

# DiscretMath Assignment1

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## 1 Q.1

a.  $r \wedge \neg q$

b.  $r \rightarrow p$

c.  $p \wedge \neg q \wedge r$

d.  $(p \wedge q) \rightarrow r$

e.  $r \leftrightarrow (p \vee q)$

## 2 Q.2

2. a

$p$	$q$	$\neg q$	$p \oplus q$	$p \wedge q$	$p \oplus \neg q$	$(p \oplus q) \rightarrow (p \wedge q)$	$(p \oplus q) \rightarrow (p \oplus \neg q)$
0	0	1	0	0	1	1	1
0	1	0	1	0	0	0	0
1	0	1	1	0	0	0	0
1	1	0	0	1	1	1	1

they are equivalent

2. b

$p$	$q$	$p \leftrightarrow q$	$p \wedge q$	$\neg p \wedge \neg q$	$(p \wedge q) \vee (\neg p \wedge \neg q)$
0	0	1	0	1	1
0	1	0	0	0	0
1	0	0	0	0	0
1	1	1	1	0	1

they are equivalent

2. c

$p$	$q$	$\neg p$	$\neg q$	$p \rightarrow q$	$\neg q \wedge \neg(p \rightarrow q)$
0	0	1	1	1	0
0	1	1	0	1	0
1	0	0	1	0	1
1	1	0	0	1	0

they are not equivalent

2. d

$p$	$q$	$r$	$p \rightarrow \neg q$	$p \vee \neg q$	$r \rightarrow (p \vee \neg q)$	$(p \rightarrow \neg q) \leftrightarrow (r \rightarrow (p \vee \neg q))$	$\neg p \wedge \neg r$	$q \vee (\neg p \wedge \neg r)$
0	0	0	1	1	1	1	1	1
0	0	1	1	1	1	1	0	0
0	1	0	1	0	1	1	1	1
0	1	1	1	0	0	0	0	1
1	0	0	1	1	1	1	0	0
1	0	1	1	1	1	1	0	0
1	1	0	0	1	1	0	0	1
1	1	1	0	1	1	0	0	1

they are not equivalent

### 3 Q.3

a.

$$\begin{aligned}
(p \wedge \neg q) \rightarrow r &\equiv \neg(p \vee \neg q) \vee r && \text{Useful} \\
&\equiv (\neg p \vee \neg \neg q) \vee r && \text{De Morgan's} \\
&\equiv (\neg p \vee q) \vee r && \text{Double Negation} \\
&\equiv \neg p \vee (q \vee r) && \text{Associative} \\
&\equiv p \rightarrow (q \vee r) && \text{Useful}
\end{aligned}$$

b.

$$\begin{aligned}
& ((p \rightarrow q) \wedge (q \rightarrow r)) \rightarrow (p \rightarrow r) \\
& \equiv \neg((\neg p \vee q) \wedge (\neg q \vee r)) \vee (\neg p \vee r) \quad \text{Useful} \\
& \equiv (\neg(\neg p \vee q) \vee \neg(\neg q \vee r)) \vee (\neg p \vee r) \quad \text{De Morgan's} \\
& \equiv (\neg\neg p \wedge \neg q) \vee (\neg\neg q \wedge \neg r) \vee (\neg p \vee r) \quad \text{De Morgan's} \\
& \equiv (p \wedge \neg q) \vee (q \wedge \neg r) \vee (\neg p \vee r) \quad \text{Double Negation} \\
& \equiv (p \vee (q \wedge \neg r) \vee (\neg p \vee r)) \wedge (\neg q \vee (q \wedge \neg r) \vee (\neg p \vee r)) \quad \text{Distributive} \\
& \equiv (p \vee (\neg p \vee r) \vee (q \wedge \neg r)) \wedge ((q \wedge \neg r) \vee \neg q \vee (\neg p \vee r)) \quad \text{Commutative} \\
& \equiv ((p \vee \neg p) \vee r \vee (q \wedge \neg r)) \wedge ((q \wedge \neg r) \vee \neg q \vee (\neg p \vee r)) \quad \text{Associative} \\
& \equiv ((p \vee \neg p) \vee r \vee (q \wedge \neg r)) \wedge ((q \vee \neg q) \vee (\neg p \vee r)) \wedge ((\neg r \vee \neg q) \vee (\neg p \vee r)) \quad \text{Distributive} \\
& \equiv ((p \vee \neg p) \vee r \vee (q \wedge \neg r)) \wedge ((q \vee \neg q) \vee (\neg p \vee r)) \wedge ((\neg r \vee r) \vee \neg q \vee \neg p)) \quad \text{Associative} \\
& \equiv (T \vee r \vee (q \wedge \neg r)) \wedge (T \vee (\neg p \vee r)) \wedge (T \vee \neg q \vee \neg p) \quad \text{Negation} \\
& \equiv T \wedge T \wedge T \quad \text{Domination} \\
& \equiv T \quad \text{Identity}
\end{aligned}$$

#### 4 Q.4

$p$	$q$	$r$	$p \rightarrow q$	$(p \rightarrow q) \rightarrow r$	$q \rightarrow r$	$p \rightarrow (q \rightarrow r)$
0	0	0	1	0	1	1
0	0	1	1	1	1	1
0	1	0	1	0	0	1
0	1	1	1	1	1	1
1	0	0	0	0	1	1
1	0	1	0	1	1	1
1	1	0	1	0	0	0
1	1	1	1	1	1	1

From the Truth table, we can get  $(p \rightarrow q) \rightarrow r$  and  $p \rightarrow (q \rightarrow r)$  are **not equivalent**.

## 5 Q.5

$$\begin{aligned}
(q \rightarrow (r \vee p)) &\rightarrow ((\neg r \vee s) \wedge \neg s) \\
&\equiv \neg(\neg q \vee (r \vee p)) \vee ((\neg r \vee s) \wedge \neg s) && \text{Useful} \\
&\equiv (q \wedge \neg(r \vee p)) \vee ((\neg r \vee s) \wedge \neg s) && \text{De Morgan's} \\
&\equiv (q \wedge \neg r \wedge \neg p) \vee ((\neg r \vee s) \wedge \neg s) && \text{De Morgan's} \\
&\equiv (q \wedge \neg r \wedge \neg p) \vee (\neg r \wedge \neg s) \vee (s \wedge \neg s) && \text{Distributive} \\
&\equiv (q \wedge \neg r \wedge \neg p) \vee (\neg r \wedge \neg s) \vee F && \text{Negation} \\
&\equiv (q \wedge \neg r \wedge \neg p) \vee (\neg r \wedge \neg s) && \text{Identity} \\
&\equiv \neg r \wedge ((q \wedge \neg p) \vee \neg s) && \text{Distributive} \\
&\rightarrow \neg r && \text{Simplication Inference}
\end{aligned}$$

## 6 Q.6

- a.  $\forall x \exists y L(x, y)$
- b.  $\exists y \forall x L(x, y)$
- c.  $\neg \exists x \forall y L(x, y)$
- d.  $\exists y \forall x \neg L(x, y)$
- e.  $\exists y \forall x (L(x, y) \wedge \forall z (L(x, z) \rightarrow y = z))$
- f.  $\exists y_1 \exists y_2 (y_1 \neq y_2 \wedge L(\text{lynn}, y_1) \wedge L(\text{lynn}, y_2) \wedge \forall z (L(\text{lynn}, z) \rightarrow (z = y_1 \vee z = y_2)))$
- g.  $\exists x (L(x, x) \wedge \forall y (L(x, y) \rightarrow x = y))$

## 7 Q.7

1.  $\forall x (I(x) \rightarrow E(x))$
2.  $\forall x \forall y (L(x, y) \rightarrow \neg Q(x, y))$
3.  $\neg \exists y \forall x L(x, y)$

## 8 Q.8

1. Domain of the  $x, y$  are both  $\mathbb{N}$ ,  $P(x, y) = (x \geq y)$ . Thus  $\exists x \forall y P(x, y)$  is **false** but  $\forall y \exists x P(x, y)$  is **true**
2. NO, as the example in 1.

## 9 Q.9

1. not equivalent. When  $P(x)$  is  $x \geq 0$  and  $Q(x)$  is  $x < 0$ . Then i is **false** but ii is **true**
2. equivalent
3. equivalent
4. equivalent

## 10 Q.10

- a. Replace  $n^3 + 6n + 5$  is odd with  $P(n)$  and  $n$  is even with  $Q(n)$ .  
Thus

$$\begin{aligned}
 \neg \forall n \in \mathbb{N} (P(n) \rightarrow Q(n)) &\equiv \exists n \in \mathbb{N} \neg (P(n) \rightarrow Q(n)) \\
 &\equiv \exists n \in \mathbb{N} \neg (\neg P(n) \vee Q(n)) \\
 &\equiv \exists n \in \mathbb{N} (\neg \neg P(n) \wedge \neg Q(n)) \\
 &\equiv \exists n \in \mathbb{N} (P(n) \wedge \neg Q(n))
 \end{aligned}$$

Replace the Predict backward, the negation is:

$$\exists n \in \mathbb{N} (n^3 + 6n + 5 \text{ is odd} \wedge n \text{ is even})$$

- b. Assume the negation is true. Then we can get  $\exists n \in \mathbb{N} n \text{ is even}$ . Thus  $n^3 + 6n + 5$  is even, which leads to the negation being false. Because of the contradiction, the negation is false. So the original statement is true.

## 11 Q.11

a. using truth table

$p$	$q$	$(q \vee \neg p)$	$\neg(p \leftrightarrow (q \vee \neg p))$	$\neg p \vee \neg q$
0	0	1	1	1
0	1	1	1	1
1	0	0	1	1
1	1	1	0	0

- b. Regardless how long  $P$  is and logical connectives  $P$  uses, the truth table of  $P$  has only  $2^2 = 4$  rows. It implies that  $P$  has totally  $2^{2^2} = 16$  truth tables.

Meanwhile, for the situation in the question, it has  $4 \times 3 \times 4$  collections. Because the connectives  $\wedge, \vee, \leftrightarrow$  are all symmetric, we can remove half of them for equivalence by the **Commutative laws**, getting  $4 \times 3 \times 4 / 2 = 24$ .

Despite the Commutative laws. Think other situations that  $A \square B$  can produce equivalence. For the reason that  $A \square B$  is the simplest form of

binary connections,  $\wedge, \vee, \leftrightarrow$  can't be equivalent in this form **unless** they are **tautology** or **contradiction**.

Enumerate all the situation that are **tautology** or **contradiction**:

Tautology	Contradiction
$p \vee \neg p$	$p \wedge \neg p$
$q \vee \neg q$	$q \wedge \neg q$
$p \leftrightarrow p$	$p \leftrightarrow \neg p$
$q \leftrightarrow q$	$q \leftrightarrow \neg q$
$\neg p \leftrightarrow \neg p$	
$\neg q \leftrightarrow \neg q$	

Just remain one tautology and one contradiction, we can get the **distinct truth table**  $A \square B$  can produce:  $2^4 - 5 - 3 = 16$ . Match the number of  $P$ 's truth tables.

Thus, we have proven that whatever  $P$ 's truth table is, we can select one  $A \square B$  that satisfy the situation.

## 12 Q.12

"All movies produced by John Sayles are wonderful. John Sayles produced a movie about coal miners. (This movie about coal miners is wonderful)"

→ **Hypothetical syllogism**

"(This movie about coal miners is wonderful) Therefore, there is a wonderful movie about coal miners." → **Existential Generalization(EG)**

## 13 Q.13

1. counterexample1:  $a = \sqrt{2}^{\sqrt{2}}$ ,  $b = \sqrt{2}$ ,  $a^b = 2$  If  $a$  is irrational number, then we have proved it. If  $a$  is rational number, then we have the  $b^b = a$  implies irrational number  $\sqrt{2}$  power irrational number  $\sqrt{2}$  is rational number  $\sqrt{2}^{\sqrt{2}}$

counterexample2:  $a = e$ ,  $b = \ln 2$ ,  $a^b = 2$

2. counterexample:  $a = -\sqrt{2}$ ,  $\sqrt{a} \notin \mathbb{R}$ , it's not an irrational number.

• I guess there miss a condition that  $a > 0$ . In this condition, proof:

Suppose  $a$  is an irrational number but  $\sqrt{a}$  is rational, Then there exist two integers  $m$  and  $n$  such that  $\gcd(m, n) = 1$ , and  $\sqrt{a} = \frac{m}{n}$ . We have then  $a = \frac{m^2}{n^2}$ . It then shows  $a$  is a rational number, which contradicts to the assumption. Thus prove the proposition.

3. For example:  $x = 2$ ,  $y = \frac{1}{\ln 2}$

proof:  $2^{\frac{1}{\ln 2}} = e$  is an irrational number.

Because  $e^{\ln 2} = 2$ , thus  $2^{\frac{1}{\ln 2}} = e$ , which is obviously an irrational number.

## 14 Q.14

Suppose  $r = \sqrt{2} + \sqrt{3}$  is rational, then  $r^2 = 5 + 2\sqrt{6}$  should be rational. It follows that  $\frac{r^2 - 5}{2} = \sqrt{6}$  should be rational. From the theorem we can get that  $\sqrt{6}$  is an irrational number, which contradicts to the assumption. Thus  $\sqrt{2} + \sqrt{3}$  is irrational.

## 15 Q.15

Suppose we have rational number  $x$  and irrational number  $y$ , we can select an irrational number  $z = \frac{x+y}{2}$ . Then we just need to proof  $z$  is irrational number. proof: Suppose  $z = \frac{x+y}{2}$  is rational, then  $y = 2z - x$  should be rational, which contradicts to the assumption. Then  $z = \frac{x+y}{2}$  is an irrational number.

Hence between every rational number and every irrational number there is an irrational number.

## 16 Q.16

Suppose  $a^2 + b^2 = 2n$ , then  $a^2 - b^2 = 2n - 2b^2$ . It follows that  $(a+b)(a-b)$  is even. Thus  $a+b$  or  $a-b$  is even. Meanwhile,  $a+b = a-b + 2b$ , which implies  $a+b$  and  $a-b$  are both even or both odd. Combine the two condition,  $a+b$  is even.

## 17 Q.17

Assume one real root  $r$  of  $f(x) = 0$  is neither integral or irrational, which means  $r$  can be expressed by  $r = \frac{m}{q}$  and  $m, q \in \mathbb{Z}$  and  $\gcd(m, q) = 1$ .

$$\text{Get } f(r) = f\left(\frac{m}{q}\right) = a_0 + a_1\left(\frac{m}{q}\right) + a_2\left(\frac{m}{q}\right)^2 + \cdots + a_{n-1}\left(\frac{m}{q}\right)^{n-1} + \left(\frac{m}{q}\right)^n = 0$$

$$\rightarrow a_0q^n + a_1mq^{n-1} + a_2m^2q^{n-2} + \cdots + a_{n-1}m^{n-1}q + m^n = 0$$

$$\rightarrow a_0q^n + a_1mq^{n-1} + a_2m^2q^{n-2} + \cdots + a_{n-1}m^{n-1}q = -m^n$$

$$\rightarrow (a_0q^{n-1} + a_1mq^{n-2} + a_2m^2q^{n-3} + \cdots + a_{n-1}m^{n-1})q = -m^n$$

Because  $q$  is factor of  $(a_0q^{n-1} + a_1mq^{n-2} + a_2m^2q^{n-3} + \cdots + a_{n-1}m^{n-1})q$

$$\rightarrow q \text{ is factor of } m^n$$

With the result in the question  $\rightarrow q$  is factor of  $m$

Denote that  $\gcd(m, q) = 1$

$$\rightarrow q \text{ is factor of } 1$$

$$\rightarrow q = 1$$

$$\rightarrow r = m \text{ (is an integer)}$$

Which is contradicts to the assumption. Thus the real root is either integral or irrational.