

## Documentation

Regular Expression:

RE =  $(ba)^* b^* a [a + ba(a+b)]^*$

Alphabet:

$\Sigma = \{a, b\}$

### 1. Explanation of the Regular Expression

Part 1:  $(ba)^*$

Means zero or more repetitions of the string ba.

Examples:  $\epsilon$ , ba, baba, bababa.

Part 2:  $b^*$

Means zero or more b.

Examples:  $\epsilon$ , b, bb, bbb.

Part 3: a

Mandatory single a.

Every valid string must contain this a.

Part 4:  $[a + ba(a+b)]^*$

Allows zero or more occurrences of:

a, baa, bab.

Overall Meaning:

A valid string starts with zero or more ba, may have extra b,  
must contain at least one a, and can end with combinations of a, baa, or bab.

## 2. Thompson's Construction (ENFA)

Thompson's construction is used to build the ENFA.

Symbols are converted to basic NFAs.

Operators:

\* for closure using epsilon loops,  
+ for union using epsilon branching,  
concatenation using epsilon transitions.

## 3. ENFA to DFA Conversion

Subset construction algorithm is used.

Each epsilon-closure becomes a DFA state.

States containing final ENFA state are accepting.

## 4. DFA Description

Start State:  $q_0$

Accepting State:  $q_3$

States:

q0 – Start, processing (ba)\*

q1 – Seen b, expecting a

q2 – Processing b\*

q3 – After mandatory a (Accepting)

q4 – After b in final part

q5 – After ba in final part

DEAD – Invalid transition

## 5. DFA Transition Table

q0: a→DEAD, b→q1

q1: a→q0, b→DEAD

q2: a→q3, b→q2

q3: a→q3, b→q4

q4: a→q5, b→DEAD

q5: a→q3, b→q3

## 6. DFA Implementation

The DFA is implemented in C++ using enumerated states and a transition function.

The input string is processed character by character.

The string is accepted if the final state is an accepting state.

## 7. Sample Test Cases

a → Accepted

bba → Accepted

bababaa → Accepted

bb → Rejected

ba → Rejected

## 8. Conclusion

The regular expression was analyzed.

Thompson's ENFA was constructed and converted to DFA.

The DFA was implemented in C++ and correctly validates strings.