Advanced Message-Passing Programming

Parallel Filesystems and Lustre











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Overview

- Lecture will cover
 - Parallel Filesystems
 - Lustre Filesystem
 - Striping
 - Simple Lustre commands
 - Bottlenecks





Parallel File Systems

Parallel computer

- constructed of many standard processors, each not particularly fast
- performance comes from using many processors at once
- requires manual distribution of data and calculation across processors

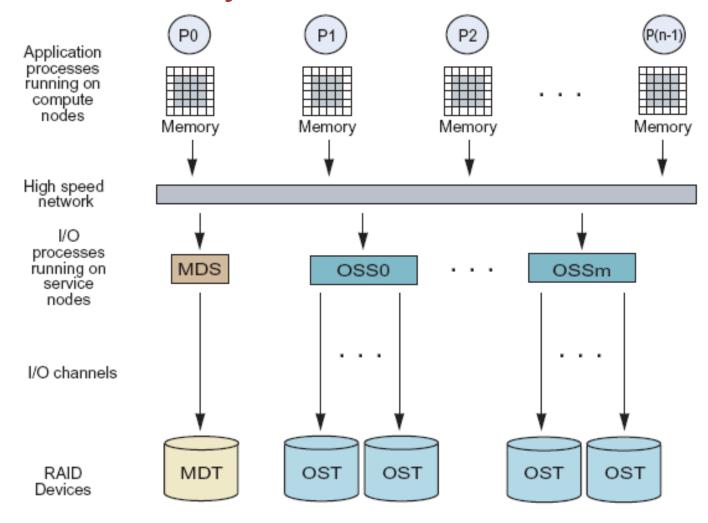
Parallel file systems

- constructed from many standard disks, each not particularly fast
- performance comes from reading / writing to many disks at once
- requires many *clients* to read / write to different disks
 - each node appears as a separate IO client
- data from a single file can be striped across many disks
- Must appear as a single file system to user
 - typically have a single *MetaData* Server (MDS)





Parallel File Systems: Lustre







ARCHER's (not ARCHER2) Cray

Sonexion Storage



SSU: Scalable Storage Unit

2 x OSSs and 8 x OSTs (Object Storage Targets)

- Contains Storage controller, Lustre server, disk controller and RAID engine
- Each unit is 2 OSSs each with 4 OSTs of 10 (8+2) disks in a RAID6 array



Multiple SSUs are combined to form storage racks



Terminology

- Lustre has many different levels and virtualisations
 - e.g. one Object Storage Server has multiple Object Storage Targets
 - a single OST has many physical disks in a RAID array
- I will refer to the following parts of Lustre
 - Meta Data Server (MDS)
 - the database that contains information on, e.g., where a file is stored
 - Object Storage Target
 - the physical device that stores your data
 - I may also call this a "disk" (although it contains multiple hard drives)
- The MDS and the OSTs are what a user interacts with





ARCHER2 hardware

- Three separate /work filesystems (work1, work2 & work3)
 - each has 12 OSTs and one MDS
 - consortia assigned to different partitions to share the load
 - multiple filesystems means the MDS is less likely to be overloaded
- One filesystem with Solid State (SSD) not spinning disks
 - still to be configured
 - expect better latency, e.g. good for small I/O transactions
- Each has around 3.3 PiB of storage
 - total of 13.2 PiB (which is 14.5 PB!)





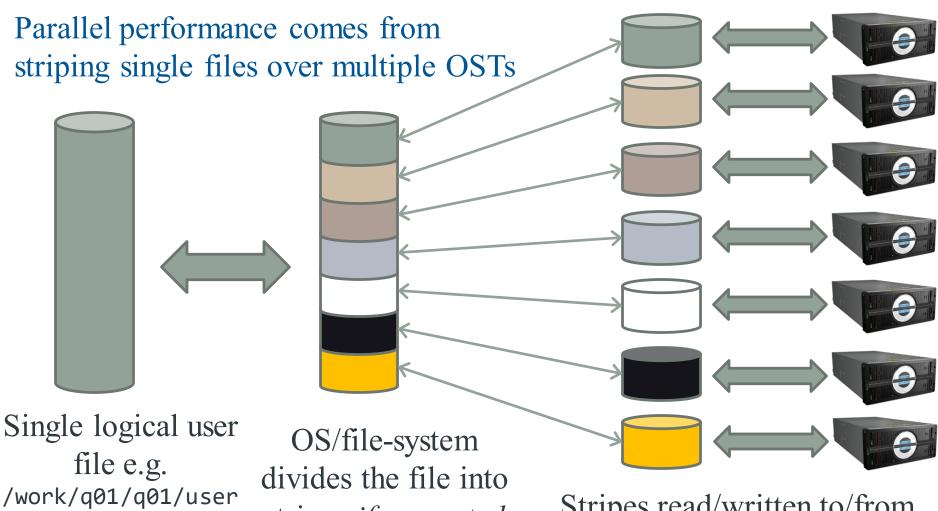
Default Configuration

- By default, each file is stored on a single OST
 - assigned when the file is created
 - distributed across all available OSTs to balance the load
 - each OST is actually a separate Linux filesystem
- This is called an "unstriped" file
- Reading and writing multiple files from multiple nodes can benefit from multiple OSTs
- Access to a single file will not benefit from the parallel nature of the filesystem





Lustre data striping



/bigfile.dat

stripes if requested by the user

Stripes read/written to/from their assigned OSTs

Striping

- Allow multiple IO processes to access same file
 - increases bandwidth as you are accessing multiple OSTs
- Typically optimised for bandwidth, not for latency
 - e.g. reading/writing small amounts of data is very inefficient
- This is called striping
 - striping of a file is fixed when it is created, under control of the user
 - fundamental parameters are the number of stripes and stripe size
- For example, if a file is created with 4 stripes
 - Lustre assigns four OTS: OST1, OST2, OST3, OST4
 - first MiB is stored on OST1, second on OST2, third on OST3, fourth on OST4, fifth on OST1, sixth on OST2,
 - i.e. round-robin with default stripe size of 1MiB





Lustre commands

To set the striping on a directory or file

```
lfs setstripe -c nstripe <dir/file>
```

- nstripe = -1 is full striping (12 on each of ARCHER2's 3 filesystems)
- Stripe size: lfs setstripe -s 4m <dir/file>
- Does not alter striping for existing files: that requires a copy
- I always use setstripe on directories
 - all files subsequently created in directory will have the same striping
- To enquire: lfs getstripe <dir/file>





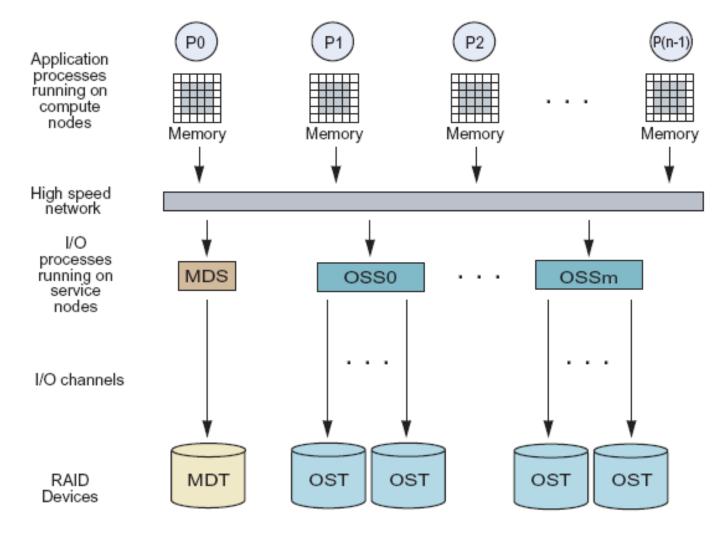
Parallel IO to a striped file

- Very complicated in practice!
 - where in the file does the local data need to be written?
 - which OSTs are the stripes located on?
 - are there write conflicts coming from different processes?
- Need to use a parallel IO library





Lustre: where are the bottlenecks?







Benchio benchmark

- Obvious questions:
 - does the MDS become overloaded for large numbers of files?
 - what is the maximum performance of a single OST?
 - can one process saturate an OST? can a node saturate an OST?
 - or is the network the limiting factor
 - how well do different IO libraries work with Lustre?
 - what are the best stripe count (and size) settings?

-

- I wrote a simple benchmark to help investigate Lustre performance characteristics and bottlenecks
 - we will use benchio for the practical examples
 - writes a large distributed 3D array of double precision numbers
 - https:/github.com/davidhenty/benchio/.





Cellular Automaton Model

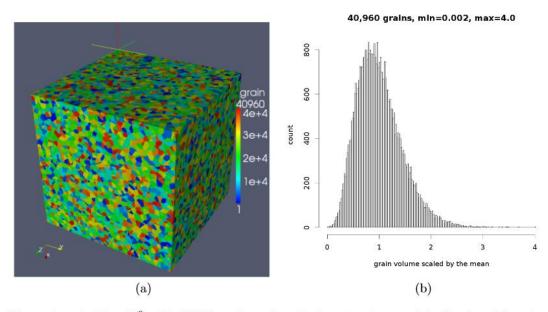


Figure 1: A 4.1×10^9 cell, 40,960 grain equiaxed microstructure model, showing (a) grain arrangement with colour denoting orientation; (b) grain size size (volume) histogram.

 Fortran coarray library for 3D cellular automata microstructure simulation, Anton Shterenlikht, proceedings of 7th International Conference on PGAS Programming Models, 3-4 October 2013, Edinburgh, UK.





Summary

- A Lustre filesystem has multiple OSTs
 - I think of these as being multiple disks
- By default on ARCHER2, each file stored on a single OST
 - i.e. an unstriped file with a stripe count of 1
 - increased performance for multiple files
 - a single user writing many files
 - multiple users each writing a single file
- Improving performance for a single file requires striping
 - fully under control of the user
 - expect parallel IO libraries to take advantage of striping



