Advanced Message-Passing Programming

Parallel Filesystems and Lustre

Reusing this material

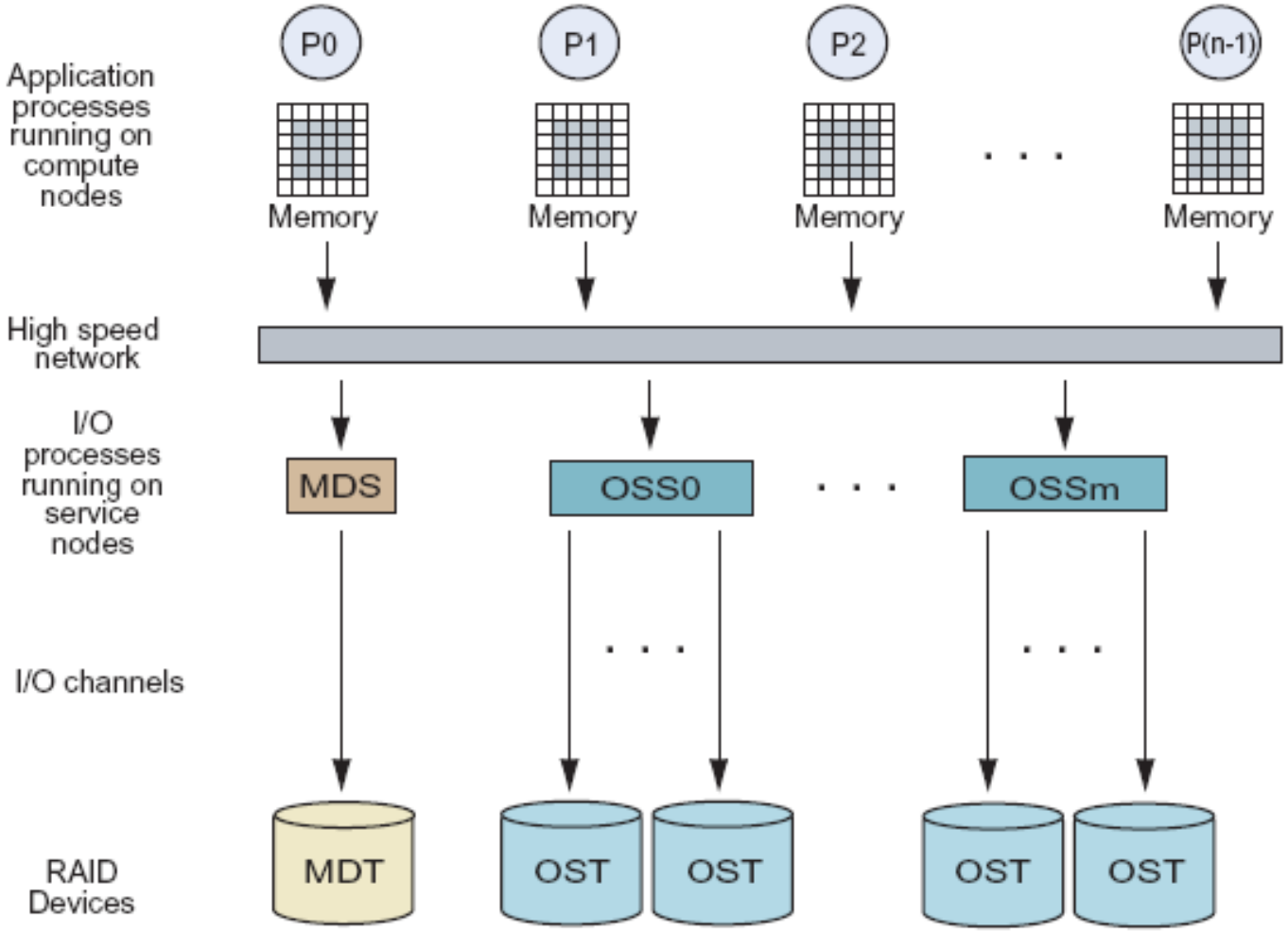
Overview

- Lecture will cover
  - Parallel Filesystems
  - Lustre Filesystem
  - Striping
  - Simple Lustre commands
  - Bottlenecks

Parallel File Systems

- Parallel computer
  - constructed of many standard processors, each not particularly fast
  - performance comes from using many processors at once
  - requires manual distribution of data and calculation across processors
- Parallel file systems
  - constructed from many standard disks, each not particularly fast
  - performance comes from reading / writing to many disks at once
  - requires many *clients* to read / write to different disks
    - each *node* appears as a separate IO client
  - data from a single file can be *striped* across many disks
- Must appear as a single file system to user
  - typically have a single *MetaData* Server (MDS)

Parallel File Systems: Lustre



ARCHER's (**not** ARCHER2) Cray Sonexion Storage

SSU: *Scalable Storage Unit*

2 x OSSs and 8 x OSTs (Object Storage Targets)

- Contains Storage controller, Lustre server, disk controller and RAID engine
- Each unit is 2 OSSs each with 4 OSTs of 10 (8+2) disks in a RAID6 array

Multiple SSUs are combined to form storage racks

Terminology

- Lustre has many different levels and virtualisations
  - e.g. one Object Storage Server has multiple Object Storage Targets
  - a single OST has many physical disks in a RAID array
- I will refer to the following parts of Lustre
  - Meta Data Server (MDS)
    - the database that contains information on, e.g., where a file is stored
  - Object Storage Target
    - the physical device that stores your data
    - I may also call this a “disk” (although it contains multiple hard drives)
- The MDS and the OSTs are what a user interacts with

ARCHER2 hardware

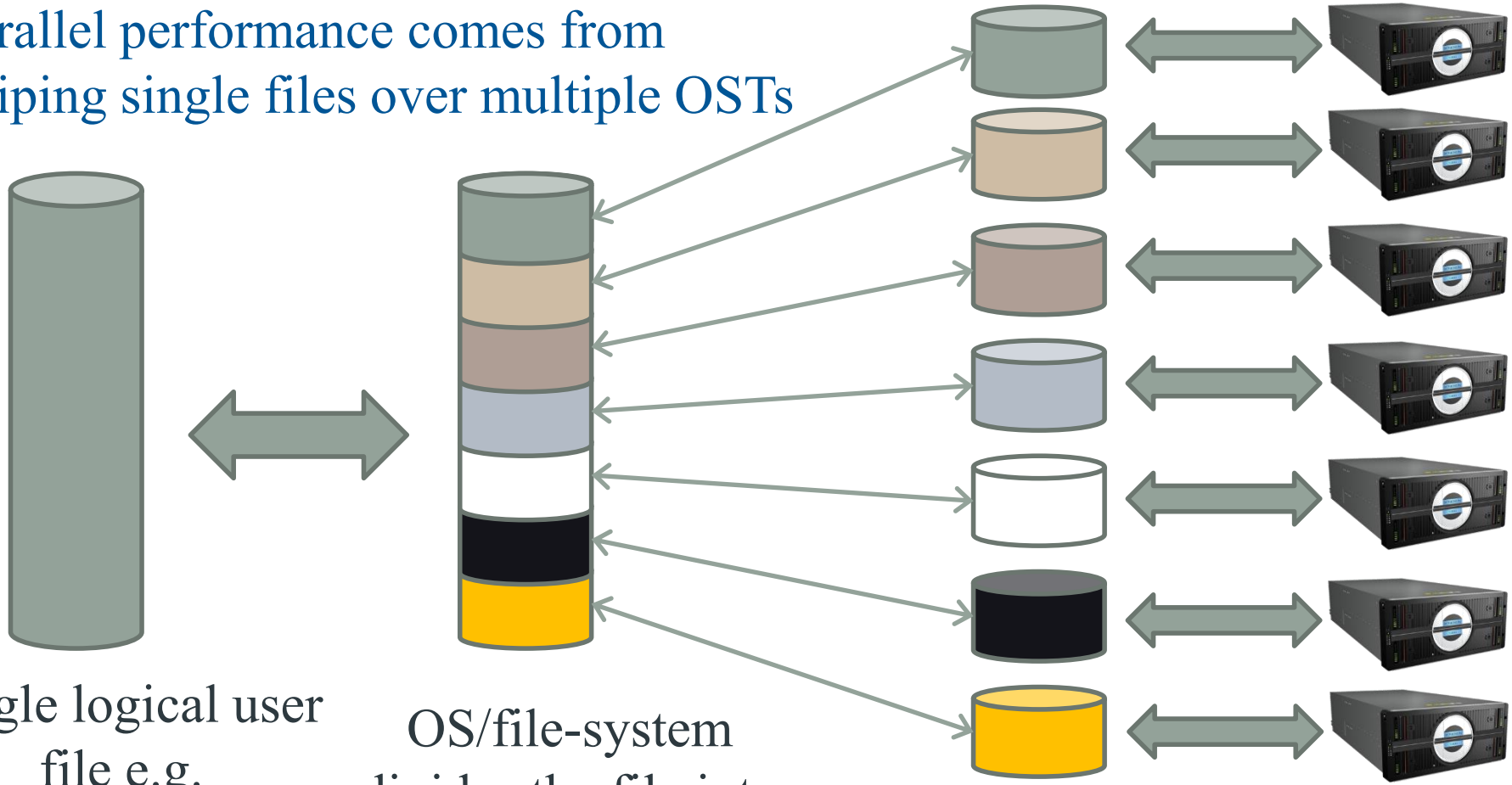
- Three separate /work filesystems (work1, work2 & work3)
  - each has 12 OSTs and one MDS
  - consortia assigned to different partitions to share the load
  - multiple filesystems means the MDS is less likely to be overloaded
- One filesystem with Solid State (SSD) not spinning disks
  - still to be configured
  - expect better latency, e.g. good for small I/O transactions
- Each has around 3.3 PiB of storage
  - total of 13.2 PiB (which is 14.5 PB !)

Default Configuration

- By default, each file is stored on a single OST
  - assigned when the file is created
  - distributed across all available OSTs to balance the load
  - each OST is actually a separate Linux filesystem
- This is called an “unstriped” file
- Reading and writing multiple files from multiple nodes can benefit from multiple OSTs
- Access to a single file will not benefit from the parallel nature of the filesystem

Lustre data striping

Parallel performance comes from striping single files over multiple OSTs



Single logical user file e.g. `/work/q01/q01/user/bigfile.dat`

OS/file-system divides the file into stripes *if requested by the user*

Stripes read/written to/from their assigned OST

epcc

Striping

- Allow multiple IO processes to access same file
  - increases bandwidth as you are accessing multiple OSTs
- Typically optimised for bandwidth, not for latency
  - e.g. reading/writing small amounts of data is very inefficient
- This is called striping
  - striping of a file is fixed when it is created, under control of the user
  - fundamental parameters are the number of stripes and stripe size
- For example, if a file is created with 4 stripes
  - Lustre assigns four OTS: OST1, OST2, OST3, OST4
  - first MiB is stored on OST1, second on OST2, third on OST3, fourth on OST4, fifth on OST1, sixth on OST2, ...
  - i.e. round-robin with default stripe size of 1MiB

Lustre commands

- To set the striping on a directory or file

  `lfs setstripe -c nstripe <dir/file>`

  - nstripe = -1 is full striping (12 on each of ARCHER2's 3 filesystems)
  - Stripe size: **`lfs setstripe -s 4m <dir/file>`**

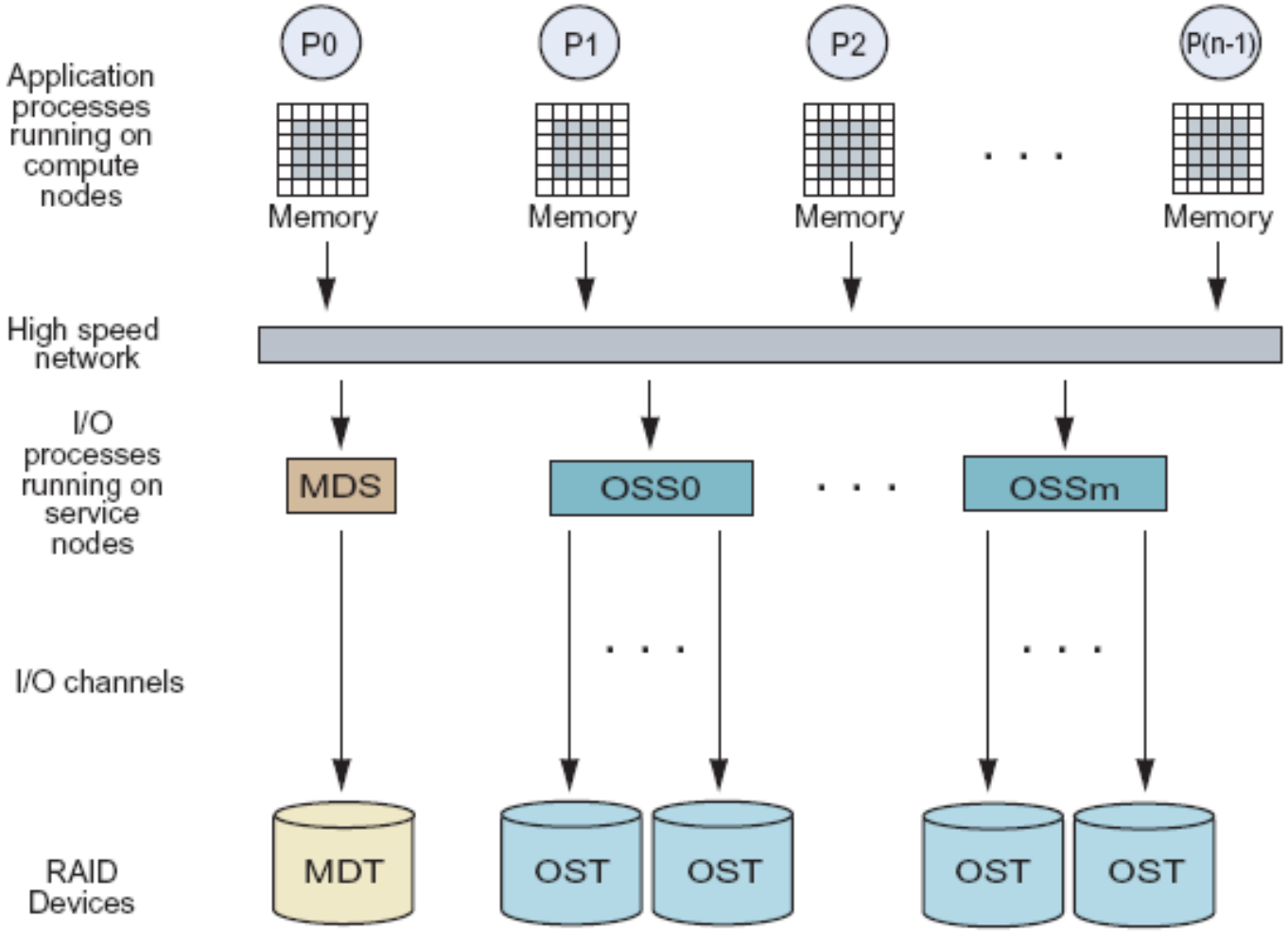
- Does **not** alter striping for existing files: that requires a copy
- I always use setstripe on directories
  - all files subsequently created in directory will have the same striping
- To enquire: **`lfs getstripe <dir/file>`**

Parallel IO to a striped file

- Very complicated in practice!
  - where in the file does the local data need to be written?
  - which OSTs are the stripes located on?
  - are there write conflicts coming from different processes?

- Need to use a parallel IO library

Lustre: where are the bottlenecks?



Benchio benchmark

- Obvious questions:
  - does the MDS become overloaded for large numbers of files?
  - what is the maximum performance of a single OST?
  - can a single process saturate an OST?
  - can a single node saturate an OST?
    - or is the network the limiting factor
  - how well do different IO libraries work with Lustre?
  - what are the best stripe count (and size) settings?
  - ....
- I wrote a simple benchmark to help investigate Lustre performance characteristics and bottlenecks
  - we will use benchio for the practical examples

Summary

- A Lustre filesystem has multiple OSTs
  - I think of these as being multiple disks
- By default on ARCHER2, each file stored on a single OST
  - i.e. an unstriped file with a stripe count of 1
  - increased performance for multiple files
    - a single user writing many files
    - multiple users each writing a single file
- Improving performance for a single file requires striping
  - fully under control of the user
  - expect parallel IO libraries to take advantage of striping