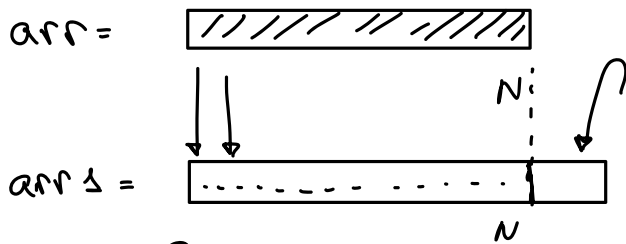


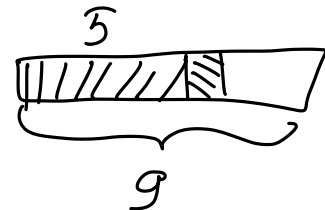
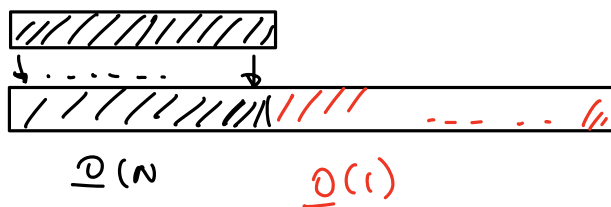
# Линейные контейнеры

① Динамически расширяющийся массив.



Асимптотика:  
 $O(N)$

for  $i = 0 \dots N$   
 $arr1[i] = arr[i]$       $O(N)$



class Dynamic Array

- 1. size - кол-во элементов в массиве
- 2. capacity - емкость контейнера
- 3. array - контейнер

```
метод is_empty ( ) ?  
return size == 0  
}
```

```
метод size ( ) ?  
return size  
}
```

```
метод resize ( ) ?
```

new\_arr = [... ] // размер capacity \* 2  
for  $i = 0 \dots \text{capacity}$

$O(N)$

$\text{new\_arr}[i] = \text{arr}[i]$

delete arr

$\text{arr} = \text{new\_arr}$

$\text{capacity} \times = 2$

}

$\text{memog push\_back}(\text{elem}) \{$

$\text{ecu size} + 1 \geq \text{capacity}$

$\hookrightarrow \text{resize}()$

$\text{arr}[\text{size}] = \text{elem}$

$\text{size} += 1$

}

$\text{memog pop\_back}() \{$

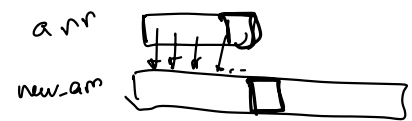
$\text{ecu size} = 0$

$\hookrightarrow \text{Queue}$

$\text{size} -= 1$

$\text{return arr}[\text{size}]$

}



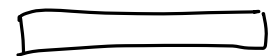
как надо

$O(N)$   $\downarrow$   $P = \frac{1}{N}$

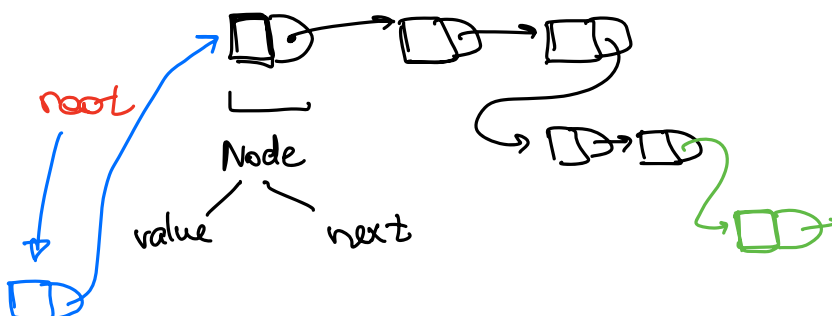
$O(1)^*$   
 $\uparrow$   
 $O(1)$   $\&$   
сложно.

Время работы push-back ?

$\hookrightarrow O^*(1)$



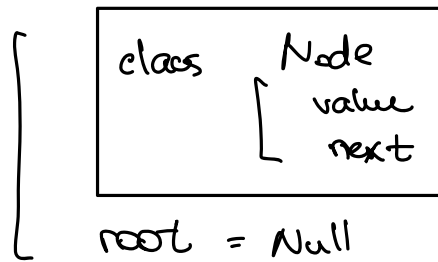
II АТД список  $\rightarrow$  АСМРК мтери тун ААтРДх  
(огнодототон список)



push-back (✓)  
pop-back (✓)  
push-front  
pop-front

```
class
{
```

List



```
push-back (elem) {
```

```
  runner = root
```

```
  while runner.next != Null
    runner = runner.next
```

```
  new-node = Node (elem)
```

```
  new-node.next = Null
```

```
  runner.next = new-node
```

```
}
```

```
pop-back () {
```

```
  runner = root
```

```
  if runner == null // ret -6
    return
```

```
  if runner.next == null // ret -7
    return
```

```
  delete runner
```

```
  root = null
```

```
  return
```

```
  while runner.next.next != null
    runner = runner.next
```

```
  delete runner.next
```

```
  runner.next = null
```

```
}
```

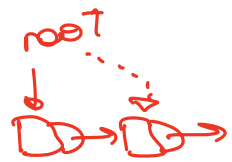


```

push-front(elem)
{
    new_elem = Node(elem)
    new_elem.next = root
    root = new_elem
}

pop-front() {
    if (root.next == null) {
        delete root
        root = null
    }
    old_root = root
    root = old_root.next
    delete old_root
}

```



Асимптотика:

push-back ||  $O(N)$   
pop-back

pop-front ||  $O(1)$   
push-front

Dynamic Array VS List

- + mem  $O^*(1)$

+ - odp. no indexy [guara]

+ - znacheniya  
sobytiy

korga integer

Ⓢ xpramye pozitsii  
↑ tunc gatax

Obor kross → 4MB

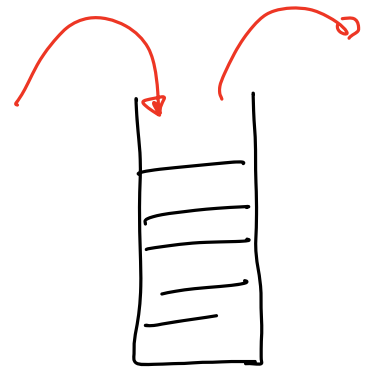
III Ha nagymértékű

1) Címek (ATL)

Címek kérés

FIFO

LIFO



Programozási Címek: (Ha Címek)

push-front  $\mathcal{O}(1)$   
pop-front  $\mathcal{O}(1)$

class Stack {

|| stack = List()

push(elem) {

stack.push\_front(elem) //  $\mathcal{O}(1)$

pop() {

stack.pop\_front() //  $\mathcal{O}(1)$

}

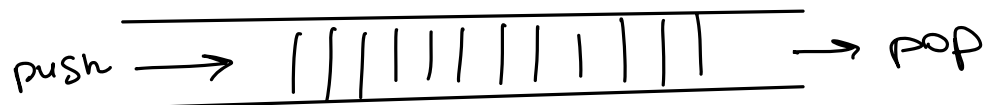
top() {

return stack.root.value //  $\text{repl} \rightarrow \text{top}$

}

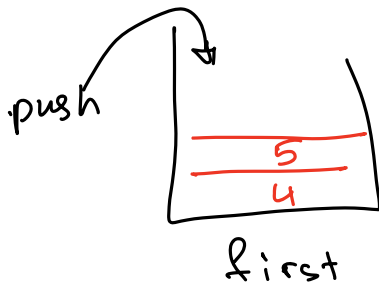
}

IV Összegzés



FIFO  
LIFO

Dragega ra 2x camera



Queue ?

|| stack first  
|| stack last

push (elem) ?  
first.push (elem)  
}

|| O(1)

pop ( ) ?

ecnu not last.is\_empty() ?  
elem = last.top()  
last.pop()  
return elem  
}

O(1)

while not first.is\_empty()  
last.push (first.top())  
first.pop()  
}

O(N)  
 $\frac{1}{N}$

ecnu last.is\_empty()

Queue OKA

unrare: elem = last.top()  
last.pop()  
return elem  
}

O(1)

}

Acharya ~~Totaka~~  $\mathcal{O}^*$ (s)