

# Threaded Programming

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## Lecture 5: Synchronisation

# Why is it required?

Recall:

- Need to synchronise actions on shared variables.
- Need to ensure correct ordering of reads and writes.
- Need to protect updates to shared variables (not atomic by default)

# BARRIER directive

- No thread can proceed past a barrier until all the other threads have arrived.
- Note that there is an implicit barrier at the end of DO/FOR, SECTIONS and SINGLE directives.

- Syntax:

Fortran: **!\$OMP BARRIER**

C/C++: **#pragma omp barrier**

- Either all threads or none must encounter the barrier: otherwise DEADLOCK!!

# BARRIER directive (cont)

Example:

```
#pragma omp parallel private(myid,neighb) shared(a,b,c)
{
    myid = omp_get_thread_num();
    neighb = myid - 1;
    if (myid.eq.0) neighb = omp_get_num_threads()-1;
    ...
    a[myid] *= 3.5;
    #pragma omp barrier
    b[myid] = a[neighb] + c;
    ...
}
```

- Barrier required to force synchronisation on **a**

# BARRIER directive (cont)

Example:

```
!$OMP PARALLEL PRIVATE(MYID,NEIGHB) SHARED(A,B,C)
  myid = omp_get_thread_num()
  neighb = myid - 1
  if (myid.eq.0) neighb = omp_get_num_threads()-1
  ...
  a(myid) = a(myid)*3.5
!$OMP BARRIER
  b(myid) = a(neighb) + c
  ...

!$OMP END PARALLEL
```

- Barrier required to force synchronisation on **a**

# Critical sections

- A critical section is a block of code which can be executed by only one thread at a time.
- Can be used to protect updates to shared variables.

# CRITICAL directive

- Syntax:

Fortran: **!\$OMP CRITICAL**  
*block*

**!\$OMP END CRITICAL**

C/C++: **#pragma omp critical**  
*structured block*

# CRITICAL directive (cont)

Example: appending to a shared list

```
#pragma omp parallel for shared(list, N) private(newitem_p)
for (int i=0; i<N; i++) {
    newitem_p = createitem(i);
#pragma omp critical
    {
        append(&list, p_newitem);
    }
}
```



# CRITICAL directive (cont)

Example: appending to a shared list

```
!$OMP PARALLEL DO SHARED(list,n) PRIVATE(newitem)
do i=1,n
    newitem = createitem(i)
    !$OMP CRITICAL
        call append(list,newitem)
    !$OMP END CRITICAL
end do
```

# CRITICAL directive (cont)

Example: pushing and popping a task stack

```
#pragma omp parallel shared(stack) private(p_next,p_new,done)
{
while (!done) {
#pragma omp critical
{
    p_next = pop(&stack);
}
    p_new = process(p_next);
#pragma omp critical
{
    if (p_new != NULL) push(p_new,&stack);
    done = isempty(&stack);
}
}
}
```

# CRITICAL directive (cont)

Example: pushing and popping a task stack

```
!$OMP PARALLEL SHARED(stack) ,PRIVATE(next,new,done)
    do while (.not. done)
!$OMP CRITICAL
        next = pop(stack)
!$OMP END CRITICAL
        new = process(next)
!$OMP CRITICAL
            if (valid(new)) call push(new,stack)
            done = isempty(stack)
!$OMP END CRITICAL
    end do
!$OMP END PARALLEL
```

# ATOMIC directive

- Used to protect a single update to a shared scalar variable of basic type.
- Applies only to a single statement.
- Syntax:

Fortran: **!\$OMP ATOMIC**  
*statement*

where *statement* must have one of these forms:

$x = x \text{ op } \text{expr}, \quad x = \text{expr op } x, \quad x = \text{intr} (x, \text{expr}) \text{ or}$   
 $x = \text{intr} (\text{expr}, x)$

*op* is one of +, \*, -, /, **.and.**, **.or.**, **.eqv.**, or **.neqv.**

*intr* is one of **MAX**, **MIN**, **IAND**, **IOR** or **IEOR**

# ATOMIC directive (cont)

C/C++: `#pragma omp atomic`  
*statement*

where *statement* must have one of the forms:

$x \text{ binop} = \text{expr}$ ,  $x++$ ,  $++x$ ,  $x--$ , or  $--x$

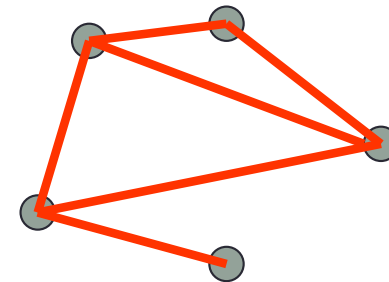
and *binop* is one of  $+$ ,  $*$ ,  $-$ ,  $/$ ,  $\&$ ,  $\wedge$ ,  $\ll$ , or  $\gg$

- Note that the evaluation of *expr* is not atomic.
- May be more efficient than using CRITICAL directives, e.g. if different array elements can be protected separately.
- No interaction with CRITICAL directives

# ATOMIC directive (cont)

Example (compute degree of each vertex in a graph):

```
#pragma omp parallel for
    for (j=0; j<nedges; j++){
#pragma omp atomic
        degree[edge[j].vertex1]++;
#pragma omp atomic
        degree[edge[j].vertex2]++;
    }
```



# Lock routines

- Occasionally we may require more flexibility than is provided by CRITICAL directive.
- A lock is a special variable that may be *set* by a thread. No other thread may *set* the lock until the thread which set the lock has *unset* it.
- Setting a lock can either be blocking or non-blocking.
- A lock must be initialised before it is used, and may be destroyed when it is no longer required.
- Lock variables should not be used for any other purpose.

# Lock routines - syntax

Fortran:

```
USE OMP_LIB
```

```
SUBROUTINE OMP_INIT_LOCK(OMP_LOCK_KIND var)
```

```
SUBROUTINE OMP_SET_LOCK(OMP_LOCK_KIND var)
```

```
LOGICAL FUNCTION OMP_TEST_LOCK(OMP_LOCK_KIND var)
```

```
SUBROUTINE OMP_UNSET_LOCK(OMP_LOCK_KIND var)
```

```
SUBROUTINE OMP_DESTROY_LOCK(OMP_LOCK_KIND var)
```

*var* should be an INTEGER of the same size as addresses (e.g. INTEGER\*8 on a 64-bit machine)

OMP\_LIB defines OMP\_LOCK\_KIND



# Lock routines - syntax

C/C++:

```
#include <omp.h>
```

```
void omp_init_lock(omp_lock_t *lock);
```

```
void omp_set_lock(omp_lock_t *lock);
```

```
int omp_test_lock(omp_lock_t *lock);
```

```
void omp_unset_lock(omp_lock_t *lock);
```

```
void omp_destroy_lock(omp_lock_t *lock);
```

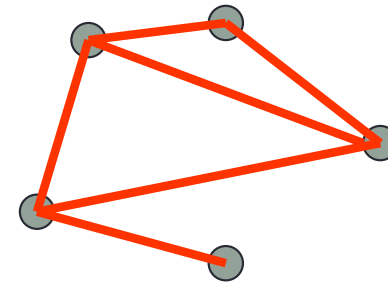
# Lock example

Example (compute degree of each vertex in a graph):

```
omp_lock_t lockvar[nvertices];
```

```
for (i=0; i<nvertexes; i++){  
    omp_init_lock(&lockvar[i]);  
}
```

```
#pragma omp parallel for  
    for (j=0; j<nedges; j++){  
        omp_set_lock(&lockvar[edge[j].vertex1]);  
        degree[edge[j].vertex1]++;  
        omp_unset_lock(&lockvar[edge[j].vertex1]);  
        omp_set_lock(&lockvar[edge[j].vertex2]);  
        degree[edge[j].vertex2]++;  
        omp_unset_lock(&lockvar[edge[j].vertex2]);  
    }
```



# Exercise: Molecular dynamics

- The code supplied is a simple molecular dynamics simulation of the melting of solid argon.
- Computation is dominated by the calculation of force pairs in subroutine **forces**.
- Parallelise this routine using a DO/FOR directive and critical sections.
  - Watch out for PRIVATE and REDUCTION variables.
  - Choose a suitable loop schedule
- Extra exercise: can you improve the performance by using locks, or atomics, or by using a reduction array.

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