Queens College

Internet and Web Technology (CSCI 355)

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1/11 Lecture07: Internet Security Notes

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Internet Security

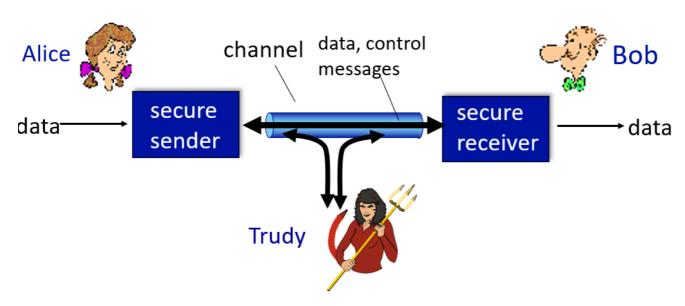
What is network security?

Goals: CIA

- Confidentiality: Do not allow other (non-intended) actors to have access to our data.
 - Not so simple in the context of networking. Analogy: driving in the highway, everyone can see you and your car.
- Integrity: ensure data is not altered in any way.
- Authentication/Availabilty: Sender, receiver confirm their identities and services must be accessible and available to users.

Goals addressed in different ways.

Scenarios:



Alice and Bob want to communicate but Trudy may intercept, delete, add messages.

Alice and Bob can represent:

- Actual people trying to communicate, exchange information.
- Web browser/server
- On-Line banking client/server
- DNS server

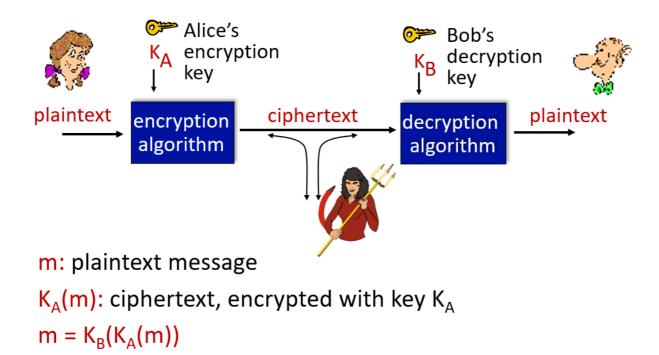
• BGP routers: Border Gateway Protocol (standard inter-AS protocol)

Trudy can perform:

- Eavesdrop: ex. using Wireshark to see what people are doing.
- Insert messages.
- Impersonate.
- Hijack
- DoS: Denial of Service, overwhelming a server with activity that it does not have time for anything else.
- DDoS: Distributed Denial of Service:
 - o more volume, generate bad traffic
 - harder to track

Answer?: Cyptography

CRYPTOGRAPHY



- m = message
- K(A) = encryption key
- K(B) = decryption
- ciphertext = plaintext packet that has been encrypted.

Question: How much of the packet do you encrypt?

• Can be super secure that it can become useless.

Shift-Cipher: Simple, shifting letters by a set amount. Ex. bump A to become D, B to E, etc.

- Encryption shifts forward
- Decryption shifts backward

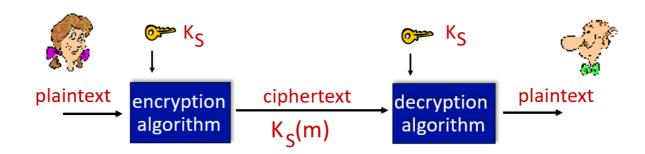
• 26 possible shifts if using english alphabet, can be more if others used.

Permutation-Cipher: "A" gets mapped to "E" but "B" gets mapped to "X". Leads to 26! possibilities.

Breaking Encryption Schemes:

- Use brute force. Shift-cipher can be easily cracked.
- Statistical Analysis.
- Trudy has the plaintext corresponding to ciphertext.

Symmetric Key Cryptography



Bob and Alice share same (symmetric) key: K

plaintext: abcdefghijklmnopqrstuvwxyz

ciphertext: mnbvcxzasdfghjklpoiuytrewq

e.g.: Plaintext: bob. i love you. alice ciphertext: nkn. s gktc wky. mgsbc

Permutation-cipher that uses an encryption key that maps from a set of 26 letters to a set of 26 letter (26!).

To make more difficult: have several permutation ciphers and alternate between them but it needs to be known/shared, this is still risky.

DES: Data Encryption Standard

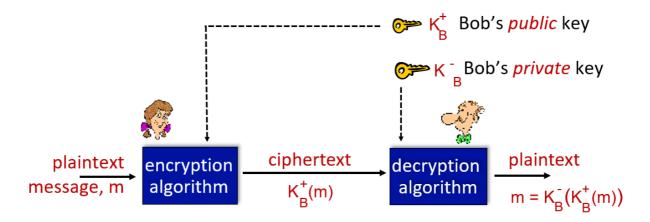
- Symmetric key.
- Chopped into chunks.
- Not that good in modern day since it can be decrypted by brute force.
- 1 sec to decrypt with modern technology.

AES: Advanced Encryption Standard

- Symmetric key.
- Harder than DES, 149 trillion years to decrypt using brute force.

Public Key Cryptography: Highlight

- How can something public be secret?
- RSA78
- public encryption key known to all, can be accessed.
- private decryption key known only to receiver.
- Not a symmetric key, just because you know how to encrypt doesn't mean you know to decrypt. Ex: A bank wants people to send it data, the bank publishes the encryption key but only the customer knows the decryption.



Alice wants to send an encrypted message to Bob: * Alice uses Bob's public key to encrypt. * But only Bob's private key can decrypt. * Bad actors do not know Bob's private key, not even Alice.

RSA: Encryption/Decryption algorithm developed by Rivest, Shamir, Adelson 0. given (n,e), public key, and (n,d), private key.

- $n = product of two large prime numbers, p and q (ex: 1024 bits, ~10^300 range).$
- e < n, e has no common factors with z
- z = (p-1)(q-1)
- d = (ed) 1 is exactly divisible by z ((ed)mod z = 1)
- 1. to encrypt message m (< n) compute $c = m^e \mod n$
 - Take message and convert to binary number
 - Take binary and convert to decimal
 - Take decimal and raise to e power.
- 2. to decrypt received bit patter c, compute $m = c^d \mod n$

Magic: $m = (m^e \mod n)^d \mod n$

- the relationship between e and d is such that they cancel each other out and get back m.
- e and d are carefully chosen
- whole world knows e, but they don't know d.
- n ~ 10^700
- nature of e and d is such that they can be interchanged and result in the same thing. Ex: $(A^B)^C = (A^C)^B$
 - useful depending on what is trying to be accomplished.

• Can be used as integrity.

The following property will be very useful later:

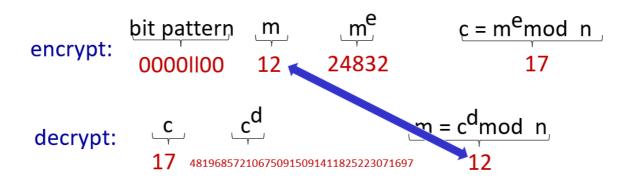
$$K_B(K_B^+(m)) = m = K_B^+(K_B^-(m))$$

use public key use private key first, followed by private key by public key

result is the same!

RSA example:

Bob chooses p=5, q=7. Then n=35, z=24. e=5 (so e, z relatively prime). d=29 (so ed-1 exactly divisible by z). encrypting 8-bit messages.



Problem: SLOW to work with. Trick: Use some RSA to share the permutation cipher then use permutation cipher for the remainder of the sessions.

Two-phased model

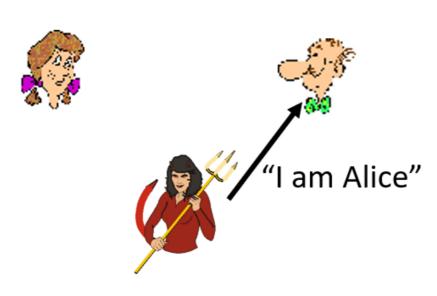
Authentication

How to use cryptography to achieve goals?





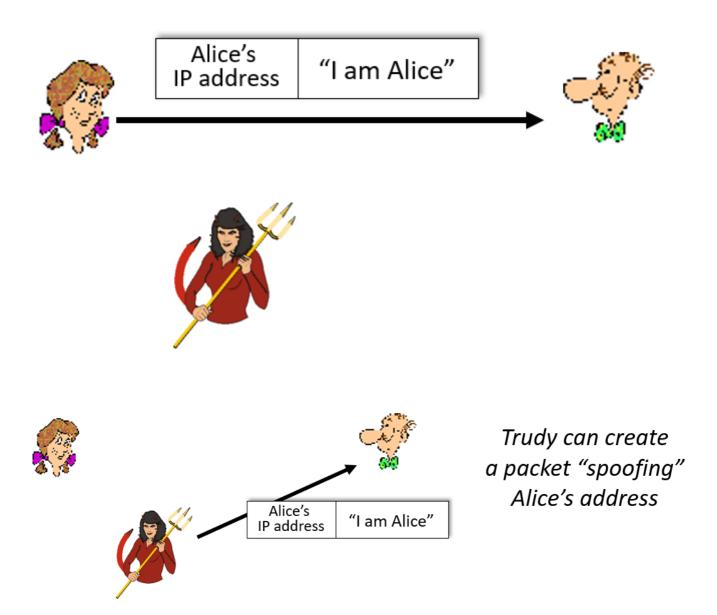
Failure Scenario 1:



in a network, Bob can not "see" Alice, so Trudy simply declares herself to be Alice

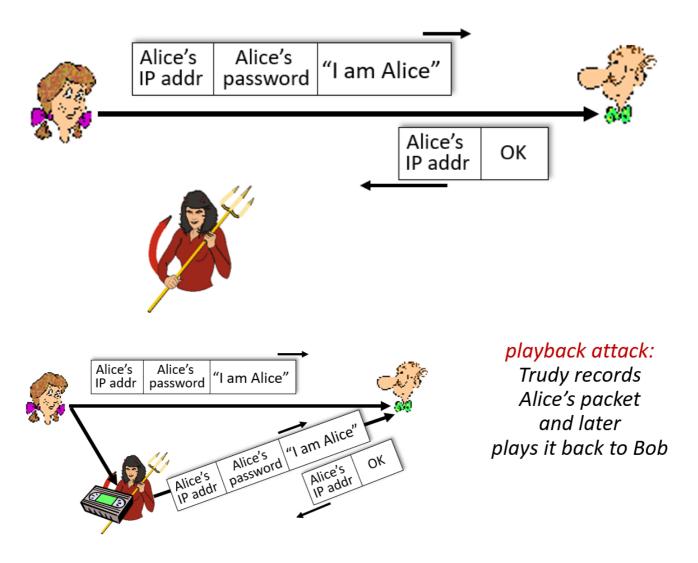
• Al can be used to replicate the "voice" tricking Bob to think it is Alice.

Failure Scenario 2:



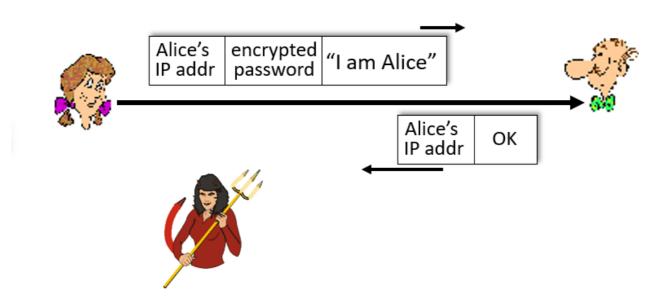
- Alice sends from her own IP address.
- Trudy spoofs Alice's IP address, Bob recognized IP address and is tricked.

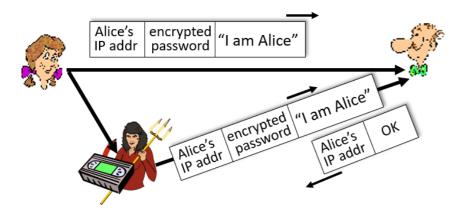
Failure Scenario 3:



- Alice uses a password in the message that Bob recognizes.
- Trudy can read channel, sees weird characters that is the PW. She records and uses the same PW to trick Bob.

Failure Scenario 4:



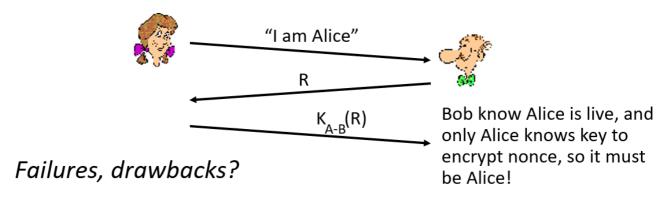


playback attack still works: Trudy records Alice's packet and later plays it back to Bob

- Alice encrypts password.
- Bob see's something encrypted and understands its Alice.
- Problem: Trudy, since she is monitoring, uses the same encrypted password.

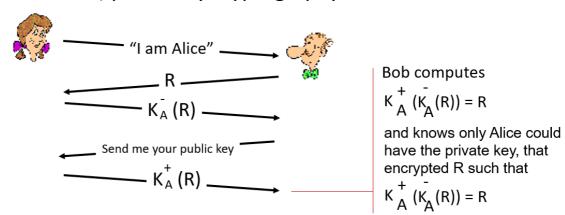
Failure Scenario 5:

Alice must return R, encrypted with shared secret key

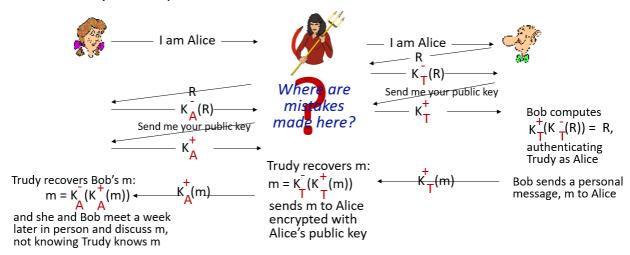


- Bob sends a **naunce** message, used once in a lifetime.
- Alice sends the naunce back to Bob encrypted.
- Will not be the same each time.

ap5.0: use nonce, public key cryptography



man (or woman) in the middle attack: Trudy poses as Alice (to Bob) and as Bob (to Alice)



Trudy sits in between and wrecks havoc.

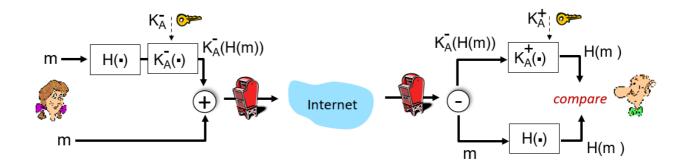
Digital Signatures

- Analogous to hand-written signatures, ex: pdf's.
- Message Digest (MD): ex MD5 available in databases, SQL.
 - Kind of a hash
 - SHA-1 also used
 - MD5 is large, must make sure there is enough space in DB to store MD5.
- Internet CheckSums:

Need for certified public keys

- ex. pizza shop asks for payment in advance.
- Certificate Authority (CA) needed to certify users. To record/validate identity for access to a public key.

Secure Email: REMEMBER FOR EXAM!!

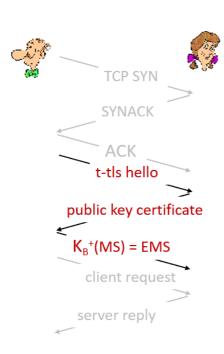


- Which side is the public key, private key in terms of confidentiality?
 - Sent with the public key. Only one who can decrypt it is the recipient.
- For message integrity and authentication, it's the other way around.
 - Use the private key to encrypt, anyone can decrypt it with the public key but only the individual can do it.

Transport-layer Security (TLS)

- A combination of message integrity, authentication and confidentiality.
- Together with HTTP makes HTTPS.
- Provides:
 - o confidentiality via symmetric encryption (NOT RSA, instead permutation ciphers types).
 - integrity via cryptographic hashing (if someone were to change one pixel or character, you will find out about it since it will be different)
 - o authentication via public key cryptography.

t-tls: initial handshake



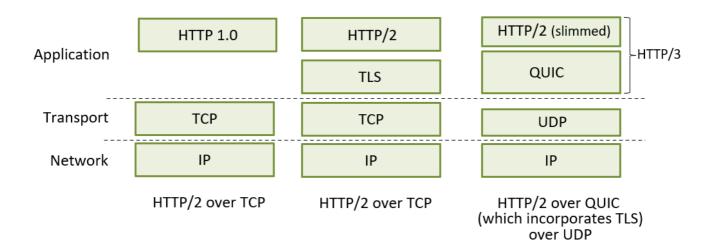
t-tls handshake phase:

- Bob establishes TCP connection with Alice
- Bob verifies that Alice is really Alice
- Bob sends Alice a master secret key (MS), used to generate all other keys for TLS session
- potential issues:
 - 3 RTT before client can start receiving data (including TCP handshake)
- top light-gray exchange is standard three-way handshake.
- Problem: 3 RTT (Round Trip Time). A lot of overhead.

T-TLS: Encrypting Data

- Possible issues:
 - Re-ordering: man-in middle intercepts TCP segments and reorders
 - Changing the sequence makes putting it back together a problem.
 - Connection closure: attacker forges TCP connection close segment. Must check if legitimate closure.

HTTP View of TLS:



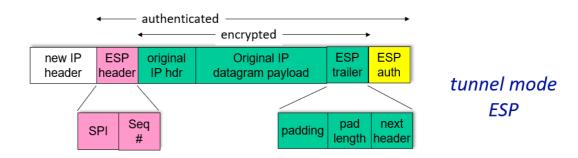
- HTTP 1.0 has no security
- HTTP/2 (HTTPS) is HTTP + TLS
- HTTP/3 uses QUIC, not using TLS works around it.

Some approaches attempt to wrap more of the cryptography earlier on in the exchange.

Where does this cryptography happen?

- Application Layer:
 - o In computer email client applications
 - ZIP, using WinZip PW protection
 - VPN essentially
 - People don't want to bother with encryption.
- IP Layer/Network Layer:

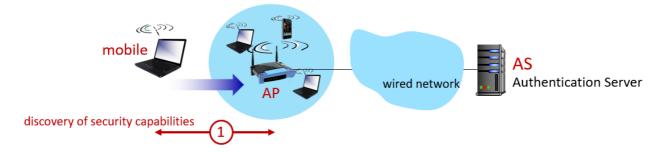
IPsec datagram



• Encrypted data sitting inside another data is sometimes done.

802.11: WIFI Authentication, Encryption

802.11: authentication, encryption



- Easier to listen to something as opposed to wired.
- Need to establish a wireless connection to an access point (AP).
- Can be unreliable.

FIREWALLS: IMPORTANT TOPIC

- Gets name from construction to slow down the spread of fire.
- In our context:
 - Keep things in: ex: company employee attempts to leak secrets to sell.
 - Company firewall looks at all communication going out.
 - Main thing, Keep things out:
 - Have a white-list of IP addresses that are allowed.
 - Doesn't necessarily prevent someone from impersonating someone else.
 - Have a black-list of bad actor IP addresses to deny entry.
 - Packet filter: if anyone sends anything to a webserver not on port 80, for example, that is a red flag.
 - Why are they sending other things? Suspicious.
 - Ex of a packet filter: **Access Control List** (ACL)

Access Control Lists

ACL: table of rules, applied top to bottom to incoming packets: (action, condition) pairs: looks like OpenFlow forwarding (Ch. 4)!

action	source address	dest address	protocol	source port	dest port	flag bit
allow	222.22/16	outside of 222.22/16	TCP	> 1023	80	any
allow	outside of 222.22/16	222.22/16	TCP	80	> 1023	ACK
allow	222.22/16	outside of 222.22/16	UDP	> 1023	53	
allow	outside of 222.22/16	222.22/16	UDP	53	> 1023	
deny	all	all	all	all	all	all

Intrusive Detection Systems

• Deep packet inspection

- o looks at data itself for virus, malware
 - virus' can be in images since it's an executable file.