

Linked Lists

03/10/2021

```
struct person
```

```
{ char name[20];  
  struct person * next;  
};
```

} node
=> a self-referential structure

```
struct person * new;
```

```
struct person * head;
```

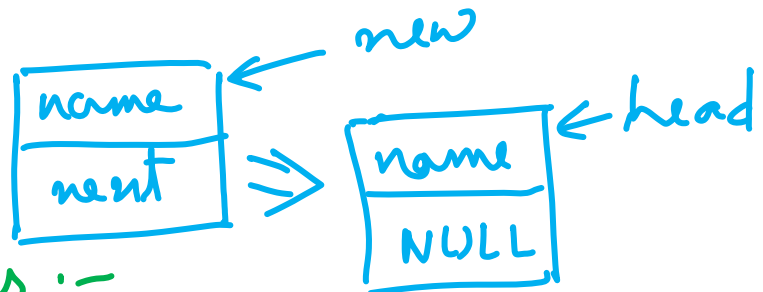
```
head = NULL;
```

Adding the first element of a linked list

```
new = (struct person *) malloc(sizeof(struct person));
```

```
new->next = head;
```

```
head = new;
```

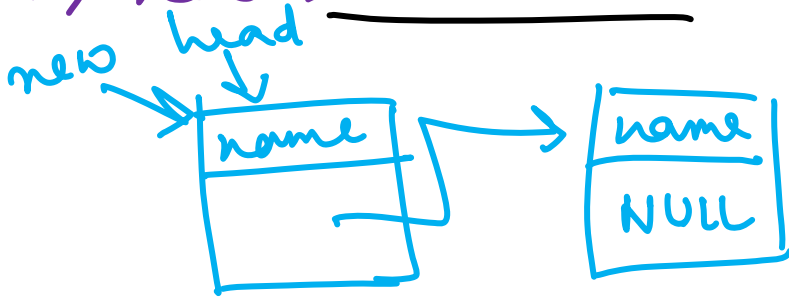


Insertion cases :-

- Adding up node at the beginning of the list
- Adding up node in an intermediate position.
- Adding up node at the last end of the linked list.

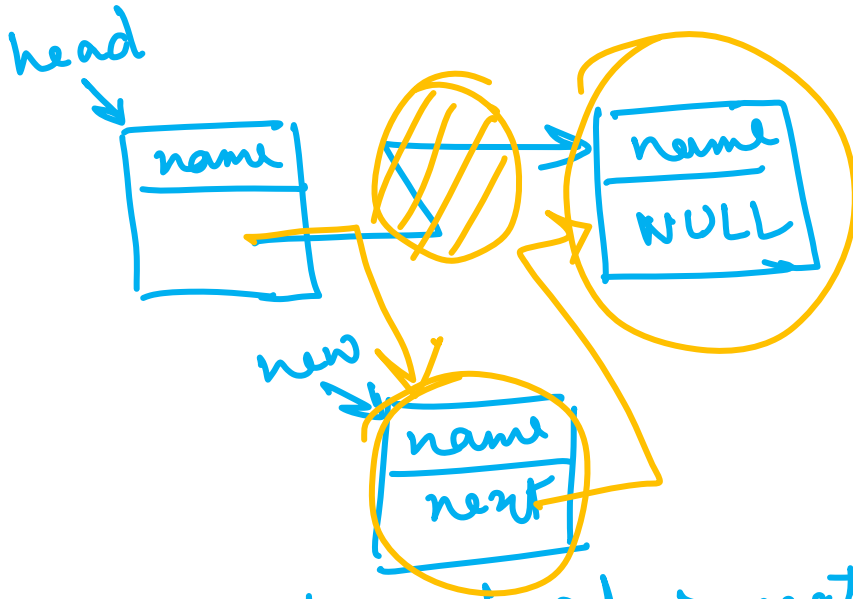
$new = (\text{struct person} *) \text{malloc}(\text{sizeof}(\text{struct person}));$

$new \rightarrow \text{next} = \text{head};$

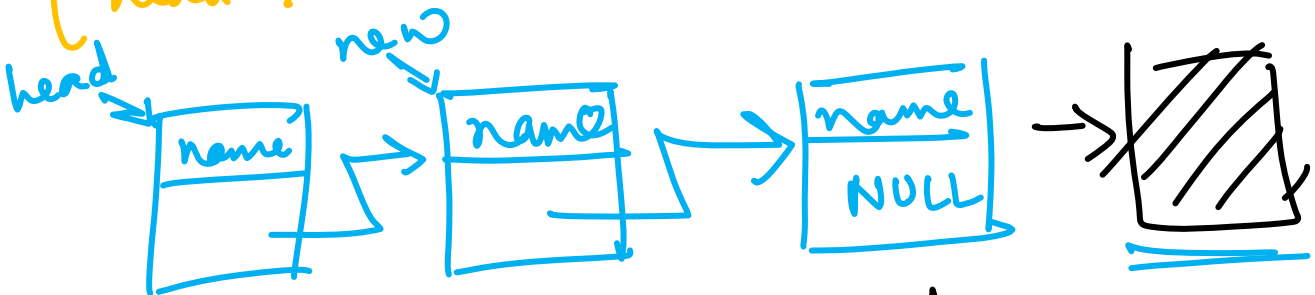


$\text{head} = \text{new};$

$new = (\text{struct person} *) \text{malloc}(\text{sizeof}(\text{struct person}));$




$\left\{ \begin{array}{l} new \rightarrow \text{next} = \text{head} \rightarrow \text{next}; \\ \text{head} \rightarrow \text{next} = \text{new}; \end{array} \right.$



pseudo code : similar to program statements but not a complete program.

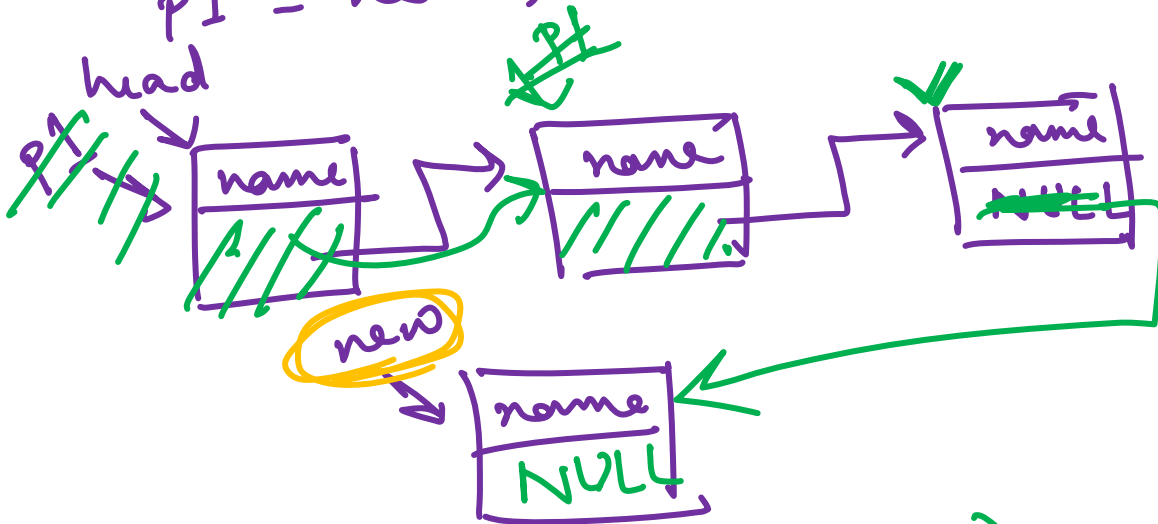
`new = (struct person*) malloc(sizeof(struct person));`



A diagram showing a single node structure. It is a rectangle divided into two horizontal sections. The top section is labeled 'name' and the bottom section is labeled 'next'. An arrow labeled 'new' points to the top section.

`struct person * p1;`

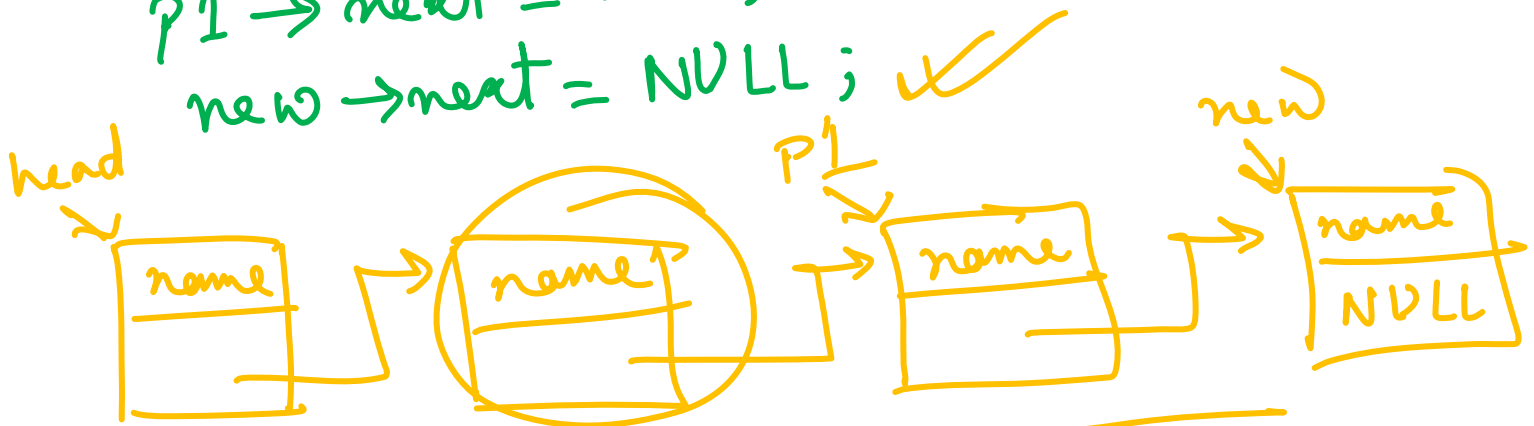
`p1 = head;`



`while (p1->next != NULL)`
`p1 = p1->next;`

`p1->next = new;`

`new->next = NULL;` ✓



We have created a linked list with
4-nodes with insertions at the

- (a) beginning
- (b) intermediate position
- (c) at the end

a single-linked list / Singly linked list.
 ⇓
 uni-directional linked list.

Deletion of a node from a linked list

07/10/2021

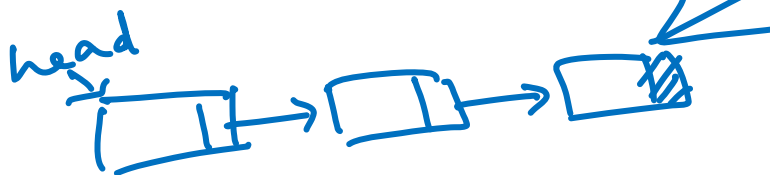
Case 1 :- Deletion of the 1st element / starting node

struct person * temp,

temp = head,

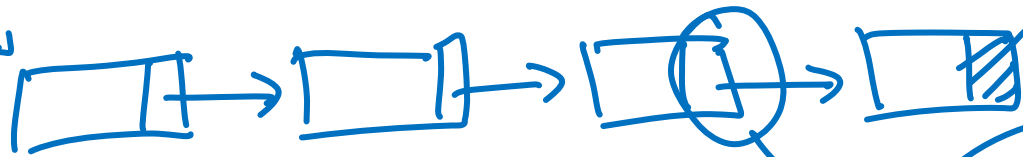
head = head → next;

free(temp); // frees the dynamically allocated memory / node pointed by temp



Case 2: — Deletion of the last node/element

head



struct person * current1;

struct person * current2;

current1 = head;

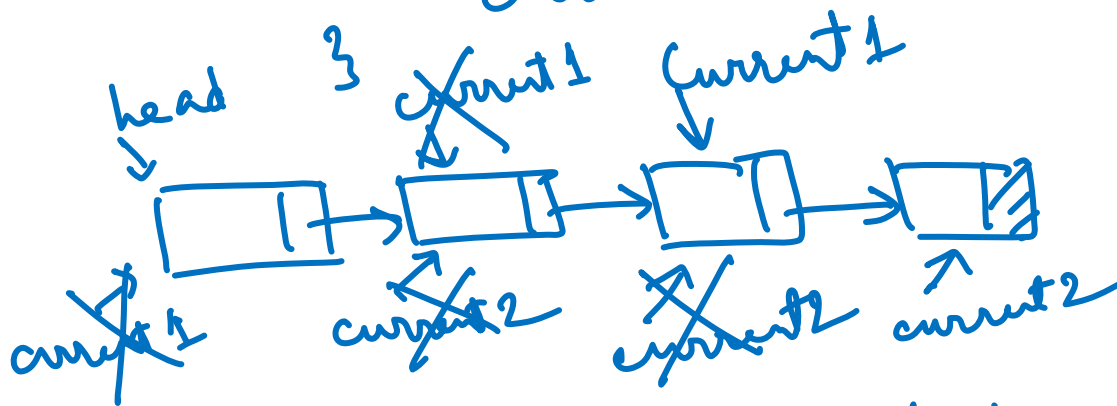
current2 = current1 -> next;

while (current2 -> next != NULL)

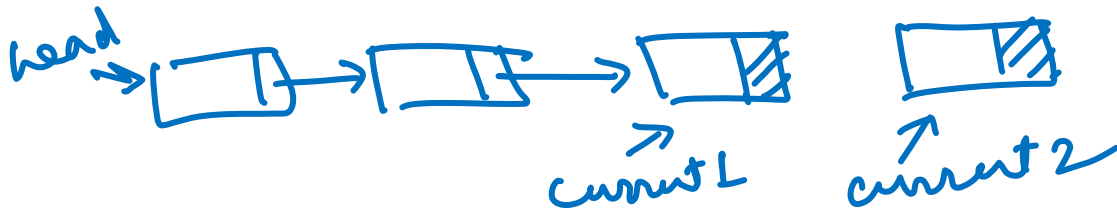
{

current1 = current2;

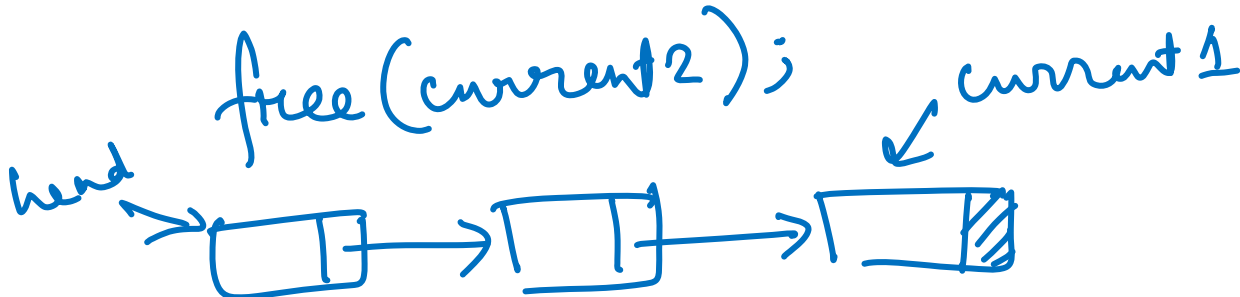
current2 = current1 -> next;



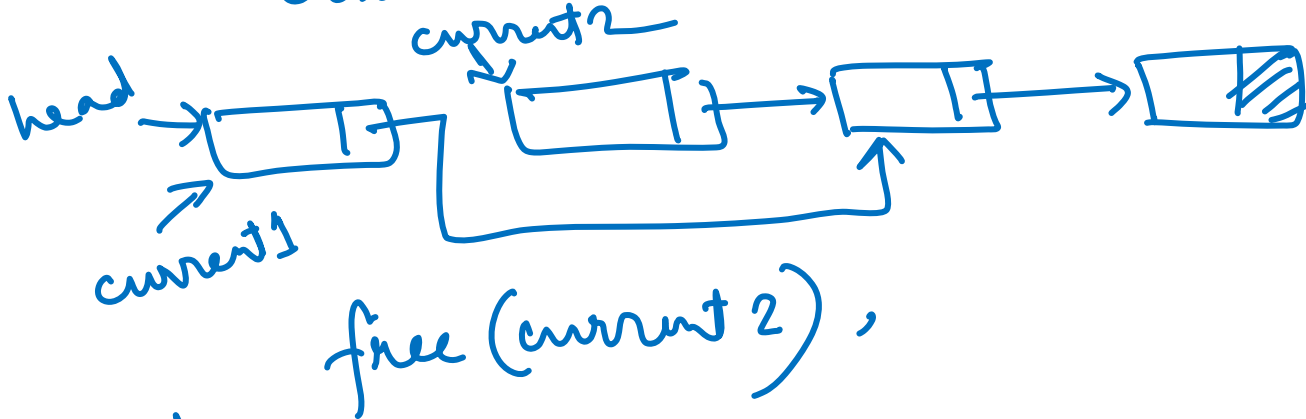
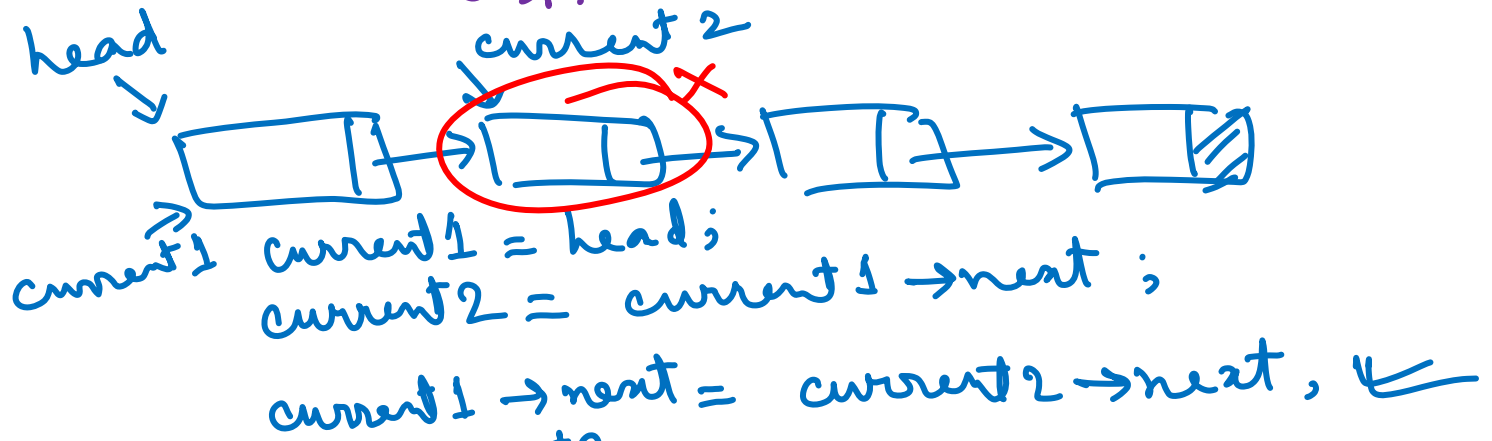
current1 -> next = NULL;



free(current2);

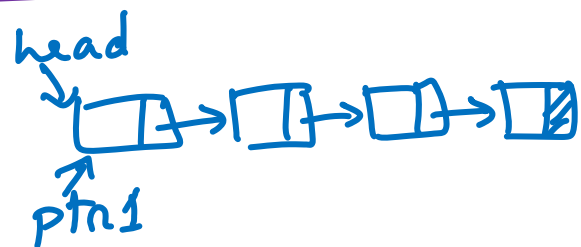


Case 3 :- Deletion of an intermediate node/element from a linked list.



Case 4 :- Deletion of the whole linkedlist

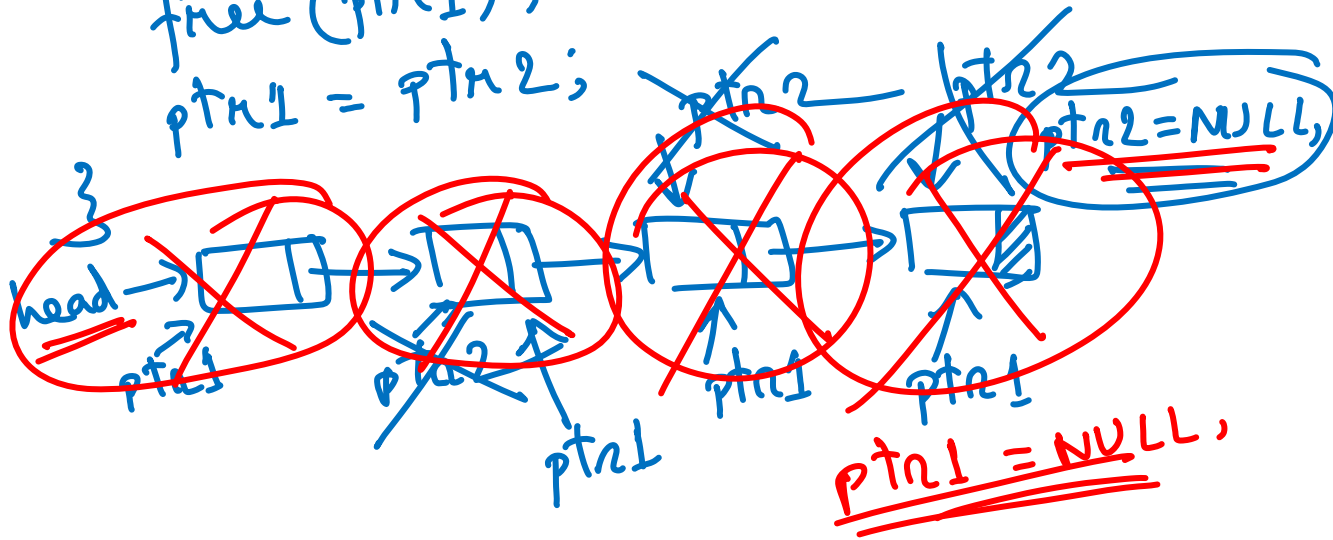
```
struct person * ptr1;  
struct person * ptr2;  
ptr1 = head;  
while (ptr1 != NULL)
```



```

{
    ptr2 = ptr1 → next,
    free(ptr1);
    ptr1 = ptr2;
}

```

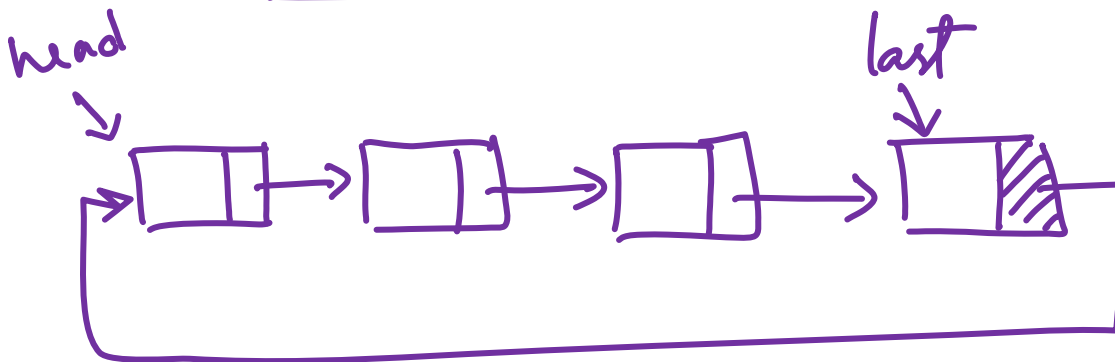


head = NULL;

— X —

Circular Linked Lists

26/10/2021



last → next = head;

Insertion Cases :

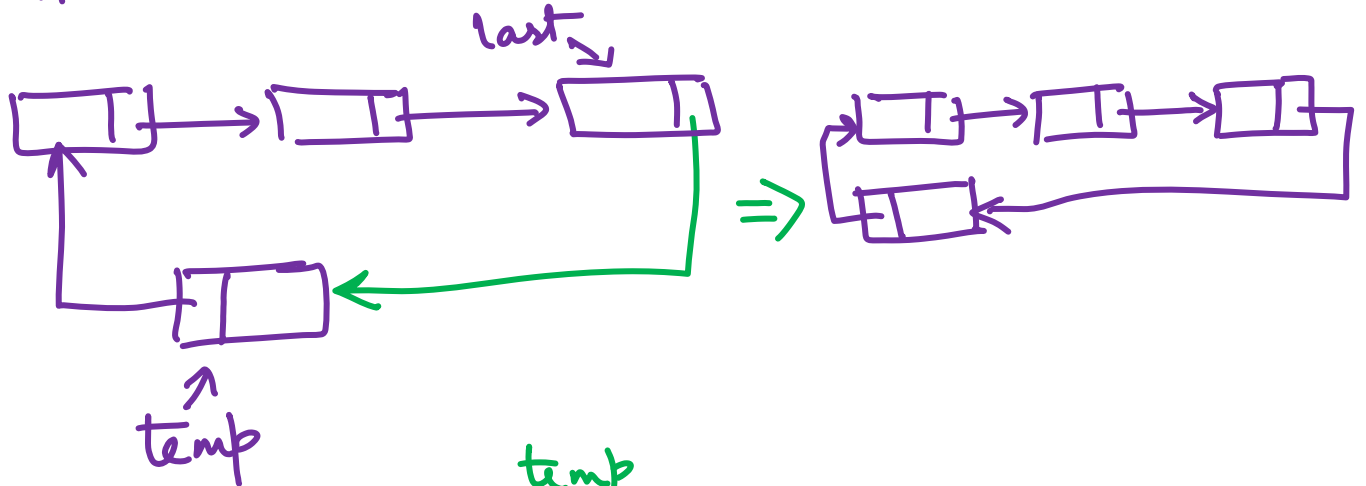
- (1) Insertion at beginning
- (2) Insertion in between

Case-1 :-

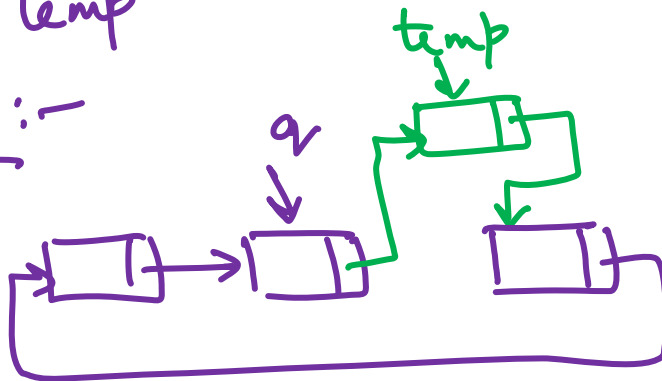
temp = (struct node *) malloc(sizeof(struct node));

temp → next = last → next;

last → next = temp;



Case 2 :-



temp = (struct node *) malloc (sizeof (struct node));

struct node

{ int data;

struct node * next;

};

temp → next = q → next

q → next = temp; ✓

Deletion within Circular Linked List

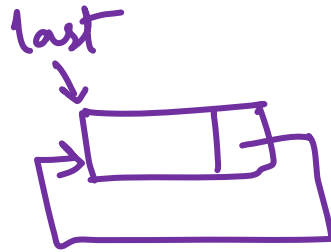
1) If the list has only one node

2) Node to be deleted is the 1st node.

3) Deletion in-between

4) Node to be deleted is the last node

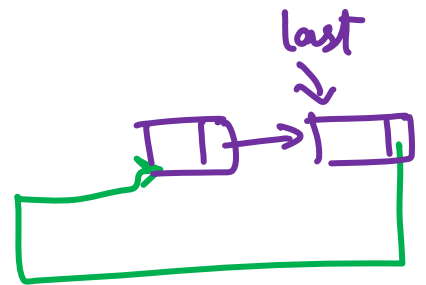
Case-1:



```
if (last->next == last)
{
    temp = last;
    last->next = NULL;
    free(temp);
}
```

Case-2 :-

$q = \text{last} \rightarrow \text{next};$

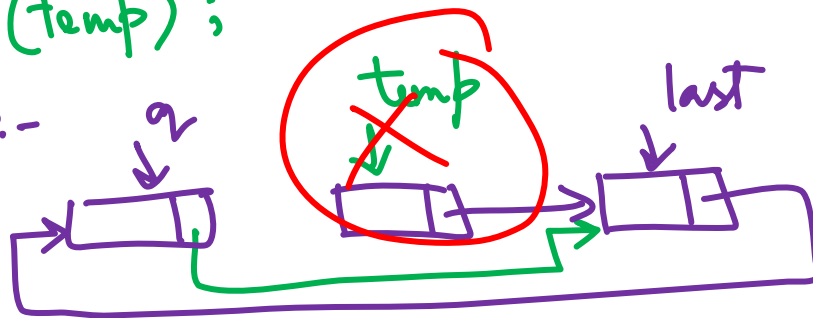


$\text{temp} = q;$

$\text{last} \rightarrow \text{next} = q \rightarrow \text{next};$ ✓

$\text{free}(\text{temp});$

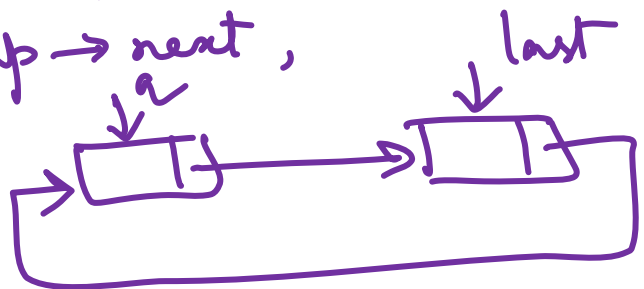
Case-3 :-



$\text{temp} = q \rightarrow \text{next};$

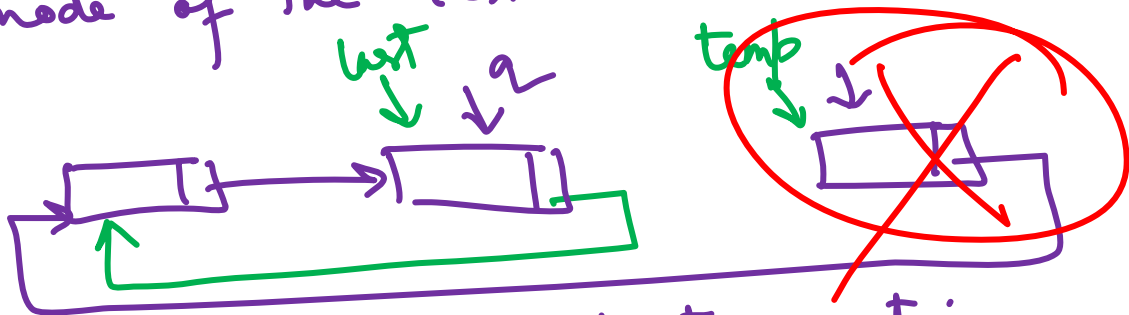
$q \rightarrow \text{next} = \text{temp} \rightarrow \text{next};$

$\text{free}(\text{temp});$



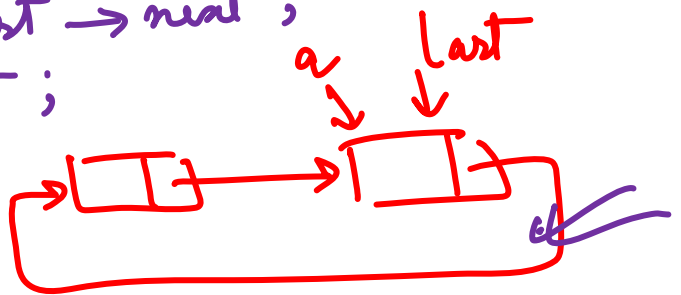
Case 4:-

Let us consider q points to the previous node of the last node



$q \rightarrow \text{next} = \text{last} \rightarrow \text{next};$
 $\text{temp} = \text{last};$

$\text{last} = q;$
 $\text{free}(\text{temp});$

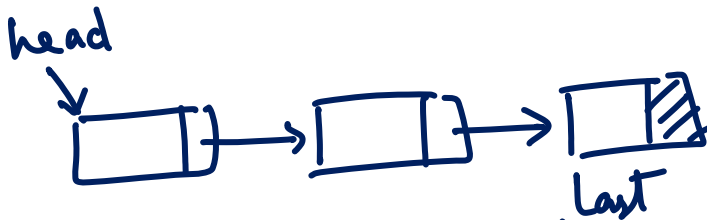


— X —

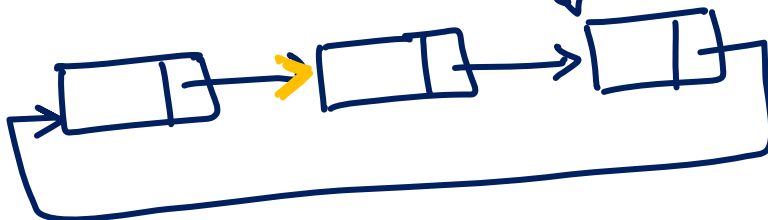
Double Linked List

28/10/2021

Type 1:



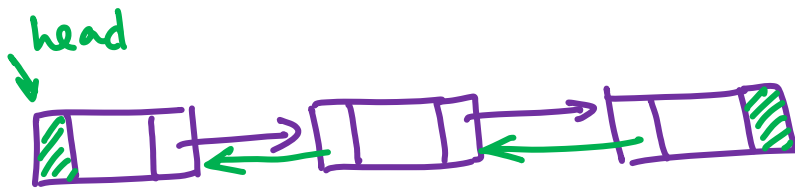
Type 2:



"Singly" linked list
 \hookrightarrow single direction

Circular Linked List (also a type of Singly linked list)

Double linked list
 or
 Doubly linked list \rightarrow Bi-directional Traversal within a linked list



struct node

```
{
    struct node * prev,
    int data;
    struct node * next;
};
```

Node becomes bulkier
More bytes needed.

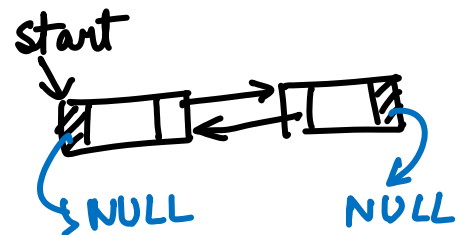


A. Insertion within a doubly linked list

- Insertion at beginning
- Insertion in between
- Insertion at the end

Case - (a) :

```
struct node * start;
struct node * temp;
```



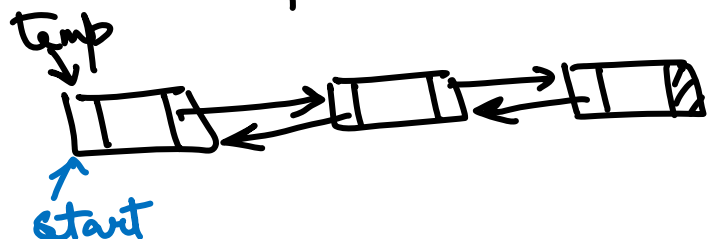
```
temp = (struct node *) malloc (sizeof(struct node ));
```

```
temp -> next = start;
```

```
start -> prev = temp;
```

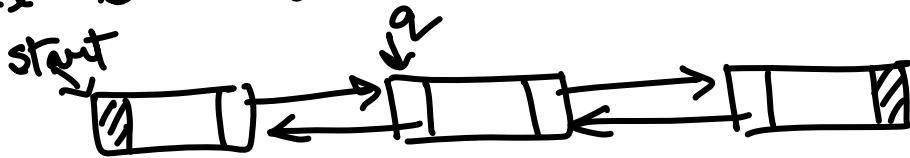
```
start = temp;
```

```
start -> prev = NULL;
```



Case - (b) :

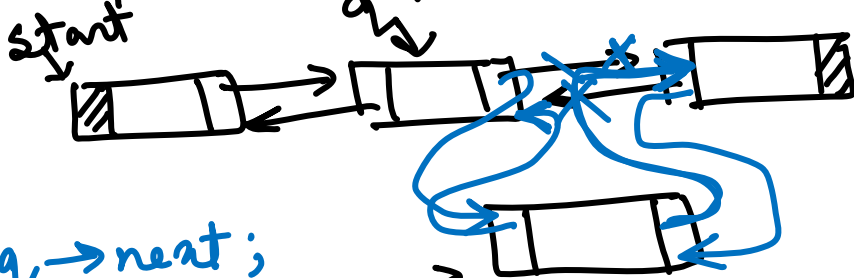
Let there be a pointer 'q' after which a new node shall be inserted then,



struct node * temp ,

temp = (struct node *) malloc (sizeof (struct node));

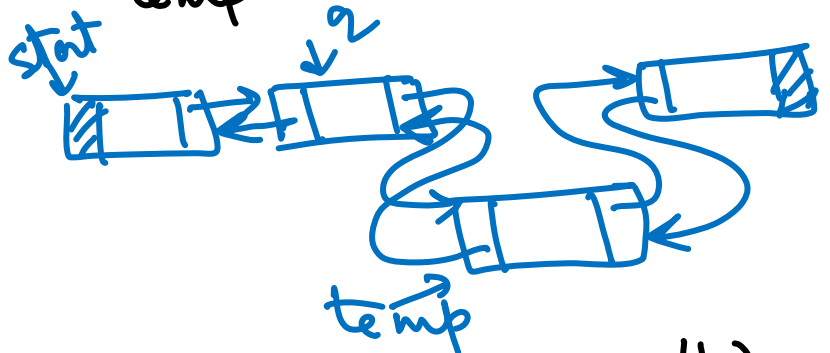
q → next → prev = temp ,



temp → next = q → next ;

temp → prev = q ;

q → next = temp ;



Case - (c) : it is a sub-case of case - (b) .

if (q → next == NULL)

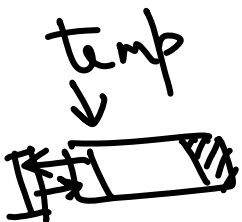
{

temp → next = q → next ,

temp → prev = q ,

q → next = temp ;

}



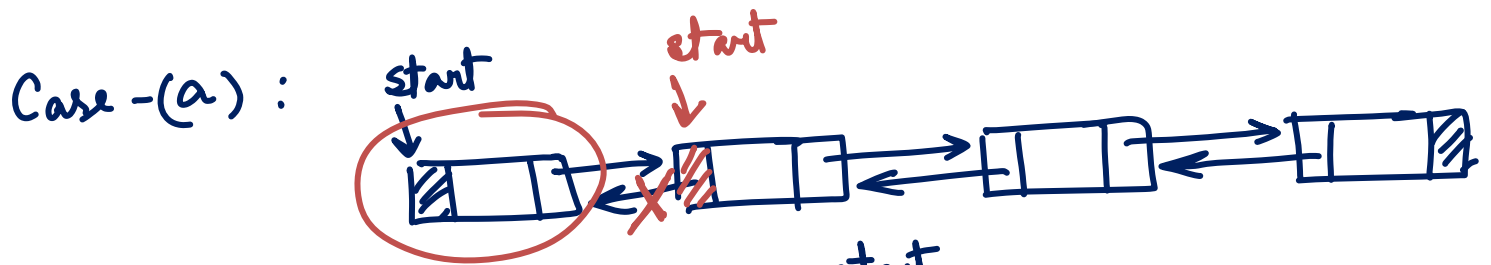
B. Deletion within a Doubly linked list

dated - 29/10/2021

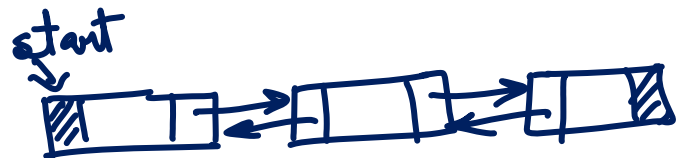
(a) Deletion at the beginning

(b) Deletion in-between

(c) Deletion at the end

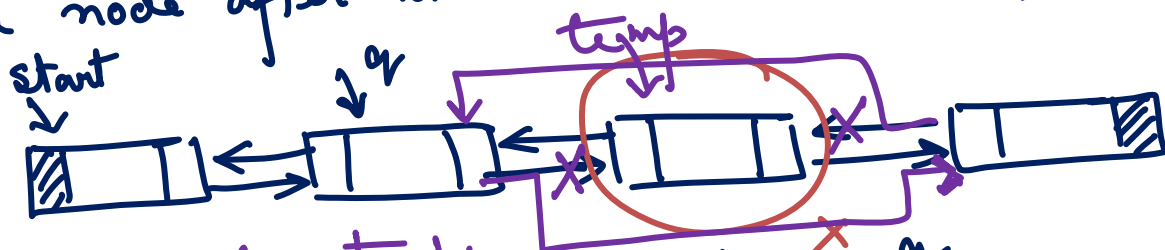


```
struct node * temp;  
temp = start;  
start = start->next;  
start->prev = NULL;  
free(temp);
```

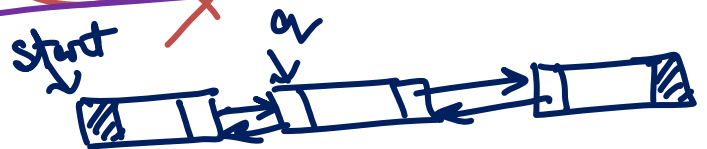


Case -(b):

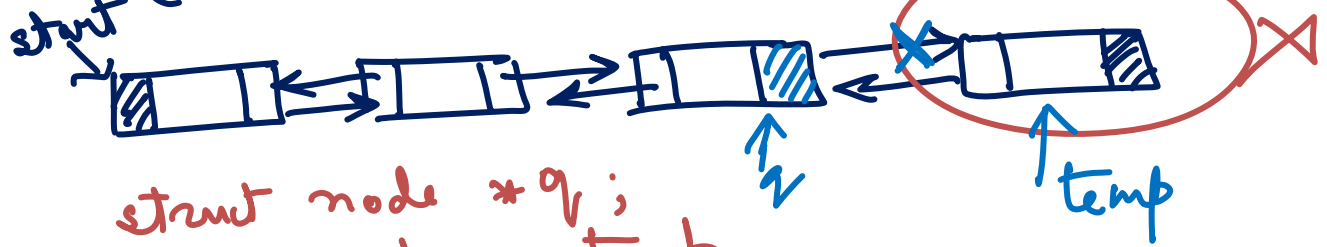
Let us assume, a struct node * q points to the node after which deletion takes place.



```
struct node * temp;  
temp = q->next;  
temp->next->prev = q;  
q->next = temp->next;  
free(temp);
```



Case - (c) :



```
struct node *q;
struct node *temp;
```

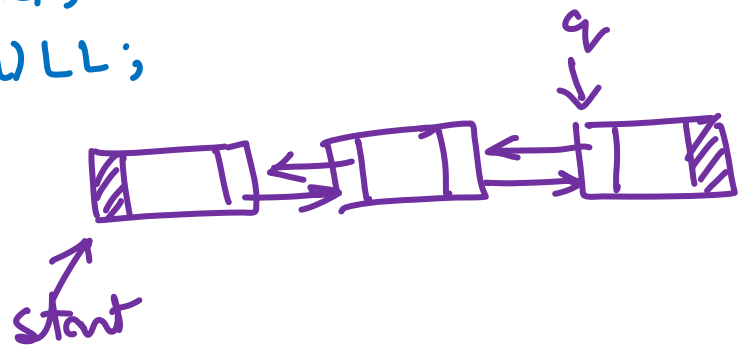
let us assume, that q points to the node after which the deletion takes place i.e. q points to the second-last node.

```
q = start;
while (q->next->next != NULL)
```

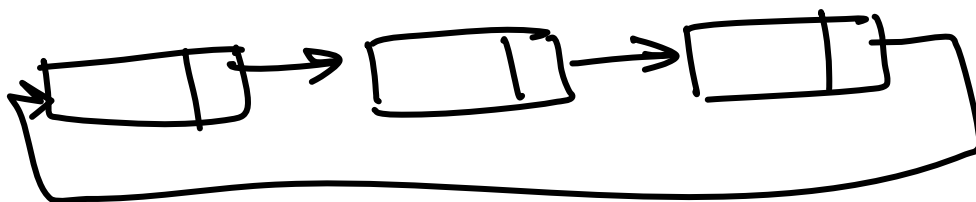
```
q = q->next;
```

```
temp = q->next;
q->next = NULL;
```

```
free(temp);
```

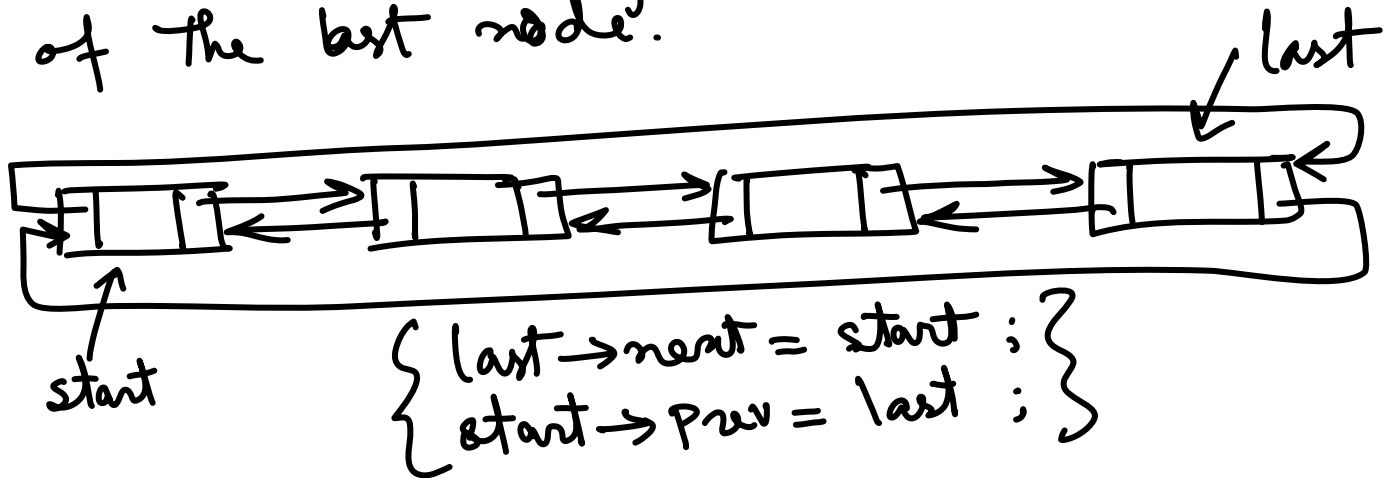


^{linked} Singly list can be circular if the last node points to the first node.



Similarly, Doubly linked list can also be converted into a Circular - Doubly linked list by making the last node's next field holding the address of the first node and the

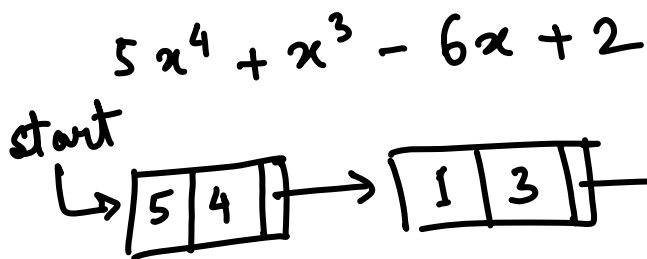
first node's prev field holding the address of the last node.



—X—

Polynomial Arithmetic With Linked List

Dated on 31/10/2021

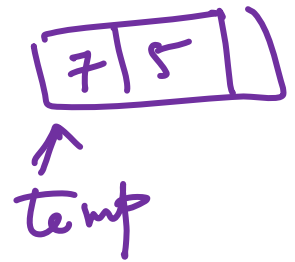
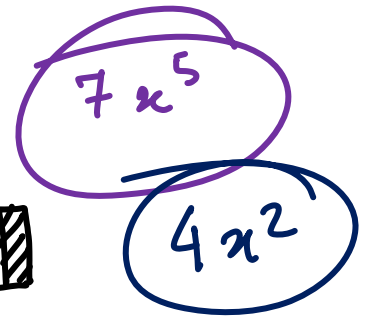


struct node

```

{
    int coeff;
    int exp;
    struct node * link;
};

```



if (start == NULL || start->exp < temp->exp)

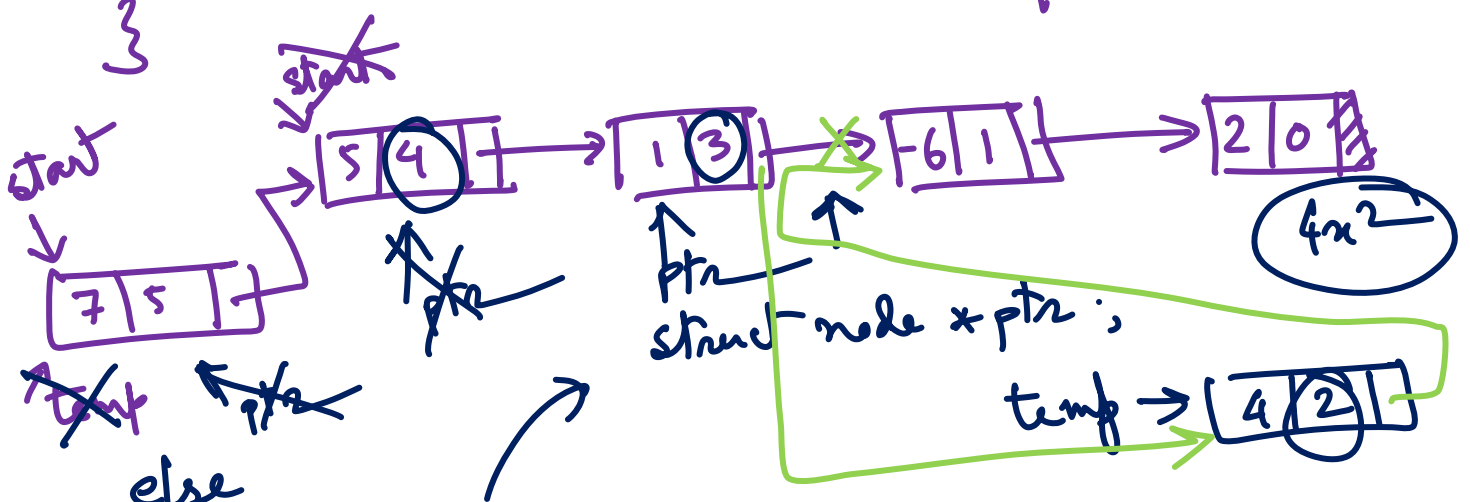
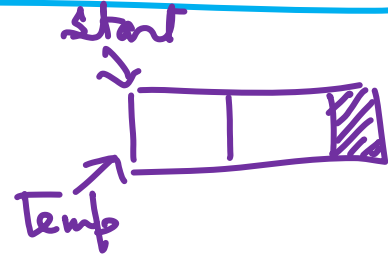
struct node *temp;

temp = (struct node*) malloc (sizeof(struct node));

scanf("... ", temp->exp);

scanf(" , temp → coeff);

```
{
    temp → link = start;
    start = temp;
}
```



```
else
{
```

ptr = start;

while (ptr → link != NULL && ptr → link → exp > temp → exp)

ptr = ptr → link;

```
temp → link = ptr → link;
ptr → link = temp;
```

$7x^3 + 2x^2 - 3x$ (circled) → $+9$ (circled)

if (ptr → link == NULL)
temp → link = NULL;



```
}
```