Department of
Computer Science and Engineering
计算机科学与工程系

CS215 DISCRETE MATH

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k-Element Permutations/Combinations of a Set

- *k-element permutation of N*: a list of k distinct elements from $\{1, 2, \ldots, n\}$. How many are there?

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- *k-element combination of N*: a set of k distinct elements from $\{1, 2, \ldots, n\}$. How many are there?

$$\binom{n}{k} = C(n, k) = \frac{P(n, k)}{k!} = \frac{n!}{k!(n-k)!}.$$

Pascal's Triangle

							1							
						1		1						
					1		2		1					
				1		3		3		1				
			1		4		6		4		1			
		1		5		10		10		5		1		
	1		6		15		20		15		6		1	
1		7		21		35		35		21		7		1

Pascal's Triangle

$$
\begin{array}{ccccccccccccccc}
&&&&&&&1&&&&&&& \\
&&&&&&1&&1&&&&&& \\
&&&&&1&&2&&1&&&&& \\
&&&&1&&3&&3&&1&&&& \\
&&&1&&4&&6&&4&&1&&& \\
&&1&&5&&10&&10&&5&&1&& \\
&1&&6&&15&&20&&15&&6&&1& \\
1&&7&&21&&35&&35&&21&&7&&1
\end{array}
$$

Pascal identity

Each (non-1) entry in Pascal's Triangle is the sum of the two entries directly above it (to left and to right).

Pascal's Triangle

$$
\begin{array}{ccccccccccccccc}
 & & & & & & & 1 & & & & & & & \\
 & & & & & & 1 & & 1 & & & & & & \\
 & & & & & 1 & & 2 & & 1 & & & & & \\
 & & & & 1 & & 3 & & 3 & & 1 & & & & \\
 & & & 1 & & 4 & & 6 & & 4 & & 1 & & & \\
 & & 1 & & 5 & & 10 & & 10 & & 5 & & 1 & & \\
 & 1 & & 6 & & 15 & & 20 & & 15 & & 6 & & 1 & \\
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$$\boxed{\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}}$$

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  We will use a combinatorial proof.

A Combinatorial Proof

- $\binom{n}{k}$ is the number of k-element subsets of an n-element set.

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$$\boxed{\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}}$$

Therefore, each term (left and right) represents the number of subsets of a particular size chosen from an appropriately sized set.

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Try to use sum principle to explain relationship among these three terms.

Example: $n = 5$, $k = 2$

$$\binom{5}{2} = \binom{4}{1} + \binom{4}{2}.$$

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Set S_1 of 2-subsets of S

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Consider $S = \{A, B, C, D, E\}$.

Set S_1 of 2-subsets of S can be partitioned into 2 disjoint parts.

S_2: the 2-subsets that contain E and

S_3: the set of 2-subsets that do not contain E.

$$S_1 = \{\{A, B\}, \{A, C\}, \{A, D\}, \{A, E\}, \{B, C\}, \{B, D\}, \{B, E\}, \{C, D\}, \{C, E\}, \{D, E\}\}.$$

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Proof: Apply sum rule.

Let S_1 be set of all k-element subsets.

To apply sum rule, partition S_1 into S_2 and S_3.

Let S_2 be set of k-element subsets that contain x_n.

Let S_3 be set of k-element subsets that don't contain x_n.

Blaise Pascal

Born 1623; Died 1662

French Mathematician

A Founder of Probability Theory

Inventor of one of the first mechanical calculating machines

Pascal Programming Language named for him

The Binomial Theorem

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$$\begin{aligned}(x+y)^3 &= x^3 + 3x^2y + 3xy^2 + y^3 \\ &= \binom{3}{0}x^3 + \binom{3}{1}x^2y + \binom{3}{2}xy^2 + \binom{3}{3}y^3\end{aligned}$$

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Proof?

Application of the Binomial Theorem

- We may use the Binomial Theorem to prove

$$\boxed{\sum_{i=0}^{n} \binom{n}{i} = 2^n}$$

Labelling and Trinomial Coefficients

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  Show that if we have k_1 labels of one kind, e.g., red, k_2 labels of a second kind, e.g., blue, and $k_3 = n - k_1 - k_2$ labels of a third kind, then there are $\frac{n!}{k_1!k_2!k_3!}$ ways to apply these labels to n objects

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  What is the coefficient of $x^{k_1}y^{k_2}z^{k_3}$ in $(x+y+z)^n$?

Labelling and Trinomial Coefficients

- There are $\binom{n}{k_1}$ ways to choose the red items There are then $\binom{n-k_1}{k_2}$ ways to choose the blue items from the remaining $n - k_1$. The remaining k_3 items get labelled a third color.

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  Using the *product rule* the total number of labellings is

$$\begin{aligned}\binom{n}{k_1}\binom{n-k_1}{k_2} &= \frac{n!}{k_1!(n-k_1)!}\frac{(n-k_1)!}{(k_2)!(n-k_1-k_2)!} \\ &= \frac{n!}{k_1!k_2!(n-k_1-k_2)!} = \frac{n!}{k_1!k_2!k_3!}\end{aligned}$$

Labelling and Trinomial Coefficients

- When $k_1 + k_2 + k_3 = n$, we call

$$\frac{n!}{k_1!k_2!k_3!}$$

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  This will be very similar to the analysis of hashing n keys into a table of size 365.

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  $\#A_n + \#B_n = 365^n$

The Birthday Paradox

n	A_n	B_n	n	A_n	B_n
1	0.00000000	1.00000000	16	0.28360400	0.71639599
2	0.00273972	0.99726027	17	0.31500766	0.68499233
3	0.00820416	0.99179583	18	0.34691141	0.65308858
4	0.01635591	0.98364408	19	0.37911852	0.62088147
5	0.02713557	0.97286442	20	0.41143838	0.58856161
6	0.04046248	0.95953751	21	0.44368833	0.55631166
7	0.05623570	0.94376429	22	0.47569530	0.52430469
8	0.07433529	0.92566470	23	0.50729723	0.49270276
9	0.09462383	0.90537616	24	0.53834425	0.46165574
10	0.11694817	0.88305182	25	0.56869970	0.43130029
11	0.14114137	0.85885862	26	0.59824082	0.40175917
12	0.16702478	0.83297521	27	0.62685928	0.37314071
13	0.19441027	0.80558972	28	0.65446147	0.34553852
14	0.22310251	0.77689748	29	0.68096853	0.31903146
15	0.25290131	0.74709868	30	0.70631624	0.29368375

"Birthday" attacks

- Event A: at least two people in the room have the same birthday

  Event B: no two people in the room have the same birthday

  $\Pr[A] = 1 - \Pr[B]$

  $$\begin{aligned}\Pr[B] &= \left(1 - \frac{1}{365}\right) \cdot \left(1 - \frac{2}{365}\right) \cdot \cdots \cdot \left(1 - \frac{n-1}{365}\right) \\ &= \prod_{i=1}^{n-1} \left(1 - \frac{i}{365}\right).\end{aligned}$$

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  $p(n; H) := 1 - \prod_{i=1}^{n-1}(1 - \frac{i}{H})$

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  Recall that $p(n; H) := 1 - \prod_{i=1}^{n-1}(1 - \frac{i}{H})$

  This probability can be approximated as
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  This probability can be approximated as
  $$p(n; H) \approx 1 - e^{-n(n-1)/2H} \approx 1 - e^{-n^2/2H}.$$
  Let $n(p; H)$ be the smallest number of values we have to choose, such that the probability for finding a collision is at least p. By inverting the expression above, we have
  $$n(p; H) \approx \sqrt{2H \ln \frac{1}{1-p}}.$$

Euclidean Algorithm

- The Euclidean algorithm in pseudocode

ALGORITHM 1 The Euclidean Algorithm.

procedure $gcd(a, b$: positive integers)
$x := a$
$y := b$
while $y \neq 0$
    $r := x$ **mod** y
    $x := y$
    $y := r$
return $x\{\gcd(a, b)$ is $x\}$

The number of divisions required to find $\gcd(a, b)$ is $O(\log b)$, where $a \geq b$. (this will be proved later.)

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Why ?

Euclidean Algorithm

- Key steps in the Euclidean algorithm

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\begin{aligned}
r_0 &= r_1 q_1 + r_2 && 0 \le r_2 < r_1, \\
r_1 &= r_2 q_2 + r_3 && 0 \le r_3 < r_2, \\
&\ \ \vdots \\
r_{n-2} &= r_{n-1} q_{n-1} + r_n && 0 \le r_n < r_{n-1}, \\
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Case (i): $r_{i+1} \le \frac{1}{2} r_i$: $r_{i+2} < r_{i+1} \le \frac{1}{2} r_i$.

Case (ii): $r_{i+1} > \frac{1}{2} r_i$: $r_{i+2} = r_i \bmod r_{i+1} = r_i - r_{i+1} < \frac{1}{2} r_i$.

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See [Theorem 1 p. 347].

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- **Definition** A *linear homogeneous relation of degree k* with constant coefficients is a recurrence relation of the form

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By induction, such a recurrence relation is uniquely determined by this recurrence relation, and k initial conditions $a_0, a_1, \ldots, a_{k-1}$.

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Examples

$$P_n = (1.11) P_{n-1}$$

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$B_n = nB_{n-1}$	coefficients are not constants

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- **Example** Consider the recurrence relation
  $a_n = 2a_{n-1} - a_{n-2}$,
  Which of the following are solutions?
  - ◇ $a_n = 3n$:
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Fact: Assume that the sequences a_n and a'_n both satisfy the recurrence, then $b_n = a_n + a'_n$, $d_n = \alpha a_n$ also satisfy the recurrence, where α is a constant.

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  So, try to find any solution of the form $a_n = r^n$ that satisfies the recurrence.

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◇ The solutions to the *characteristic equation* can yield an explicit formula for the sequence.

$$(r^k - c_1 r^{k-1} - \cdots - c_k) = 0$$

Recall: Problem IV

- **Fibonacci number**

  $F_0 = 0,\ F_1 = 1,\ F_n = F_{n-1} + F_{n-2}$ for $n \geq 2$

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  ◇ What is the closed-form expression of F_n?

  Consider $x^n = x^{n-1} + x^{n-2}$, with $x \neq 0$. There are two different roots

  $$\phi = \frac{1+\sqrt{5}}{2}, \quad \psi = \frac{1-\sqrt{5}}{2}$$

  Then F_n can be the form of $a\phi^n + b\psi^n$. By $F_0 = 0$ and $F_1 = 1$, we have $a + b = 0$ and $\phi a + \psi b = 1$, leading to $a = \frac{1}{\sqrt{5}}$, $b = -a$. Therefore,

  $$F_n = \frac{\phi^n - \psi^n}{\sqrt{5}}$$

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  See [Theorem 1 p. 515].

Department of Computer Science and Engineering 南方科技大学计算机科学与工程系

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  $a_n = 3 \cdot 2^n - (-1)^n$

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- **Example 2** $a_n = 7a_{n-1} - 10a_{n-2}$, with $a_0 = 2$, $a_1 = 1$

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  Example $a_n = 2a_{n-1} + 5a_{n-2} - 6a_{n-3}$

The Case of Degenerate Roots

- **Theorem** If the CE $r^2 - c_1 r - c_2 = 0$ has only 1 root r_0, then

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  Exercise.

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The Case of Degenerate Roots in General

- **Theorem** [Theorem 4, p.519] Suppose that there are t roots $r_1, \ldots, r_t$ with multiplicities $m_1, \ldots, m_t$. Then

$$a_n = \sum_{i=1}^{t} \left(\sum_{j=0}^{m_i-1} \alpha_{i,j} n^j \right) r_i^n,$$

for all $n \geq 0$ and constants $\alpha_{i,j}$.

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Example

$a_n = -3a_{n-1} - 3a_{n-2} - a_{n-3}$ with $a_0 = 1$, $a_1 = -2$, $a_2 = -1$

Linear Nonhomogeneous Recurrence Relations

- **Definition** A *linear nonhomogeneous relation* with constant coefficients may contain some terms $F(n)$ that depend only on n

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  $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \cdots + c_k a_{n-k}$ is called the *associated homogeneous recurrence relation*.

  Fact: Assume that the sequence b_n satisfies the recurrence. Then another sequece a_n satisfies the non-homogeneous recurrence if and only if $h_n = a_n - b_n$ is a sequence that satisfies the associated homogeneous recurrence.

  Idea: We already know how to find h_n. For many common $f(n)$, a solution b_n to the non-homogeneous recurrence is similar to $f(n)$. We then need find $a_n = b_n + h_n$ to the non-homogeneous recurrence that satisfies both recurrence and initial conditions.

- **Theorem** If $a_n = p(n)$ is any particular solution to the linear nonhomogeneous relation with constant coefficients,

  $$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \cdots + c_k a_{n-k} + F(n),$$

  Then all its solutions are of the form

  $$a_n = p(n) + h(n),$$

  where $a_n = h(n)$ is any solution to the associated homogeneous recurrence relation

  $$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \cdots + c_k a_{n-k}.$$

- **Example** $a_n = 3a_{n-1} + 2n$. Which solution has $a_1 = 3$?

Solving Linear Nonhomogeneous Recurrence Relations

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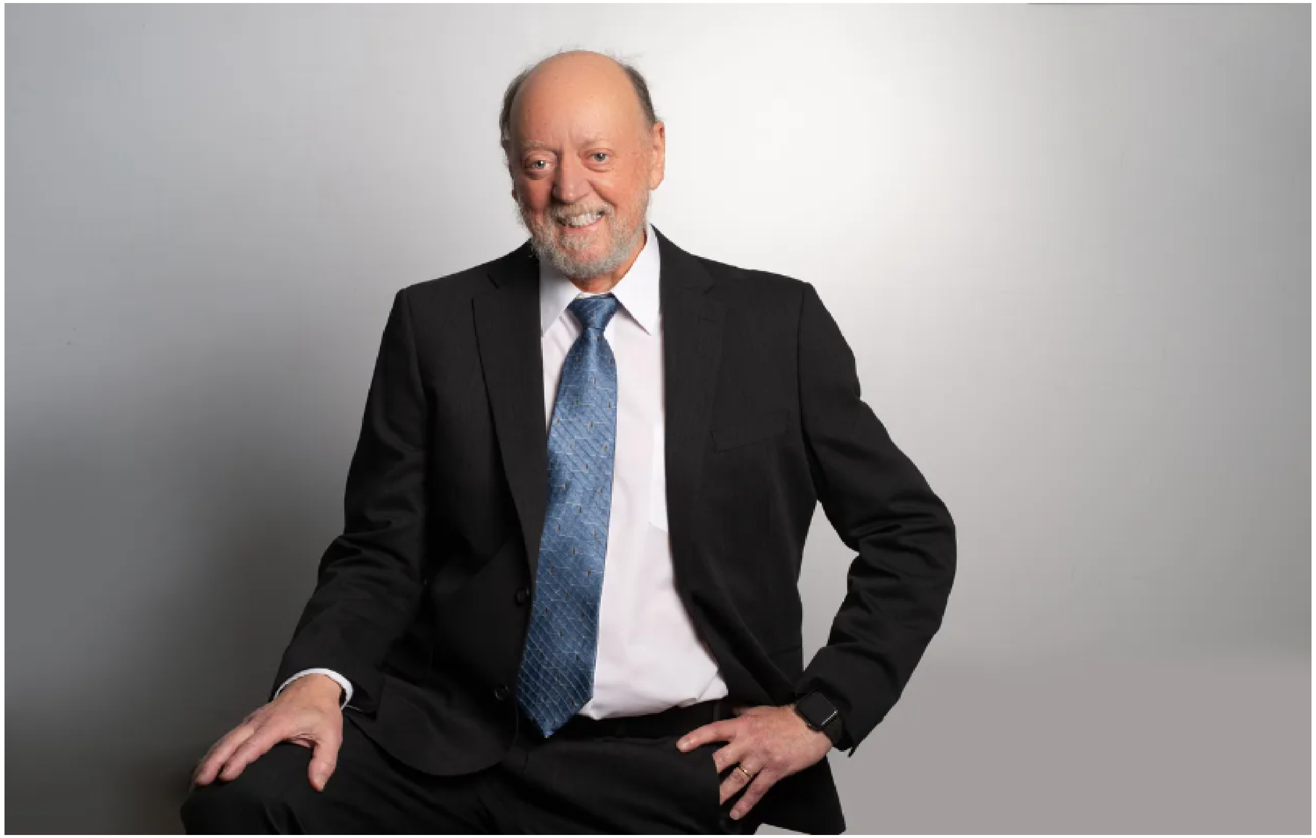
  We get $c = -1$ and $d = -3/2$. Thus,
  $$p(n) = -n - 3/2$$

Solving Linear Nonhomogeneous Recurrence Relations



"Science is driven by simulation," says Dongarra. "It's that match between the hardware capability, and the necessity of the simulations to use that hardware, where my software fits in."

"I'm a mathematician, to me, everything is linear algebra, but the world is seeing that as well," he said. "It's a fabric on which we build other things." Most problems in machine learning and AI, he said, go back to an "eternal computational component" in linear algebra.

Next Lecture

- generating function, relation ...