

# RANDOMIZED LOW-RANK APPROXIMATION WITH ADAPTIVE SAMPLING

**Anonymous authors**

Paper under double-blind review

## ABSTRACT

The randomized singular value decomposition (rSVD) is an efficient algorithm for computing low-rank approximations to matrices through matrix-vector products with standard Gaussian vectors. Recently, Boullé and Townsend generalized the rSVD to leverage information on the right singular vectors of the target matrix  $A$  using Gaussian sketches with correlated entries. One open problem is whether this approach allows for the adaptive selection of the covariance matrix. We develop an adaptive sampling randomized singular value decomposition algorithm (AS-rSVD), which iteratively updates the covariance matrix after batches of rSVD applications. After each batch, the distribution of the random sampling vectors is updated to incorporate knowledge of the dominant right singular vectors of  $A$  obtained from the intermediate low-rank approximations. Our approach is demonstrated by a theoretical analysis and numerical experiments, and yields lower approximation errors than the rSVD and generalized rSVD while using the same number of matrix-vector products.

**NB: Check and cite papers by Chris and Cameron Musco on randomized SVD, such as (Musco & Musco, 2015). Make the distinction clear between our approach and the power method.**

## 1 INTRODUCTION

Computing low-rank matrix approximations through sketching has been a problem of interest at the intersection of computational linear algebra and machine learning for the past two decades (Drineas et al., 2006; Halko et al., 2011; Liberty et al., 2007; Martinsson et al., 2011; Musco & Musco, 2015; Nakatsukasa, 2020; Woodruff, 2014). In many real-world applications, it is often not possible to run experiments in parallel. Consider the following setting, there are a set of  $n$  inputs and  $m$  outputs, and there exists a PDE such it maps any set of inputs in  $\mathbb{C}^m \rightarrow \mathbb{C}^n$ . However, PDE experiments have intensive time and monetary requirements, e.g. aerodynamics (Fan et al., 2005), fluid dynamics (Lomax et al., 2001). Thus, after each experimental run, we want to sample a function such that in expectation, we will be exploring an area of the PDE, which we have the least knowledge of. The randomized SVD is a popular algorithm for low rank matrix approximation in the matrix vector product query model (Halko et al., 2011). The rSVD utilizes Gaussian sketching to approximate the range of  $A$  and then constructs the matrix  $QQ^*A$  where  $Q$  is an orthonormal basis of the range of the sketch  $AG$  for a standard normal matrix  $G$ . This method has been generalized recently by Boullé et al. (2022a) to allow for Gaussian input vectors with a general positive semi-definite covariance matrix. The current state-of-the-art algorithms for low-rank matrix approximation in the matrix-vector product model used a fixed covariance matrix structure.

Adaptively sampling vectors for matrix problems have been studied in detail by Sun et al. (2021). The theoretical properties of adaptively sampled matrix-vector queries for estimating the minimum eigenvalue of a Wishart matrix have been studied by Braverman et al. (2020). Their results are used by Bakshi et al. (2022) to develop lower bounds for rank-one low-rank matrix approximation with adaptive matrix-vector product queries. Adaptive sampling techniques for low-rank matrix approximation first appeared in finding a CUR matrix approximation (Frieze et al., 2004). Optimal column sampling for a CUR matrix decomposition received much attention in the literature (Deshpande et al., 2006; Deshpande & Vempala, 2006; Har-Peled, 2014). More recently, Paul et al. (2015) gave an adaptive algorithm for the column subset selection problem and proved it is possible to improve upon any relative-error column subset selection algorithm by adaptive sampling.

In this paper, we consider the adaptive setting where the algorithm  $\mathcal{A}$  chooses a vector  $\omega^{(k)}$  with access to the previous query vectors  $\omega^{(1)}, \dots, \omega^{(k-1)}$ , the matrix-vector products  $A\omega^{(1)}, \dots, A\omega^{(k-1)}$ , and the intermediate low-rank matrix approximations,  $Q_k Q_k^* A$ , where  $Q_k R_k$  is the economized QR decomposition of  $A\Omega^{(k)}$ , where  $\Omega^{(k)}$  is the concatenation of vectors  $\omega^{(1)}, \dots, \omega^{(k)}$ . The basis of our analysis is the idea of sampling vectors in the null space of the low-rank approximation. This idea has been introduced recently in Machine Learning in Wang et al. (2021) for training neural networks for sequential tasks. In a Bayesian sense, we want to maximize the expected information gain of the PDE in each iteration by sampling in the space where we have no information. This leads to the formulation of our iterative algorithm for sampling vectors for the low-rank approximation. Our algorithm utilizes the SVD computation of the low-rank approximation at each step to sample the next vector. Although there are runtime limitations, both in theory under certain conditions and many real-world matrices, our algorithm obtains a closer estimate of the optimal in the Frobenius norm.

### Main Contributions.

1. We develop a novel adaptive sampling algorithm for the low-rank matrix approximation problem in the matrix-vector product model which does not utilize prior information of  $A$ .
2. We provide a novel theoretical analysis for adaptive sampling in the matrix-vector product query model.
3. We derive improved relative-error bounds for the generalized randomized SVD for both the spectral and Frobenius norms (Boullé & Townsend, 2022, Thm. 2) and show the spectral-norm bound is also an improvement for the rSVD (Halko et al., 2011, Thm. 10.8).
4. We perform numerical experiments on real-world and synthetic matrices that confirm our theoretical claims.

## 2 BACKGROUND AND RELATED WORKS

The randomized SVD is a method to find an orthonormal matrix that captures the range of the top left singular space of an  $m \times n$  matrix  $A$  by multiplying it with a matrix with standard normal entries (Martinsson & Tropp, 2020). One first samples  $A$  at  $k + p$  Gaussian vectors  $\mathbf{x}_i \sim \mathcal{N}(\mathbf{0}, I)$  for  $1 \leq i \leq k + p$ , where  $k$  is the target rank,  $p$  is the oversampling parameter, and  $I$  denotes the  $n \times n$  identity matrix. One then calculates an orthonormal basis for the range by the economized QR decomposition and obtains  $QR = AX$ . Halko et al. (2011) that  $Q$  is a good approximation to the range of the top right singular space of  $A$  and thus the approximation  $QQ^* A$  is close to  $A$  in both the spectral and Frobenius norms. The analysis has also been extended to subsampled randomized trigonometric transform (SRTT) matrices (Boutsidis & Gittens, 2013). More recently, a generalized of the randomized SVD was proposed by Boullé & Townsend (2022) to allow for the columns of the Gaussian matrix to be sampled from a multivariate normal distribution with correlated covariance. Their results indicate that there exist covariance matrices that are able to obtain better approximation bounds than the standard rSVD. Correlated Gaussian sketching has also been studied for performing Nyström approximations (Persson et al., 2024).

## 3 RANDOMIZED SVD WITH ADAPTIVE SAMPLES

The adaptive range finder algorithm was proposed by Halko et al. (2011) as a method to guarantee a high accuracy low-rank approximation by sampling vectors one at a time until a stop condition is reached, which is not fixed as the sampled vectors are Gaussian, hence the algorithm is adaptive. It does not modify the distribution of the sampling vectors to reduce sample complexity or error. To clarify the distinction in the meaning of the term *adaptive*, we provide a formal definition of the problem we study.

**Definition 1** (Adaptive sampling). *With access to  $r\ell$  right matrix-vector product queries. Sample  $A\omega_i$  for  $1 \leq i \leq r\ell$  such that, after each matrix-vector query, one has access to the low-rank approximation  $Q_t Q_t^* A$  for  $t \in [r]$ . Obtain a matrix  $Q \in \mathbb{O}_{m, r\ell}$  using maximally  $r\ell$  left and right matrix-vector product queries to  $A$  to find  $Q$  such that*

$$\|A - QQ^* A\|_F \leq (1 + \epsilon) \min_{U: U^* U = I} \|A - UU^* A\|_F.$$

Our definition of *adaptive* is standard in the theoretical computer science literature (see e.g., Deshpande & Vempala (2006); Paul et al. (2015)), in that we are using our previous samples to inform how we sample our future matrix vector product queries. We define *continued* sampling as sampling all vectors at once.

### 3.1 DESCRIPTION OF THE ALGORITHM

**NB:** Rename the number of samples per round as batch size throughout the paper, and rename round as iteration. Define the two terms in this section.

**NB:** We need a more extensive description of the algorithm before the pseudo-code.

We follow the convention in Paul et al. (2015) to perform adaptive sampling in rounds. Algorithm 1 runs in multiple rounds. In each round, we use the information we have learned from the matrix to update the covariance matrix.

**Covariance Update in Algorithm 1.** In each round, we update our covariance matrix to be the projection matrix of the singular space of the low-rank approximation. It can be calculated either from an SVD calculation or from a pseudoinverse, as we have from an expansion of the SVD,

$$\widehat{V}_{t(k+p)} \widehat{V}_{t(k+p)}^* = \widehat{V}^{(t)} \widehat{\Sigma}^{(t)} + \widehat{U}^{(t)*} \widehat{U}^{(t)} \widehat{\Sigma}^{(t)} \widehat{V}^{(t)*} = \widehat{A}^{(t)} + \widehat{A}^{(t)} \quad (1)$$

Therefore, one can call a pseudoinverse procedure of SVD procedure to update the covariance matrix.

The pseudo-code for the optimal function sampling is given in Algorithm 1. Algorithm 1 is developed with the goal of minimizing the number of matrix-vector products necessary to obtain the same accuracy as the randomized SVD or generalized randomized SVD.

---

#### Algorithm 1 AS-rSVD, Adaptive sampling for low-rank matrix approximation

---

**input:** Target matrix  $A \in \mathbb{C}^{m \times n}$ , batch size  $\ell$ , number of iterations  $r$

**output:** Low-rank approximation  $\widehat{A}$  of  $A$  to minimize  $\|\widehat{A} - A\|_F$

- 1:  $\widetilde{A}_0 \leftarrow 0$
- 2: **for**  $t \in [r]$  **do**
- 3:   Form  $\Omega_t \in \mathbb{C}^{n \times \ell}$  with columns sampled i.i.d from  $\mathcal{N}(\mathbf{0}, C_t)$
- 4:    $Y \leftarrow [Y, A\Omega_t] \in \mathbb{C}^{m \times t\ell}$  and obtain the economized QR decomposition  $Y = QR$
- 5:   Update the low-rank approximation:  $\widehat{A}^{(t)} = \widehat{A}^{t-1} + Q_-^{t\ell} Q_-^{t\ell*} A$
- 6:   Calculate the SVD:  $\widehat{U} \widehat{\Sigma} \widehat{V}^* = Q^* A$
- 7:   Update the covariance matrix:  $C_t \leftarrow \widehat{V}_{k(t-1), \ell t} \widehat{V}_{k(t-1), \ell t}^*$

**return:**  $Q_t Q_t^* A$

---

**Complexity analysis.** In total, we sample  $k+p$  Gaussian vectors from  $t$  different Gaussian Distributions and therefore the matrix-matrix product  $A\Omega$  scales in  $O(tmn\ell)$ . We perform  $t$  QR factorizations which scale in  $O(tmn\ell)$ . We then must perform the SVD decomposition on the  $Q^* A \in \mathbb{C}^{m \times n}$  a total of  $t$  times which can be done on the order of  $O(tmn\ell)$ . Thus the total complexity can be observed as  $O(tmn\ell)$ . We note that the right singular vectors of  $QQ^* A$  are equivalent to the right singular vectors of  $Q^* A$ .

## 4 THEORETICAL ANALYSIS

In this section, we will first give the mathematical setup for the theoretical analysis. We next derive improvements to the Generalized Randomized SVD Approximation Relative Error Bounds. We utilize our improved approximation bounds to derive relative error bounds for adaptive sampling. The proofs of all results presented in this section are deferred to Appendix A.

For any integer  $t$ , We define  $[t]$  as the set of integers  $\{1, \dots, t\}$ . Let  $a \wedge b = \min(a, b)$  and  $a \vee b = \max(a, b)$  for any two real numbers  $a, b$ . We define  $\mathbb{O}_{n,k}$  as the set of all  $n \times k$  matrices with orthonormal columns, i.e.  $\{V \in \mathbb{C}^{n \times k} : V^* V = I_{k \times k}\}$ . For any vector  $\mathbf{x} \in \mathbb{C}^d$ , we denote the  $\ell_2$

norm as  $\|\mathbf{x}\|_2^2 = \sum_{i=1}^d |x_i|^2$ . We define  $\mathbb{S}^{d-1} = \{\mathbf{x} \in \mathbb{C}^d : \|\mathbf{x}\|_2 = 1\}$ . We use Big-O notation,  $y = O(x)$ , to denote there exists a positive constant  $C$  such that  $y \leq Cx$ . We define  $\mathbb{E}[X]$  as the expectation of the random variable  $X$  and  $\Pr\{A\}$  as the probability of event  $A$  occurring.

#### 4.1 LINEAR ALGEBRA PRELIMINARIES

**NB: This should be much shorter.** The Moore–Penrose Inverse Moore (1920); Penrose (1955) of  $X \in \mathbb{R}^{m \times n}$  is denoted as  $X^+ \in \mathbb{R}^{n \times m}$ . The matrix  $X^+$  satisfies the following four conditions,

$$XX^+X = X, \quad X^+XX^+ = X^+, \quad (XX^+)^* = XX^+, \quad (X^+X)^* = X^+X.$$

If  $X$  has full row-rank, then the pseudo-inverse can be given explicitly as  $X^+ = (X^*X)^{-1}X^*$ . For any matrix  $X$  with economized singular value decomposition  $X = U\Sigma V^*$ , the pseudoinverse is given as  $X^+ = V\Sigma^{-1}U^*$ . An orthogonal projector matrix is a Hermitian Matrix and satisfies  $P^2 = P$ . This property of the orthogonal projector matrix implies  $\mathbf{0} \preceq P \preceq I$ . Suppose  $Q$  represents an orthonormal basis of the column space of a matrix  $Y$ , then  $P_Y = QQ^*$  represents the unique projection matrix on to the range of  $Y$ . We follow a similar setup as previous literature. We factorize  $A$  as follows,

$$A = \begin{bmatrix} k & n-k \\ U_k & U_{k,\perp} \end{bmatrix} \begin{bmatrix} k & n-k \\ \Sigma_k & \Sigma_{k,\perp} \end{bmatrix} \begin{bmatrix} V_k^* \\ V_{k,\perp}^* \end{bmatrix} \begin{bmatrix} k \\ n-k \end{bmatrix} = \sum_{i=1}^n \sigma_i \mathbf{u}_i \mathbf{v}_i^*.$$

The trace of a matrix is given as  $\text{Tr}(A) = \sum_{i \in [n]} \sigma_i(A)$ . We furthermore define for  $\xi \in \{2, F\}$ ,

$$A_k = \min_{B: \text{rank}(B)=k} \|A - B\|_\xi.$$

The Eckart-Young-Mirsky Theorem Eckart & Young (1936); Mirsky (1960) tells us for the spectral and Frobenius norms that  $A_k = \sum_{i=1}^k \sigma_i \mathbf{u}_i \mathbf{v}_i^*$ . The spectral norm for a matrix  $A \in \mathbb{C}^{m \times n}$  and  $m \geq n$  is given as  $\|A\|_2^2 = \|A^*A\|_2 = \sigma_1^2(A)$ . The Frobenius Norm is given as  $\|A\|_F^2 = \text{Tr}(A^*A) = \sum_{i=1}^n \sigma_i^2(A)$ . Our theoretical arguments rely on the following proposition given by Halko et al. Halko et al. (2011).

**Proposition 2** (Conjugation Rule). *Suppose that  $M \succeq \mathbf{0}$ . Then for any conformal matrix  $A$ ,*

$$M \preceq N \implies A^*MA \preceq A^*NA$$

We also utilize a corollary of Weyl’s Inequality, a classical result in Perturbation Theory.

**Lemma 3** (Weyl’s Inequality). *Suppose  $A, B$  have real eigenvalues (e.g. they are Hermitian), then for any  $i \in [n]$ ,*

$$\lambda_i(A + B) \leq \lambda_i(A) + \lambda_1(B)$$

#### 4.2 GENERALIZED RANDOMIZED SVD

In this subsection we present our result for the Frobenius norm approximation bounds for the Generalized Randomized SVD presented by Boullé and Townsend Boullé & Townsend (2022). To our knowledge these are the first relative-error upper bounds for the spectral and Frobenius norms for the generalized rSVD. **NB: We need to remove the dependence on the dimension here using** (Persson et al., 2024, Lem. 2.2). **This gives you a tighter bound than Prop.. 19 that you can replace simply later on in the proof. To control the expectation term in the Lem. 2.2, you can use** (Persson et al., 2024, Lem. A.2).

**Theorem 4.** *Let  $A \in \mathbb{C}^{m \times n}$ , set  $k \geq 2$  an integer, an oversampling parameter  $p \geq 4$ . Let  $\Omega \in \mathbb{R}^{n \times k+p}$  represent the test matrix with columns sampled from  $\mathcal{N}(\mathbf{0}, C)$ . Suppose  $QR = A\Omega$  is the economized QR decomposition of  $A\Omega$ , then for all  $u > 1$ ,*

$$\|A - QQ^*A\|_F \leq \|\Sigma_{k,\perp}\|_F \left( 1 + \sqrt{6ut \log((n-k)(k+p))} \|K_{22}\|_2^{1/2} \sqrt{\text{Tr}(K_{11}^{-1})} \sqrt{\frac{1}{p+1}} \right),$$

with probability at least  $1 - \sqrt{2/\pi}(nm)^{1-u} - t^{-p}$  where

$$K = V^*CV = \begin{bmatrix} V_k^*CV_k & V_k^*CV_{k,\perp} \\ V_{k,\perp}^*CV_k & V_{k,\perp}^*CV_{k,\perp} \end{bmatrix} = \begin{bmatrix} K_{11} & K_{12} \\ K_{21} & K_{22} \end{bmatrix}.$$

**Remark 1.** A similar spectral norm bound can be developed similarly to the Frobenius norm in Theorem 4 and is derived in Appendix A.

The proof of Theorem 4 can be derived with the deterministic error bound of Theorem 9.1 in Halko et al. (2011) and combining Lemma 11 with both relations of Proposition 20. From Theorem 4, we infer that we obtain a better theoretical bound than the randomized SVD when  $C$  has better alignment with the top right singular space of  $A$  than the identity. As example, choosing  $C = V_k V_k^*$ , we obtain  $K_{22} = 0$  and we then obtain the optimal  $\|\Sigma_{k,\perp}\|_F^2$  Frobenius norm approximation error. The weakness in the bound when compared to the probabilistic bound in Halko et al. (2011) is the  $O(\log((n-k)(k+p)))$  term that is a result of Proposition 20. Our bounds presented in Theorem 4 are the strongest relative spectral and Frobenius norm error approximations to the generalized Randomized SVD presented in the literature to our knowledge. The difference between our bounds and the bounds derived in Theorem 2.4 of Persson et al. (2024) is that we lose the additive constant that has no inverse dependence on the oversampling parameter  $p$ , giving us the desired  $(1 + \epsilon)$  approximation results for low-rank approximation by choosing a sufficiently large oversampling parameter,  $p$ . We formalize our relative error bound in the following corollary.

**Corollary 5.** Let  $A \in \mathbb{C}^{m \times n}$ , set  $k \geq 1$  an integer, an oversampling parameter  $p \geq 4$ . Let  $\Omega \in \mathbb{R}^{n \times (k+p)}$  represent the test matrix with columns sampled from  $\mathcal{N}(\mathbf{0}, C)$ . Then, let  $QR = A\Omega$  represent the economized QR decomposition of  $A\Omega$ . Then if

$$p \geq (12/\epsilon) \log(n^2/\delta) \|K_{22}\|_2 \text{Tr}(K_{11}^{-1})$$

With probability exceeding  $1 - \delta - 0.0625$ ,

$$\|A - QQ^*A\|_F^2 \leq (1 + \epsilon) \|A - A_k\|_F^2$$

Corollary 5 develops a relative Frobenius norm error bound for the Generalized Randomized SVD and can be derived from elementary algebraic manipulations after an application of the first relation in Theorem 4. Since (Persson et al., 2024, Rem. 2.5) does not obtain a  $(1 + \epsilon)$  relative error bound, there exists a sufficiently large oversampling parameter  $p$  such that our bound gives a stronger error bound. Furthermore, our relative-error bound for the spectral norm is stronger than the relative-error spectral norm bound given by Halko et al. (Halko et al., 2011, Equation (1.9)). We are able to reduce the  $O(\sqrt{n-k})$  upper bound to  $O(\sqrt{\log(n-k)})$ . We formalize our improvement and give numerical illustrations in Appendix C.

#### 4.3 ADAPTIVE SAMPLING THEORETICAL ANALYSIS

In this section, we make theoretical connections from Adaptive sampling in the implicit matrix model to the more well-studied problem of Adaptive sampling in the column subselection problem. Adaptive sampling in works such as Deshpande & Vempala (2006) and Paul et al. (2015) for the column subselection problem (see e.g. Civil (2014) § 1 for a problem definition) sample columns from the residual matrix,  $A - Q_t Q_t^* A$ , at each iteration  $t$  instead of directly from  $A$ . Our first result is to formalize this statement and prove that indeed, adaptive low-rank matrix approximation in the implicit matrix model performs low-rank matrix approximation at each round. We require the following Lemma, which gives us theoretical insight into where we can derive advantages with adaptive sampling.

**Lemma 6.** Let  $\Omega = [\Omega_1, \Omega_2]$  and  $Y = A\Omega$ . Suppose  $QR$  is the economized QR decomposition of  $Y$ . Finally, for shorthand, define  $P_\perp = I - QQ^*$ , then for  $\xi \in \{2, F\}$ ,

$$\|A - P_Y A\|_\xi = \|P_\perp A - Q_- Q_-^* P_\perp A\|_\xi. \quad (2)$$

From Lemma 6, to minimize the RHS of Equation (2) we observe that  $Q_-$  is equal to the dominant left singular space of  $P_\perp A$ . From this result, we then see that at each iteration, the optimal vector query is the top right singular vector of  $P_\perp A$ . Thus, if one has knowledge of the matrix  $(I - P_{Y_t})A$  at any iteration, then performing a power iteration by sampling  $X$  as a Gaussian matrix, then calculate

$$X_q = (I - P_{Y_t})((I - P_{Y_t})^*(I - P_{Y_t}))^q X$$

Then, sample  $X_q$  as the sampling vectors for the next round.

**Theorem 7.** Let  $A \in \mathbb{C}^{m \times n}$  and let  $C \in \mathbb{C}^{n \times n}$  be a PSD covariance matrix. Let the sampling matrix be decomposed as  $\Omega = [\Omega, \Omega_-]$ , the matrix-matrix products  $Y = A\Omega$ ,  $Y_- = A\Omega_-$ , then let  $Q_+$  be the orthonormal matrix from an economized QR decomposition of  $[A\Omega_-, Y]$ . Let  $p \geq 4$  be the oversampling parameter, then with probability exceeding  $1 - \delta - t^{-p}$ ,

$$\|A - Q_+ Q_+^* A\|_F^2 \leq \epsilon(1 + \epsilon)\|A - A_k\|_F^2 + (1 + \epsilon)\|A - A_{2k}\|_F^2$$

From Theorem 7, we observe that our choice of covariance matrix reduces the number of matrix-vector products required to obtain the same accuracy as the Randomized SVD or even generalized randomized SVD.

**Comparison with continued sampling.** Consider the case where  $t = 2$ . In the continued sampling model, we sample all the  $2k$  vectors initially. Continued sampling incurs Frobenius Norm Error

$$\|A - Q_+ Q_+ A\|_F^2 \leq \left(1 + \frac{\epsilon}{2}\right)\|A - A_k\|_F^2$$

Next, we consider the bound given in Theorem 7. From some simple algebraic manipulations, we obtain when

$$\|A - A_{2k}\|_F^2 \leq \left(\frac{1 - \epsilon/2 - \epsilon^2}{1 + \epsilon}\right)\|A - A_k\|_F^2 \leq (1 - \epsilon)\|A - A_k\|_F^2$$

The adaptive bound is stronger than the classical bound given sufficient singular decay when the goal is a  $(1 + \epsilon)\|A - A_k\|_F^2$  approximation. Even stronger, is we can note there exists an  $\epsilon$  for which the adaptive bound is always stronger than the classical bound,

$$\sum_{i=2k+1}^n \sigma_i^2 \leq (1 - \epsilon) \sum_{i=k+1}^n \sigma_i^2 \iff \epsilon \sum_{i=k+1}^n \sigma_i^2 \leq \sum_{i=k+1}^{2k} \sigma_i^2$$

Therefore, we see there exists a sufficiently small  $\epsilon$  that gives a stronger  $(1 + \epsilon)\|A - A_k\|_F$  approximation. We now consider the analysis for  $T$  steps.

**Theorem 8.** Let  $A \in \mathbb{C}^{m \times n}$  and let  $C_i \in \mathbb{C}^{n \times n}$  for all  $i \in [T]$  be PSD covariance matrices. The sampling matrix is given as  $\Omega = [C_1^{1/2}G_1, \dots, C_t^{1/2}G_t]$ . Let  $QR$  be the economized QR decomposition of  $A\Omega$ , then

$$\|A - QQ^* A\|_F^2 \leq \|A - A_{tk}\|_F^2 + \epsilon \sum_{i=1}^{T-1} (1 + \epsilon)^{T-i} \|A - A_{ik}\|_F^2,$$

with probability exceeding  $1 - T(\delta + t^{-p})$ .

The theorem can be proved by induction, see Theorem 1 in Paul et al. (2015). Interestingly, without any algorithmic modifications to the randomized SVD, the fact that  $Q$  is an orthonormal basis of  $A\Omega$  gives us the same result as in Paul et al. (2015) without explicit knowledge of the matrix  $A - YY^+ A$ .

#### 4.4 ANALYSIS OF AS-RSVD

Our covariance update in Algorithm 1 of Algorithm 1 does not align with the theory we have developed in Section 4.3. Consider Lemma 6, it is optimal to sample in the dominant right singular subspace of  $A - P_Y A$ . However, in the implicit matrix problem, we do not have access to  $A - P_Y A$ . We first show for one round of our adaptive procedure in Algorithm 1, it is sufficient to sample in our estimation of the dominant right singular subspace of  $A$ .

**Theorem 9.** Consider one round of Algorithm 1 with a covariance matrix  $C \in \mathbb{C}^{n \times n}$ . Suppose  $G_1, G_2 \in \mathbb{C}^{n \times (k+p)}$  are Gaussian Matrices. If  $\widehat{V}_{k+p}^* G_1$  and  $V_k^* C^{1/2} G_2$  have full-row rank, then for  $\xi \in \{2, F\}$ , with probability at least 0.95,

$$\|A - QQ^* A\|_\xi \lesssim \left(1 + \sqrt{8 \log(10(k+p))} \|V_{k,\perp}^* C^{1/2}\|_2 \|(V_k^* C^{1/2})^+\|_2 \frac{e\sqrt{k+p}}{p+1}\right) \|A - A_{2k}\|_\xi$$



We can now note that  $(1 + \epsilon)^2 = 1 + O(\epsilon)$ . Then, from ideas in subspace perturbation theory (see e.g. O’Rourke et al. (2018)), the sample complexity to obtain a  $(1 + \epsilon)\|A - A_{2k}\|_F$  is reduced. We will now extend [Theorem 9](#) to multiple rounds.

**Theorem 10.** *Consider  $t$  rounds of [Algorithm 1](#) with covariance matrices  $C_i \in \mathbb{C}^{n \times n}$  for  $i \in [t]$ . Suppose  $\Psi_i \in \mathbb{C}^{n \times (k+p)}$  for  $i \in [t]$  are Gaussian Matrices. If  $\widehat{V}_{k+p}^* G_1$  and  $V_k^* C^{1/2} G_2$  have full-row rank, then for  $\xi \in \{2, F\}$ , with probability at least 0.95,*

$$\|A - QQ^*A\|_\xi \lesssim \prod_{i=1}^t \left( 1 + \sqrt{8 \log(10n)} \|V_{\mathcal{I}_i, \perp}^* C_i^{1/2}\|_2 \|(V_{\mathcal{I}_i}^* C_i^{1/2})^+\|_2 \frac{e\sqrt{k+p}}{p+1} \right) \|A - A_{tk}\|_\xi$$

[Theorem 10](#) implies we can first sample  $k + p$  standard Gaussian vectors. Then, at the  $i$ th iteration, we can sample vectors that align with  $V_{\mathcal{I}_i}$ . When given the intermediate low-rank matrix approximations, from obtaining the SVD, we can then sample our approximation of  $V_{\mathcal{I}_i}$ . We have nearly arrived at the covariance update given in [Algorithm 1](#). We next note in practice minimizing  $\|A - QQ^*A\|_\xi$  is not equivalent to minimizing  $\epsilon$  for  $(1 + \epsilon)\|A - A_k\|_\xi$ . Therefore, it is often the case even if we know the optimal covariance matrix  $V_k V_k^*$  it might not be optimal since  $\|A - A_k\|_\xi \geq \Omega(\|A - A_{k+s}\|_\xi)$  where  $s$  is an integer. Hence at each iteration it is not optimal to choose a rank- $k$  covariance matrix.

**Remark 2** (Randomized Nyström approximation with adaptive samples). *When the matrix  $A \in \mathbb{R}^{n \times n}$  is symmetric positive semi-definite (SPSD), the randomized Nyström approximation is more suitable than the randomized SVD (Tropp & Webber, 2023). It is given as  $\hat{A} := A\Omega(\Omega^*A\Omega)^+ \Omega^*A \approx A$ . While analyzing adaptive sampling for computing Nyström approximations is notoriously challenging (Gittens & Mahoney, 2016), Persson et al. (2024) recently generalized the randomized Nyström approximation to allow for a general covariance matrix. Then, our algorithm and analysis should extend naturally to improve Nyström approximations with adaptive samples.*

## 5 NUMERICAL EXPERIMENTS

In this section, we evaluate the performance of our adaptive sampling algorithm on performing low-rank approximations to matrices (issued from the discretization of differential operators or standard benchmarks) and compare it against the non-adaptive rSVD and generalized rSVD. Each point represents one round of sampling. We define

$$\text{OPT} = \min_{Z \in \mathbb{C}^{m \times t(k+p)}} \|A - ZZ^*A\|_F$$

From an application of the Eckart-Young-Mirsky Theorem for the Frobenius Norm Eckart & Young (1936); Mirsky (1960), we find that the minimizing  $Z = U_{t(k+p)}$  and therefore,  $\text{OPT} = \sqrt{\sum_{i=t(k+p)+1}^n \sigma_i^2(A)}$ .

**NB:** This needs more discussion and should be organized as “Experiment 1: Description of the figure and discussion of the results, Experiment 2:...”

**Experiment 1: Learning an Inverse Differential Operator.** In our first experiment we attempt to learn the discretized  $10^3 \times 10^3$  matrix of the inverse of the following differential operator:

$$\mathcal{L}u = \frac{\partial^2 u}{\partial x^2} - 100 \sin(5\pi x) u, \quad x \in [0, 1] \quad (3)$$

Learning the inverse operator of a PDE is equivalent to learning the Green’s Function of a PDE. This has been theoretically proven for certain classes of PDEs (Linear Parabolic Boullé et al. (2022b); Boullé & Townsend (2023)) as the inverse Differential operator is compact. Furthermore, there are multiple works suggesting the learning of inverse differential operators is provably efficient Boullé et al. (2022a; 2023). We observe in [Fig. 1\(a\)](#), even without prior knowledge of the dominant right singular space of  $A$ , after approximately 280 matrix-vector products, AS-rSVD learns a better low-rank approximation with respect to the Frobenius norm than the generalized rSVD which utilizes prior information on the Green’s function for inverse Laplacian PDE.

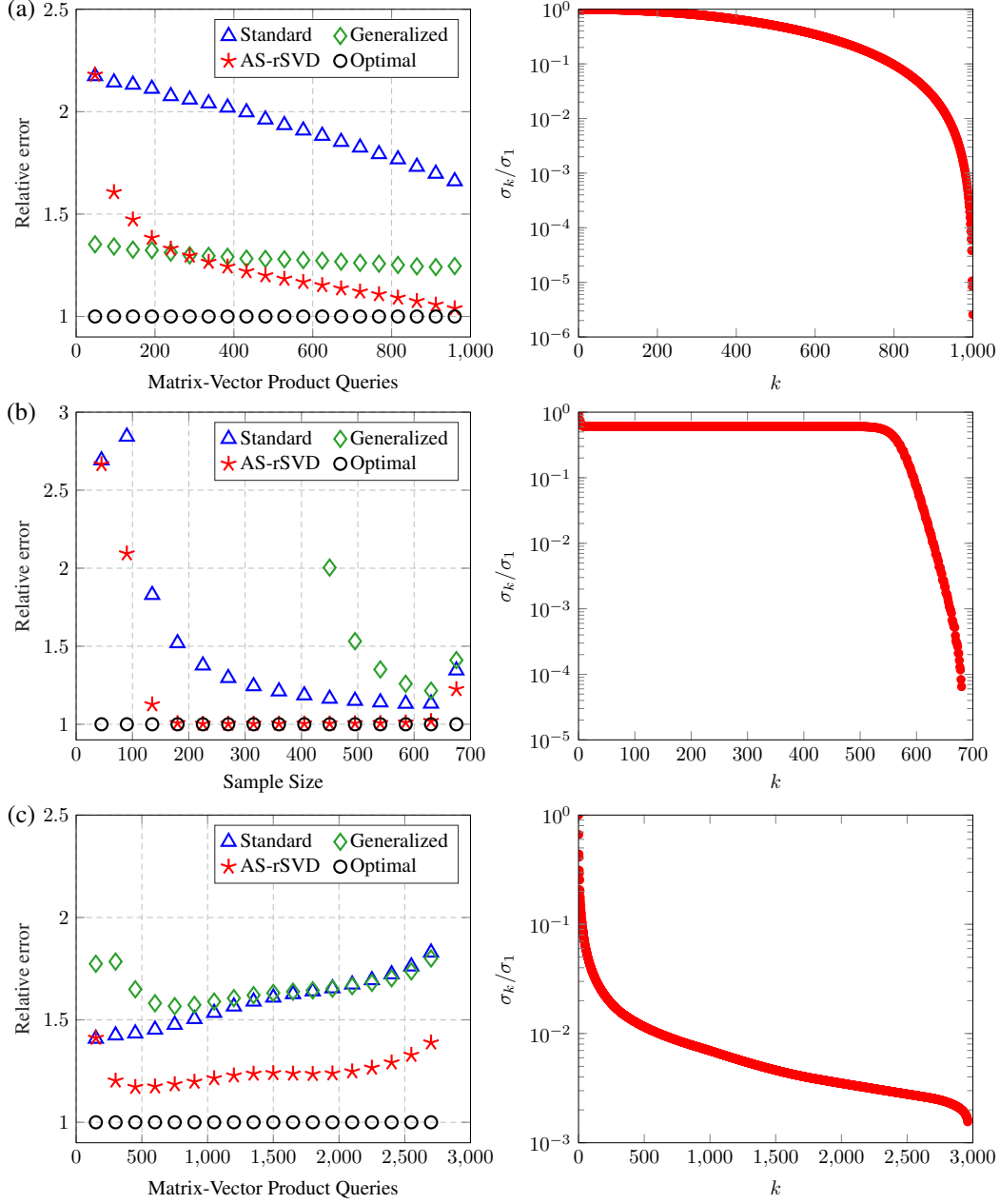


Figure 1: Left panels: Relative error (in the Frobenius norm) of the low-rank matrix approximation with respect to the number of matrix-vector products. Standard denotes the rSVD with i.i.d. Gaussian input vectors (Halko et al., 2011) while generalized uses a prior covariance matrix (Boullé & Townsend, 2022). Right panels: Singular value decay of the matrix. (a) Inverse of a discretized Laplacian differential operator (see Eq. (3) and (Boullé & Townsend, 2022, Fig. 2)),  $k = 8, p = 16, t = 20$ . (b) Matrix fs-680-1 from Davis & Hu (2011),  $k = 16, p = 32, t = 20$ . (c) Matrix pde2961 from Davis & Hu (2011),  $k = 50, p = 100, t = 18$ . All points in the plots for Fig. 1 are the average over 10 randomized runs.



**Experiment 2: Learning an Inverse Differential Operator in the Real World.** We attempt to learn the inverse of the discretized differential operator given in matrix `pde2961`, sourced from the TAMU Matrix Suite Davis & Hu (2011). `pde2961` is the matrix associated with a Model Partial Differential Equations Problem. In Fig. 1(b), we find that Algorithm 1 obtains nearly-optimal Frobenius norm error from 175 to 650 matrix-vector product queries.

**Experiment 3: Learning a forward matrix in the Real World.** We learn the a low rank matrix approximation of `fs-680-1`, also sourced from the TAMU Sparse Matrix Suite Davis & Hu (2011). `fs-680-1` is the associated matrix for a chemical kinetics problem. Algorithm 1 performs well in large matrices after considering the results in Fig. 1(c). The `pde2961` matrix is of size  $2961 \times 2961$  and Algorithm 1 outperforms rSVD and the generalized rSVD.

We find that in general, Algorithm 1 works well in a wide variety of singular value spectra. This can be concluded from observing that Fig. 1 display a diverse range of singular value spectra and cover both real and synthetic scenarios. Our experiments on real-world matrices are promising and indicate that our algorithm and implementation can be used in real-world applications of learning low-rank approximations of matrices that are only accessible via matrix-vector products.

**NB: Need a timing comparison in the Appendix (with a comment in the main text) and the following sentence:**

## 6 CONCLUSIONS

We have theoretically and empirically analyzed a novel Covariance Update to iteratively construct the sampling matrix,  $\Omega$  in the Randomized SVD algorithm. We introduce a new adaptive sampling framework for low-rank matrix approximation when the matrix is only accessible by matrix-vector products by giving the algorithm access to intermediate low-rank matrix approximations. Our covariance update for generating sampling vectors and functions find applications in PDE learning (Boullé et al., 2022a; Brunton et al., 2020). Numerical Experiments indicate without prior knowledge of the matrix, we are able to obtain superior performance to the rSVD, and the generalized rSVD with covariance matrix utilizing prior information of the right singular vectors of  $A$ . Theoretically, we provide an analysis of our update extended to  $k$ -steps and show under certain singular value decay conditions, we obtain better performance expectation.

## REFERENCES

- Ainesh Bakshi, Kenneth L Clarkson, and David P Woodruff. Low-rank approximation with  $1/\epsilon^{1/3}$  matrix-vector products. In *Proceedings of the 54th Annual ACM SIGACT Symposium on Theory of Computing*, pp. 1130–1143, 2022.
- Roger L Berger and George Casella. *Statistical inference*. Duxbury, 2001.
- Nicolas Boullé and Alex Townsend. A generalization of the randomized singular value decomposition. In *International Conference on Learning Representations*, 2022.
- Nicolas Boullé and Alex Townsend. Learning elliptic partial differential equations with randomized linear algebra. *Found. Comput. Math.*, 23(2):709–739, Apr 2023.
- Nicolas Boullé, Christopher J. Earls, and Alex Townsend. Data-driven discovery of Green’s functions with human-understandable deep learning. *Sci. Rep.*, 12(1):4824, 2022a.
- Nicolas Boullé, Seick Kim, Tianyi Shi, and Alex Townsend. Learning Green’s functions associated with time-dependent partial differential equations. *J. Mach. Learn. Res.*, 23(1):9797–9830, 2022b.
- Nicolas Boullé, Diana Halikias, and Alex Townsend. Elliptic PDE learning is provably data-efficient. *Proc. Natl. Acad. Sci. U.S.A.*, 120(39):e2303904120, 2023.
- Christos Boutsidis and Alex Gittens. Improved matrix algorithms via the subsampled randomized Hadamard transform. *SIAM J. Matrix Anal. Appl.*, 34(3):1301–1340, 2013.
- Christos Boutsidis, Petros Drineas, and Malik Magdon-Ismail. Near-optimal column-based matrix reconstruction. *SIAM J. Comput.*, 43(2):687–717, 2014.

- Mark Braverman, Elad Hazan, Max Simchowitz, and Blake Woodworth. The gradient complexity of linear regression. In *Conference on Learning Theory*, pp. 627–647, 2020.
- Steven L Brunton, Bernd R Noack, and Petros Koumoutsakos. Machine learning for fluid mechanics. *Annu. Rev. Fluid Mech.*, 52:477–508, 2020.
- Ali Civril. Column subset selection problem is ug-hard. *J. Comput. Syst. Sci.*, 80(4):849–859, 2014.
- Timothy A Davis and Yifan Hu. The University of Florida sparse matrix collection. *ACM Trans. Math. Softw.*, 38(1):1–25, 2011.
- Amit Deshpande and Santosh Vempala. Adaptive sampling and fast low-rank matrix approximation. In *International Workshop on Approximation Algorithms for Combinatorial Optimization*, pp. 292–303. Springer, 2006.
- Amit Deshpande, Luis Rademacher, Santosh S Vempala, and Grant Wang. Matrix approximation and projective clustering via volume sampling. *Theory Comput.*, 2(1):225–247, 2006.
- Petros Drineas, Ravi Kannan, and Michael W Mahoney. Fast monte carlo algorithms for matrices ii: Computing a low-rank approximation to a matrix. *SIAM Journal on computing*, 36(1):158–183, 2006.
- Carl Eckart and Gale Young. The approximation of one matrix by another of lower rank. *Psychometrika*, 1(3):211–218, 1936.
- Hui-Yuan Fan, George S Dulikravich, and Zhen-Xue Han. Aerodynamic data modeling using support vector machines. *Inverse Probl. Sci. En.*, 13(3):261–278, 2005.
- Alan Frieze, Ravi Kannan, and Santosh Vempala. Fast monte-carlo algorithms for finding low-rank approximations. *J. ACM*, 51(6):1025–1041, 2004.
- Alex Gittens and Michael W Mahoney. Revisiting the nyström method for improved large-scale machine learning. *J. Mach. Learn. Res.*, 17(1):3977–4041, 2016.
- Yehoram Gordon. Some inequalities for gaussian processes and applications. *Israel Journal of Mathematics*, 50:265–289, 1985.
- Nathan Halko, Per-Gunnar Martinsson, and Joel A. Tropp. Finding structure with randomness: Probabilistic algorithms for constructing approximate matrix decompositions. *SIAM Rev.*, 53(2): 217–288, 2011.
- Sariel Har-Peled. Low rank matrix approximation in linear time. *arXiv preprint arXiv:1410.8802*, 2014.
- Edo Liberty, Franco Woolfe, Per-Gunnar Martinsson, Vladimir Rokhlin, and Mark Tygert. Randomized algorithms for the low-rank approximation of matrices. *Proceedings of the National Academy of Sciences*, 104(51):20167–20172, 2007.
- Harvard Lomax, Thomas H Pulliam, David W Zingg, Thomas H Pulliam, and David W Zingg. *Fundamentals of computational fluid dynamics*. Springer, 2001.
- Per-Gunnar Martinsson and Joel A Tropp. Randomized numerical linear algebra: Foundations and algorithms. *Acta Numerica*, 29:403–572, 2020.
- Per-Gunnar Martinsson, Vladimir Rokhlin, and Mark Tygert. A randomized algorithm for the decomposition of matrices. *Applied and Computational Harmonic Analysis*, 30(1):47–68, 2011.
- Leonid Mirsky. Symmetric gauge functions and unitarily invariant norms. *Q. J. Math.*, 11(1):50–59, 1960.
- Eliakim H Moore. On the reciprocal of the general algebraic matrix. *Bulletin of the american mathematical society*, 26:294–295, 1920.

- Cameron Musco and Christopher Musco. Randomized block krylov methods for stronger and faster approximate singular value decomposition. In *Advances in neural information processing systems*, volume 28, 2015.
- Yuji Nakatsukasa. Fast and stable randomized low-rank matrix approximation. *arXiv preprint arXiv:2009.11392*, 2020.
- Sean O’Rourke, Van Vu, and Ke Wang. Random perturbation of low rank matrices: Improving classical bounds. *Linear Algebra and its Applications*, 540:26–59, 2018. doi: 10.1016/j.laa.2017.11.014. URL <https://doi.org/10.1016/j.laa.2017.11.014>.
- Saurabh Paul, Malik Magdon-Ismail, and Petros Drineas. Column selection via adaptive sampling. In *Advances in Neural Information Processing Systems*, volume 28, 2015.
- Roger Penrose. A generalized inverse for matrices. *Mathematical proceedings of the Cambridge philosophical society*, 51(3):406–413, 1955.
- David Persson, Nicolas Boullé, and Daniel Kressner. Randomized Nyström approximation of non-negative self-adjoint operators. *arXiv preprint arXiv:2404.00960*, 2024.
- Philippe Rigollet and Jan-Christian Hütter. High-dimensional statistics. *arXiv preprint arXiv:2310.19244*, 2023.
- Erhard Schmidt. Zur theorie der linearen und nichtlinearen integralgleichungen. *Math. Ann.*, 63(4): 433–476, 1907.
- Xiaoming Sun, David P Woodruff, Guang Yang, and Jialin Zhang. Querying a matrix through matrix-vector products. *ACM Transactions on Algorithms (TALG)*, 17(4):1–19, 2021.
- Joel A Tropp and Robert J Webber. Randomized algorithms for low-rank matrix approximation: Design, analysis, and applications. *arXiv preprint arXiv:2306.12418*, 2023.
- Shipeng Wang, Xiaorong Li, Jian Sun, and Zongben Xu. Training networks in null space of feature covariance for continual learning. In *Conference on Computer Vision and Pattern Recognition*, pp. 184–193, 2021.
- David P. Woodruff. Sketching as a tool for numerical linear algebra. In *Foundations and Trends in Theoretical Computer Science*, volume 10, pp. 1–157. Now Publishers, Inc., 2014.

## A PROOFS

**NB:** This section should be better organized and not just a succession of proofs.

In this section, we prove the results described in the main text.

### A.1 DISTRIBUTION OF $V^*\Omega$

In this section we will derive the distribution of  $V^*\Omega$ .

**Lemma 11.** *Let  $\Omega = [\omega_1, \dots, \omega_\ell] \in \mathbb{C}^{n \times \ell}$  such that  $\omega_i \sim \mathcal{N}(\mathbf{0}, C)$  for all  $i \in [\ell]$  and  $C \succeq 0$  is symmetric. For  $V \in \mathbb{O}_{n,k}$ , it follows that the columns of  $V^*\Omega$  are sampled from a centered multivariate Gaussian distribution with second-moment matrix  $K = V^*CV$ .*

**Proof.** The matrix  $V^*\Omega$  can be decomposed as follows,

$$V^*\Omega = \begin{bmatrix} \mathbf{v}_1^* \omega_1 & \cdots & \mathbf{v}_1^* \omega_\ell \\ \vdots & \ddots & \vdots \\ \mathbf{v}_k^* \omega_1 & \cdots & \mathbf{v}_k^* \omega_\ell \end{bmatrix}.$$

Let  $\mathcal{D} = \mathcal{N}(\mathbf{0}, I)$ . We will first show that the entries of each column of  $V^*\Omega$  are Gaussian. From the fact that  $C$  is symmetric, we have that  $C = U\Sigma U$  for a unitary  $U$  and diagonal  $\Sigma \succeq 0$ . Then for any  $i \in [k]$  and  $j \in [\ell]$ , we have for  $\mathbf{x} \sim \mathcal{D}$ ,

$$\mathbf{v}_i^* \omega_j = \mathbf{v}_i^* C^{1/2} \mathbf{x} = \mathbf{v}_i^* U \Sigma^{1/2} \mathbf{x} = \sum_{k \in [n]} \mathbf{v}_i^* \mathbf{u}_k \sqrt{\lambda_k(C)} x_k.$$

In the above, we have that each  $[V^*\Omega]_{i,j}$  is Gaussian for  $(i, j) \in [k] \times [\ell]$  as a linear combination of Gaussians is Gaussian. We will calculate the mean and covariance. We first calculate the mean of a column of  $V^*\Omega$ . For any  $(i, j) \in [k] \times [\ell]$ ,

$$\begin{aligned} \mathbf{E}_{\omega_j \sim \mathcal{N}(\mathbf{0}, C)} [\mathbf{v}_i^* \omega_j] &= \mathbf{E}_{\mathbf{x} \sim \mathcal{D}} [\mathbf{v}_i^* C^{1/2} \mathbf{x}] = \mathbf{v}_i^* C^{1/2} \mathbf{E}_{\mathbf{x} \sim \mathcal{D}} [\mathbf{x}] \\ &= \sum_{p \in [n]} \mathbf{v}_i^* \mathbf{u}_p \sqrt{\lambda_p(C)} \mathbf{E}_{\mathbf{x} \sim \mathcal{D}} [x_p] = 0. \end{aligned}$$

Now we calculate the covariance matrix. Let  $\mathbf{v} \in \mathbb{S}^{n-1}$ , then for any  $(i, i', j) \in [k] \times [k] \times [\ell]$ , and  $i \neq i'$ , we have

$$\begin{aligned} \mathbf{E}_{\omega_j \sim \mathcal{N}(\mathbf{0}, C)} \left[ (\mathbf{v}_i^* \omega_j - \mathbf{E}_{\omega_j \sim \mathcal{N}(\mathbf{0}, C)} [\mathbf{v}_i^* \omega_j]) (\mathbf{v}_{i'}^* \omega_j - \mathbf{E}_{\omega_j \sim \mathcal{N}(\mathbf{0}, C)} [\mathbf{v}_{i'}^* \omega_j]) \right] \\ = \mathbf{E}_{\omega_j \sim \mathcal{N}(\mathbf{0}, C)} [\mathbf{v}_i^* \omega_j \omega_j^* \mathbf{v}_{i'}] = \mathbf{v}_i^* C \mathbf{v}_{i'}. \end{aligned} \quad (4)$$

For the diagonal covariance elements, we have

$$\mathbf{E}_{\omega_j \sim \mathcal{N}(\mathbf{0}, C)} [(\mathbf{v}_i^* \omega_j - \mathbf{E}_{\omega_j \sim \mathcal{N}(\mathbf{0}, C)} [\mathbf{v}_i^* \omega_j])^2] = \mathbf{E}_{\omega_j \sim \mathcal{N}(\mathbf{0}, C)} [\mathbf{v}_i^* \omega_j \omega_j^* \mathbf{v}_i] = \mathbf{v}_i^* C \mathbf{v}_i. \quad (5)$$

Then combining Equations (4) and (5), we have

$$K = V^*CV.$$

Our proof is complete. ■

### A.2 IMPROVEMENTS TO THE GENERALIZED RSVD

We now give an improvement to the deterministic error bound given in Theorem 9.1 of Halko et al. (2011) by Boutsidis et al. (2014).

**Lemma 12** (Lemma 3.2 in Boutsidis et al. (2014)). *Let  $A \in \mathbb{R}^{m \times n}$ , and let  $Q$  be an orthonormal basis of  $A\Omega$  for  $\Omega \in \mathbb{R}^{n \times (k+p)}$ , then*

$$\|A - Q(Q^*A)_k\|_F^2 \leq \|A - A_k\|_F^2 + \|\Sigma_{k,\perp} V_{k,\perp} \Omega (V_k \Omega)^+\|_F^2.$$

Before we give our improvement to the Generalized Randomized SVD, we present the following necessary lemma on the concentration of  $\|(V^* \Omega)^+\|_F$  for  $\Omega$  with columns sampled from  $\mathcal{N}(\mathbf{0}, C)$ .

**Lemma 13** (Lemma 3 in Boullé & Townsend (2023)). *Suppose  $V_k \in \mathbb{O}^{n,k}$  and  $\Omega \in \mathbb{C}^{n \times k+p}$  with columns sampled i.i.d from  $\mathcal{N}(\mathbf{0}, K)$  and  $p \geq 4$ . Then,*

$$\|(V_k^* \Omega)^+\|_F \leq \sqrt{\frac{3 \operatorname{Tr}(K^{-1})}{p+1}} \cdot t^2,$$

with probability at least  $1 - t^{-p}$ .

For our improvement to the Generalized Randomized SVD for the spectral norm, we present the following necessary lemma.

**Lemma 14** (Proposition 10.4 in Halko et al. (2011)). *Let  $G \in \mathbb{R}^{k \times (k+p)}$  have elements sampled i.i.d from  $\mathcal{N}(0, 1)$  for  $p \geq 4$ . Then for  $t \geq 1$ , with probability exceeding  $1 - t^{-(p+1)}$ ,*

$$\|G^+\|_2 \leq \frac{e\sqrt{k+p}}{p+1} \cdot t.$$

We are now ready to prove our relative Spectral and Frobenius Error Norm Bounds for the Generalized Randomized SVD originally presented in Boullé & Townsend (2022).

**Proof.**[Theorem 4] Recall  $A = U\Sigma V^*$ ,  $\Omega_k = V_k^* \Omega$  and  $\Omega_{k,\perp} = V_{k,\perp}^* \Omega$  where the columns of  $\Omega \in \mathbb{R}^{n \times (k+p)}$  are sampled from  $\mathcal{N}(\mathbf{0}, C)$ . Let

$$K = \begin{bmatrix} V_k^* C V_k & V_k^* C V_{k,\perp} \\ V_{k,\perp}^* C V_k & V_{k,\perp}^* C V_{k,\perp} \end{bmatrix} = \begin{bmatrix} K_{11} & K_{12} \\ K_{21} & K_{22} \end{bmatrix}.$$

We then have the following manipulations, from Lemma 12,

$$\begin{aligned} \|A - Q(Q^* A)_k\|_F &\leq (\|\Sigma_{k,\perp}\|_F^2 + \|\Sigma_{k,\perp} \Omega_{k,\perp} \Omega_k^+\|_F^2)^{1/2} \\ &= \left( \|\Sigma_{k,\perp}\|_F^2 + \|\Sigma_{k,\perp} V_{k,\perp}^* C^{1/2} X (V_k^* C^{1/2} X)^+\|_F^2 \right)^{1/2} \\ &= \left( \|\Sigma_{k,\perp}\|_F^2 + \|\Sigma_{k,\perp} V_{k,\perp}^* K_{22}^{1/2} X_2 (K_{11}^{1/2} X_1)^+\|_F^2 \right)^{1/2} \\ &\leq \left( \|\Sigma_{k,\perp}\|_F^2 + 2u \log((n-k)(k+p)) \|\Sigma_{k,\perp} K_{22}^{1/2}\|_F^2 \|(K_{11}^{1/2} X_1)^+\|_F^2 \right)^{1/2} \\ &\leq \|\Sigma_{k,\perp}\|_F \left( 1 + t^2 \sqrt{6u \log((n-k)(k+p))} \|K_{22}\|_2^{1/2} \sqrt{\operatorname{Tr}(K_{11}^{-1})} \sqrt{\frac{1}{p+1}} \right) \end{aligned}$$

In the above, the first relation follows from the deterministic error bound in Lemma 12. The second relation follows from noting that the columns of  $\Omega$  are sampled from  $\mathcal{N}(\mathbf{0}, C)$ , thus  $\Omega = C^{1/2} X$  where  $X \in \mathbb{R}^{n \times (k+p)}$  is a standard Gaussian matrix. In the third relation, we apply Lemma 11 to note that  $V_{k,\perp}^* C^{1/2} X = K_{22}^{1/2} X_2$  and  $V_k^* C^{1/2} X = K_{11}^{1/2} X_1$  for standard Gaussian matrices  $X_1 \in \mathbb{R}^{k \times (k+p)}$  and  $X_2 \in \mathbb{R}^{(n-k) \times (k+p)}$ . The fourth relation follows from Proposition 20 for the Frobenius norm with failure probability at most  $\delta$ , and the final inequality follows from noting by the sub-multiplicativity of the Frobenius Norm, which gives us

$$\|\Sigma_{k,\perp} V_{k,\perp}^* C^{1/2}\|_F^2 \leq \|\Sigma_{k,\perp}\|_F^2 \|V_{k,\perp}^* C^{1/2}\|_2^2 = \|\Sigma_{k,\perp}\|_F^2 \|K_{22}\|_2$$

and Lemma 13 which holds with failure probability at most  $t^{-p}$ . We then apply the elementary inequality  $\sqrt{a+b} \leq \sqrt{a} + \sqrt{b}$ , to complete our proof for the probabilistic bound for the Frobenius norm. ■

**Theorem 15.** *Let  $A \in \mathbb{C}^{m \times n}$ , set  $k \geq 2$  an integer, an oversampling parameter  $p \geq 4$ . Let  $\Omega \in \mathbb{R}^{n \times (k+p)}$  represent the test matrix with columns sampled from  $\mathcal{N}(\mathbf{0}, C)$ . Suppose  $QR = A\Omega$  is the economized QR decomposition of  $A\Omega$ , then for all  $u \geq 1$ ,*

$$\|A - QQ^* A\|_2 \leq \|A - A_k\|_2 \left( 1 + \sqrt{2ut \log((n-k)(k+p))} \|K_{22}\|_2^{1/2} \|K_{11}^{-1}\|_2 \frac{e\sqrt{k+p}}{p+1} \right)$$

**Proof.** We follow the same notation as in the proof of [Theorem 4](#). From the deterministic error bound ([Halko et al., 2011, Thm. 9.1](#)),

$$\begin{aligned}
\|A - QQ^*A\|_2 &\leq (\|\Sigma_{k,\perp}\|_2^2 + \|\Sigma_{k,\perp}\Omega_{k,\perp}\Omega_k^+\|_2^2)^{1/2} \\
&= \left(\|\Sigma_{k,\perp}\|_2^2 + \|\Sigma_{k,\perp}V_{k,\perp}^*C^{1/2}X(V_k^*C^{1/2}X)^+\|_2^2\right)^{1/2} \\
&= \left(\|\Sigma_{k,\perp}\|_2^2 + \|\Sigma_{k,\perp}K_{22}^{1/2}X_2(K_{11}^{1/2}X_1)^+\|_2^2\right)^{1/2} \\
&\leq \left(\|\Sigma_{k,\perp}\|_2^2 + 2\log((n-k)(k+p)/\delta)\|\Sigma_{k,\perp}K_{22}^{1/2}\|_2^2\|(K_{11}^{1/2}X_1)^+\|_2^2\right)^{1/2} \\
&\leq \|\Sigma_{k,\perp}\|_2 \left(1 + \sqrt{2t\log((n-k)(k+p)/\delta)}\|K_{22}\|_2^{1/2}\|K_{11}^{-1}\|_2^{1/2}\frac{e\sqrt{k+p}}{p+1}\right).
\end{aligned}$$

The first inequality follows from the deterministic error bound in [Theorem 9.1](#) of [Halko et al. \(2011\)](#). The second relation follows from noting  $\Omega = C^{1/2}X$  where  $X$  is a Gaussian matrix. The third relation follows from [Lemma 11](#) which gives Gaussian matrices  $X_1 \in \mathbb{R}^{k \times (k+p)}$  and  $X_2 \in \mathbb{R}^{(n-k) \times (k+p)}$ . The final relation follows from noting by Cauchy–Schwarz,

$$\|\Sigma_{k,\perp}V_{k,\perp}^*C^{1/2}\|_2^2 \leq \|\Sigma_{k,\perp}\|_2^2\|V_{k,\perp}^*C^{1/2}\|_2^2 = \|\Sigma_{k,\perp}\|_2^2\|K_{22}\|_2$$

Furthermore, from noting that for any rank- $k$  matrix  $\|B^+\|_2 = \sigma_k^{-1}(B)$ , we have

$$\|(V_k^*C^{1/2}X)^+\|_2^2 = \sigma_k^{-2}(V_k^*C^{1/2}X_1) \leq \lambda_k^{-1}(K_{11})\sigma_k^{-2}(X_1) = \|K_{11}^{-1}\|_2\|X_1^+\|_2^2$$

Then, applying the bound in [Lemma 14](#) with failure probability at most  $t^{-p+1}$  completes the proof.  $\blacksquare$

### A.3 GENERAL ADAPTIVE SAMPLING

In this section, we will prove deferred results on general adaptive sampling. We first will prove [Lemmas 6](#) and [16](#) which we then use to prove [Theorem 7](#).

**Proof.** [\[Lemma 6\]](#) Recall that  $Q_{12}$  is an orthonormal basis of  $[A\Omega_1, A\Omega_2]$ . Let  $Q_1$  be an orthonormal basis of  $A\Omega_1$ , it thus follows that  $Q_{12}$  is also an orthonormal basis of  $[A\Omega_1, (I - Q_1Q_1^*)A\Omega_2]$ . This can be seen from the Classical Gram-Schmidt Procedure ([Schmidt, 1907](#)). Thus, we have

$$Q_{12}Q_{12}^* = Q_1Q_1^* + Q_{2|1}Q_{2|1}^*$$

Then, expanding out the residual matrix, we have

$$A - Q_{12}Q_{12}^*A = A - Q_1Q_1^*A - Q_{2|1}Q_{2|1}^*A$$

Then from noting,  $Q_1^*Q_{2|1} = 0$ , we have

$$A - Q_1Q_1^*A - Q_{2|1}Q_{2|1}^*A = (A - Q_1Q_1^*A) - Q_{2|1}Q_{2|1}^*(A - Q_1Q_1^*A)$$

Our proof is complete.  $\blacksquare$

We now require the following linear algebraic lemma.

**Lemma 16.** *Let  $A \in \mathbb{C}^{m \times n}$  and  $B \in \mathbb{C}^{m \times n}$  and  $Q \in \mathbb{O}_{m,\ell}$  for  $\ell > k$ . Then for any  $i \in [n]$ ,*

$$\sigma_i^2(A - QQ^*A) \leq \sigma_i^2(A - Q(Q^*A)_k)$$

**Proof.** We can note that  $\sigma_i(QQ^*A) \leq \sigma_i(A)$  for all  $i \in [n]$  by noting  $QQ^* \preceq I$  as  $Q \in \mathbb{O}_{m,k}$ . We first note that for any matrix  $B \in \mathbb{C}^{m \times n}$ , the singular values and eigenvalues are related by  $\sigma_i^2(B) = \lambda_i(B^*B)$  for any  $i \in [n]$ . Then let us consider any  $i \in [n]$ .

$$\lambda_i((A - QQ^*A)^*(A - QQ^*A)) = \lambda_i(A^*A - A^*QQ^*A).$$

Let us now consider the matrix in the RHS of the lemma statement,

$$\begin{aligned}
\sigma_i^2(A - Q(Q^*A)_k) &= \lambda_i((A - Q(Q^*A)_k)^*(A - Q(Q^*A)_k)) \\
&= \lambda_i(A^*A - A^*Q(Q^*A)_k - (Q^*A)_k^*Q^*A + (Q^*A)_k^*Q^*Q(Q^*A)_k) \\
&= \lambda_i(A^*A - (A^*QQ^*A)_k - (A^*QQ^*A)_k + (A^*QQ^*A)_k) \\
&= \lambda_i(A^*A - (A^*QQ^*A)_k).
\end{aligned} \tag{6}$$

In the above, the final relation follows noting  $Q^*Q = I$  because  $Q \in \mathbb{O}_{m,\ell}$ , and for any matrix  $B$ , we have that  $B_k^*B = (B^*B)_k$  by an expansion of the SVD. We can now conclude our Lemma,

$$\begin{aligned}
\sigma_i^2(A - QQ^*A) &= \lambda_i(A^*A - A^*QQ^*A) \\
&= \lambda_i(A^*A - (A^*QQ^*A)_k - (A^*QQ^*A)_{k,\perp}) \\
&\leq \lambda_i(A^*A - (A^*QQ^*A)_k) && \text{Lemma 3} \\
&= \sigma_i^2(A - A(Q^*A)_k) && \text{Equation (6)}
\end{aligned}$$

In the above, the third relation follows from noting  $(A^*QQ^*A)_{k,\perp} \geq 0$ . Our proof is complete. ■  
We are now ready to prove our general adaptive sampling error bound for one round.

**Proof.**[Theorem 7] From Lemma 6, we obtain

$$\|A - Q_{t+1}Q_{t+1}^*A\|_F = \|(I - Q_tQ_t^*)A - \tilde{Q}_{t+1}\tilde{Q}_{t+1}^*(I - Q_tQ_t^*)A\|_F.$$

The above equation follows from the same argument as Lemma 6. Then from our relative-error accuracy bound in Corollary 5, we have

$$\|(I - Q_-Q_-^*)A - QQ^*(I - Q_-Q_-^*)A\|_F^2 \leq (1 + \epsilon)\|(A - Q_-Q_-^*A)_{k,\perp}\|_F^2$$

We then have from Lemma 16,

$$\begin{aligned}
\|(A - Q_-Q_-^*A)_{k,\perp}\|_F^2 &= \sum_{i \in [n] \setminus [k]} \sigma_i^2(A - Q_-Q_-^*A) && \text{Frobenius Norm Defn.} \\
&\leq \sum_{i \in [n] \setminus [k]} \sigma_i^2(A - Q_-(Q_-^*A)_k) && \text{Lemma 16} \\
&= \|(A - Q_-(Q_-^*A)_k)_{k,\perp}\|_F^2 && \text{Frobenius Norm Defn.}
\end{aligned}$$

Then, from expanding out the Frobenius Norm, we have

$$\|(A - Q_-(Q_-^*A)_k)_{k,\perp}\|_F^2 = \|A - Q_-(Q_-^*A)_k\|_F^2 - \sum_{i \in [k]} \sigma_i^2(A - Q_-(Q_-^*A)_k)$$

We now leverage the Eckart-Young-Mirsky Theorem Eckart & Young (1936); Mirsky (1960), and obtain

$$\sum_{i \in [k]} \sigma_i^2(A - Q_-(Q_-^*A)_k) \geq \inf_{\substack{B \in \mathbb{C}^{m \times n} \\ \text{rank}(B)=k}} \sum_{i \in [k]} \sigma_i^2(A - B) = \|A_k - A_{2k}\|_F^2 \tag{7}$$

Then, from our improved Generalized Randomized SVD error bound in Theorem 4, we have with probability at least  $0.99 - t^{-p}$ ,

$$\|A - Q_-(Q_-^*A)_k\|_F^2 \leq \|A - A_k\|_F^2 \left( 1 + 6t^2 \log(100n\ell) \|V_{k,\perp}^* C^{1/2}\|_2^2 \frac{\text{Tr}((V_k^* C V_k)^{-1})}{p+1} \right) \tag{8}$$

Then, from choosing our oversampling parameter sufficiently large in accordance to Corollary 5,

$$p \geq (6t^2/\epsilon) \|V_{k,\perp}^* C^{1/2}\|_2^2 \text{Tr}((V_k^* C V_k)^{-1})$$

We can then combine our results in Equations (7) and (8) to obtain

$$\begin{aligned}
\|A - Q_+Q_+^*A\|_F^2 &\leq (1 + \epsilon)^2 \|A - A_k\|_F^2 - (1 + \epsilon) \|A_k - A_{2k}\|_F^2 \\
&= \epsilon(1 + \epsilon) \|A - A_k\|_F^2 + (1 + \epsilon) \|A - A_{2k}\|_F^2
\end{aligned}$$

Our proof is complete. ■



#### A.4 ADAPTIVE SAMPLING

We first give necessary lemmas for our proof.

**Lemma 17** (Lemma 5.3 in Boutsidis & Gittens (2013)). *Let  $A \in \mathbb{R}^{m \times n}$  and  $Y \in \mathbb{R}^{m \times \ell}$ , then for all  $X \in \mathbb{R}^{\ell \times n}$  and  $\xi \in \{2, F\}$ ,*

$$\|A - YY^+A\|_\xi \leq \|A - YX\|_\xi.$$

**Lemma 18** (Matrix Pythagoras). *Let  $X, Y$  be conformal matrices, then if  $X^*Y = \mathbf{0}$  or  $Y^*X = \mathbf{0}$ ,*

$$\|X + Y\|_2^2 \leq \|X\|_2^2 + \|Y\|_2^2 \quad \text{and} \quad \|X + Y\|_F^2 = \|X\|_F^2 + \|Y\|_F^2.$$

**Proof.** The proof for the Spectral norm follows directly from the definition.

$$\|X + Y\|_2^2 = \|X^*X + X^*Y + Y^*X + Y^*Y\|_2 = \|X^*X + Y^*Y\|_2 \leq \|X\|_2^2 + \|Y\|_2^2$$

In the above, the final relation follows from the sub-additivity of the spectral norm and noting for any matrix  $B$ , it holds  $\|B^*B\|_2 = \|B\|_2^2$ . The proof for the Frobenius norm follows directly from the definition of the Frobenius Norm.

$$\|X + Y\|_F^2 = \text{Tr}(X^*X + X^*Y + Y^*X + Y^*Y) = \text{Tr}(X^*X) + \text{Tr}(Y^*Y) = \|X\|_F^2 + \|Y\|_F^2$$

In the above, we use the fact that  $\text{Tr}(A + B) = \text{Tr}(A) + \text{Tr}(B)$  for conformal matrices,  $A, B$ . ■

**Proof.**[Theorem 9] Our argument relies on an equivalence for the range of  $Q$ . Since we have that  $Q$  is an orthonormal basis of  $[Y_1 \ Y_2]$ . Then we have that  $Q$  is also the unique orthonormal basis of  $[Y_1 Y_1^+ Y_1 \ (I - Y_1 Y_1^+) Y_2]$  by the definition of the Moore-Penrose Inverse (see the first property of the pseudoinverse in Section 4.1) and the classical Gram-Schmidt orthonormalization procedure (Schmidt, 1907). Let  $X \in \mathbb{R}^{2\ell \times n}$ , then we have from Lemma 17,

$$\begin{aligned} \|A - YY^+A\|_\xi^2 &\leq \left\| A - \begin{bmatrix} P_{Y_1} A \Psi_1 & (I - P_{Y_1}) A C^{1/2} \Psi_2 \end{bmatrix} \begin{bmatrix} X_1 \\ X_2 \end{bmatrix} \right\|_\xi^2 \\ &\leq \underbrace{\|P_{Y_1} A - P_{Y_1} A \Psi_1 X_1\|_\xi^2}_I + \underbrace{\|(A - P_{Y_1} A) - (A - P_{Y_1} A) C^{1/2} \Psi_2 X_2\|_\xi^2}_{II} \end{aligned}$$

In the above, the second relation follows from Matrix Pythagoras (see Lemma 18). Let the SVD of  $P_{Y_1} A$  be  $\widehat{U} \widehat{\Sigma} \widehat{V}^*$ , we then choose  $X_1 = (\widehat{V}_{k+p}^* \Psi_1)^+ \widehat{V}_{k+p}^*$ , then we have  $\widehat{V}_{k+p}^* \Psi_1 (\widehat{V}_{k+p}^* \Psi_1)^+ \widehat{V}_{k+p}^* = \widehat{V}_{k+p}^*$  from our assumption that  $\widehat{V}_{k+p}^* \Psi_1$  is full-rank. Then we can note that  $\widehat{V}_{k+p}^* \Psi_1$  is invertible almost surely with probability 1. We then obtain,

$$I = \|P_{Y_1} A - P_{Y_1} A \Psi_1 X_1\|_\xi^2 = \|P_{Y_1} A - P_{Y_1} A\|_\xi^2 = 0.$$

We then have for the second term,

$$\begin{aligned} II &= \|(A - P_{Y_1} A) - (A - P_{Y_1} A) C^{1/2} \Psi_2 X_2\|_\xi^2 \\ &= \|(A - A C^{1/2} \Psi_2 X_2) - (P_{Y_1} A - P_{Y_1} A C^{1/2} \Psi_2 X_2)\|_\xi^2. \end{aligned}$$

Then, setting  $X_2 = (V_k^* C^{1/2} \Psi_2)^+ V_k^*$ , we obtain

$$\begin{aligned} II^{1/2} &\leq \|A_{k,\perp} - P_{Y_1} A_{k,\perp}\|_\xi + \|(A_{k,\perp} - P_{Y_1} A_{k,\perp}) C^{1/2} \Psi_2 (V_k^* C^{1/2} \Psi_2)^+\|_\xi \\ &\leq \|A_{k,\perp} - P_{A_{k,\perp} \Psi_1} A_{k,\perp}\|_\xi + \|(A_{k,\perp} - P_{A_{k,\perp} \Psi_1} A_{k,\perp}) C^{1/2} \Psi_2 (V_k^* C^{1/2} \Psi_2)^+\|_\xi \\ &\leq \|A_{k,\perp} - P_{A_{k,\perp} \Psi_1} A_{k,\perp}\|_\xi \left( 1 + \|V_{k,\perp}^* C^{1/2} \Psi_2 (V_k^* C^{1/2} \Psi_2)^+\|_2 \right) \\ &\leq (1 + \epsilon) \|A - A_{2k}\|_\xi \left( 1 + \sqrt{4 \log(10(n-k)(k+p))} \|V_{k,\perp}^* C^{1/2}\|_2 \|(V_k^* C^{1/2})^+\|_2 \frac{e\sqrt{k+p}}{p+1} \right) \end{aligned}$$

In the above, the second inequality follows from the Conjugation Rule (see Proposition 2) with  $P_{A_{k,\perp} \Psi_1} \preceq P_{A \Psi_1}$ . Suppose for two conformal matrices,  $Z_1, Z_2$  such that  $\text{range}(Z_1) \subset \text{range}(Z_2)$ , it then follows that  $P_{Z_1} \preceq P_{Z_2}$ , and we have for the Frobenius norm,

$$\|(I - P_{Z_2})A\|_F^2 = \sum_{i \in [n]} \|(I - P_{Z_2})A e_i\|_2^2 = \sum_{i \in [n]} e_i^* A^* (I - P_{Z_2}) A e_i$$

$$\leq \sum_{i \in [n]} \mathbf{e}_i^* A^* (I - P_{Z_1}) A \mathbf{e}_i = \|(I - P_{Z_1}) A\|_F^2$$

In the above, the first and final relations follow from an alternative definition of the Frobenius Norm, the second relation follows from the definition of the spectral norm, and the third relation follows from the Conjugation Rule. Then, noting that  $\text{range}(A_{k,\perp} \Psi_1) \subset \text{range}(A \Psi_1)$ , we obtain the desired result. A similar result for the spectral norm follows from a simpler argument. The final inequality follows from the standard rSVD Frobenius norm error bound (see e.g. Theorem 10.6 in Halko et al. (2011)) for  $k + p$  samples on the matrix  $A - A_k$ , we thus obtain

$$\|A_{k,\perp} - P_{A_{k,\perp} \Psi_1} A_{k,\perp}\|_\xi \leq (1 + \epsilon) \|A - A_{2k}\|_\xi$$

and the second term follows from the manipulations in the Spectral Norm bound proof given in Theorem 15 with failure probability less than  $\delta + t^{-p+1} \leq 0.01 + 0.0325 \leq 0.05$ . ■

**Theorem 19.** Consider Algorithm 1 for  $t$  steps, then

$$\|A - QQ^* A\|_F \lesssim \prod_{i=1}^t \left(1 + \sqrt{8 \log((n - ki)(ki + pi))} \|V_{\mathcal{I},\perp}^* C_i^{1/2}\|_2 \|(V_{\mathcal{I},i}^* C_i^{1/2})^+ \|_2\right) \|A - A_{tk}\|_\xi$$

**Proof.** The proof follows similarly to the above theorem. Let  $\tilde{\Omega} = [\Omega \quad \Omega_t]$  be all the sampled vectors. We furthermore decompose  $\Omega = [\Omega_1 \quad \dots \quad \Omega_{t-1}]$ . Note  $\Omega_i \in \mathbb{R}^{n \times (k+p)}$  and let  $Y = A\Omega$  and  $Y_i = A\Omega_i$  for all  $i \in [t]$ . We then obtain,

$$\begin{aligned} \|A - QQ^* A\|_F^2 &\leq \left\| A - \begin{bmatrix} P_Y A \Omega & (I - P_Y) A \Omega_t \end{bmatrix} \begin{bmatrix} X \\ X_t \end{bmatrix} \right\|_\xi^2 \\ &= \sum_{i=1}^{t-1} \underbrace{\|P_{Y_i} A - P_{Y_i} A C_i^{1/2} \Psi_i X_i\|_\xi^2}_{I_i} + \underbrace{\|(A - P_Y A) - (A - P_Y A) C^{1/2} \Psi_t X_t\|_\xi^2}_{II} \end{aligned}$$

We can then note that by choosing  $C_i^{1/2} = (P_{Y_i} A)^+ (P_{Y_i} A)$ , we have that  $P_{Y_i} A C_i^{1/2} = P_{Y_i} A$  and then choosing  $X_i = (\hat{V}_{k+p}^* C_i^{1/2} \Psi)^+ \hat{V}_{k+p}^*$ , with probability almost surely unitary, we obtain

$$\sum_{i=1}^{t-1} I_i \stackrel{\text{def}}{=} \sum_{i=1}^{t-1} \|P_{Y_i} A - P_{Y_i} A C_i^{1/2} \Psi_i X_i\|_\xi^2 = \sum_{i=1}^{t-1} \|P_{Y_i} A - P_{Y_i} A\|_\xi^2 = 0$$

Then, considering the second term, we have

$$\begin{aligned} II &= \|(A - P_Y A) - (A - P_Y A) C_t^{1/2} \Psi_t X_t\|_\xi^2 \\ &= \|(A - A C_t^{1/2} \Psi_t X_t) - (P_Y A - P_Y A C_t^{1/2} \Psi_t X_t)\|_\xi^2 \end{aligned}$$

Let  $\mathcal{I} = \{kt, \dots, k(t+1)\}$ . We now set  $X_t = (V_{\mathcal{I}}^* C_t^{1/2} \Psi_t)^+ V_{\mathcal{I}}^*$ . We then obtain,

$$\begin{aligned} \sqrt{II} &\leq \|A_{\mathcal{I}} - P_Y A_{\mathcal{I}}\|_\xi \left(1 + \|V_{\mathcal{I},\perp}^* C_t^{1/2} \Psi_t (V_{\mathcal{I}}^* C_t^{1/2} \Psi_t)^+ \|_2\right) \\ &\leq \|A_{\mathcal{I},\perp} - P_{A_{\mathcal{I},\perp} \Omega} A_{\mathcal{I},\perp}\|_\xi \left(1 + \sqrt{8 \log(n/\delta)} \|V_{\mathcal{I},\perp}^* C_t^{1/2}\|_2 \|(V_{\mathcal{I}}^* C_t^{1/2})^+ \|_2 \frac{e\sqrt{k+p}}{p+1}\right) \end{aligned}$$

In the above, we observe when  $C_t$  is a good approximation of  $V_{\mathcal{I}}$  we obtain an improvement. Let us now consider the term  $\|A_{\mathcal{I}} - P_{A_{\mathcal{I}} \Omega} A_{\mathcal{I}}\|_\xi$ .

$$\begin{aligned} \|A_{\mathcal{I}} - P_{A_{\mathcal{I}} \Omega} A_{\mathcal{I}}\|_\xi^2 &\leq \left\| A_{\mathcal{I}} - \begin{bmatrix} P_{A_{\mathcal{I}} \Omega} A_{\mathcal{I}} \Omega & (I - P_{A_{\mathcal{I}} \Omega}) A_{\mathcal{I}} \Omega_{t-1} \end{bmatrix} \begin{bmatrix} X \\ X_{t-1} \end{bmatrix} \right\|_\xi^2 \\ &\leq \underbrace{\|P_{A_{\mathcal{I}} \Omega} A_{\mathcal{I}} - P_{A_{\mathcal{I}} \Omega} A_{\mathcal{I}} \Omega X\|_\xi^2}_I + \underbrace{\|(A_{\mathcal{I}} - P_{A_{\mathcal{I}} \Omega} A_{\mathcal{I}}) - (A_{\mathcal{I}} - P_{A_{\mathcal{I}} \Omega} A_{\mathcal{I}}) C_{t-1}^{1/2} \Psi_{t-1} X_{t-1}\|_\xi^2}_{II} \end{aligned}$$

We will now consider  $I$ . Let the SVD of  $P_{A_{\mathcal{I}} \Omega} A_{\mathcal{I}} = U_{\mathcal{I},i} \Sigma_{\mathcal{I},i} V_{\mathcal{I},i}^*$ . Then we choose

$$X_i = \left( (V_{\mathcal{I},i}^*)_{k+p} C_i^{1/2} \Psi_i \right)^+ (V_{\mathcal{I},i}^*)_{k+p}$$

From which we obtain,

$$I = \sum_{i=1}^{t-1} \|P_{A_{\mathcal{I}}\Omega} A_{\mathcal{I}} - P_{A_{\mathcal{I}}\Omega} A_{\mathcal{I}} C_i^{1/2} \Psi_i X_i\|_{\xi}^2 = \sum_{i=1}^{t-1} \|P_{A_{\mathcal{I}}\Omega_i} A_{\mathcal{I}} - P_{A_{\mathcal{I}}\Omega_i} A_{\mathcal{I}}\|_{\xi}^2 = 0$$

Let  $\mathcal{F} = \{k(t-1), \dots, k(t+1)\}$ . Next, we set  $X_{t-1} = (V_{\mathcal{I}}^* C_{t-1}^{1/2} \Psi_{t-1})^+ V_{\mathcal{I}}^*$ , we then obtain

$$\begin{aligned} \sqrt{II} &\leq \|A_{\mathcal{F},\perp} - P_{A_{\mathcal{F},\perp}\Omega} A_{\mathcal{F},\perp}\|_{\xi} \left(1 + \|V_{\mathcal{I},\perp}^* C_t^{1/2} \Psi_{t-1} (V_{\mathcal{I}}^* C_{t-1}^{1/2} \Psi_{t-1})^+ V_{\mathcal{I}}^*\|_2\right) \\ &\leq \|A_{\mathcal{F},\perp} - P_{A_{\mathcal{F},\perp}\Omega} A_{\mathcal{F},\perp}\|_{\xi} \left(1 + \sqrt{8 \log(n/\delta)} \|V_{\mathcal{I},\perp}^* C_{t-1}^{1/2}\|_2 \|(V_{\mathcal{I}}^* C_{t-1}^{1/2})^+\|_2 \frac{e\sqrt{k+p}}{p+1}\right) \end{aligned}$$

We can then complete the proof with induction, which gives us the desired result.  $\blacksquare$

## B PROBABILITY THEORY

In this section, we present and prove an interesting property of Gaussian matrices that gives us the improved bounds for the generalized rSVD.

**Proposition 20.** Fix matrices  $S \in \mathbb{R}^{k \times n}$  and  $T \in \mathbb{R}^{m \times \ell}$ , then for a conformal matrix  $G$  with elements sampled i.i.d from  $\mathcal{N}(0, 1)$ , then for all  $\xi \in \{2, F\}$  and  $u \geq 1$ ,

$$\|SGT\|_{\xi} \leq \|S\|_{\xi} \|T\|_{\xi} \sqrt{2u \log(nm)},$$

with failure probability at most  $\sqrt{2/\pi}(nm)^{1-u}$ .

**Proof.** The proof for the Frobenius norm follows from a brute-force calculation followed by a maximal tail bound on a sample of Gaussians.

$$\begin{aligned} \|SGT\|_F^2 &= \sum_{i \in [k]} \sum_{j \in [\ell]} \sum_{(k_1, k_2) \in [n] \times [m]} S_{i, k_1}^2 T_{k_2, j}^2 G_{k_1, k_2}^2 \\ &\leq \sum_{i \in [k]} \sum_{j \in [\ell]} \sum_{(k_1, k_2) \in [n] \times [m]} S_{i, k_1}^2 T_{k_2, j}^2 \max_{(k_1, k_2) \in [n] \times [m]} G_{k_1, k_2}^2 \\ &= \|S\|_F^2 \|T\|_F^2 \max_{(k_1, k_2) \in [n] \times [m]} G_{k_1, k_2}^2. \end{aligned} \quad (9)$$

We now show an analogous result for the spectral norm with just a few more steps to recover the result. We will use classical techniques in probability theory to obtain a similar proof structure as the Frobenius norm result. From Rigollet & Hütter (2023), we have the following relation for any matrix  $B \in \mathbb{R}^{m \times n}$ ,

$$\|B\|_2 = \sup_{\mathbf{v} \in \mathbb{S}^{n-1}} \|B\mathbf{v}\|_2 = \sup_{\mathbf{u} \in \mathbb{S}^{m-1}} \sup_{\mathbf{v} \in \mathbb{S}^{n-1}} |\mathbf{u}^* B \mathbf{v}|. \quad (10)$$

With the relation given in Equation (10) in hand, we have

$$\begin{aligned} \|SGT\|_2^2 &= \sup_{\mathbf{u} \in \mathbb{S}^{k-1}} \sup_{\mathbf{v} \in \mathbb{S}^{\ell-1}} (\mathbf{u}^* SGT\mathbf{v})^2 \\ &= \sup_{\mathbf{u} \in \mathbb{S}^{k-1}} \sup_{\mathbf{v} \in \mathbb{S}^{\ell-1}} \sum_{(i,j) \in [n] \times [m]} (\mathbf{u}^* S)_i^2 G_{i,j}^2 (T\mathbf{v})_j^2 \\ &\leq \sup_{\mathbf{u} \in \mathbb{S}^{k-1}} \sup_{\mathbf{v} \in \mathbb{S}^{\ell-1}} \sum_{(i,j) \in [n] \times [m]} (\mathbf{u}^* S)_i^2 (T\mathbf{v})_j^2 \max_{(i,j) \in [n] \times [m]} G_{i,j}^2 \\ &\leq \sup_{\mathbf{u} \in \mathbb{S}^{k-1}} \sup_{\mathbf{v} \in \mathbb{S}^{\ell-1}} \|S^* \mathbf{u}\|_2^2 \|T\mathbf{v}\|_2^2 \max_{(i,j) \in [n] \times [m]} G_{i,j}^2 \\ &= \|S\|_2^2 \|T\|_2^2 \max_{(i,j) \in [n] \times [m]} G_{i,j}^2. \end{aligned} \quad (11)$$

Combining our results from Equations (9) and (11), we obtain for all  $\xi \in \{2, F\}$ ,

$$\|SGT\|_{\xi}^2 \leq \|S\|_{\xi}^2 \|T\|_{\xi}^2 \max_{(i,j) \in [n] \times [m]} G_{i,j}^2.$$

We now bound the maximum Gaussian over a finite sample.

$$\begin{aligned} \Pr\left\{\max_{(k_1, k_2) \in [n] \times [m]} G_{k_1, k_2}^2 \geq ut\right\} &= \Pr\left\{\max_{(k_1, k_2) \in [n] \times [m]} |G_{k_1, k_2}| \geq \sqrt{ut}\right\} \\ &\leq \frac{\sqrt{2nm}}{\sqrt{\pi}} \int_{\sqrt{ut}}^{\infty} e^{-x^2/2} dx \leq \frac{\sqrt{2nm}}{\sqrt{\pi}} \int_{\sqrt{ut}}^{\infty} \frac{xe^{-x^2/2}}{\sqrt{ut}} dx = \frac{\sqrt{2nm}}{\sqrt{\pi}} e^{-ut/2}. \end{aligned}$$

In the above, the first inequality follows from a union bound over  $[n] \times [m]$ , and integrating over the PDF of a standard normal Gaussian (Berger & Casella, 2001). Then, from a change of variables  $t = 2 \log(nm)$ ,

$$\|SGT\|_{\xi}^2 \leq 2u \log(nm) \|S\|_{\xi}^2 \|T\|_{\xi}^2,$$

with failure probability at most  $\sqrt{2/\pi}(nm)^{1-u}$ . Taking the square root of both sides completes the proof.  $\blacksquare$

## C THEORETICAL IMPROVEMENTS TO RSVD

In this section, we take a closer look at our spectral norm bound given in [Theorem 15](#). We show that our bound implies a stronger relative spectral norm error bound than (Halko et al., 2011, Thm. 10.8). We next will show the conditions where our Frobenius norm error bound is stronger than the bound derived by Persson et al. (2024). We will first restate the implication of (Halko et al., 2011, Thm. 10.8).

**Theorem 21** (Implication of HMT Theorem 10.8). *Let  $A \in \mathbb{C}^{m \times n}$ , choose target rank  $k \geq 2$  and oversampling parameter  $p \geq 4$ . Then sample  $\Omega \in \mathbb{R}^{m \times n}$  with standard normal entries. Construct the sample matrix  $Y = A\Omega$  and let  $QR$  be the economized  $QR$  decomposition of  $Y$ . Then,*

$$\|A - QQ^*A\|_2 \leq \|A - A_k\|_2 \left(1 + t\sqrt{12k/p} + t \cdot \frac{e\sqrt{k+p}}{p+1} \sqrt{n-k} + ut \cdot \frac{\sqrt{k+p}}{p+1}\right),$$

with failure probability at most  $5t^{-p} + e^{-u^2/2}$ .

The dependence on the  $O(\sqrt{n-k})$  factor arises from Chevet’s inequality (Gordon, 1985).

**Lemma 22** (Chevet’s inequality). *Fix matrices  $S \in \mathbb{R}^{k \times n}$  and  $T \in \mathbb{R}^{m \times \ell}$ , then for a Gaussian matrix  $G \in \mathbb{R}^{n \times \ell}$  with elements sampled i.i.d from  $\mathcal{N}(0, 1)$ ,*

$$\mathbb{E}\|SGT\|_2 \leq \|S\|_2 \|T\|_F + \|S\|_F \|T\|_2.$$

In the rSVD analysis (see e.g. (Halko et al., 2011, Thm. 10.7)),  $S = A - A_k$  and thus to obtain a relative error upper bound, we can note  $A - A_k$  is maximally rank  $n - k$  and obtain,

$$\|A - A_k\|_F \leq \sqrt{n-k} \|A - A_k\|_2.$$

Although this upper bound is loose with strong singular value decay, we show we can give a stronger bound without any assumptions on the singular value spectra. We now will formalize our comparison by giving a corollary where we set  $C = I$  and use our result from [Theorem 15](#).

**Corollary 23.** *Consider the same hypothesis as [Theorem 21](#), then for all  $u \geq 1$ ,*

$$\|A - QQ^*A\|_2 \leq \|A - A_k\|_2 \left(1 + \sqrt{2ut \log((n-k)(k+p))} \frac{e\sqrt{k+p}}{p+1}\right),$$

with failure probability at most  $\sqrt{2/\pi}(nm)^{1-u} + t^{-p}$ .

**Proof.** Consider the spectral norm result of [Theorem 15](#) and set  $C = I$ .  $\blacksquare$

Our analysis allows us to remove the  $O(\sqrt{n-k})$  term and reduce it to a  $O(\log(n-k))$  term, giving us a stronger bound. We showcase the improved relative error bound in the following synthetic experiment. We observe our bounds offer a significant improvement for relative-error spectral norm  $\|A - Q(Q^*A)_k\|_2 / \|A_{k,\perp}\|_2$  compared to [Theorem 21](#) and also are bounds are a proper upper bound.

We now give the Frobenius norm error bound with tools derived in Persson et al. (2024).

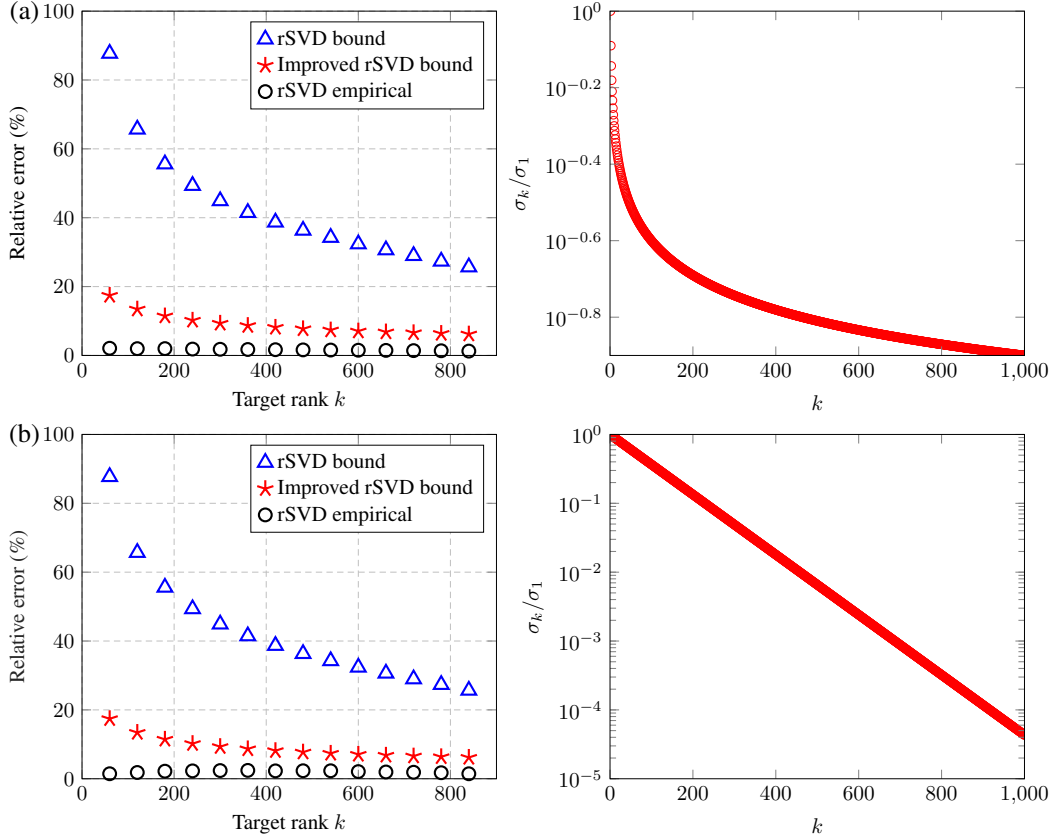


Figure 2: Left panels: Relative error (in the spectral norm) of the low-rank matrix approximation with respect to the number of matrix-vector products. At each target rank  $k$ , we set the oversampling parameter  $p = \lceil k/4 \rceil$ . Right panels: Singular value decay of the matrix. (a) Polynomial singular value decay (see Equation (12) with  $p = 0.3$ ). (b) Exponential singular value decay (see Equation (13) with  $\delta = 0.05$ ). Empirical Mean Error points in Figure 2 are averaged over 10 randomized runs.

**Theorem 24.** Let  $A \in \mathbb{C}^{m \times n}$ , let  $k \geq 2$  be a target rank, and  $p \geq 4$  be an oversampling parameter. Let  $\Omega$  be a random sketch matrix with columns sampled i.i.d from  $\mathcal{N}(0, C)$ , then

$$\|A - QQ^*A\|_F \leq \Xi$$

**Proof.** Following (Persson et al., 2024, Lem. 2.2) and Jensen’s inequality, for  $u \geq 1$ , we have

$$\|SGT\|_F = \|SGT\|_{(2)} \leq (\mathbf{E}[\|SGT\|_{(2)}^2])^{1/2} + u\|S\|_2\|T\|_2,$$

with probability  $\geq 1 - e^{-u^2}$ . Then, by (Persson et al., 2024, Lem. A.2),

$$\|SGT\|_F \leq \|S\|_F\|T\|_F + u\|S\|_2\|T\|_2.$$

If we do the change of variables  $\delta = e^{-u^2}$  and bound  $\|\cdot\|_2 \leq \|\cdot\|_F$ , we obtain

$$\|SGT\|_F \leq (1 + \sqrt{\log(1/\delta)})\|S\|_F\|T\|_F.$$

■

## D ADDITIONAL DETAILS ON THE NUMERICAL EXPERIMENTS

**NB: This requires a lot more details and discussion.** In this section, we perform additional experiments on synthetic matrices with polynomial and exponential singular value decay.

**Experiment 4: Matrices with Polynomial and Exponential Singular Value Decay.** We sample  $U$  uniformly randomly from the Haar Measure over  $\mathbb{O}_{m,m}$  and  $V$  uniformly randomly from the Haar Measure over  $\mathbb{O}_{n,n}$ . We then form the matrices  $A$  in the scheme described in Equation (12) and Equation (13), respectively.

$$A = \sum_{i \in [n]} i^{-p} \cdot \mathbf{u}_i \mathbf{v}_i^*, \quad U \in \mathbb{O}_{m,m}, V \in \mathbb{O}_{n,n} \quad (12)$$

$$A = \sum_{i \in [n]} (1 - \delta)^i \cdot \mathbf{u}_i \mathbf{v}_i^*, \quad U \in \mathbb{O}_{m,m}, V \in \mathbb{O}_{n,n}. \quad (13)$$

**Observation 4:** We observe Algorithm 1 is stronger than the rSVD and generalized rSVD in both experiments. **NB: We might also want to show the variance of all methods here.**

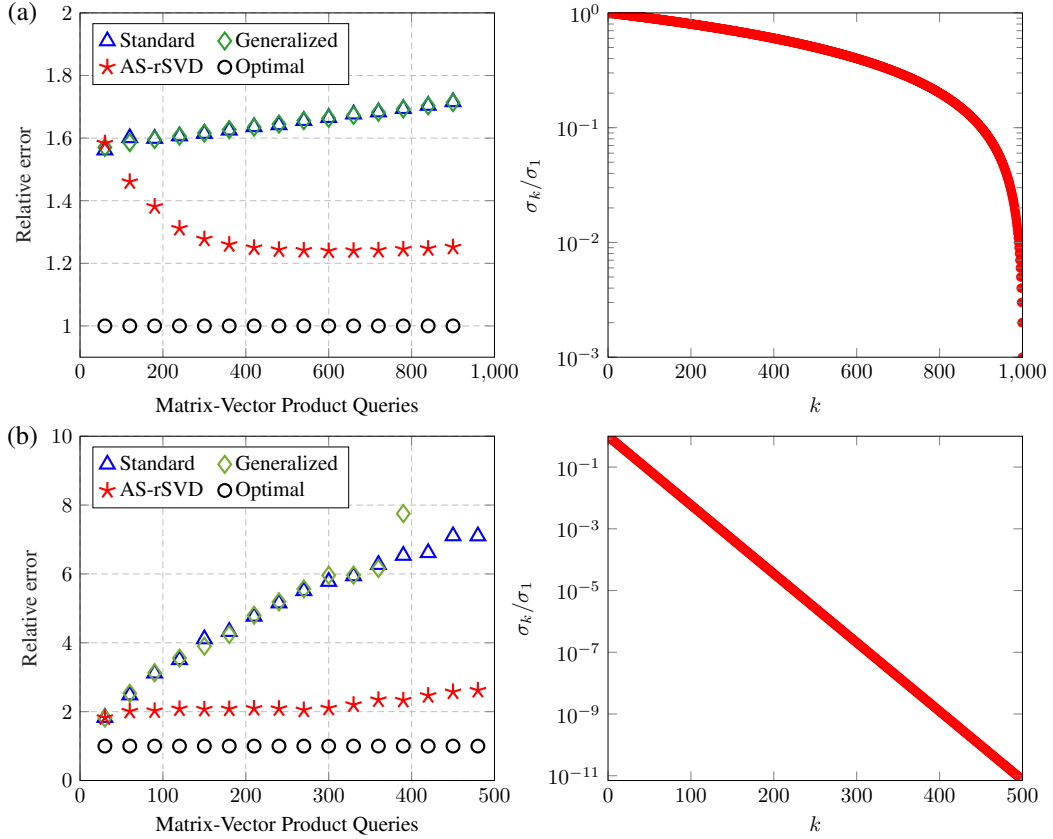


Figure 3: Left panels: Relative error (in the Frobenius norm) of the low-rank matrix approximation with respect to the number of matrix-vector products. Right panels: Singular value decay of the matrix. (a) Polynomial singular value decay (see Equation (12) with  $p = 1$ ). (b) Exponential singular value decay (see Equation (13) with  $\delta = 0.05$ ).