

# Introduction to AVL Trees

Unit 6

CC4 Data Structures and Algorithms

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# Introduction to AVL Trees

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What is an AVL tree?



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# Review: Binary Search Trees

- Each parent node has a maximum of two child nodes
  - Left child: lesser value than the parent node
  - Right child: higher value than the parent node
- Each tree has a tree height
  - Number of levels from the root node
  - Shorter tree height is usually more efficient in searching

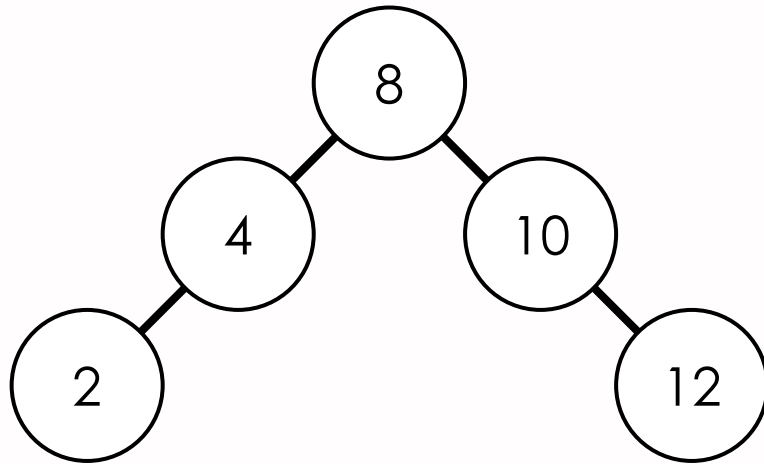


# Review: Binary Search Trees

**BST values: 8, 4, 2, 10, 12**

Tree height: 2

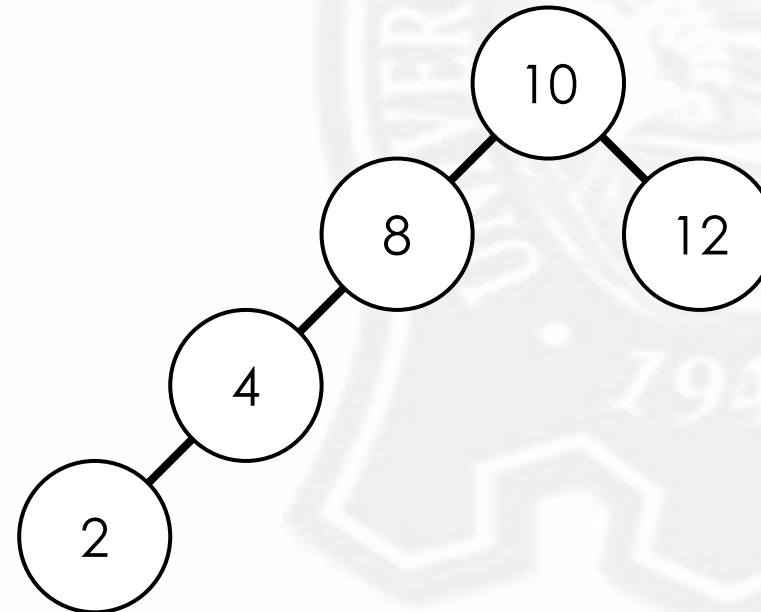
Type of tree: balanced



**BST values: 10, 8, 4, 2, 12**

Tree height: 3

Type of tree: unbalanced / skewed



# The Problem with Binary Search Trees

- Unbalanced / skewed BSTs can be less efficient:
  - Binary search trees are made to make searching more efficient
  - The higher the tree height, the more levels
- Balanced BSTs are usually more efficient:
  - Tends to have a lesser tree height

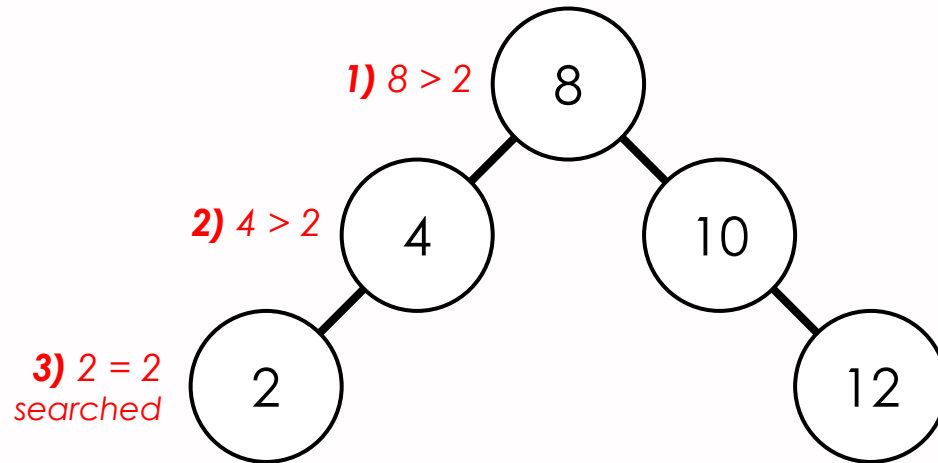


# The Problem with Binary Search Trees

## Example: Search '2' from the tree

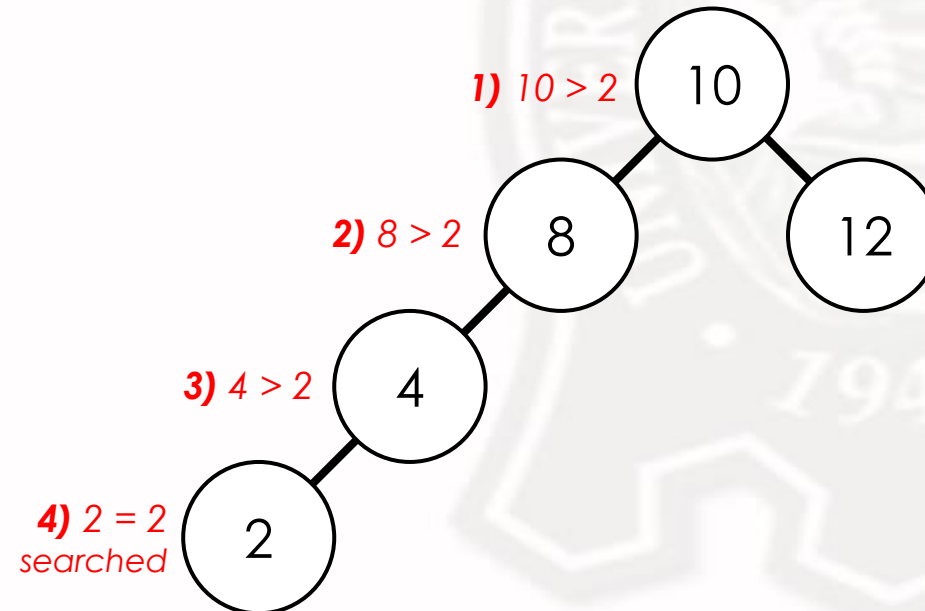
Tree height: 2

Number of comparisons: 3



Tree height: 3

Number of comparisons: 4



# Introduction to AVL Trees

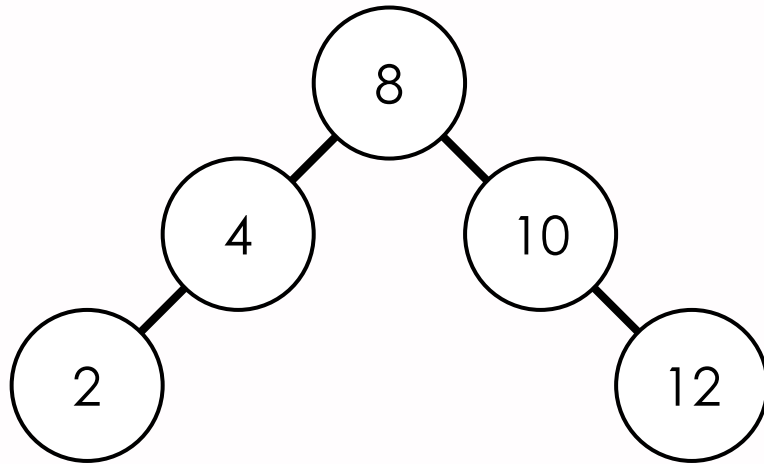
- Short for Adelson-Velsky and Landis tree
- Self-balancing binary search tree
- Determined by tree
  - Difference of tree height of the left subtree and the right subtree must be 0 or 1
  - When subtracting the heights, get the absolute value
  - If difference is higher than 1, then it is subject for rotation



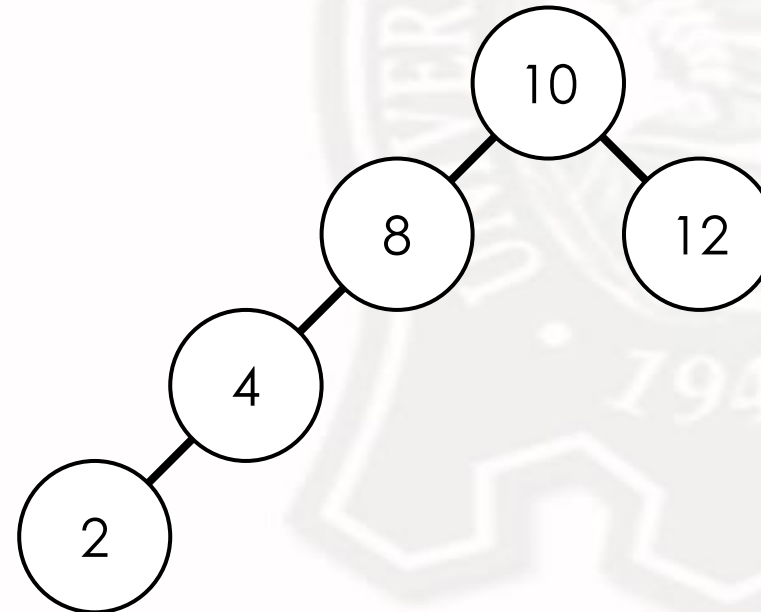


# Introduction to AVL Trees

Left subtree height: 2  
Right subtree height: 2  
 $2 - 2 = /0/ = \mathbf{0} \therefore \text{an AVL tree}$

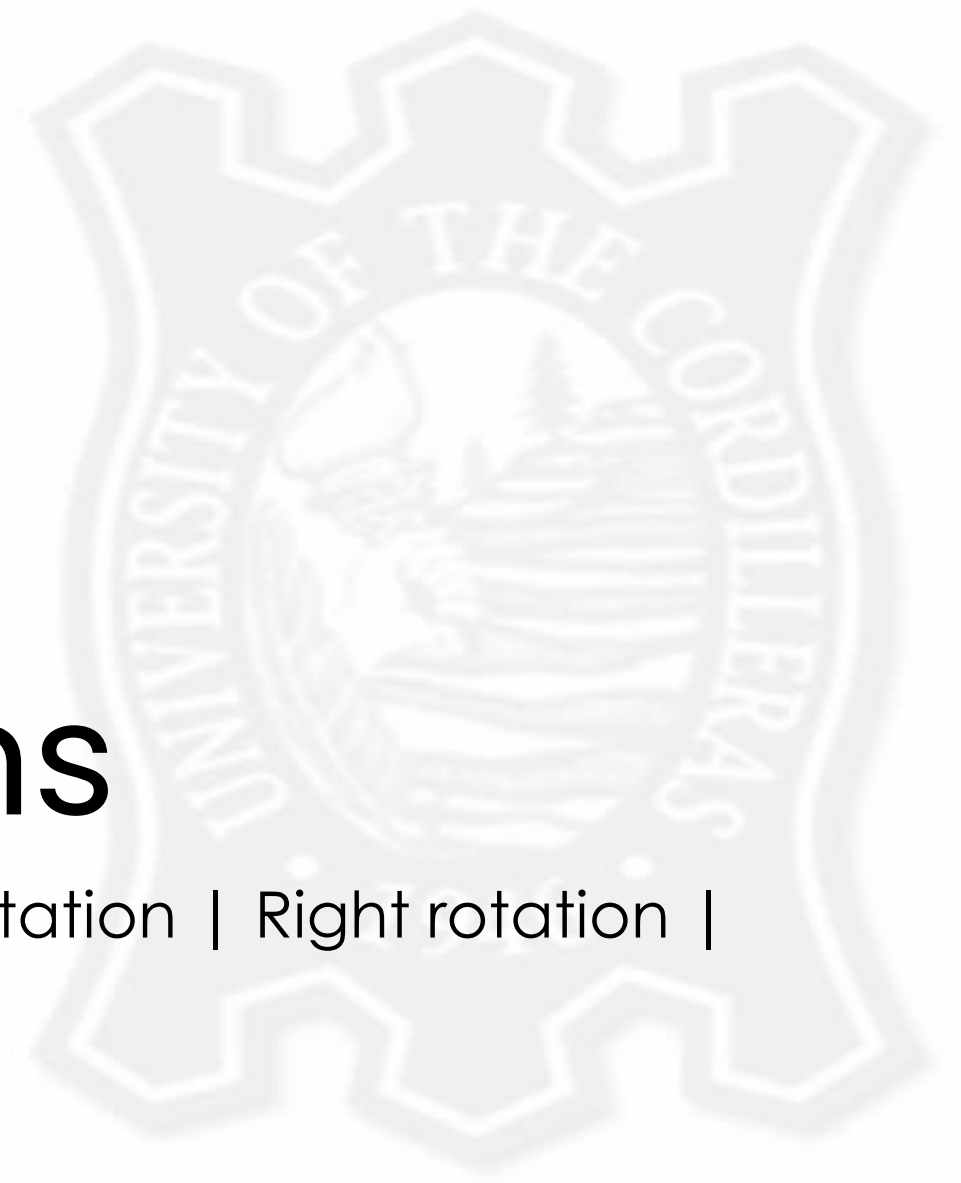


Left subtree height: 3  
Right subtree height: 1  
 $3 - 1 = /2/ = \mathbf{2} \therefore \text{not an AVL tree}$



# Actions in AVL Trees

- Searching, Insertion, and Deletion
  - Steps are the same as a binary search tree
  - Rotation may be done after insertion or deletion
- Rotation
  - Done if the tree is considered as unbalanced
  - Ensure that when rotating, values still follow the rules of BSTs



# AVL Tree Rotations

Identifying problems and violations | Left rotation | Right rotation |  
Left-right rotation | Right-left rotation



# Identifying Problems and Violations

- When doing rotations, it is necessary to determine the items that are the cause of the problem and the steps necessary to fix it
- Problem
  - Node that causes the imbalance of the AVL tree
  - Usually a leaf node
- Violation
  - Node that is subjected to rotation
  - Usually the grandfather of the problem node



# Identifying Problems and Violations

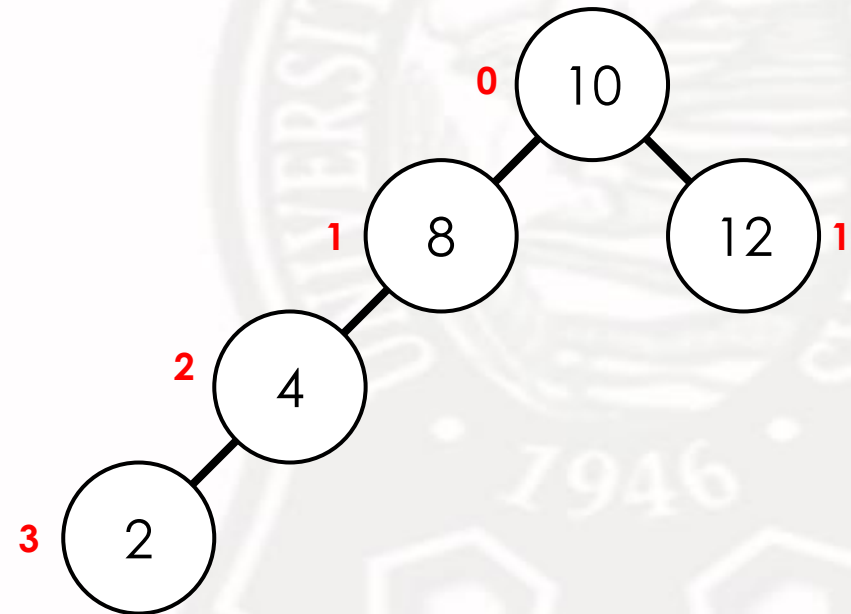
- Identify left and right subtree height
- Compute for the difference in subtree heights
- Determine the subtree with the longer tree height
- Identify the problem
- Identify the violation
- Determine the type of rotation necessary



# Identifying Problems and Violations

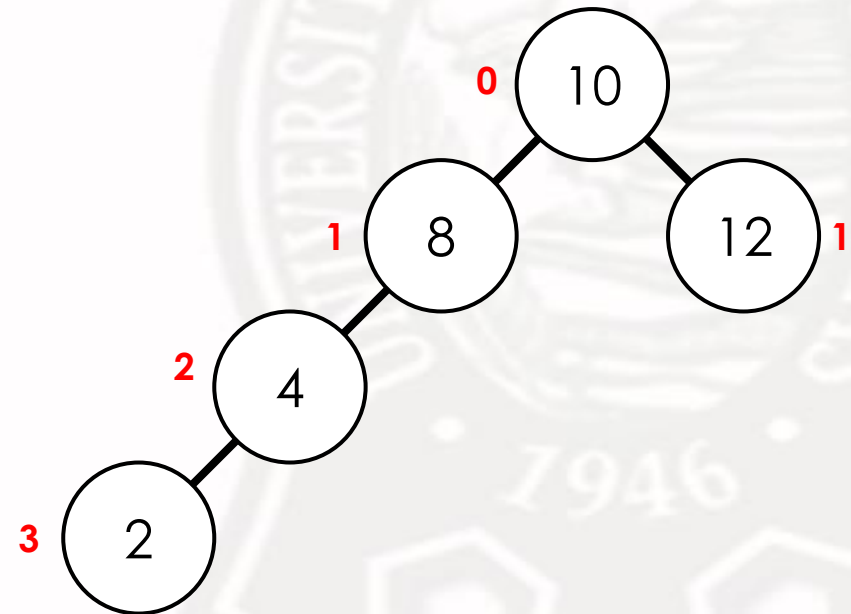
- ***Identify left and right subtree height***

- Root node is counted as 0
- Left subtree height (LSH) = 3
- Right subtree height (RSH) = 1



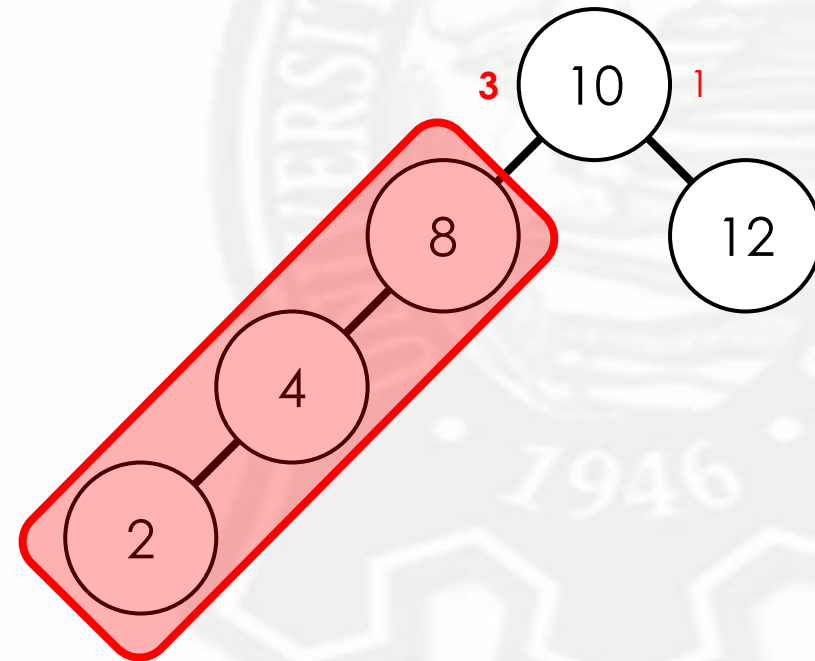
# Identifying Problems and Violations

- **Compute for the difference in subtree heights**
  - Difference =  $/LSH - RSH/$
  - $3 - 1 = /2/ = 2$
  - 2 is greater than 1 or 0; therefore it is unbalanced
  - Unbalanced BSTs are subject to rotation



# Identifying Problems and Violations

- Determine the subtree with the longer tree height
  - **Left subtree height = 3**
  - Right subtree height = 1

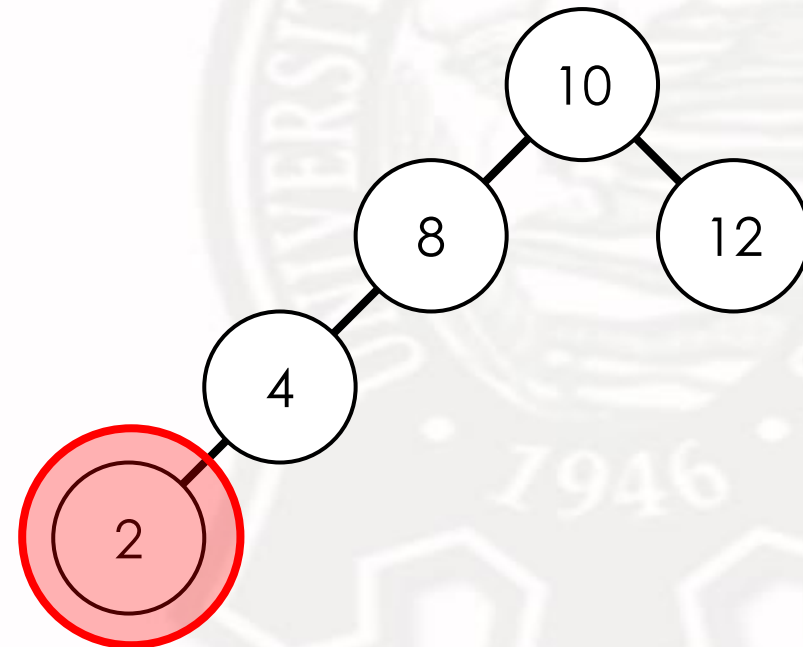




# Identifying Problems and Violations

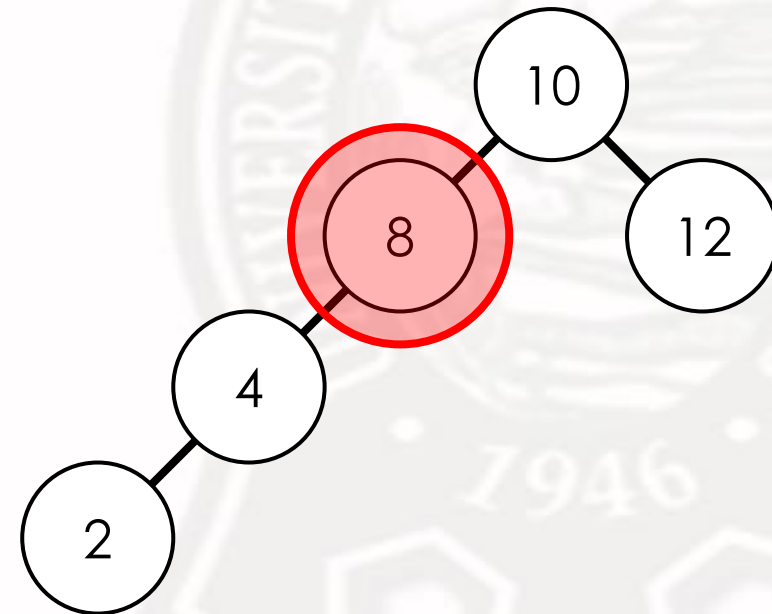
- ***Identify the problem***

- Problem should be on the subtree with longer tree height (left subtree)
- Problem is usually the leaf node
- Problem node = 2



# Identifying Problems and Violations

- Identify the violation
  - Violation is usually 2 nodes away from the leaf node
  - AKA “grandparent node”
  - Violation node = 8
- Determine the type of rotation necessary



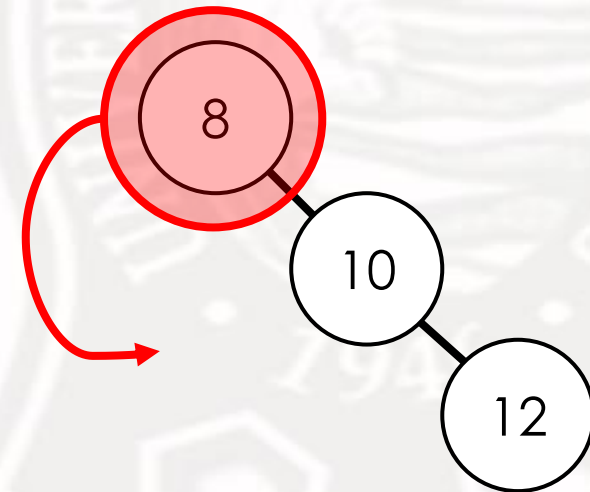
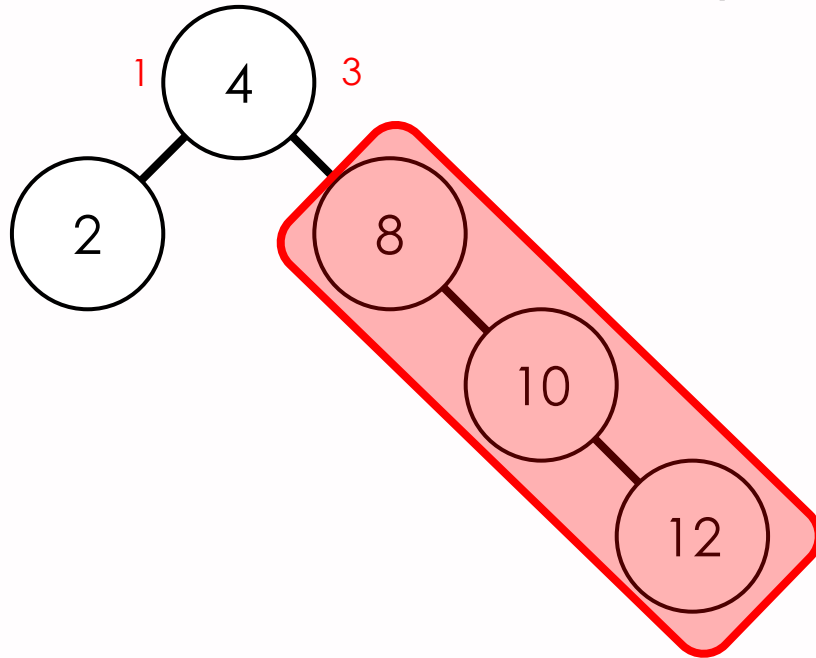
# Common AVL Tree Rotations

- Left rotation (LL)
- Right rotation (RR)
- Left-right rotation (LR)
- Right-left rotation (RL)



# Left Rotation

- Rotate the grandparent node to the left, becoming the left child node of the parent node



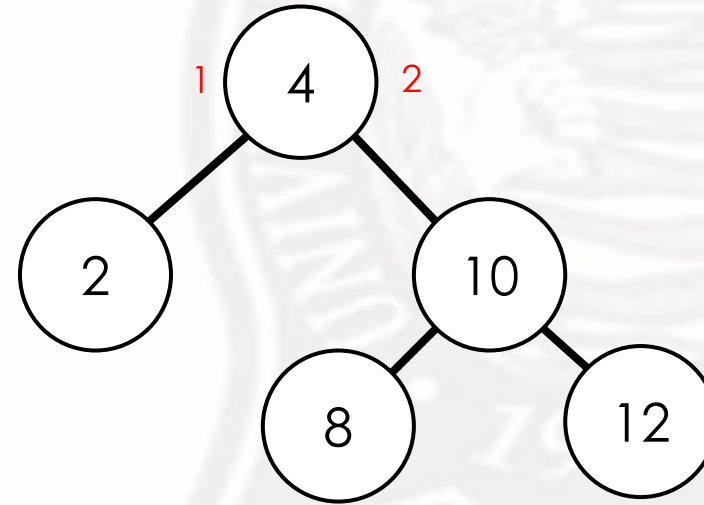
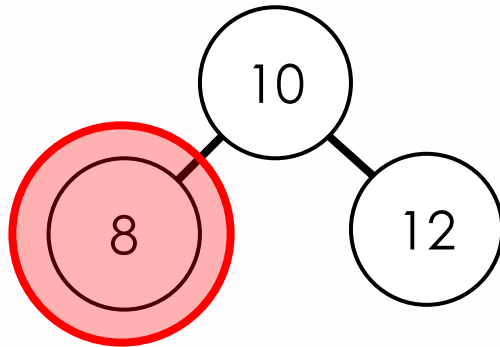
- Left subtree height: 1
- Right subtree height: 3

•  $1 - 3 = -2 \Rightarrow 2 \therefore$  NOT an AVL tree



# Left Rotation

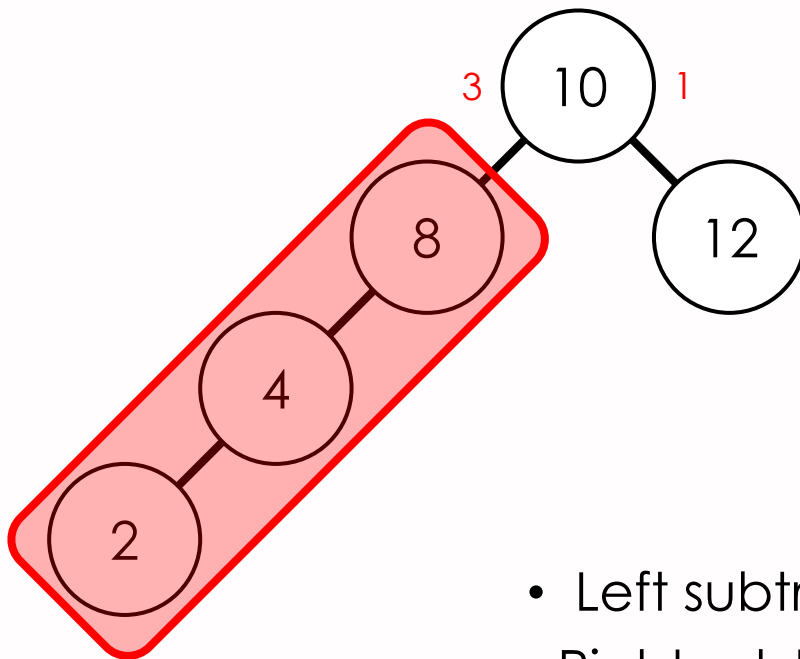
- Rotate the grandparent node to the left, becoming the left child node of the parent node



- Left subtree height: 1
- Right subtree height: 2
- $1 - 2 = -1 \neq 1 \therefore$  an AVL tree

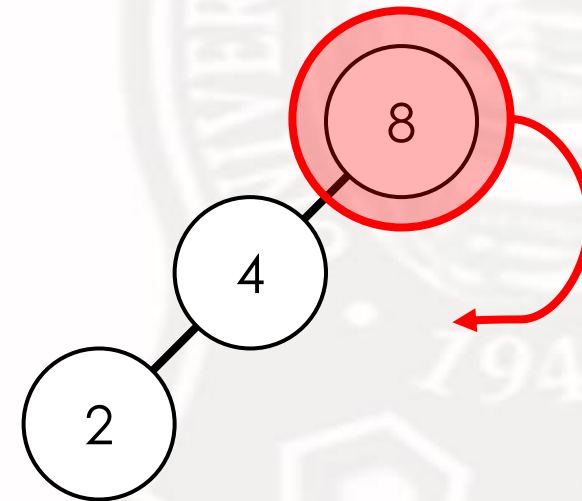
# Right Rotation

- Rotate the grandparent node to the right, becoming the right child node of the parent node



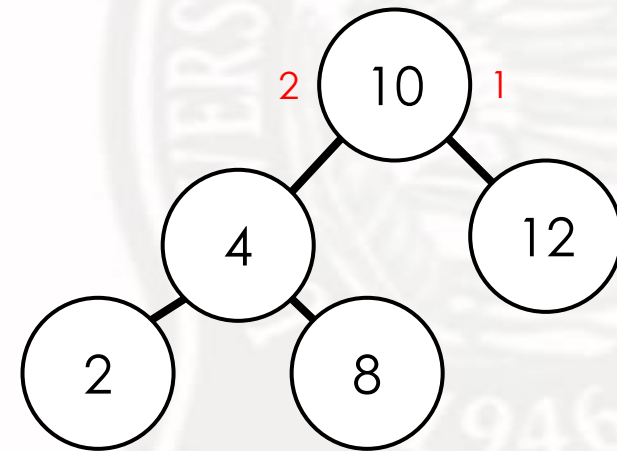
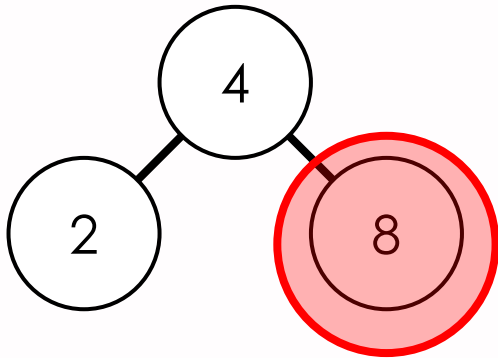
- Left subtree height: 3
- Right subtree height: 1

$3 - 1 = /2/ = 2 \therefore$  NOT an AVL tree



# Right Rotation

- Rotate the grandparent node to the right, becoming the right child node of the parent node

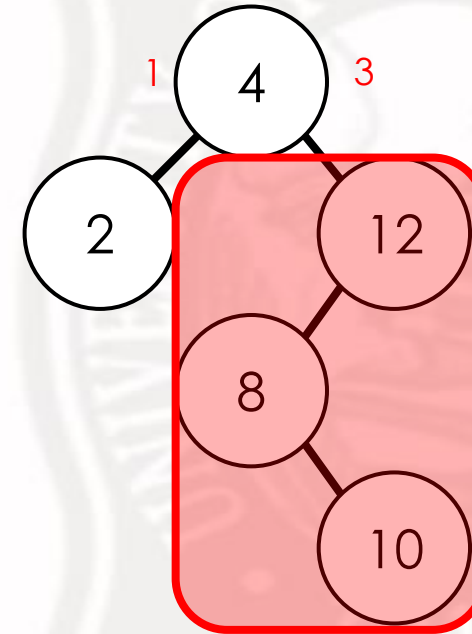


- Left subtree height: 2
- Right subtree height: 1
- $2 - 1 = 1 \therefore$  an AVL tree



# Left-Right Rotation

- Rotate the parent node to the left; rotate the problem node (grandchild node) to the right
- Perform right rotation
  - Rotate the grandparent node to the right, becoming the right child node of the parent node



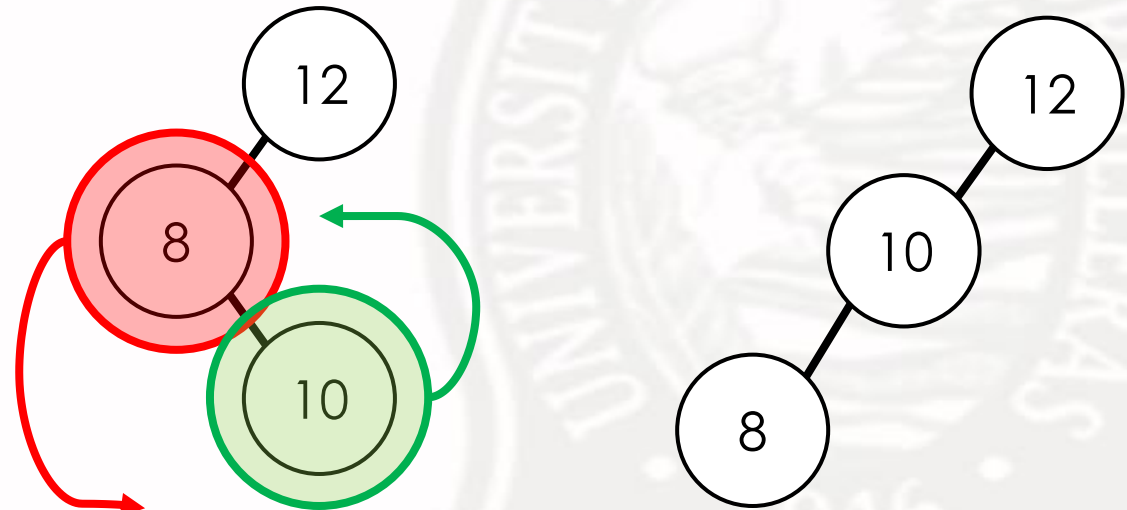
- Left subtree height: 1
- Right subtree height: 3
- $1 - 3 = -2 \neq 0 \therefore$  NOT an AVL tree





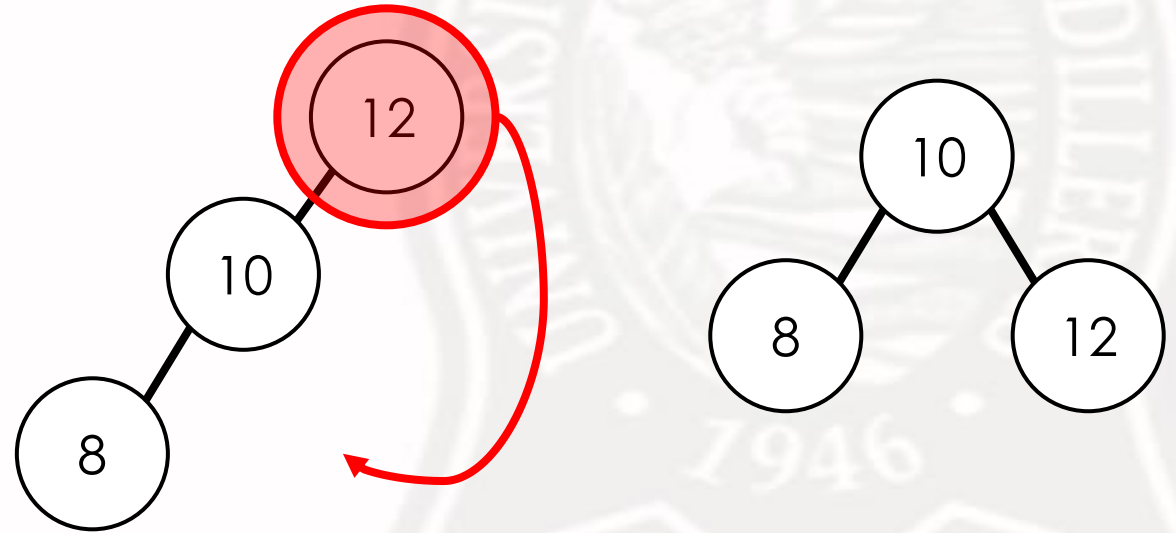
# Left-Right Rotation

- **Rotate the parent node to the left; rotate the problem node (grandchild node) to the right)**
- Perform right rotation
  - Rotate the grandparent node to the right, becoming the right child node of the parent node



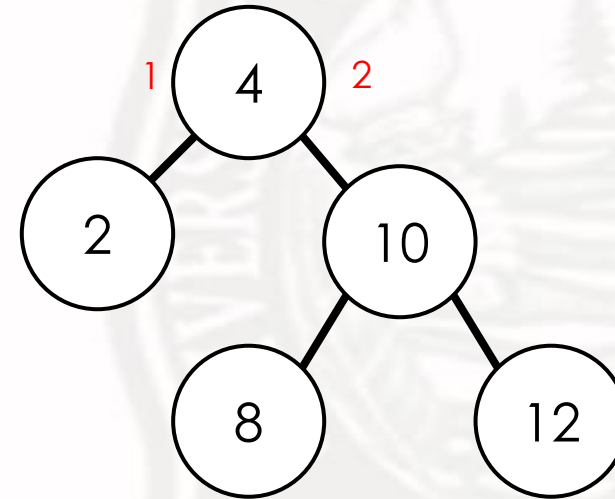
# Left-Right Rotation

- Rotate the parent node to the left; rotate the problem node (grandchild node) to the right
- **Perform right rotation**
  - Rotate the grandparent node to the right, becoming the right child node of the parent node



# Left-Right Rotation

- Rotate the parent node to the left; rotate the problem node (grandchild node) to the right
- Perform right rotation
  - Rotate the grandparent node to the right, becoming the right child node of the parent node

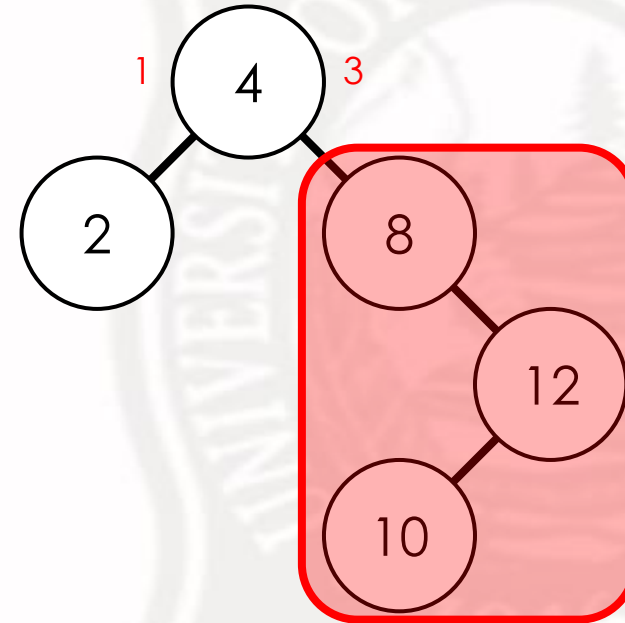


- Left subtree height: 1
- Right subtree height: 2
- $1 - 2 = -1 \neq \pm 1 \therefore$  an AVL tree



# Right-Left Rotation

- Rotate the parent node to the right; rotate the problem node (grandchild node) to the left
- Perform left rotation
  - Rotate the grandparent node to the left, becoming the left child node of the parent node

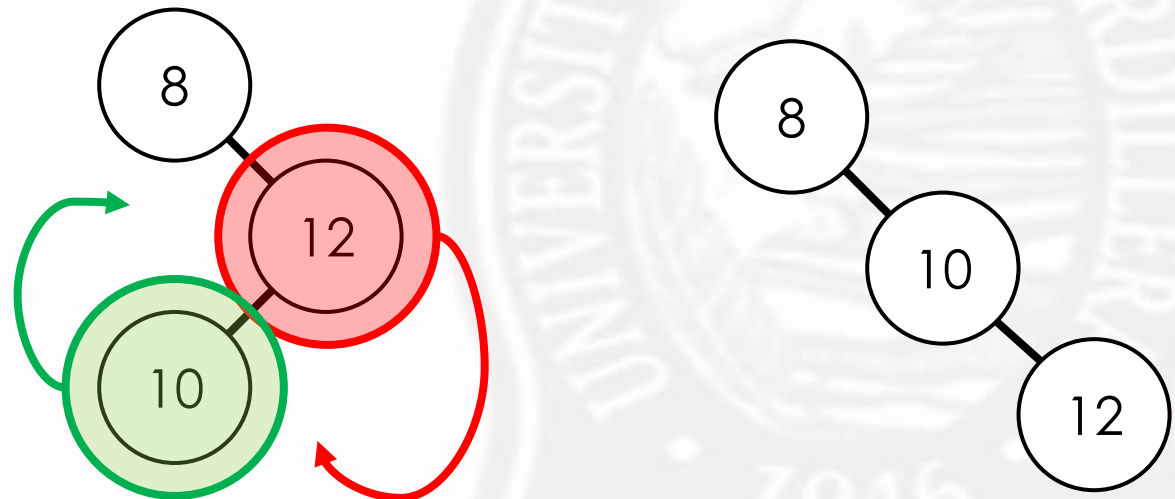


- Left subtree height: 1
- Right subtree height: 3
- $1 - 3 = -2 \Rightarrow \text{NOT an AVL tree}$



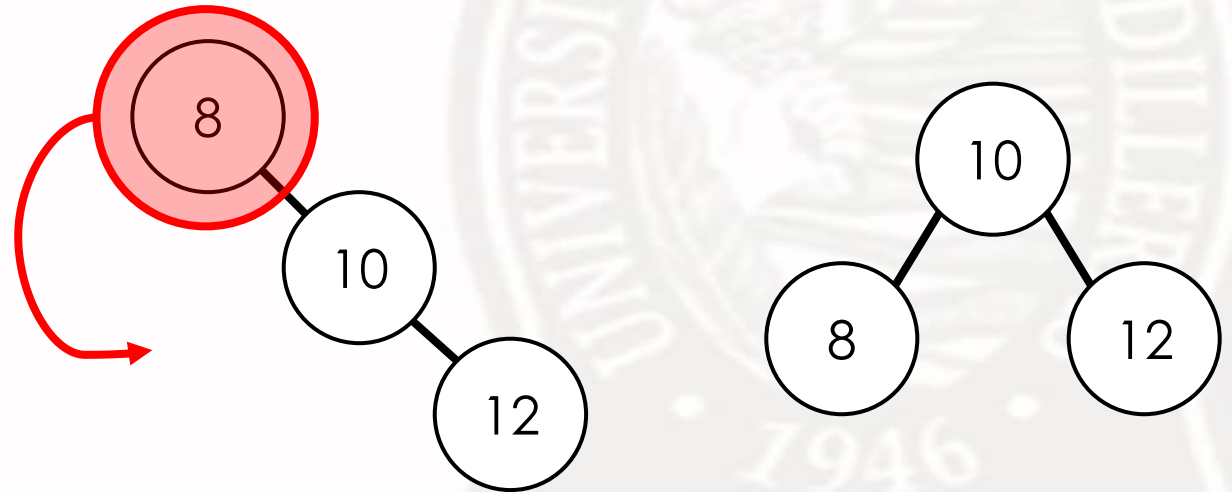
# Right-Left Rotation

- **Rotate the parent node to the right; rotate the problem node (grandchild node) to the left**
- Perform left rotation
  - Rotate the grandparent node to the left, becoming the left child node of the parent node



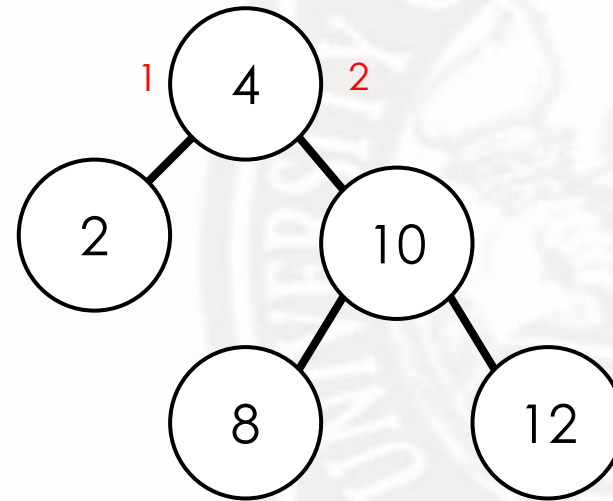
# Right-Left Rotation

- Rotate the parent node to the right; rotate the problem node (grandchild node) to the left
- Perform left rotation
  - Rotate the grandparent node to the left, becoming the left child node of the parent node



# Right-Left Rotation

- Rotate the parent node to the right; rotate the problem node (grandchild node) to the left
- Perform left rotation
  - Rotate the grandparent node to the left, becoming the left child node of the parent node



- Left subtree height: 1
- Right subtree height: 2
- $1 - 2 = -1 \Rightarrow$  an AVL tree

