

## Question 2

### 2.1

Number of customers  $n = 3$ .

For  $n = 1$ ,  $s_1 = 1$ ,  $e_1 = 100$ ,  $a_1 = 1$ ,  $b_1 = 1$  and  $c_1 = 1$ .

For  $n = 2$ ,  $s_2 = 2$ ,  $e_2 = 100$ ,  $a_2 = 1$ ,  $b_2 = 1$  and  $c_2 = 1$ .

For  $n = 3$ ,  $s_3 = 3$ ,  $e_3 = 100$ ,  $a_3 = 1$ ,  $b_3 = 1$  and  $c_3 = 100000$ .

Running the greedy algorithm, we process each customer in order of arrival.

For customer 1,  $a_1 = b_1$  so we assign that customer to Alice.

For customer 2,  $a_2 = b_2$ , so we assign that customer to Bob.

For customer 3, there is no mechanic available, as both Alice and Bob are busy until they finish their current job, which ends at  $e_1 = e_2 = 100$ . So, customer 3 will not be served.

Customer 1 and customer 2 will make 1 dollar each having been served by Alice and Bob respectively. Customer 3 will lose the store 100000 dollars having not been served at all. Thus, the greedy algorithm will lose the store 99998 dollars total.

A higher net figure can be made by cancelling the job Alice/Bob is on and serving the third customer. When doing this, the total earned will remain the same (2 dollars) while the total lost will reduce to 1 dollar. Thus, the net earnings will be 1 dollar, which is higher than the greedy algorithm's net earnings of negative 99998 dollars.

### 2.2

#### Subproblem:

For  $0 \leq i, j \leq n$ , let  $P(i, j)$  be the problem of determining  $\text{opt}(i, j)$ , the maximum total earnings that can be achieved using up to  $i$  customers and using only the first  $j$  customers.

#### Recurrence:

$$\text{opt}(i, j) = \max(\text{opt}(i, j - 1) - c_i, \text{opt}(i, j - 1) + a_i, \text{opt}(i, j - 1) + b_i, \text{opt}(i - 1, j - 1) + a_i - (a_k + c_k), \text{opt}(i - 1, j - 1) + b_i - (b_k + c_k)).$$

There are five options when processing a new  $\text{opt}(i)$ :

1. Nobody is available and you don't serve the new customer, losing  $c_i$  dollars.
2. Alice is available and serves the new customer, making  $a_i$  dollars.
3. Bob is available and serves the new customer, making  $b_i$  dollars.
4. Alice doesn't serve her current customer and serves the new customer, losing  $(a_k + c_k)$  dollars and making  $a_i$  dollars.
5. Bob doesn't serve his current customer and serves the new customer, losing  $(b_k + c_k)$  dollars and making  $b_i$  dollars.

#### Base case:

$\text{opt}(0, 0) = 0$ , Alice and Bob don't have any current customers.

Order of computation:

Ascending order of  $j$  then  $i$ . This guarantees subproblems are completed when we reach  $P(i,j)$ .

Final answer:

$\text{opt}(n,n)$ .

Time Complexity:

$O(n^2)$  since for each  $i$ , order  $O(1)$  steps are carried out  $j$  times.