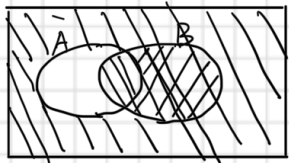


Q.1 (a) False. counter example:  $A = \{1, 2, 3\}$   $B = \{4, 5\}$ .

(b) True  $A \cap B \cap C = \{x \mid x \in A \wedge x \in B \wedge x \in C\}$   
 $\subseteq \{x \mid x \in A \wedge x \in B\}$   
 $\subseteq \{x \mid x \in A \vee x \in B\}$   
 $= A \cup B$

(c) False

$$\overline{(A - B)} \cap (B - A) = \{x \mid (x \in A \wedge x \notin B) \wedge (x \in B \wedge x \notin A)\}$$



$$= \{x \mid (x \notin A \vee x \in B) \wedge (x \in B \wedge x \notin A)\}$$

$$= \{x \mid x \notin A \wedge x \in B\}$$

$$= B - A$$

Q.2.

(1) Disproof: counterexample:

$$A = \{1, 2, 3, 4\} \quad B = \{2, 3, 5, 6\} \quad C = \{3, 4, 7, 8\}$$

$$C - (A \cap B) = \{4, 7, 8\}$$

$$(C - A) \cap (C - B) = \{7, 8\}$$

(2) Proof:

$X \in P(A) \cap P(B)$  means  $X \in P(A)$  and  $X \in P(B)$ , implies that  $X \subseteq A$  and  $X \subseteq B$ , then  $X \subseteq A \cap B$ , thus  $X \in P(A \cap B)$ . So  $P(A) \cap P(B) \subseteq P(A \cap B)$

$X \in P(A \cap B)$  means that  $X$  is a subset of both  $A$  and  $B$ , which implies  $X \in P(A)$  and  $X \in P(B)$

Then we get  $X \in P(A) \cap P(B)$ . So  $P(A \cap B) \subseteq P(A) \cap P(B)$

Since two statements above, we have  $P(A) \cap P(B) = P(A \cap B)$

(3) False  $A = \{1\}$   $B = \{2\}$

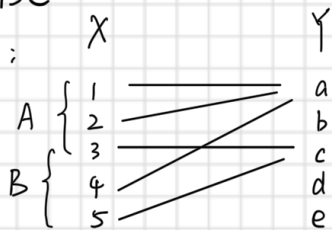
$$P(A) \cup P(B) = \{\{1\}, \{2\}\}$$

$$P(A \cup B) = \{\{1\}, \{2\}, \{1, 2\}\}$$

$$P(A) \cup P(B) \neq P(A \cup B)$$

(4) False

Counter:



$$f(A \cap B) = \{c\}$$

$$f(A) \cap f(B) = \{a, c\}$$

Q.3 (a)

A	B	C	$A \oplus B$	$B \oplus C$	$A \oplus (B \oplus C)$	$(A \oplus B) \oplus C$
0	0	0	0	0	0	0
0	0	1	0	1	1	1
0	1	0	1	1	0	0
0	1	1	1	0	0	0
1	0	0	1	0	1	1
1	0	1	1	1	0	0
1	1	0	0	1	0	0
1	1	1	0	0	1	1

From the truth table, we can get  $A \oplus (B \oplus C) = (A \oplus B) \oplus C$ .

(b) Suppose  $A \neq B$ , then  $\exists x, x \in A \wedge x \notin B$  or  $\exists x, x \notin A \wedge x \in B$ .

when  $\exists x, x \in A \wedge x \notin B$   $A \oplus C = 1 \oplus C = \neg C$ ,  $B \oplus C = 0 \oplus C = C$ .

$$A \oplus C \neq B \oplus C$$

when  $\exists x, x \notin A \wedge x \in B$   $A \oplus C = 0 \oplus C = C$ ,  $B \oplus C = 1 \oplus C = \neg C$

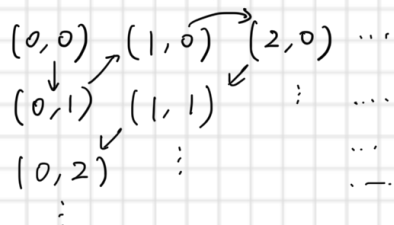
$$A \oplus C \neq B \oplus C$$

By contradiction,  $A = B$

Q.4. (a) finity  $\Rightarrow$  countable.

(b) countable

constructing the list: first list  $(a, b)$  with  $a+b=0$ , then list  $(a, b)$  with  $a+b=1$ , and so on



(c) uncountable.

Because for any  $a_0 \in \mathbb{N}$ , it has a subset  $\{(a_0, b) \mid b \in \mathbb{R}\}$  which is an uncountable set.

Q.5

(a)  $A = B = \mathbb{R}$

(b)  $A = \mathbb{R} \quad B = \mathbb{R} - \mathbb{N}$

(c)  $A = \mathbb{R} \quad B = (0, 1)$

Q.6. only if: Because  $A \subseteq P(A)$ ,  $P(A) \subseteq P(B)$ , we have  $A \subseteq P(B)$ .  
which means  $A$  is a subset of  $B$ , Thus  $A \subseteq B$

if: Let  $X \in P(A)$  which means  $X \subseteq A$ . Since  $X \subseteq A$ ,  $A \subseteq B$   
Therefore  $X \subseteq B$ . Thus  $X \in P(B)$   
Since for all element  $X$  in  $P(A)$  we can show that  $X \in P(B)$   
Hence,  $P(A) \subseteq P(B)$

Q.7.

$f$  is one-to-one  $\Rightarrow$  Proof every  $f(x) = f(y)$  implies  $x = y$  for  $x, y \in A$ .

Proof: Assume  $f(x) = f(y)$ ,  $x, y \in A$

since  $g \circ f = 1_A$ , we have  $g \circ f(x) = g(f(x)) = x$

$$g \circ f(y) = g(f(y)) = y$$

since  $f(x) = f(y)$ , thus  $g(f(x)) = g(f(y))$

Hence,  $x = y$ . So  $f$  is one-to-one.

$g$  is onto: for every  $a \in A$ , there exists an  $b \in B$  that  $g(b) = a$

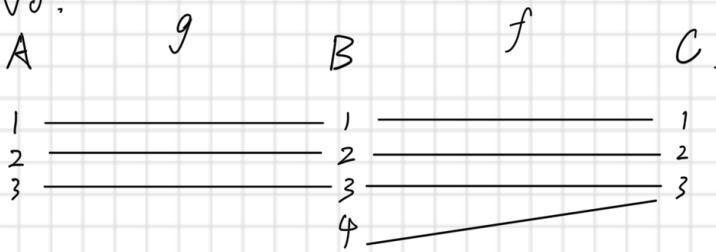
$$\text{since } g \circ f(a) = g(f(a)) = a$$

Thus for every  $a \in A$ , there exists  $f(a) \in B$  that  $g(f(a)) = a$

Hence  $g$  is onto.

Q.8. (a) No.

counter  
example:



(b)  $g$  must be one-to-one.

Suppose  $g$  is not one-to-one, then there exists  $a_1 \neq a_2 \in A$  and  $g(a_1) = g(a_2)$ . Thus  $f(g(a_1)) = f(g(a_2))$  which means that  $\exists a_1 \neq a_2 \in A$   $f \circ g(a_1) = f \circ g(a_2)$ . Contradict to  $f \circ g$  is one-to-one. Hence  $g$  must be one-to-one.

(c)  $g$  must be one-to-one.

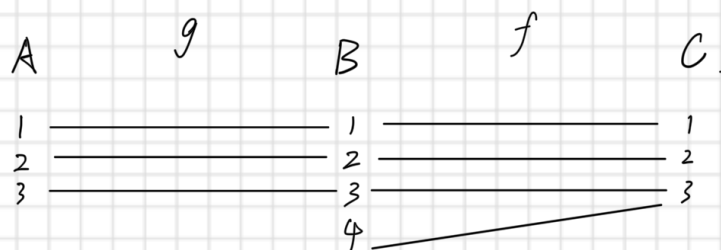
Let  $a_1, a_2 \in A$  that  $g(a_1) = g(a_2)$  then  $f(g(a_1)) = f(g(a_2))$ . Because  $f \circ g$  is one-to-one, thus  $a_1 = a_2$ . Hence  $g$  is one-to-one.

(d) Yes.

If  $f \circ g$  is onto, then for every  $c \in C$ , there exist an  $a \in A$  that  $f \circ g(a) = c$ . Since for every  $a$ , we have a  $b \in B$  that  $b \in B$ . Thus for every  $c \in C$ , we have a  $b \in g(b)$  let  $f(b) = c$ .

Hence  $f$  is onto

(e) No.



Q.9.

$$n^3 - (n-1)^3 = 3n^2 - 3n + 1$$

$$(n-1)^3 - (n-2)^3 = 3(n-1)^2 - 3(n-1) + 1$$

⋮

$$2^3 - 1^3 = 3 \times 2^2 - 3 \times 2 + 1$$

$$1^3 - 0^3 = 3 \times 1^2 - 3 \times 1 + 1$$

$$n^3 = 3(1^2 + 2^2 + \dots + n^2) - 3(1 + 2 + \dots + n) + n$$

$$= 3 \sum_{k=1}^n k^2 - 3 \times \frac{n(n+1)}{2} + n$$

$$\sum_{k=1}^n k^2 = \frac{n^3}{3} + \frac{n(n+1)}{2} - \frac{n}{3}$$

$$= \frac{2n^3 + 3n^2 + n}{6}$$

$$= \frac{n(2n^2 + 3n + 1)}{6} = \frac{n(2n+1)(n+1)}{6}$$

Q.10. 
$$n^4 - (n-1)^4 = (n^2 + (n-1)^2)(n^2 - (n-1)^2)$$

$$= (2n^2 - 2n + 1)(2n - 1)$$

$$= 4n^3 - 6n^2 + 4n - 1$$

$$(n-1)^4 - (n-2)^4 = 4(n-1)^3 - 6(n-1)^2 + 4(n-1) - 1$$

⋮

$$1^4 - 0^4 = 4 \times 1^3 - 6 \times 1^2 + 4 \times 1 - 1$$

$$n^4 = 4 \sum_{k=1}^n k^3 - 6 \sum_{k=1}^n k^2 + 4 \sum_{k=1}^n k - n$$

$$\sum_{k=1}^n k^3 = \frac{1}{4} n^4 - \frac{1}{4} n(2n+1)(n+1) + \frac{n(n+1)}{2} - \frac{n}{4}$$

$$= \frac{n^4 + 2n^3 + 3n^2 + n - 2n^2 - 2n + n}{4}$$

$$= \frac{n^4 + 2n^3 + n^2}{4} = \left[ \frac{n(n+1)}{2} \right]^2$$

Q.11. we have  $(n+1)^2 - n^2$  numbers of  $n$

thus if  $\lfloor \sqrt{m} \rfloor = n+1$   $n = \lfloor \sqrt{m} \rfloor - 1$

$$\begin{aligned} \sum_{k=1}^m \lfloor k \rfloor &= \sum_{i=1}^n i \cdot [(i+1)^2 - i^2] + \sum_{k=(n+1)^2}^m \lfloor k \rfloor \\ &= \sum_{i=1}^n i \cdot (2i+1) + [m - (n+1)^2 + 1] (n+1) \\ &= \frac{n(n+1)(2n+1)}{3} + \frac{n(n+1)}{2} + (n+1) [m - (n+1)^2 + 1] \\ &= \frac{4n^3 + 9n^2 + 5n}{6} + (m+1)(n+1) - (n+1)^3 \\ &= \frac{n(n+1)(4(n+1)+1)}{6} + (m+1)(n+1) - (n+1)^3 \\ &= \frac{(\lfloor \sqrt{m} \rfloor - 1) \lfloor \sqrt{m} \rfloor (4\lfloor \sqrt{m} \rfloor + 1)}{6} + (m+1) \lfloor \sqrt{m} \rfloor - (\lfloor \sqrt{m} \rfloor)^3 \end{aligned}$$

Q.12.  $|A \times C| = |A| \times |C|$

$|B \times D| = |B| \times |D|$

Since  $|A| = |B|$ ,  $|C| = |D|$

we have  $|A| \times |C| = |B| \times |D|$  Hence  $|A \times C| = |B \times D|$

\*Q.13 If  $|A| = |B|$  we can find a bijection  $f: A \rightarrow B$

Thus we have an injection  $f: A \rightarrow B$ , so  $|A| \leq |B|$ .

Meanwhile we have another injection  $f^{-1}: B \rightarrow A$ . so  $|B| \leq |A|$ .

Q.14. We prove a lemma first: If  $f$  and  $g$  are both onto function, then  $f \circ g$  is also an onto function.

Proof: Assume  $g: A \rightarrow B$   $f: B \rightarrow C$ .

Let  $y \in C$ . Because  $f$  is onto,  $y = f(b)$  for some  $b \in B$ .

And Because  $g$  is onto  $b = g(a)$  for some  $a \in A$ .

Hence  $y = f(b) = f(g(a)) = f \circ g(a)$ . which means  $f \circ g$  is onto.

Solution:

If  $A$  is finite, then  $|A| = k$ . And because there is an onto function from  $A$  to  $B$ , then  $|B| \leq |A| = k$   $B$  is finite.

Hence  $B$  is countable.

If  $A$  is countably infinite, Let  $f$  be the onto function from  $A$  to  $B$  and  $g$  be the  $\mathbb{Z}^+ \rightarrow A$  onto function. Thus we can find an onto function  $fo g: \mathbb{Z}^+ \rightarrow B$ .  
Hence  $B$  is countable.

Q.15. If we fix  $m+n=x$ , then the value of  $f(m,n)$  is

$$\left[ \frac{(x-2)(x-1)}{2} + 1, \frac{(x-2)(x-1)}{2} + x - 1 \right] \text{ when } \begin{cases} m=1 \\ n=x-1 \end{cases} \text{ and } \begin{cases} m=x-1 \\ n=1 \end{cases} \text{ get "="}$$

Thus the images of fixed  $m+n$  are continuous integers.

Because  $g(x) = \frac{(x-2)(x-1)}{2}$  covers  $\mathbb{Z}^+$ , to prove  $f$  is one-to-one and onto only need to show the values of  $x$  and  $x+1$  is precisely connected. Which means show  $\frac{(x-2)(x-1)}{2} + x - 1 + 1 = \frac{(x+1-2)(x+1-1)}{2} + 1$

$$\text{left: } \frac{(x-2)(x-1)}{2} + x - 1 + 1 = \frac{x^2 - 3x + 2 + 2x}{2} = \frac{x^2 - x + 2}{2}$$

$$\text{right: } \frac{(x+1-2)(x+1-1)}{2} + 1 = \frac{(x-1)x + 2}{2} = \frac{x^2 - x + 2}{2}$$

$$\text{left} = \text{right}.$$

Proof Done.

Q.16. Define  $f: (0,1) \rightarrow [0,1]$  by  $f(x)=x$  thus  $|(0,1)| \leq |[0,1]|$ .  
Define  $g: [0,1] \rightarrow (0,1)$  by  $g(x) = \frac{x}{2} + \frac{1}{4}$  thus  $|[0,1]| \leq |(0,1)|$ .  
 $\therefore |(0,1)| = |[0,1]|$

Q.17.  $f(x)$  is  $\Theta(g(x)) \rightarrow c_1 g(x) \leq f(x) \leq c_2 g(x) \quad c_1 > 0$ .  
 $g(x)$  is  $\Theta(h(x)) \rightarrow c_3 h(x) \leq g(x) \leq c_4 h(x) \quad c_3 > 0$

$$\text{Then } c_1 c_3 h(x) \leq f(x) \leq c_2 c_4 h(x) \quad c_1 c_3 > 0.$$

$$\text{Thus } f(x) = \Theta(h(x)).$$

Q.18. Counter example:  $f_1(x) = x^2 + x$   $f_2(x) = x^2$ ,  $g(x) = x^2$   
 $(f_1 - f_2)x = x = \Theta(x) \neq \Theta(x^2)$ .

Q.19.

when  $x \geq 1$ .

$$\begin{aligned} |f(x)| &= |a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0| \\ &\leq |a_n| x^n + |a_{n-1}| x^{n-1} + \dots + |a_1| x + |a_0| \\ &\leq |a_n| x^n + |a_{n-1}| x^n + \dots + |a_1| x^n + |a_0| x^n \end{aligned}$$

$$\begin{aligned} \text{Let } M &= \max \{|a_n|, |a_{n-1}|, \dots, |a_0|\} \\ &\leq M x^n + M x^n + \dots + M x^n + M x^n \\ &= (n+1) M x^n = C_1 x^n. \end{aligned}$$

Besides:

$$\begin{aligned} |f(x)| &= |a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0| \\ &= x^n \left| a_n + \frac{a_{n-1}}{x} + \dots + \frac{a_1}{x} + \frac{a_0}{x^n} \right| \end{aligned}$$

$$\text{let } \left| \frac{a_{n-1}}{x} + \dots + \frac{a_0}{x^n} \right| < \left| \frac{a_n}{2} \right|.$$

$$\text{we should have } \left| \frac{nM}{x} \right| < \left| \frac{a_n}{2} \right|.$$

$$\begin{aligned} \text{Therefore: when } x &> \frac{2nM}{a_n} \\ |f(x)| &= x^n \left| a_n + \frac{a_{n-1}}{x} + \dots + \frac{a_1}{x} + \frac{a_0}{x^n} \right| \\ &\geq \left| \frac{a_n}{2} \right| x^n = C_2 x^n \end{aligned}$$

$$\text{Hence } f(x) = \Theta(x^n).$$

$$\text{Q.20. Definition: } C_1 \log_a n \leq \Theta(\log_a n) \leq C_2 \log_a n \quad (a > 1)$$

$$C_1 \cdot \frac{\ln 2}{\ln a} \cdot \frac{\ln n}{\ln a} \cdot \frac{\ln a}{\ln 2} \leq \Theta(\log_a n) \leq C_2 \cdot \frac{\ln 2}{\ln a} \cdot \frac{\ln n}{\ln a} \cdot \frac{\ln a}{\ln 2}$$

$$\underbrace{C_1 \cdot \frac{\ln 2}{\ln a}}_{C_3} \cdot \log_2 n \leq \Theta(\log_a n) \leq \underbrace{C_2 \cdot \frac{\ln 2}{\ln a}}_{C_4} \cdot \log_2 n$$

$$C_3 \log_2 n \leq \Theta(\log_a n) \leq C_4 \log_2 n$$

$$\text{Thus, } \Theta(\log_a n) = \Theta(\log_2 n).$$

$$\text{Q.21. multiplication: } 2n$$

$$\text{addition: } n$$

$$\text{total: } 2n + n = 3n$$



Q22. multiplication :  $n$   
addition :  $n$   
total :  $n+n=2n$

Q23. (1)  $2^{(\log n)^{\log(\log n)}} = 2^{\log(\log n) \cdot \log n} = n^{\log(\log n)} \leq 2^{\log(n^n)} = n^n$  when  $n > 1$ .  
since  $2^n$  is monotonous increasing

Thus  $(\log n)^{\log(\log n)} \leq \log(n^n)$ .

Hence  $(\log n)^{\log(\log n)} = O(\log(n^n))$

(2)  $f_1(n) = \log n$

$f_2(n) = (\log n)^{\log \log n}$

$f_3(n) = (\log \log n)^{\log n}$

$f_4(n) = (\log n)^{\log n}$

$f_5(n) = n \log n$

$f_6(n) = (\log n)^n$

$f_7(n) = n^2$

$f_8(n) = n^{\log n}$

$f_9(n) = 3^{\frac{n}{2}}$

Q24. A, C, E.