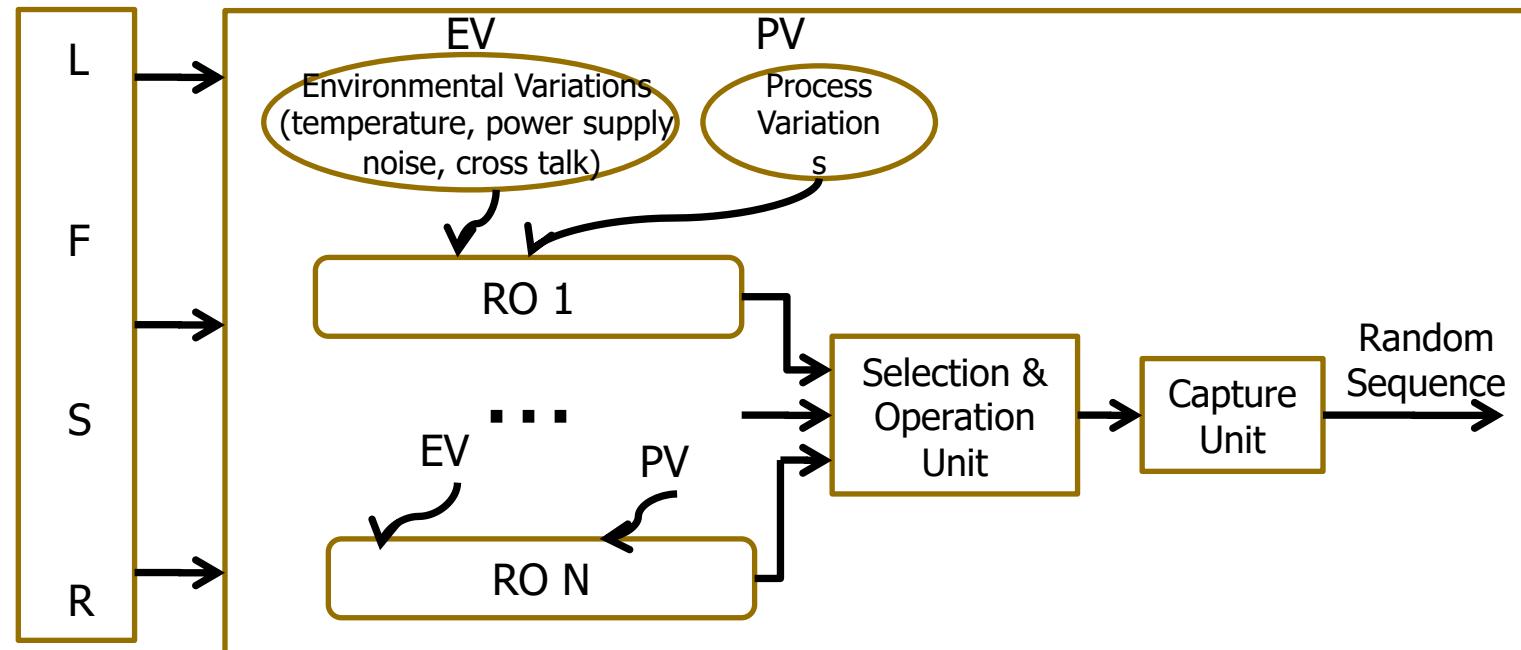
## □ Ring Oscillators

- Process variations & environmental variations
- Random phase jitter



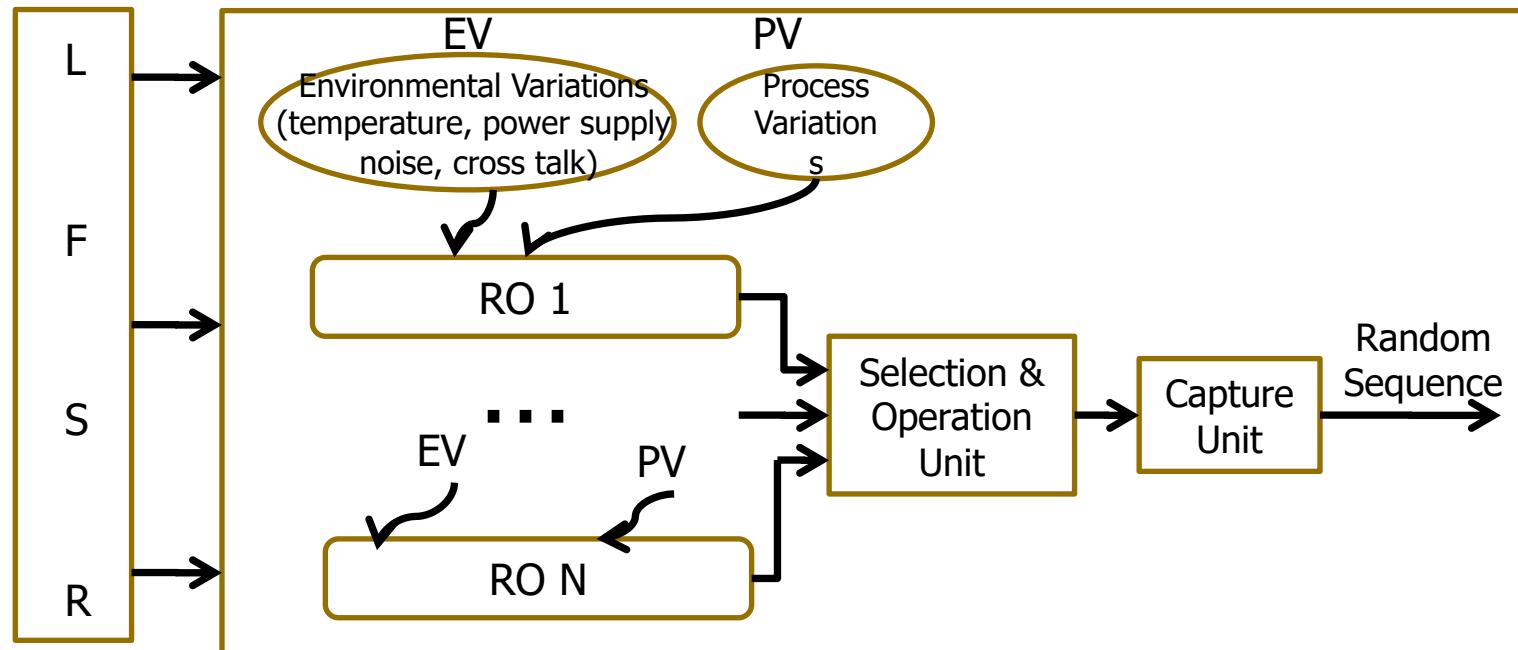
# TRNG Structure

- **Selection & Operation Unit:** The random phase of ring oscillators could be translated into digital values by this unit, such as XOR operation

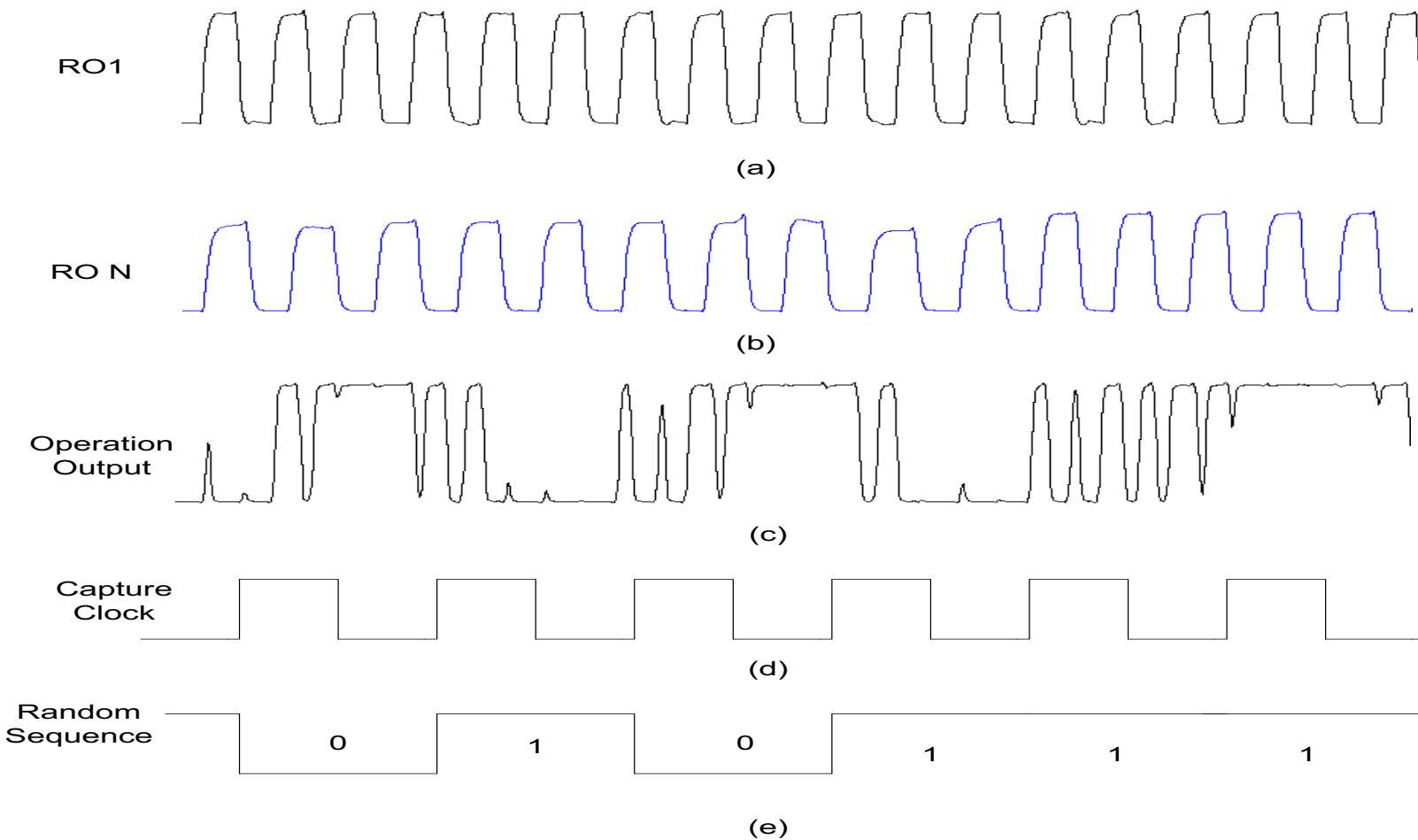


# TRNG Structure

- **Capture Unit:** Make sure the digital value is sampled with the frequency of the required true random number.



# TRNG Output



# References

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- [2]Gassend, B., Lim, D., Clarke, D., van Dijk, M., Devadas, S.: Identification and au-thentication of integrated circuits: Research articles. Concurr. Comput. : Pract. Exper.16(11), 1077-1098.
- [3]Gassend, B., Clarke, D., van Dijk, M., Devadas, S.: Controlled physical random functions. In: ACSAC '02: Proceedings of the 18th Annual Computer Security Applications Conference, p. 149. IEEE Computer Society, Washington, DC, USA (2002)
- [4]B. Gassend, D. Clarke, M. van Dijk, and S. Devadas. Silicon physical random functions. In Proceedings of the Computer and Communication Security Conference , November 2002.
- [5] Dinesh Ganta, Vignesh Vivekraja, Kanu Priya and Leyla Nazhandali, “A Highly Stable Leakage-Based Silicon Physical Unclonable Functions”

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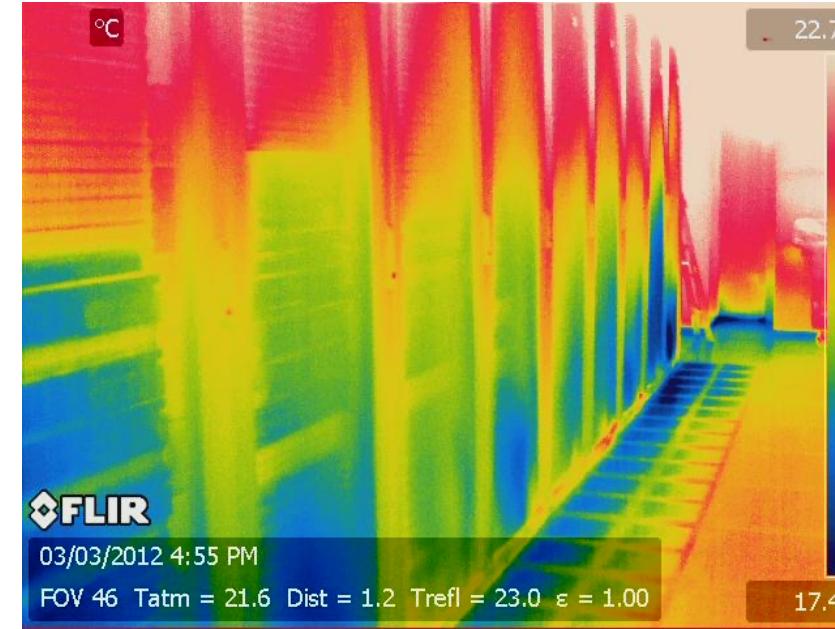
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- [6] A. Maiti and P. Schaumont, “Improved ring oscillator puf: An fpga-friendly secure primitive,” *J. Cryptology*, vol. 24, no. 2, pp. 375–397.,2011.
- [7] B. Sunar, W. J. Martin, D. R. Stinson. A Provably Secure True Random Number Generator with Built-in Tolerance to Active Attacks. *IEEE Transactions on Computers*, vol 58, no 1, pages 109-119, January 2007.

# More on PUF Applications

# Motivation: Securely Sensing without Key

- Physical features should be closely monitored
  - The temperature of data center, Facebook, Google, Amazon, etc



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- Physical features should be closely monitored
  - The temperature of data center, Facebook, Google, Amazon, etc
- Some physical features might be meaningful for security/privacy
  - Relative location between bank cards and automated teller machines (ATMs), card (not) present withdrawal

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    - The temperature of data center, Facebook, Google, Amazon, etc
  - Some physical features might be meaningful for security/privacy
    - Relative location between bank cards and automated teller machines (ATMs), card (not) present withdrawal
  - Some physical features are hard to sense but favoring many applications
    - Digital rights management, physically/irreversibly canceling membership?
  - Secure sensors are needed:
    - Operation safety, secrecy of sensitive data
-

# Key Is a Target

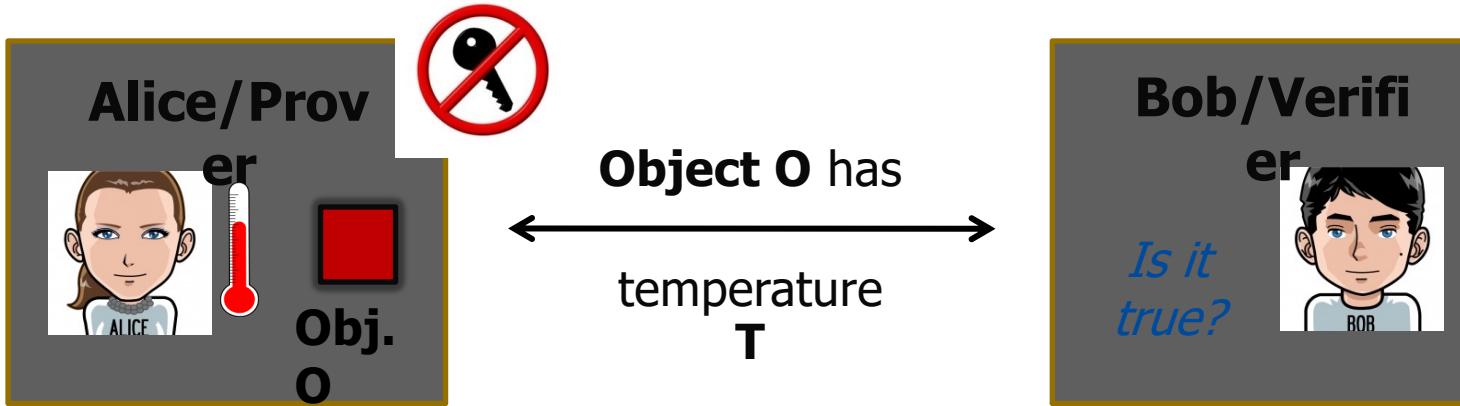
- Modern security protocols are commonly based on secret keys.
- A robust key enhances the robustness of security systems, but also announces itself as an interested target for attackers. [1]



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[1] Rivest, Ronald L. Illegitimi non carborundum. CRYPTO (2011).

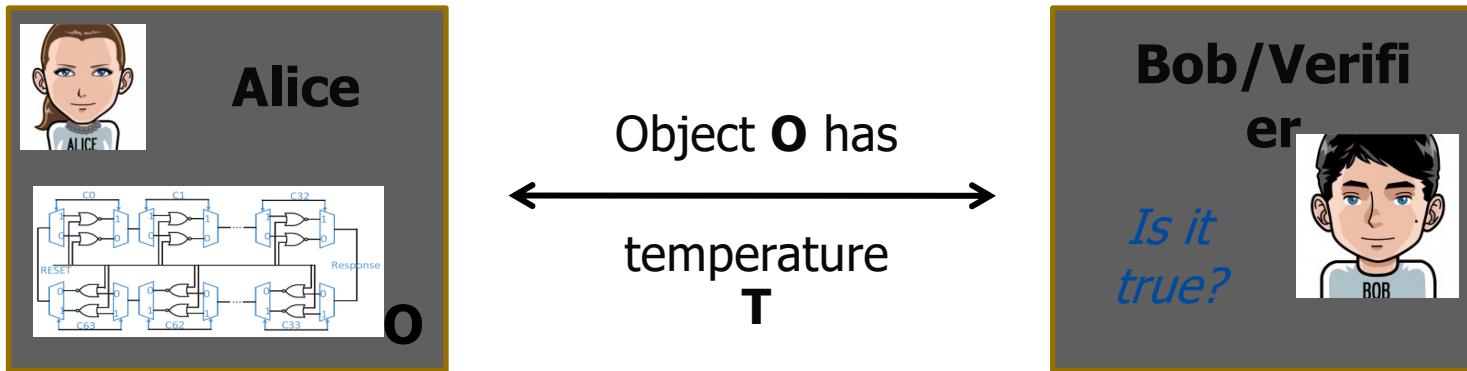
# Virtual Proofs of Temperature (Or: Keyless Temperature Sensors)



- Alice wants to prove to Bob that
  - a certain **object O**
  - is at temperature **T**
  - at the time of the execution of the VP
- **Classical approach:** Secure sensors with key known to Bob...
  - But, of course, we want no keys...

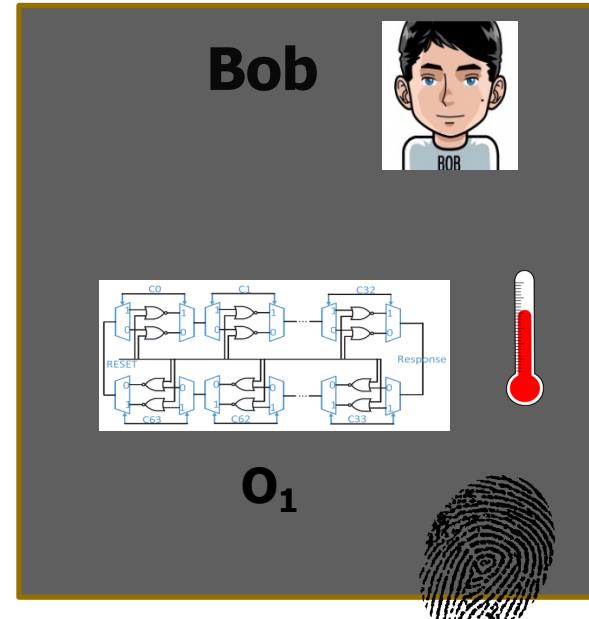
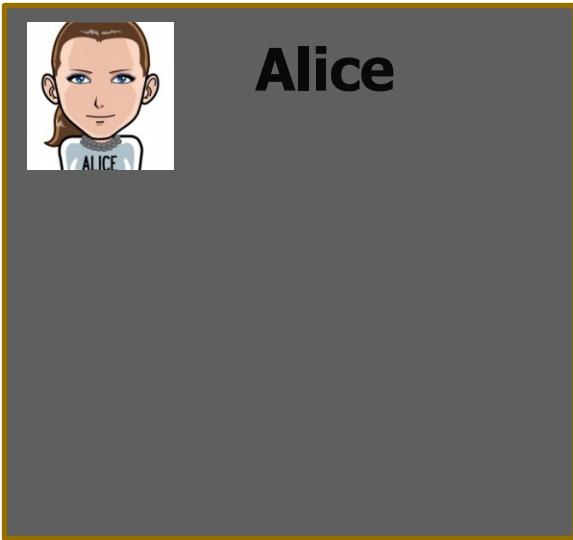
# VPs of Temperature via Noise Sensitivity

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- Key Idea: Use a ***temperature-sensitive Strong PUF*** as Object O
    - ie., an integrated circuit (IC) / PUF whose output **intentionally** depends upon the ambient temperature T
    - Output of the IC/PUF to Bob's random challenges proves temperature
-

# VP of Temperature $T_i \in \{T_1, \dots, T_k\}$



## Private set-up phase:

- For each temperature  $T_i \in \{T_1, \dots, T_k\}$ :
  - Bob puts the objects  $O_1$  at temperature  $T_i$   
Chooses  $n$  random challenges  $C_1^i, \dots, C_n^i$   
Measures the  $n$  resulting responses:  $r_1^i, \dots, r_n^i$
  - Bob creates&stores private list  $\text{List}(T_i) = (C_1^i, r_1^i), \dots, (C_n^i, r_n^i)$

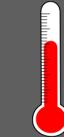
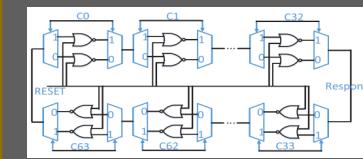
# VP of Temperature $T_i \in \{T_1, \dots, T_k\}$

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Alice

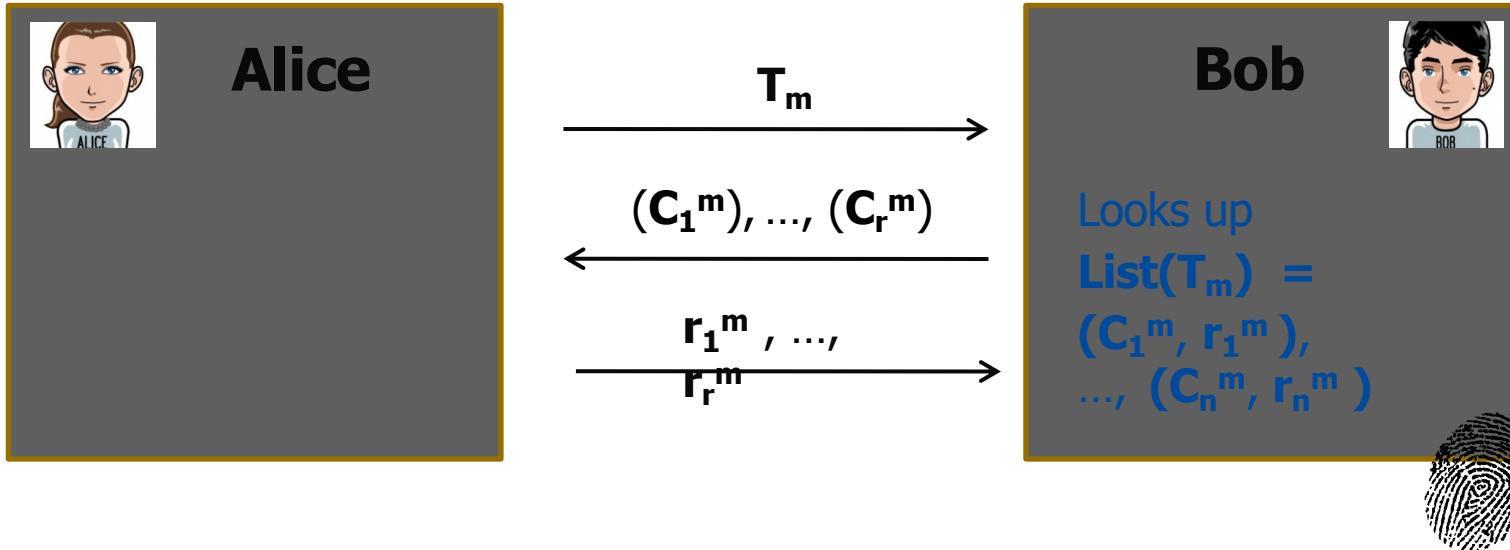
Bob



- Set-up phase closes, and the objects are transferred to **Alice**
- Some time may pass...  
(time for Alice to attack the system...)
- **VP starts!!**



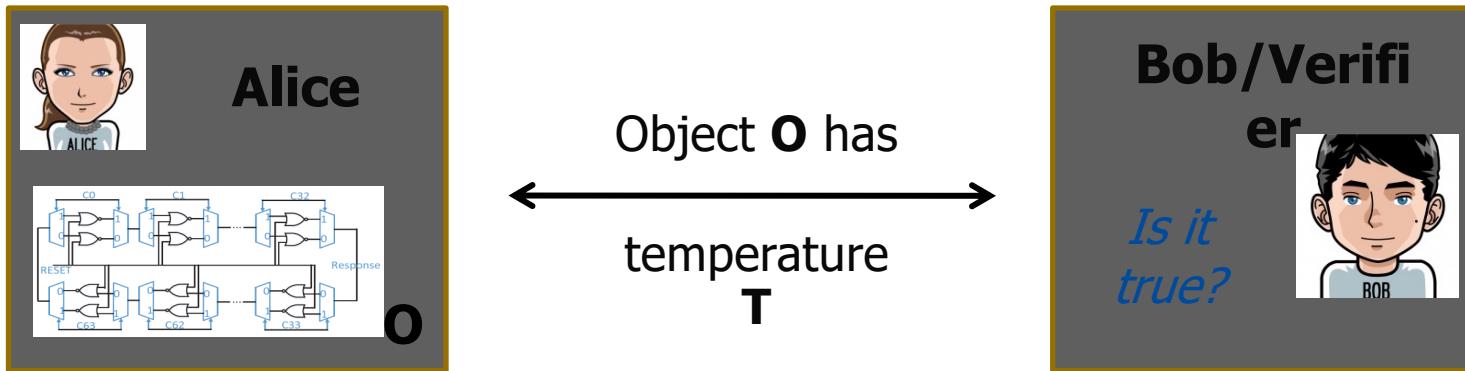
# VP of Temperature $T_i \in \{T_1, \dots, T_k\}$



- If pre-recorded values  $I_i^m$  in the **match** the values that **Alice** sent then **Bob** will **accept** the **VP**



# VPs of Temperature via Noise Sensitivity



- Leads to „keyfree temperature sensors“
  - Practical outcome of our new theoretical concept!!!
- Temp.-dependent behavior in PUFs usually **unwanted**
  - Turning known weakness of PUF into strength!
- **Again**, all standard properties of PUFs needed
  - Unclonability, no modeling, many challenge-response pairs