

1. Bottom-Up-Cut-Rod(p, n)

Let $r[0 \dots n]$ be a new array

$r[0] = 0$

for $j=1$ to n do

$q = -\infty$

for $i=1$ to j do

$q = \max(q, p[i] + r[j-i] - c)$

$r[j] = q$

return $r[n]$

2. MEMOIZED-CUT-ROD(p, n)

let $r[0:n]$ be a new array

let $s[0:n]$ be a new array

for $i=0$ to n

$r[i] = -\infty$

$s[i] = 0$

return MEMOIZED-CUT-AUX(p, n, r, s)

MEMOIZED-CUT-ROD-AUX(p, n, r, s)

if $r[n] \geq 0$

return $r[n], s[n]$

if $n == 0$

$q = 0$

else $q = -\infty$

cut-position = 0

for $i=1$ to n

$tmp = p[i] + \text{MEMOIZED-CUT-ROD-AUX}(p, n-i, r)[0]$

if $tmp > q$

$q = tmp$

cut-position = i

$r[n] = q$

$s[n] = \text{cut-position}$

return $q, s[n]$.

GET - CUTS (s, n)

cuts = []

while $n \geq 0$:

cut.append $s[n]$

$n = s[n]$

return cuts

3. let $F[0:n]$ be a new array

for $i = 0$ to n

$F[i] = 0$

$F[1] = 1$

index = 1

FIBONACCI(n)

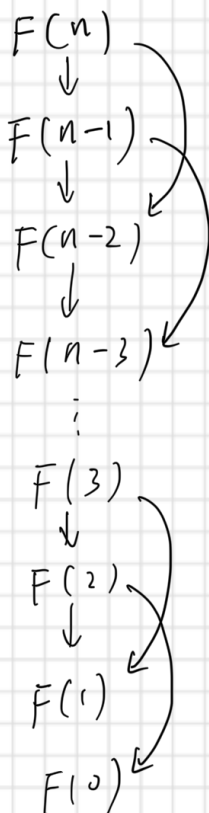
if $n \leq \text{index}$

return $F[n]$

for $i = \text{index} + 1$ to n

$F[i] = F[i-1] + F[i-2]$

return $F[n]$.



$$V = n + 1$$

$$E = n + n - 1 = 2n - 2$$

4. a)

$$A[k] = \begin{cases} A[k-1] & \text{not contain the shares} \\ B[k-1] + a[k] & \text{contain the shares} \end{cases}$$

$$B[k] = \begin{cases} B[k-1] + a[k] & \text{not share} \\ a[k] & \text{buy the share on } k. \end{cases}$$

MAX-RETURN (a)

if $n \leq 1$ return 0

let $A[0, \dots, n]$, $B[0, \dots, n]$ be new array

$$A[0] = 0$$

$$B[0] = -\infty$$

let $C[0, \dots, n]$ be a new array

for i from 1 to n

if $B[i-1] > 0$

$$B[i] = B[i-1] + a[i]$$

else $B[i] = a[i]$; $C[i] = i$.

buy-day = sell-day = 0.

for i from 1 to n

if $A[i-1] > B[i-1] + a[i]$

$$a[i] = A[i-1]$$

else $A[i] = B[i-1] + a[i]$

buy-day = $C[i-1]$.

sell-day = i

return $A[n]$, buy-day, sell-day.

$$b) T(n) = O(1) + O(1) + \dots + O(n) + O(n) = O(n)$$