

## Abstract Interpretation

Software Quality Assurance - Static Code Analysis, II | Florian Sihler | December 9, 2024

# The Why

```
public static void main(String[] args) {  
    int a = 1;  
    double r = Math.random() * 10;  
    if (r > 5) {  
        a = 2;  
    }  
    System.out.println(a);  
}
```

# The Why

```
public static void main(String[] args) {  
    int a = 1;  
    double r = Math.random() * 10;  
    if (r > 5) {  
        a = 2;  
    }  
    System.out.println(a);  
}
```

# The Why

```
public static void main(String[] args) {  
    int a = 1;  
    double r = Math.random() * 10;    {a ∈ {1}}  
    if (r > 5) {  
        a = 2;  
    }  
    System.out.println(a);  
}
```

# The Why

```
public static void main(String[] args) {  
    int a = 1;  
    double r = Math.random() * 10;    {  $a \in \{1\}$  }  
    if (r > 5) {                        {  $r \in [0..10)$  }  
        a = 2;  
    }  
    System.out.println(a);  
}
```

# The Why

```
public static void main(String[] args) {  
    int a = 1;  
    double r = Math.random() * 10;    { a ∈ {1} }  
    if (r > 5) {                        { r ∈ [0..10) }  
        a = 2;  
    }                                  { a ∈ {2} }  
    System.out.println(a);  
}
```

# The Why

```
public static void main(String[] args) {  
    int a = 1;  
    double r = Math.random() * 10;    {  $a \in \{1\}$  }  
    if (r > 5) {                       {  $r \in [0..10)$  }  
        a = 2;  
    }                                  {  $a \in \{2\}$  }  
    System.out.println(a);  
}
```

# The Why

```
public static void main(String[] args) {  
    int a = 1;  
    double r = Math.random() * 10;    {  $a \in \{1\}$  }  
    if (r > 5) {                        {  $r \in [0..10)$  }  
        a = 2;  
    }                                  {  $a \in \{2\}$  }  
    System.out.println(a);  
}
```

{  $a \in \{1,2\}$  } → Valid? Ok? Safe?



# The Why

```
public static void main(String[] args) {  
    int a = 1;  
    double r = Math.random() * 10;    {  $a \in \{1\}$  }  
    if (r > 5) {                        {  $r \in [0..10)$  }  
        a = 2;  
    }                                  {  $a \in \{2\}$  }  
    System.out.println(a);  
}                                     {  $a \in \{1,2\}$  } → Valid? Ok? Safe?
```

- We want to proof, that a program satisfies certain properties

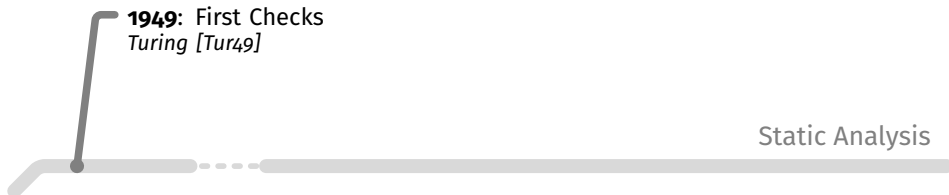
# Origins



Static Analysis

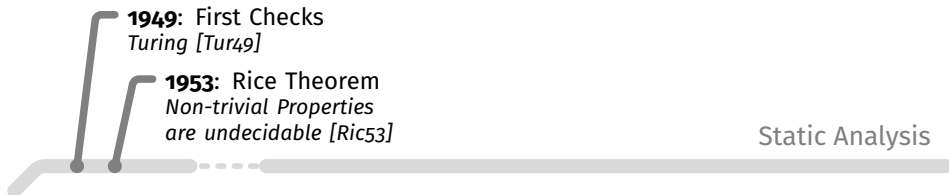
Based on the amazing “Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation” by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

# Origins



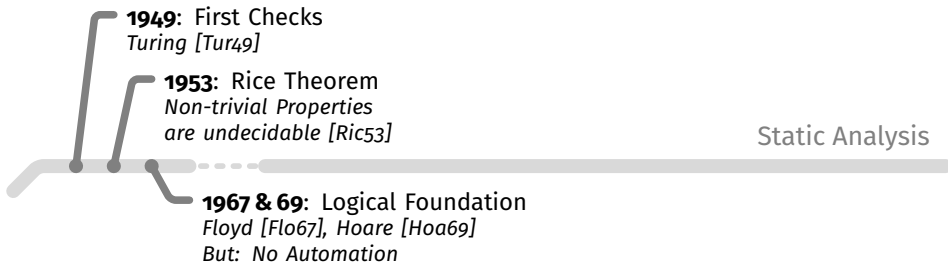
Based on the amazing "Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation" by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

# Origins



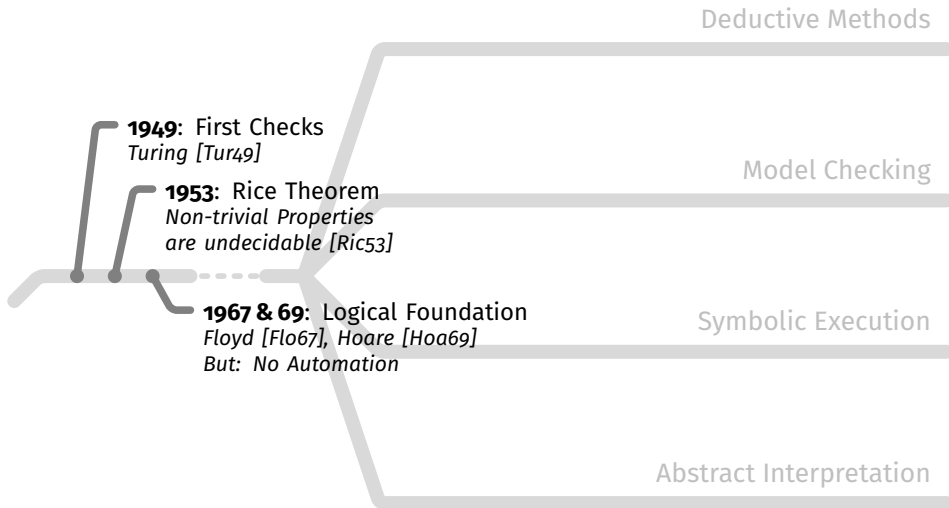
Based on the amazing "Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation" by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

# Origins



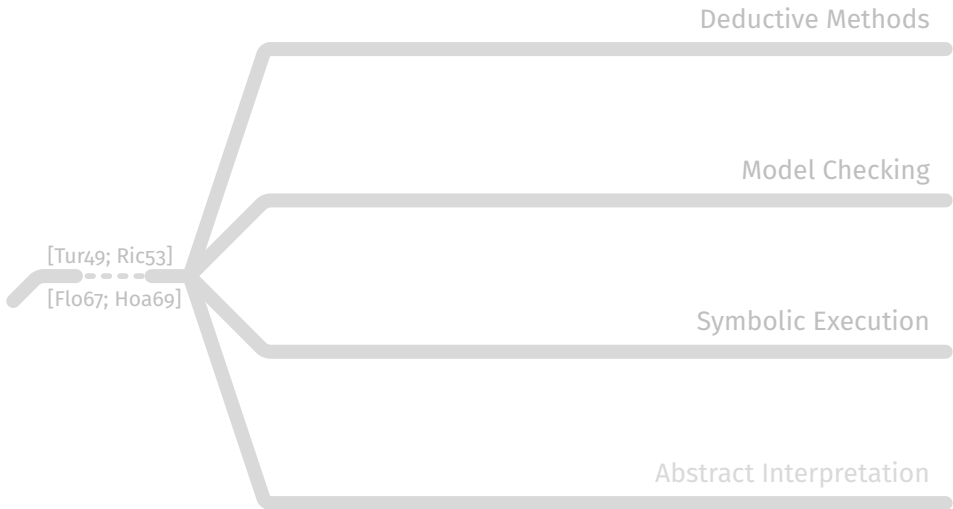
Based on the amazing "Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation" by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

# Origins



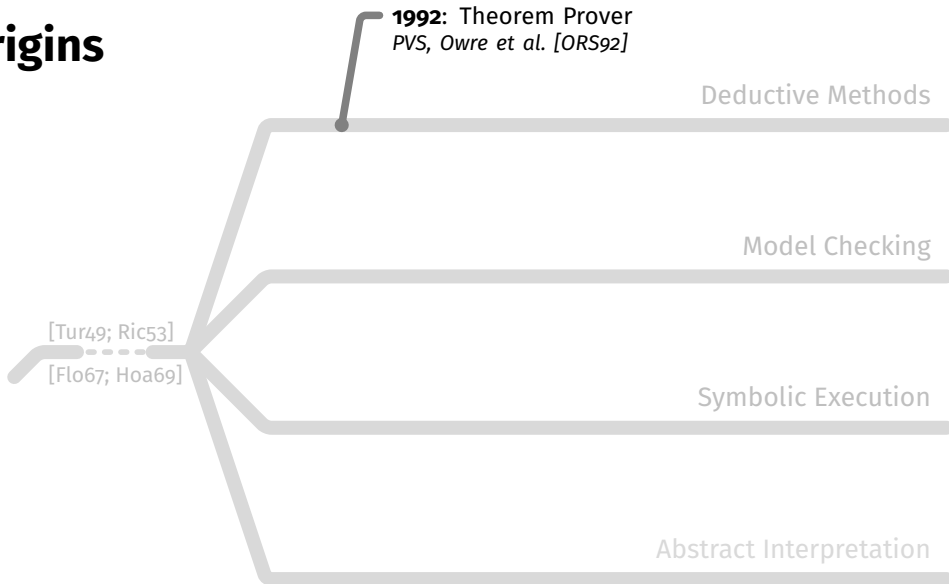
Based on the amazing "Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation" by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

# Origins



Based on the amazing "Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation" by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

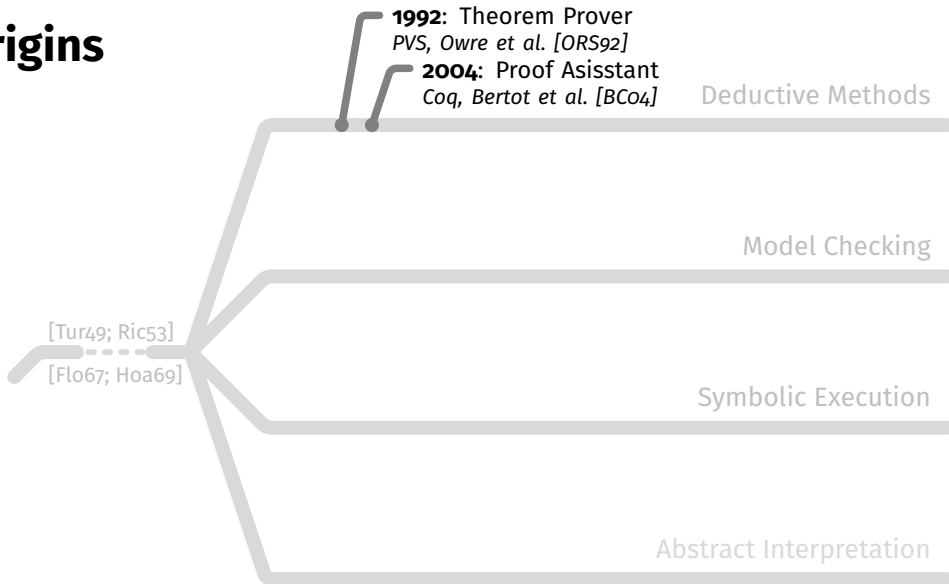
# Origins



Based on the amazing "Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation" by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

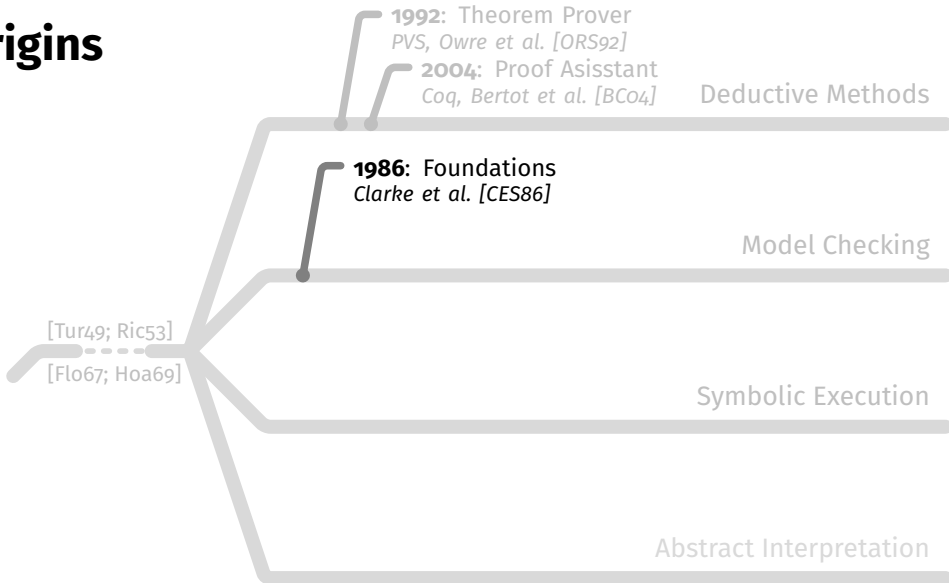


# Origins



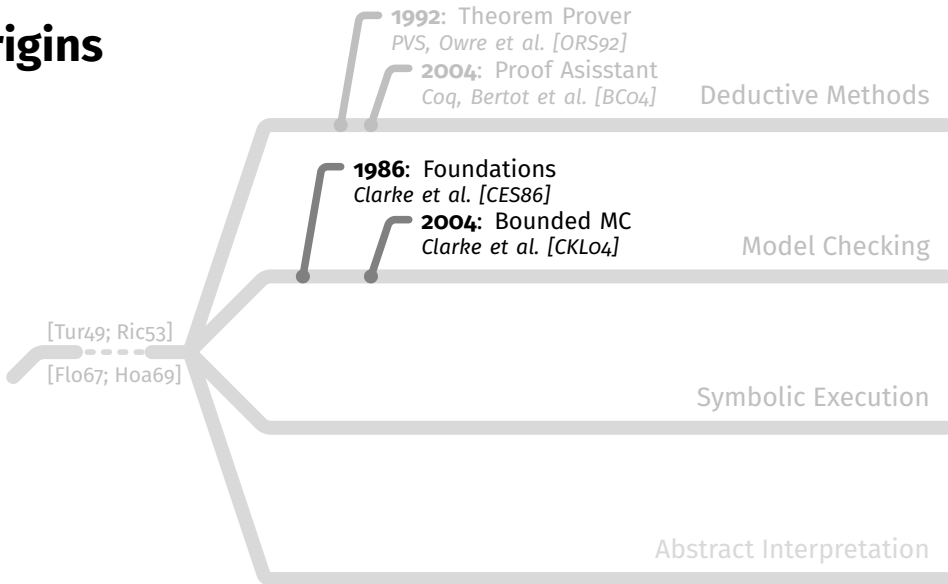
Based on the amazing "Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation" by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

# Origins



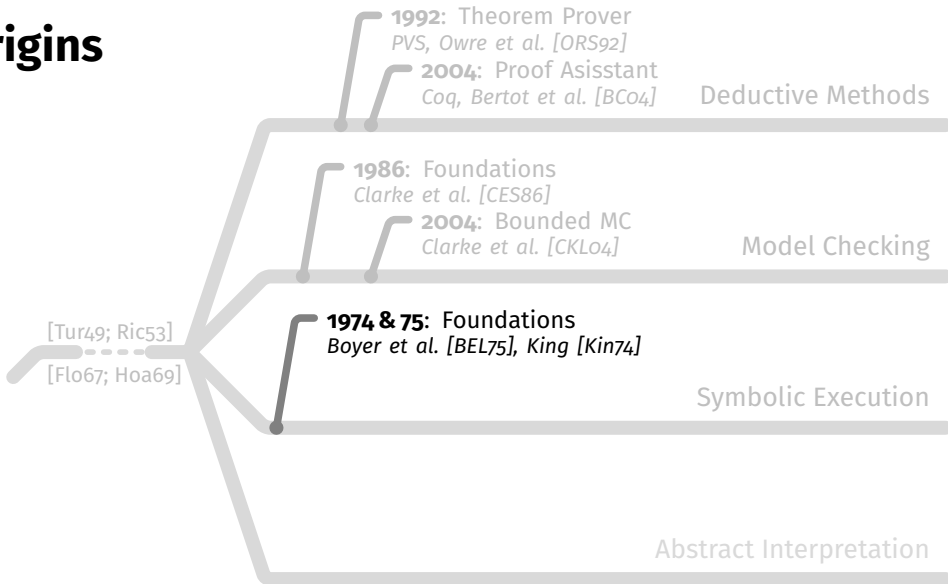
Based on the amazing "Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation" by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

# Origins



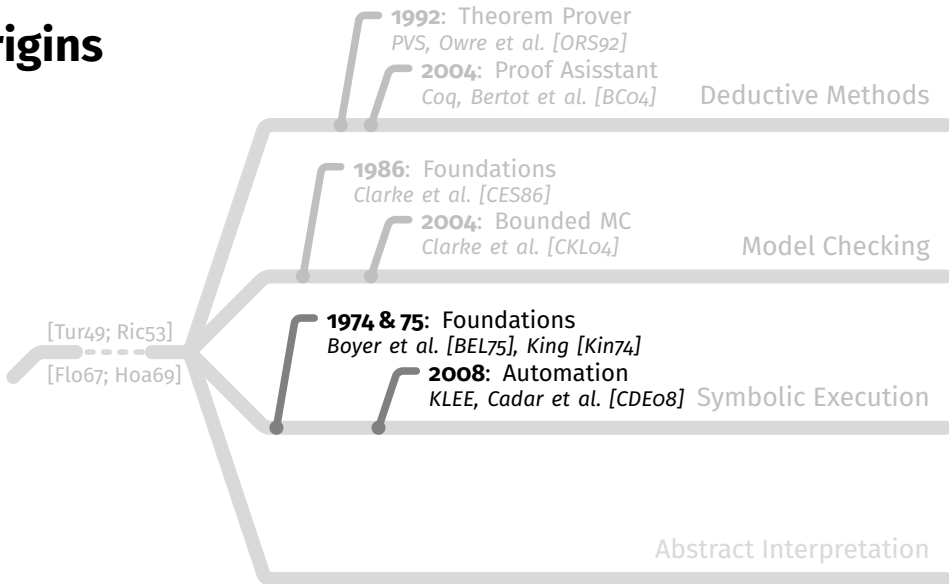
Based on the amazing "Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation" by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

# Origins



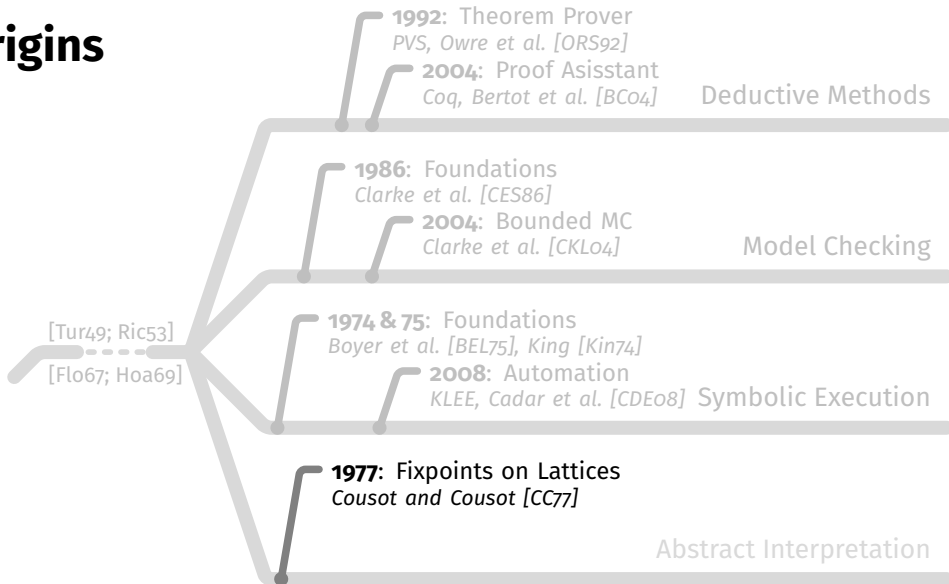
Based on the amazing "Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation" by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

# Origins



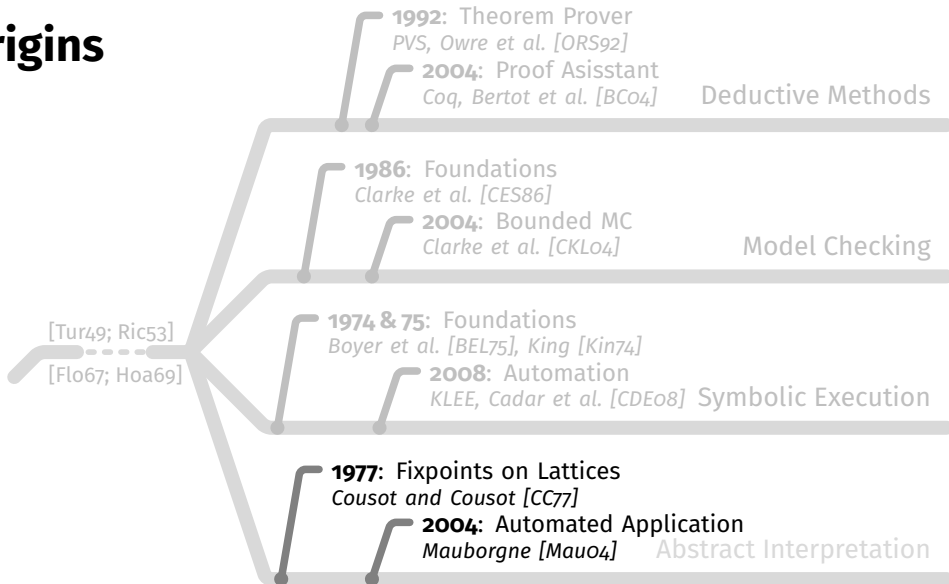
Based on the amazing "Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation" by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

# Origins



Based on the amazing "Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation" by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

# Origins



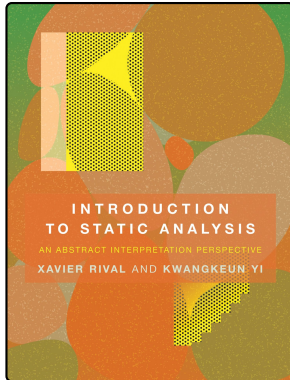
Based on the amazing "Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation" by Miné [Min17], <https://www.di.ens.fr/~cousot/AI/>, and [Bal+18; GR22]

# Recommended Resources



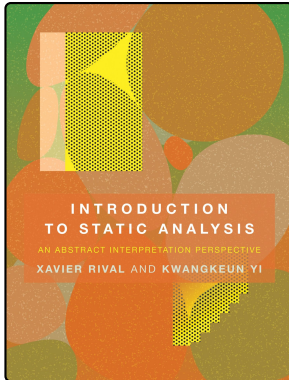
# Recommended Resources

Using Analyses [RY20]

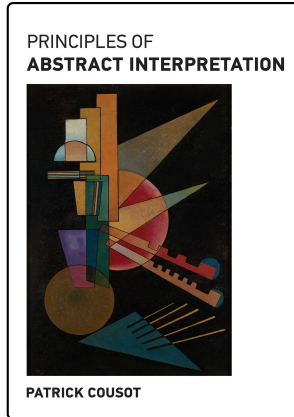


# Recommended Resources

Using Analyses [RY20]

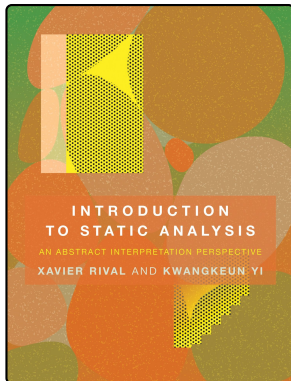


Formal Foundations [Cou21]

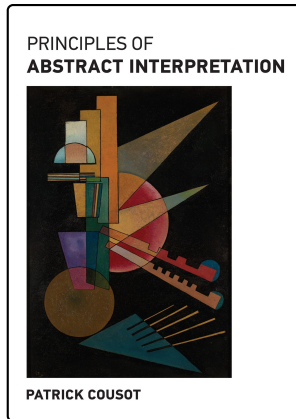


# Recommended Resources

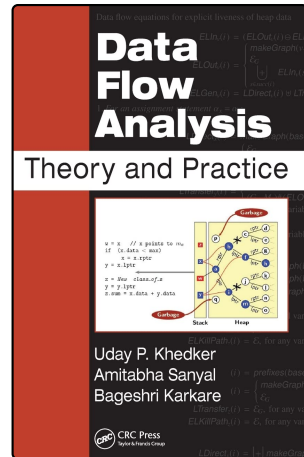
## Using Analyses [RY20]



## Formal Foundations [Cou21]



## Dataflow Perspective [RY20]



# Abstract Interpretation

# Abstract Interpretation

```
public static void main(String[] args) {  
    int a = 1;           {  $a \in \{1\}$  }  
    double r = Math.random() * 10; {  $r \in [0..10)$  }  
    if (r > 5) {  
        a = 2;           {  $a \in \{2\}$  }  
    }  
    System.out.println(a); {  $a \in \{1,2\}$  } → Valid? Ok? Safe?  
}
```

# Abstract Interpretation

```
public static void main(String[] args) {  
    int a = 1;           {  $a \in \{1\}$  }  
    double r = Math.random() * 10; {  $r \in [0..10)$  }  
    if (r > 5) {  
        a = 2;           {  $a \in \{2\}$  }  
    }  
    System.out.println(a); {  $a \in \{1,2\}$  } → Valid? Ok? Safe?  
}
```

# Abstract Interpretation

- We want to proof interesting properties of programs

```
public static void main(String[] args) {  
    int a = 1;           { a ∈ {1} }  
    double r = Math.random() * 10; { r ∈ [0..10) }  
    if (r > 5) {  
        a = 2;           { a ∈ {2} }  
    }  
    System.out.println(a); { a ∈ {1,2} } → Valid? Ok? Safe?  
}
```

# Abstract Interpretation

- We want to proof interesting properties of programs

- *Dataflow Properties*

Liveness, Fainting, Reaching Definitions, ...

```
public static void main(String[] args) {  
    int a = 1;           { a ∈ {1} }  
    double r = Math.random() * 10; { r ∈ [0, 10) }  
    if (r > 5) {  
        a = 2;           { a ∈ {2} }  
    }  
    System.out.println(a); { a ∈ {1, 2} } → Valid? Ok? Safe?  
}
```



# Abstract Interpretation

```
public static void main(String[] args) {  
    int a = 1;           { a ∈ {1} }  
    double r = Math.Random() * 10; { r ∈ [0, 10) }  
    if (r > 5) {  
        a = 2;           { a ∈ {2} }  
    }  
    System.out.println(a); { a ∈ {1, 2} }  
}
```

- We want to proof interesting properties of programs
- *Dataflow Properties*  
Liveness, Fainting, Reaching Definitions, ...
- *Safety Properties*  
No Null Dereference, No Division by Zero, ...

→ Valid? Ok? Safe?

# Abstract Interpretation

```
public static void main(String[] args) {  
    int a = 1;           { a ∈ {1} }  
    double r = Math.random() * 10; { r ∈ [0, 10) }  
    if (r > 5) {  
        a = 2;           { a ∈ {2} }  
    }  
    System.out.println(a); { a ∈ {1, 2} }  
}
```

- We want to proof interesting properties of programs
- *Dataflow Properties*  
Liveness, Fainting, Reaching Definitions, ...
- *Safety Properties*  
No Null Dereference, No Division by Zero, ...
- *Numerical Properties*  
Signs, Intervals, Octagons, Polyhedra, ...

$\{a \in \{1, 2\}\} \rightarrow \text{Valid? Ok? Safe?}$

# Abstract Interpretation

- ```
public static void main(String[] args) {  
    int a = 1;           { a ∈ {1} }  
    double r = Math.Random() * 10; { r ∈ [0, 10) }  
    if (r > 5) {  
        a = 2;           { a ∈ {2} }  
    }  
    System.out.println(a); { a ∈ {1, 2} }  
}
```
- We want to proof interesting properties of programs
    - *Dataflow Properties*  
Liveness, Fainting, Reaching Definitions, ...
    - *Safety Properties*  
No Null Dereference, No Division by Zero, ...
    - *Numerical Properties*  
Signs, Intervals, Octagons, Polyhedra, ...
    - ...

# Abstract Interpretation



See "A casual introduction to Abstract Interpretation" [Cou12]

# Concrete Interpretation



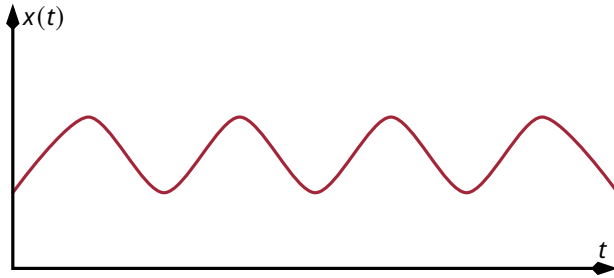
See "A casual introduction to Abstract Interpretation" [Cou12]

# Concrete Interpretation



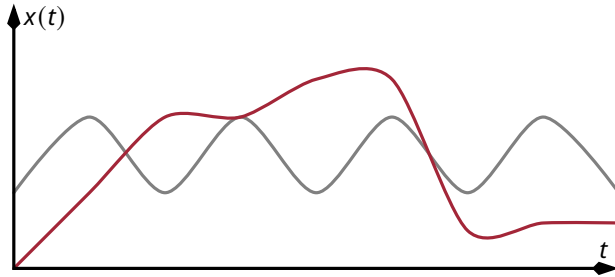
See "A casual introduction to Abstract Interpretation" [Cou12]

# Concrete Interpretation



See "A casual introduction to Abstract Interpretation" [Cou12]

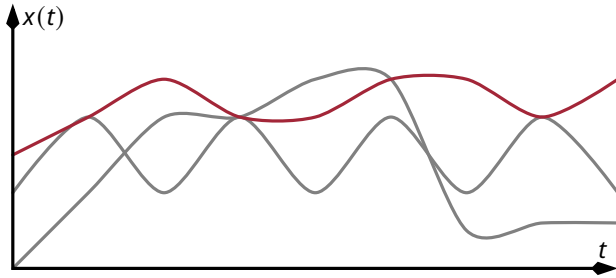
# Concrete Interpretation



See "A casual introduction to Abstract Interpretation" [Cou12]

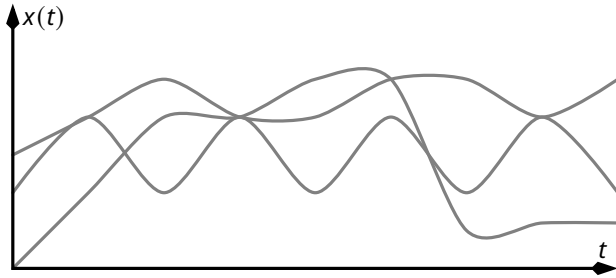


# Concrete Interpretation



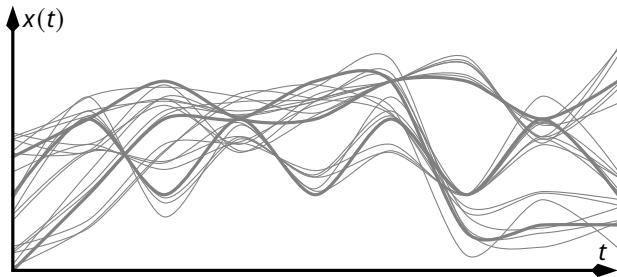
See "A casual introduction to Abstract Interpretation" [Cou12]

# Concrete Interpretation



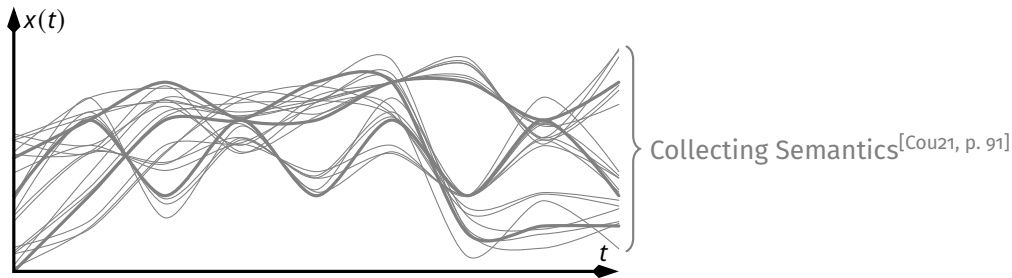
See "A casual introduction to Abstract Interpretation" [Cou12]

# Concrete Interpretation



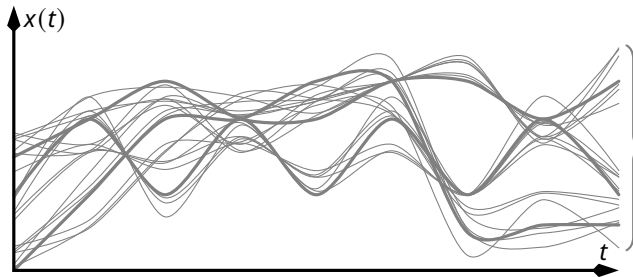
See "A casual introduction to Abstract Interpretation" [Cou12]

# Concrete Interpretation



See "A casual introduction to Abstract Interpretation" [Cou12]

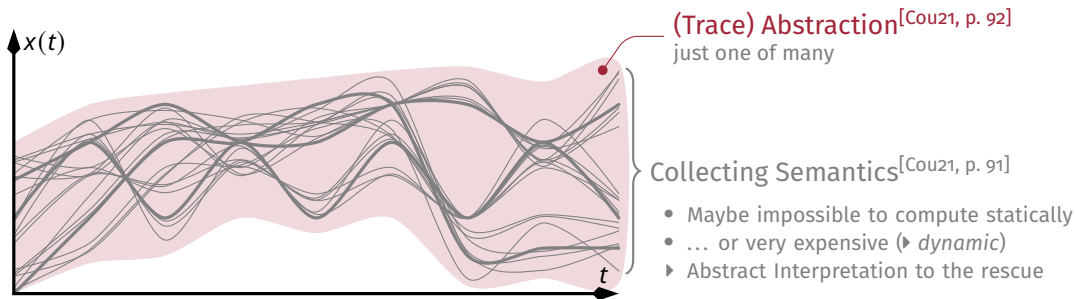
# Concrete Interpretation



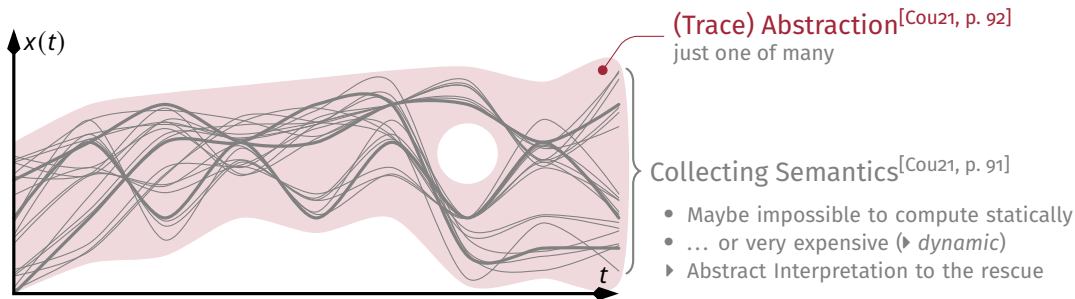
Collecting Semantics<sup>[Cou21, p. 91]</sup>

- Maybe impossible to compute statically
- ... or very expensive (► *dynamic*)
- Abstract Interpretation to the rescue

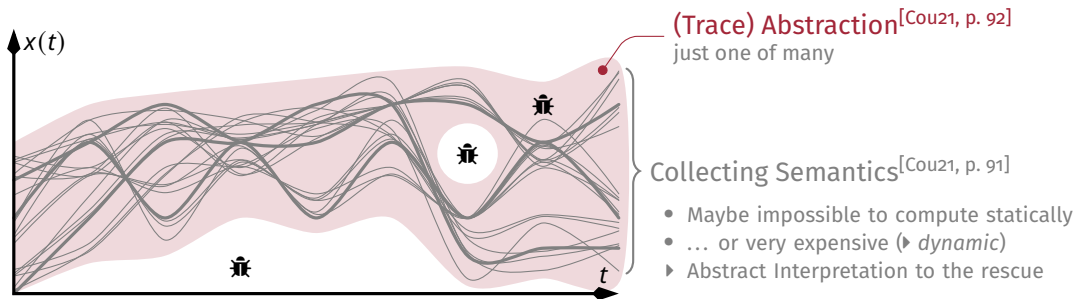
# Abstract Interpretation



# Abstract Interpretation

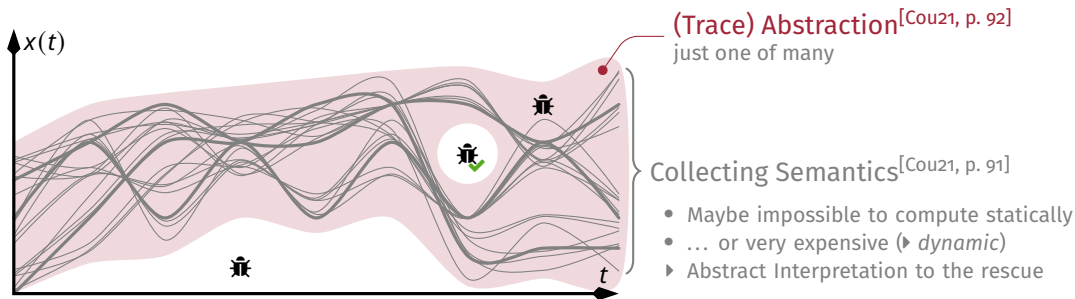


# Abstract Interpretation

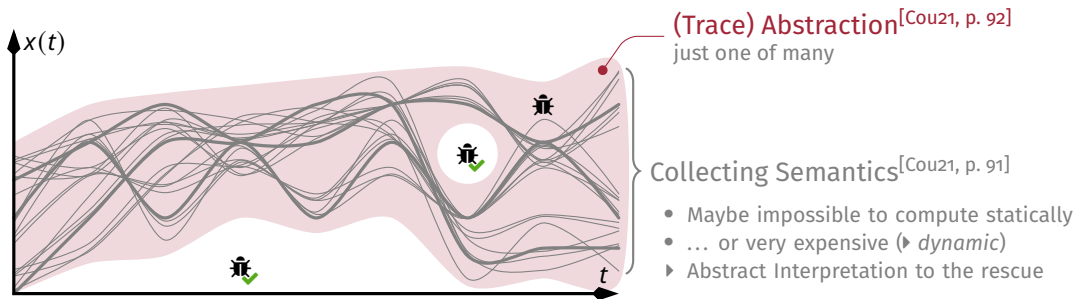




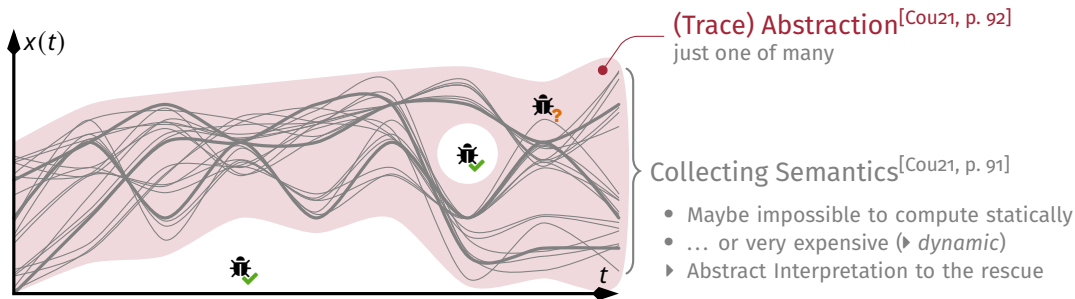
# Abstract Interpretation



# Abstract Interpretation



# Abstract Interpretation



# Terminology

- **Property** — Set of states/traces that satisfy that property

Even integers:  $P = \{z \in \mathbb{Z} \mid \exists k \in \mathbb{Z} : z = 2k\} = \{0, 2, 4, 6, \dots\} \subseteq \mathcal{P}(\mathbb{Z})$

Strongest  $\nearrow \emptyset \subseteq P_1 \subseteq P_2 \subseteq \mathbb{U}$

What is a Property? Set basis Poset etc. I have to abstract! Galois, Semantics Principles of Abstract Interpretation book





# References I

- [Bal+18] Roberto Baldoni et al. “A Survey of Symbolic Execution Techniques”. In: *ACM Comput. Surv.* 51.3 (2018), 50:1–50:39. DOI: 10.1145/3182657. URL: <https://doi.org/10.1145/3182657>.
- [BCo4] Yves Bertot and Pierre Castéran. *Interactive Theorem Proving and Program Development - Coq’Art: The Calculus of Inductive Constructions*. Texts in Theoretical Computer Science. An EATCS Series. Springer, 2004. ISBN: 978-3-642-05880-6. DOI: 10.1007/978-3-662-07964-5. URL: <https://doi.org/10.1007/978-3-662-07964-5>.
- [BEL75] Robert S. Boyer, Bernard Elspas, and Karl N. Levitt. “SELECT - a formal system for testing and debugging programs by symbolic execution”. In: *Proceedings of the International Conference on Reliable Software 1975, Los Angeles, California, USA, April 21-23, 1975*. Ed. by Martin L. Shooman and Raymond T. Yeh. ACM, 1975, pp. 234–245. DOI: 10.1145/800027.808445. URL: <https://doi.org/10.1145/800027.808445>.

## References II

- [CC77] Patrick Cousot and Radhia Cousot. “Abstract Interpretation: A Unified Lattice Model for Static Analysis of Programs by Construction or Approximation of Fixpoints”. In: *Conference Record of the Fourth ACM Symposium on Principles of Programming Languages, Los Angeles, California, USA, January 1977*. Ed. by Robert M. Graham, Michael A. Harrison, and Ravi Sethi. ACM, 1977, pp. 238–252. DOI: 10.1145/512950.512973. URL: <https://doi.org/10.1145/512950.512973>.
- [CDEo8] Cristian Cadar, Daniel Dunbar, and Dawson R. Engler. “KLEE: Unassisted and Automatic Generation of High-Coverage Tests for Complex Systems Programs”. In: *8th USENIX Symposium on Operating Systems Design and Implementation, OSDI 2008, December 8-10, 2008, San Diego, California, USA, Proceedings*. Ed. by Richard Draves and Robbert van Renesse. USENIX Association, 2008, pp. 209–224. URL: [http://www.usenix.org/events/osdi08/tech/full%5C\\_papers/cadar/cadar.pdf](http://www.usenix.org/events/osdi08/tech/full%5C_papers/cadar/cadar.pdf).



# References III

- [CES86] Edmund M. Clarke, E. Allen Emerson, and A. Prasad Sistla. “Automatic Verification of Finite-State Concurrent Systems Using Temporal Logic Specifications”. In: *ACM Trans. Program. Lang. Syst.* 8.2 (1986), pp. 244–263. DOI: 10.1145/5397.5399. URL: <https://doi.org/10.1145/5397.5399>.
- [CKLo4] Edmund Clarke, Daniel Kroening, and Flavio Lerda. “A tool for checking ANSI-C programs”. In: *Tools and Algorithms for the Construction and Analysis of Systems: 10th International Conference, TACAS 2004, Held as Part of the Joint European Conferences on Theory and Practice of Software, ETAPS 2004, Barcelona, Spain, March 29–April 2, 2004. Proceedings 10*. Springer. 2004, pp. 168–176.
- [Cou12] Patrick Cousot. “A casual introduction to Abstract Interpretation”. In: *CMACS Workshop on Systems Biology and Formals Methods (SBFM’12)* (2012). URL: <https://pcousot.github.io/talks/PCousot-SBFM-2012-1-1.pdf> (visited on 12/09/2024).
- [Cou21] Patrick Cousot. “Principles of Abstract Interpretation”. In: (2021).

## References IV

- [Flo67] Robert W. Floyd. “Assigning Meanings to Programs”. In: *Proc. of the American Mathematical Society Symposia on Applied Mathematics*. Vol. 19. 1967, pp. 19–32.
- [GR22] Roberto Giacobazzi and Francesco Ranzato. “History of Abstract Interpretation”. In: *IEEE Ann. Hist. Comput.* 44.2 (2022), pp. 33–43. DOI: 10.1109/MAHC.2021.3133136. URL: <https://doi.org/10.1109/MAHC.2021.3133136>.
- [Hoa69] C. A. R. Hoare. “An Axiomatic Basis for Computer Programming”. In: *Commun. ACM* 12.10 (1969), pp. 576–580. DOI: 10.1145/363235.363259. URL: <https://doi.org/10.1145/363235.363259>.
- [Kin74] James C. King. “A New Approach to Program Testing”. In: *Programming Methodology, 4th Informatik Symposium, IBM Germany, Wildbad, September 25-27, 1974*. Ed. by Clemens Hackl. Vol. 23. Lecture Notes in Computer Science. Springer, 1974, pp. 278–290. DOI: 10.1007/3-540-07131-8\_30. URL: [https://doi.org/10.1007/3-540-07131-8\\_30](https://doi.org/10.1007/3-540-07131-8_30).

# References V

- [Mau04] Laurent Mauborgne. “Astrée: verification of absence of run-time error”. In: *Building the Information Society, IFIP 18th World Computer Congress, Topical Sessions, 22-27 August 2004, Toulouse, France*. Ed. by René Jacquart. Vol. 156. IFIP. Kluwer/Springer, 2004, pp. 385–392. DOI: [10.1007/978-1-4020-8157-6\\_30](https://doi.org/10.1007/978-1-4020-8157-6_30). URL: [https://doi.org/10.1007/978-1-4020-8157-6%5C\\_30](https://doi.org/10.1007/978-1-4020-8157-6%5C_30).
- [Min17] Antoine Miné. “Tutorial on Static Inference of Numeric Invariants by Abstract Interpretation”. In: *Found. Trends Program. Lang.* 4.3-4 (2017), pp. 120–372. DOI: [10.1561/25000000034](https://doi.org/10.1561/25000000034). URL: <https://doi.org/10.1561/25000000034>.
- [ORS92] Sam Owre, John M. Rushby, and Natarajan Shankar. “PVS: A Prototype Verification System”. In: *Automated Deduction - CADE-11, 11th International Conference on Automated Deduction, Saratoga Springs, NY, USA, June 15-18, 1992, Proceedings*. Ed. by Deepak Kapur. Vol. 607. Lecture Notes in Computer Science. Springer, 1992, pp. 748–752. DOI: [10.1007/3-540-55602-8\\_217](https://doi.org/10.1007/3-540-55602-8_217). URL: [https://doi.org/10.1007/3-540-55602-8%5C\\_217](https://doi.org/10.1007/3-540-55602-8%5C_217).

# References VI

- [Ric53] Henry Gordon Rice. “Classes of recursively enumerable sets and their decision problems”. In: *Transactions of the American Mathematical society* 74.2 (1953), pp. 358–366.
- [RY20] Xavier Rival and Kwangkeun Yi. “Introduction to Static Analysis: An Abstract Interpretation Perspective”. In: (2020).
- [Tur49] Alan Turing. “Checking a large routine”. In: *Report of a Conference on High Speed Automatic Calculating Machines*. 1949, pp. 67–69.