

A1

For any set $A = \{a_1, a_2, a_3, a_4\}$ of four distinct positive integers with sum $s_A = a_1 + a_2 + a_3 + a_4$, let p_A denote the number of pairs (i, j) with $1 \leq i < j \leq 4$ for which $a_i + a_j$ divides s_A . Among all sets of four distinct positive integers, determine those sets A for which p_A is maximal.

Answer. The sets A for which p_A is maximal are the sets the form $\{d, 5d, 7d, 11d\}$ and $\{d, 11d, 19d, 29d\}$, where d is any positive integer. For all these sets p_A is 4.

Solution. Firstly, we will prove that the maximum value of p_A is at most 4. Without loss of generality, we may assume that $a_1 < a_2 < a_3 < a_4$. We observe that for each pair of indices (i, j) with $1 \leq i < j \leq 4$, the sum $a_i + a_j$ divides s_A if and only if $a_i + a_j$ divides $s_A - (a_i + a_j) = a_k + a_l$, where k and l are the other two indices. Since there are 6 distinct pairs, we have to prove that at least two of them do not satisfy the previous condition. We claim that two such pairs are (a_2, a_4) and (a_3, a_4) . Indeed, note that $a_2 + a_4 > a_1 + a_3$ and $a_3 + a_4 > a_1 + a_2$. Hence $a_2 + a_4$ and $a_3 + a_4$ do not divide s_A . This proves $p_A \leq 4$.

Now suppose $p_A = 4$. By the previous argument we have

$$\begin{array}{lll} a_1 + a_4 \mid a_2 + a_3 & \text{and} & a_2 + a_3 \mid a_1 + a_4, \\ a_1 + a_2 \mid a_3 + a_4 & \text{and} & a_3 + a_4 \nmid a_1 + a_2, \\ a_1 + a_3 \mid a_2 + a_4 & \text{and} & a_2 + a_4 \nmid a_1 + a_3. \end{array}$$

Hence, there exist positive integers m and n with $m > n \geq 2$ such that

$$\begin{cases} a_1 + a_4 = a_2 + a_3 \\ m(a_1 + a_2) = a_3 + a_4 \\ n(a_1 + a_3) = a_2 + a_4. \end{cases}$$

Adding up the first equation and the third one, we get $n(a_1 + a_3) = 2a_2 + a_3 - a_1$. If $n \geq 3$, then $n(a_1 + a_3) > 3a_3 > 2a_2 + a_3 > 2a_2 + a_3 - a_1$. This is a contradiction. Therefore $n = 2$. If we multiply by 2 the sum of the first equation and the third one, we obtain

$$6a_1 + 2a_3 = 4a_2,$$

while the sum of the first one and the second one is

$$(m+1)a_1 + (m-1)a_2 = 2a_3.$$

Adding up the last two equations we get

$$(m+7)a_1 = (5-m)a_2.$$

It follows that $5 - m \geq 1$, because the left-hand side of the last equation and a_2 are positive. Since we have $m > n = 2$, the integer m can be equal only to either 3 or 4. Substituting $(3, 2)$ and $(4, 2)$ for (m, n) and solving the previous system of equations, we find the families of solutions $\{d, 5d, 7d, 11d\}$ and $\{d, 11d, 19d, 29d\}$, where d is any positive integer.

A2

Determine all sequences $(x_1, x_2, \dots, x_{2011})$ of positive integers such that for every positive integer n there is an integer a with

$$x_1^n + 2x_2^n + \dots + 2011x_{2011}^n = a^{n+1} + 1.$$

Answer. The only sequence that satisfies the condition is

$$(x_1, \dots, x_{2011}) = (1, k, \dots, k) \quad \text{with } k = 2 + 3 + \dots + 2011 = 2023065.$$

Solution. Throughout this solution, the set of positive integers will be denoted by \mathbb{Z}_+ .

Put $k = 2 + 3 + \dots + 2011 = 2023065$. We have

$$1^n + 2k^n + \dots + 2011k^n = 1 + k \cdot k^n = k^{n+1} + 1$$

for all n , so $(1, k, \dots, k)$ is a valid sequence. We shall prove that it is the only one.

Let a valid sequence (x_1, \dots, x_{2011}) be given. For each $n \in \mathbb{Z}_+$ we have some $y_n \in \mathbb{Z}_+$ with

$$x_1^n + 2x_2^n + \dots + 2011x_{2011}^n = y_n^{n+1} + 1.$$

Note that $x_1^n + 2x_2^n + \dots + 2011x_{2011}^n < (x_1 + 2x_2 + \dots + 2011x_{2011})^{n+1}$, which implies that the sequence (y_n) is bounded. In particular, there is some $y \in \mathbb{Z}_+$ with $y_n = y$ for infinitely many n .

Let m be the maximum of all the x_i . Grouping terms with equal x_i together, the sum $x_1^n + 2x_2^n + \dots + 2011x_{2011}^n$ can be written as

$$x_1^n + 2x_2^n + \dots + x_{2011}^n = a_m m^n + a_{m-1} (m-1)^n + \dots + a_1$$

with $a_i \geq 0$ for all i and $a_1 + \dots + a_m = 1 + 2 + \dots + 2011$. So there exist arbitrarily large values of n , for which

$$a_m m^n + \dots + a_1 - 1 - y \cdot y^n = 0. \tag{1}$$

The following lemma will help us to determine the a_i and y :

Lemma. Let integers b_1, \dots, b_N be given and assume that there are arbitrarily large positive integers n with $b_1 + b_2 2^n + \dots + b_N N^n = 0$. Then $b_i = 0$ for all i .

Proof. Suppose that not all b_i are zero. We may assume without loss of generality that $b_N \neq 0$.

Dividing through by N^n gives

$$|b_N| = \left| b_{N-1} \left(\frac{N-1}{N} \right)^n + \cdots + b_1 \left(\frac{1}{N} \right)^n \right| \leq (|b_{N-1}| + \cdots + |b_1|) \left(\frac{N-1}{N} \right)^n.$$

The expression $\left(\frac{N-1}{N} \right)^n$ can be made arbitrarily small for n large enough, contradicting the assumption that b_N be non-zero. \square

We obviously have $y > 1$. Applying the lemma to (1) we see that $a_m = y = m$, $a_1 = 1$, and all the other a_i are zero. This implies $(x_1, \dots, x_{2011}) = (1, m, \dots, m)$. But we also have $1 + m = a_1 + \cdots + a_m = 1 + \cdots + 2011 = 1 + k$ so $m = k$, which is what we wanted to show.

A3

Determine all pairs (f, g) of functions from the set of real numbers to itself that satisfy

$$g(f(x+y)) = f(x) + (2x+y)g(y)$$

for all real numbers x and y .

Answer. Either both f and g vanish identically, or there exists a real number C such that $f(x) = x^2 + C$ and $g(x) = x$ for all real numbers x .

Solution. Clearly all these pairs of functions satisfy the functional equation in question, so it suffices to verify that there cannot be any further ones. Substituting $-2x$ for y in the given functional equation we obtain

$$g(f(-x)) = f(x). \quad (1)$$

Using this equation for $-x-y$ in place of x we obtain

$$f(-x-y) = g(f(x+y)) = f(x) + (2x+y)g(y). \quad (2)$$

Now for any two real numbers a and b , setting $x = -b$ and $y = a+b$ we get

$$f(-a) = f(-b) + (a-b)g(a+b).$$

If c denotes another arbitrary real number we have similarly

$$f(-b) = f(-c) + (b-c)g(b+c)$$

as well as

$$f(-c) = f(-a) + (c-a)g(c+a).$$

Adding all these equations up, we obtain

$$((a+c) - (b+c))g(a+b) + ((a+b) - (a+c))g(b+c) + ((b+c) - (a+b))g(a+c) = 0.$$

Now given any three real numbers x , y , and z one may determine three reals a , b , and c such that $x = b+c$, $y = c+a$, and $z = a+b$, so that we get

$$(y-x)g(z) + (z-y)g(x) + (x-z)g(y) = 0.$$

This implies that the three points $(x, g(x))$, $(y, g(y))$, and $(z, g(z))$ from the graph of g are collinear. Hence that graph is a line, i.e., g is either a constant or a linear function.

Let us write $g(x) = Ax + B$, where A and B are two real numbers. Substituting $(0, -y)$ for (x, y) in (2) and denoting $C = f(0)$, we have $f(y) = Ay^2 - By + C$. Now, comparing the coefficients of x^2 in (1) we see that $A^2 = A$, so $A = 0$ or $A = 1$.

If $A = 0$, then (1) becomes $B = -Bx + C$ and thus $B = C = 0$, which provides the first of the two solutions mentioned above.

Now suppose $A = 1$. Then (1) becomes $x^2 - Bx + C + B = x^2 - Bx + C$, so $B = 0$. Thus, $g(x) = x$ and $f(x) = x^2 + C$, which is the second solution from above.

Comment. Another way to show that $g(x)$ is either a constant or a linear function is the following. If we interchange x and y in the given functional equation and subtract this new equation from the given one, we obtain

$$f(x) - f(y) = (2y + x)g(x) - (2x + y)g(y).$$

Substituting $(x, 0)$, $(1, x)$, and $(0, 1)$ for (x, y) , we get

$$\begin{aligned} f(x) - f(0) &= xg(x) - 2xg(0), \\ f(1) - f(x) &= (2x + 1)g(1) - (x + 2)g(x), \\ f(0) - f(1) &= 2g(0) - g(1). \end{aligned}$$

Taking the sum of these three equations and dividing by 2, we obtain

$$g(x) = x(g(1) - g(0)) + g(0).$$

This proves that $g(x)$ is either a constant or a linear function.

A4

Determine all pairs (f, g) of functions from the set of positive integers to itself that satisfy

$$f^{g(n)+1}(n) + g^{f(n)}(n) = f(n+1) - g(n+1) + 1$$

for every positive integer n . Here, $f^k(n)$ means $\underbrace{f(f(\dots f(n)\dots))}_k$.

Answer. The only pair (f, g) of functions that satisfies the equation is given by $f(n) = n$ and $g(n) = 1$ for all n .

Solution. The given relation implies

$$f(f^{g(n)}(n)) < f(n+1) \quad \text{for all } n, \quad (1)$$

which will turn out to be sufficient to determine f .

Let $y_1 < y_2 < \dots$ be all the values attained by f (this sequence might be either finite or infinite). We will prove that for every positive n the function f attains at least n values, and we have (i)_n: $f(x) = y_n$ if and only if $x = n$, and (ii)_n: $y_n = n$. The proof will follow the scheme

$$(i)_1, (ii)_1, (i)_2, (ii)_2, \dots, (i)_n, (ii)_n, \dots \quad (2)$$

To start, consider any x such that $f(x) = y_1$. If $x > 1$, then (1) reads $f(f^{g(x-1)}(x-1)) < y_1$, contradicting the minimality of y_1 . So we have that $f(x) = y_1$ is equivalent to $x = 1$, establishing (i)₁.

Next, assume that for some n statement (i)_n is established, as well as all the previous statements in (2). Note that these statements imply that for all $k \geq 1$ and $a < n$ we have $f^k(x) = a$ if and only if $x = a$.

Now, each value y_i with $1 \leq i \leq n$ is attained at the unique integer i , so y_{n+1} exists. Choose an arbitrary x such that $f(x) = y_{n+1}$; we necessarily have $x > n$. Substituting $x-1$ into (1) we have $f(f^{g(x-1)}(x-1)) < y_{n+1}$, which implies

$$f^{g(x-1)}(x-1) \in \{1, \dots, n\} \quad (3)$$

Set $b = f^{g(x-1)}(x-1)$. If $b < n$ then we would have $x-1 = b$ which contradicts $x > n$. So $b = n$, and hence $y_n = n$, which proves (ii)_n. Next, from (i)_n we now get $f(k) = n \iff k = n$, so removing all the iterations of f in (3) we obtain $x-1 = b = n$, which proves (i)_{n+1}.

So, all the statements in (2) are valid and hence $f(n) = n$ for all n . The given relation between f and g now reads $n + g^n(n) = n + 1 - g(n+1) + 1$ or $g^n(n) + g(n+1) = 2$, from which it

immediately follows that we have $g(n) = 1$ for all n .

Comment. Several variations of the above solution are possible. For instance, one may first prove by induction that the smallest n values of f are exactly $f(1) < \dots < f(n)$ and proceed as follows. We certainly have $f(n) \geq n$ for all n . If there is an n with $f(n) > n$, then $f(x) > x$ for all $x \geq n$. From this we conclude $f^{g(n)+1}(n) > f^{g(n)}(n) > \dots > f(n)$. But we also have $f^{g(n)+1} < f(n+1)$. Having squeezed in a function value between $f(n)$ and $f(n+1)$, we arrive at a contradiction.

In any case, the inequality (1) plays an essential rôle.

A5

Prove that for every positive integer n , the set $\{2, 3, 4, \dots, 3n + 1\}$ can be partitioned into n triples in such a way that the numbers from each triple are the lengths of the sides of some obtuse triangle.

Solution. Throughout the solution, we denote by $[a, b]$ the set $\{a, a + 1, \dots, b\}$. We say that $\{a, b, c\}$ is an *obtuse triple* if a, b, c are the sides of some obtuse triangle.

We prove by induction on n that there exists a partition of $[2, 3n + 1]$ into n obtuse triples A_i ($2 \leq i \leq n + 1$) having the form $A_i = \{i, a_i, b_i\}$. For the base case $n = 1$, one can simply set $A_2 = \{2, 3, 4\}$. For the induction step, we need the following simple lemma.

Lemma. Suppose that the numbers $a < b < c$ form an obtuse triple, and let x be any positive number. Then the triple $\{a, b + x, c + x\}$ is also obtuse.

Proof. The numbers $a < b + x < c + x$ are the sides of a triangle because $(c + x) - (b + x) = c - b < a$. This triangle is obtuse since $(c + x)^2 - (b + x)^2 = (c - b)(c + b + 2x) > (c - b)(c + b) > a^2$. \square

Now we turn to the induction step. Let $n > 1$ and put $t = \lfloor n/2 \rfloor < n$. By the induction hypothesis, there exists a partition of the set $[2, 3t + 1]$ into t obtuse triples $A'_i = \{i, a'_i, b'_i\}$ ($i \in [2, t + 1]$). For the same values of i , define $A_i = \{i, a'_i + (n - t), b'_i + (n - t)\}$. The constructed triples are obviously disjoint, and they are obtuse by the lemma. Moreover, we have

$$\bigcup_{i=2}^{t+1} A_i = [2, t + 1] \cup [n + 2, n + 2t + 1].$$

Next, for each $i \in [t + 2, n + 1]$, define $A_i = \{i, n + t + i, 2n + i\}$. All these sets are disjoint, and

$$\bigcup_{i=t+2}^{n+1} A_i = [t + 2, n + 1] \cup [n + 2t + 2, 2n + t + 1] \cup [2n + t + 2, 3n + 1],$$

so

$$\bigcup_{i=2}^{n+1} A_i = [2, 3n + 1].$$

Thus, we are left to prove that the triple A_i is obtuse for each $i \in [t + 2, n + 1]$.

Since $(2n + i) - (n + t + i) = n - t < t + 2 \leq i$, the elements of A_i are the sides of a triangle. Next, we have

$$(2n + i)^2 - (n + t + i)^2 = (n - t)(3n + t + 2i) \geq \frac{n}{2} \cdot (3n + 3(t + 1) + 1) > \frac{n}{2} \cdot \frac{9n}{2} \geq (n + 1)^2 \geq i^2,$$

so this triangle is obtuse. The proof is completed.

A6

Let f be a function from the set of real numbers to itself that satisfies

$$f(x + y) \leq yf(x) + f(f(x)) \quad (1)$$

for all real numbers x and y . Prove that $f(x) = 0$ for all $x \leq 0$.

Solution 1. Substituting $y = t - x$, we rewrite (1) as

$$f(t) \leq tf(x) - xf(x) + f(f(x)). \quad (2)$$

Consider now some real numbers a, b and use (2) with $t = f(a)$, $x = b$ as well as with $t = f(b)$, $x = a$. We get

$$\begin{aligned} f(f(a)) - f(f(b)) &\leq f(a)f(b) - bf(b), \\ f(f(b)) - f(f(a)) &\leq f(a)f(b) - af(a). \end{aligned}$$

Adding these two inequalities yields

$$2f(a)f(b) \geq af(a) + bf(b).$$

Now, substitute $b = 2f(a)$ to obtain $2f(a)f(b) \geq af(a) + 2f(a)f(b)$, or $af(a) \leq 0$. So, we get

$$f(a) \geq 0 \quad \text{for all } a < 0. \quad (3)$$

Now suppose $f(x) > 0$ for some real number x . From (2) we immediately get that for every $t < \frac{xf(x) - f(f(x))}{f(x)}$ we have $f(t) < 0$. This contradicts (3); therefore

$$f(x) \leq 0 \quad \text{for all real } x, \quad (4)$$

and by (3) again we get $f(x) = 0$ for all $x < 0$.

We are left to find $f(0)$. Setting $t = x < 0$ in (2) we get

$$0 \leq 0 - 0 + f(0),$$

so $f(0) \geq 0$. Combining this with (4) we obtain $f(0) = 0$.

Solution 2. We will also use the condition of the problem in form (2). For clarity we divide the argument into four steps.

Step 1. We begin by proving that f attains nonpositive values only. Assume that there exist some real number z with $f(z) > 0$. Substituting $x = z$ into (2) and setting $A = f(z)$, $B = -zf(z) - f(f(z))$ we get $f(t) \leq At + B$ for all real t . Hence, if for any positive real number t we substitute $x = -t$, $y = t$ into (1), we get

$$\begin{aligned} f(0) &\leq tf(-t) + f(f(-t)) \leq t(-At + B) + Af(-t) + B \\ &\leq -t(At - B) + A(-At + B) + B = -At^2 - (A^2 - B)t + (A + 1)B. \end{aligned}$$

But surely this cannot be true if we take t to be large enough. This contradiction proves that we have indeed $f(x) \leq 0$ for all real numbers x . Note that for this reason (1) entails

$$f(x + y) \leq yf(x) \tag{5}$$

for all real numbers x and y .

Step 2. We proceed by proving that f has at least one zero. If $f(0) = 0$, we are done. Otherwise, in view of Step 1 we get $f(0) < 0$. Observe that (5) tells us now $f(y) \leq yf(0)$ for all real numbers y . Thus we can specify a positive real number a that is so large that $f(a)^2 > -f(0)$. Put $b = f(a)$ and substitute $x = b$ and $y = -b$ into (5); we learn $-b^2 < f(0) \leq -bf(b)$, i.e. $b < f(b)$. Now we apply (2) to $x = b$ and $t = f(b)$, which yields

$$f(f(b)) \leq (f(b) - b)f(b) + f(f(b)),$$

i.e. $f(b) \geq 0$. So in view of Step 1, b is a zero of f .

Step 3. Next we show that if $f(a) = 0$ and $b < a$, then $f(b) = 0$ as well. To see this, we just substitute $x = b$ and $y = a - b$ into (5), thus getting $f(b) \geq 0$, which suffices by Step 1.

Step 4. By Step 3, the solution of the problem is reduced to showing $f(0) = 0$. Pick any zero r of f and substitute $x = r$ and $y = -1$ into (1). Because of $f(r) = f(r - 1) = 0$ this gives $f(0) \geq 0$ and hence $f(0) = 0$ by Step 1 again.

Comment 1. Both of these solutions also show $f(x) \leq 0$ for all real numbers x . As one can see from Solution 1, this task gets much easier if one already knows that f takes nonnegative values for sufficiently small arguments. Another way of arriving at this statement, suggested by the proposer, is as follows:

Put $a = f(0)$ and substitute $x = 0$ into (1). This gives $f(y) \leq ay + f(a)$ for all real numbers y . Thus if for any real number x we plug $y = a - x$ into (1), we obtain

$$f(a) \leq (a - x)f(x) + f(f(x)) \leq (a - x)f(x) + af(x) + f(a)$$

and hence $0 \leq (2a - x)f(x)$. In particular, if $x < 2a$, then $f(x) \geq 0$.

Having reached this point, one may proceed almost exactly as in the first solution to deduce $f(x) \leq 0$ for all x . Afterwards the problem can be solved in a few lines as shown in steps 3 and 4 of the second

solution.

Comment 2. The original problem also contained the question whether a nonzero function satisfying the problem condition exists. Here we present a family of such functions.

Notice first that if $g : (0, \infty) \rightarrow [0, \infty)$ denotes any function such that

$$g(x+y) \geq yg(x) \tag{6}$$

for all positive real numbers x and y , then the function f given by

$$f(x) = \begin{cases} -g(x) & \text{if } x > 0 \\ 0 & \text{if } x \leq 0 \end{cases} \tag{7}$$

automatically satisfies (1). Indeed, we have $f(x) \leq 0$ and hence also $f(f(x)) = 0$ for all real numbers x . So (1) reduces to (5); moreover, this inequality is nontrivial only if x and y are positive. In this last case it is provided by (6).

Now it is not hard to come up with a nonzero function g obeying (6). E.g. $g(z) = Ce^z$ (where C is a positive constant) fits since the inequality $e^y > y$ holds for all (positive) real numbers y . One may also consider the function $g(z) = e^z - 1$; in this case, we even have that f is continuous.

A7

Let a , b , and c be positive real numbers satisfying $\min(a+b, b+c, c+a) > \sqrt{2}$ and $a^2 + b^2 + c^2 = 3$. Prove that

$$\frac{a}{(b+c-a)^2} + \frac{b}{(c+a-b)^2} + \frac{c}{(a+b-c)^2} \geq \frac{3}{(abc)^2}. \quad (1)$$

Throughout both solutions, we denote the sums of the form $f(a, b, c) + f(b, c, a) + f(c, a, b)$ by $\sum f(a, b, c)$.

Solution 1. The condition $b+c > \sqrt{2}$ implies $b^2 + c^2 > 1$, so $a^2 = 3 - (b^2 + c^2) < 2$, i.e. $a < \sqrt{2} < b+c$. Hence we have $b+c-a > 0$, and also $c+a-b > 0$ and $a+b-c > 0$ for similar reasons.

We will use the variant of HÖLDER's inequality

$$\frac{x_1^{p+1}}{y_1^p} + \frac{x_2^{p+1}}{y_2^p} + \dots + \frac{x_n^{p+1}}{y_n^p} \geq \frac{(x_1 + x_2 + \dots + x_n)^{p+1}}{(y_1 + y_2 + \dots + y_n)^p},$$

which holds for all positive real numbers $p, x_1, x_2, \dots, x_n, y_1, y_2, \dots, y_n$. Applying it to the left-hand side of (1) with $p = 2$ and $n = 3$, we get

$$\sum \frac{a}{(b+c-a)^2} = \sum \frac{(a^2)^3}{a^5(b+c-a)^2} \geq \frac{(a^2 + b^2 + c^2)^3}{(\sum a^{5/2}(b+c-a))^2} = \frac{27}{(\sum a^{5/2}(b+c-a))^2}. \quad (2)$$

To estimate the denominator of the right-hand part, we use an instance of SCHUR's inequality, namely

$$\sum a^{3/2}(a-b)(a-c) \geq 0,$$

which can be rewritten as

$$\sum a^{5/2}(b+c-a) \leq abc(\sqrt{a} + \sqrt{b} + \sqrt{c}).$$

Moreover, by the inequality between the arithmetic mean and the fourth power mean we also have

$$\left(\frac{\sqrt{a} + \sqrt{b} + \sqrt{c}}{3} \right)^4 \leq \frac{a^2 + b^2 + c^2}{3} = 1,$$

i.e., $\sqrt{a} + \sqrt{b} + \sqrt{c} \leq 3$. Hence, (2) yields

$$\sum \frac{a}{(b+c-a)^2} \geq \frac{27}{(abc(\sqrt{a} + \sqrt{b} + \sqrt{c}))^2} \geq \frac{3}{a^2 b^2 c^2},$$

thus solving the problem.

Comment. In this solution, one may also start from the following version of HÖLDER's inequality

$$\left(\sum_{i=1}^n a_i^3\right) \left(\sum_{i=1}^n b_i^3\right) \left(\sum_{i=1}^n c_i^3\right) \geq \left(\sum_{i=1}^n a_i b_i c_i\right)^3$$

applied as

$$\sum \frac{a}{(b+c-a)^2} \cdot \sum a^3(b+c-a) \cdot \sum a^2(b+c-a) \geq 27.$$

After doing that, one only needs the slightly better known instances

$$\sum a^3(b+c-a) \leq (a+b+c)abc \quad \text{and} \quad \sum a^2(b+c-a) \leq 3abc$$

of SCHUR's Inequality.

Solution 2. As in Solution 1, we mention that all the numbers $b+c-a$, $a+c-b$, $a+b-c$ are positive. We will use only this restriction and the condition

$$a^5 + b^5 + c^5 \geq 3, \tag{3}$$

which is weaker than the given one. Due to the symmetry we may assume that $a \geq b \geq c$.

In view of (3), it suffices to prove the inequality

$$\sum \frac{a^3 b^2 c^2}{(b+c-a)^2} \geq \sum a^5,$$

or, moving all the terms into the left-hand part,

$$\sum \frac{a^3}{(b+c-a)^2} ((bc)^2 - (a(b+c-a))^2) \geq 0. \tag{4}$$

Note that the signs of the expressions $(yz)^2 - (x(y+z-x))^2$ and $yz - x(y+z-x) = (x-y)(x-z)$ are the same for every positive x, y, z satisfying the triangle inequality. So the terms in (4) corresponding to a and c are nonnegative, and hence it is sufficient to prove that the sum of the terms corresponding to a and b is nonnegative. Equivalently, we need the relation

$$\frac{a^3}{(b+c-a)^2} (a-b)(a-c)(bc + a(b+c-a)) \geq \frac{b^3}{(a+c-b)^2} (a-b)(b-c)(ac + b(a+c-b)).$$

Obviously, we have

$$a^3 \geq b^3 \geq 0, \quad 0 < b+c-a \leq a+c-b, \quad \text{and} \quad a-c \geq b-c \geq 0,$$

hence it suffices to prove that

$$\frac{ab + ac + bc - a^2}{b+c-a} \geq \frac{ab + ac + bc - b^2}{c+a-b}.$$

Since all the denominators are positive, it is equivalent to

$$(c + a - b)(ab + ac + bc - a^2) - (ab + ac + bc - b^2)(b + c - a) \geq 0,$$

or

$$(a - b)(2ab - a^2 - b^2 + ac + bc) \geq 0.$$

Since $a \geq b$, the last inequality follows from

$$c(a + b) > (a - b)^2$$

which holds since $c > a - b \geq 0$ and $a + b > a - b \geq 0$.

Claim 3. $f(n) \neq f(1)$ if and only if $a \mid n$.

Proof. Since $f(1) = \dots = f(a-1) < f(a)$, the claim follows from the fact that

$$f(n) = f(1) \iff f(n+a) = f(1).$$

So it suffices to prove this fact.

Assume that $f(n) = f(1)$. Then $f(n+a) \mid f(a) - f(-n) = f(a) - f(n) > 0$, so $f(n+a) \leq f(a) - f(n) < f(a)$; in particular the difference $f(n+a) - f(n)$ is strictly smaller than $f(a)$. Furthermore, this difference is divisible by $f(a)$ and nonnegative since $f(n) = f(1)$ is the least value attained by f . So we have $f(n+a) - f(n) = 0$, as desired. For the converse direction we only need to remark that $f(n+a) = f(1)$ entails $f(-n-a) = f(1)$, and hence $f(n) = f(-n) = f(1)$ by the forward implication. \square

We return to the induction step. So let us take two arbitrary integers m and n with $f(m) \leq f(n)$. If $a \nmid m$, then we have $f(m) = f(1) \mid f(n)$. On the other hand, suppose that $a \mid m$; then by Claim 3 $a \mid n$ as well. Now define the function $g(x) = f(ax)$. Clearly, g satisfies the conditions of the problem, but $N_g < N_f - 1$, since g does not attain $f(1)$. Hence, by the induction hypothesis, $f(m) = g(m/a) \mid g(n/a) = f(n)$, as desired.

Comment. After the fact that f attains a finite number of values has been established, there are several ways of finishing the solution. For instance, let $f(0) = b_1 > b_2 > \dots > b_k$ be all these values. One may show (essentially in the same way as in Claim 3) that the set $S_i = \{n : f(n) \geq b_i\}$ consists exactly of all numbers divisible by some integer $a_i \geq 0$. One obviously has $a_i \mid a_{i-1}$, which implies $f(a_i) \mid f(a_{i-1})$ by Claim 1. So, $b_k \mid b_{k-1} \mid \dots \mid b_1$, thus proving the problem statement.

Moreover, now it is easy to describe all functions satisfying the conditions of the problem. Namely, all these functions can be constructed as follows. Consider a sequence of nonnegative integers a_1, a_2, \dots, a_k and another sequence of positive integers b_1, b_2, \dots, b_k such that $|a_k| = 1$, $a_i \neq a_j$ and $b_i \neq b_j$ for all $1 \leq i < j \leq k$, and $a_i \mid a_{i-1}$ and $b_i \mid b_{i-1}$ for all $i = 2, \dots, k$. Then one may introduce the function

$$f(n) = b_{i(n)}, \quad \text{where } i(n) = \min\{i : a_i \mid n\}.$$

These are all the functions which satisfy the conditions of the problem.

N6

Let $P(x)$ and $Q(x)$ be two polynomials with integer coefficients such that no nonconstant polynomial with rational coefficients divides both $P(x)$ and $Q(x)$. Suppose that for every positive integer n the integers $P(n)$ and $Q(n)$ are positive, and $2^{Q(n)} - 1$ divides $3^{P(n)} - 1$. Prove that $Q(x)$ is a constant polynomial.

Solution. First we show that there exists an integer d such that for all positive integers n we have $\gcd(P(n), Q(n)) \leq d$.

Since $P(x)$ and $Q(x)$ are coprime (over the polynomials with rational coefficients), EUCLID's algorithm provides some polynomials $R_0(x), S_0(x)$ with rational coefficients such that $P(x)R_0(x) - Q(x)S_0(x) = 1$. Multiplying by a suitable positive integer d , we obtain polynomials $R(x) = d \cdot R_0(x)$ and $S(x) = d \cdot S_0(x)$ with integer coefficients for which $P(x)R(x) - Q(x)S(x) = d$. Then we have $\gcd(P(n), Q(n)) \leq d$ for any integer n .

To prove the problem statement, suppose that $Q(x)$ is not constant. Then the sequence $Q(n)$ is not bounded and we can choose a positive integer m for which

$$M = 2^{Q(m)} - 1 \geq 3^{\max\{P(1), P(2), \dots, P(d)\}}. \quad (1)$$

Since $M = 2^{Q(m)} - 1 \mid 3^{P(m)} - 1$, we have $2, 3 \nmid M$. Let a and b be the multiplicative orders of 2 and 3 modulo M , respectively. Obviously, $a = Q(m)$ since the lower powers of 2 do not reach M . Since M divides $3^{P(m)} - 1$, we have $b \mid P(m)$. Then $\gcd(a, b) \leq \gcd(P(m), Q(m)) \leq d$. Since the expression $ax - by$ attains all integer values divisible by $\gcd(a, b)$ when x and y run over all nonnegative integer values, there exist some nonnegative integers x, y such that $1 \leq m + ax - by \leq d$.

By $Q(m + ax) \equiv Q(m) \pmod{a}$ we have

$$2^{Q(m+ax)} \equiv 2^{Q(m)} \equiv 1 \pmod{M}$$

and therefore

$$M \mid 2^{Q(m+ax)} - 1 \mid 3^{P(m+ax)} - 1.$$

Then, by $P(m + ax - by) \equiv P(m + ax) \pmod{b}$ we have

$$3^{P(m+ax-by)} \equiv 3^{P(m+ax)} \equiv 1 \pmod{M}.$$

Since $P(m + ax - by) > 0$ this implies $M \leq 3^{P(m+ax-by)} - 1$. But $P(m + ax - by)$ is listed among $P(1), P(2), \dots, P(d)$, so

$$M < 3^{P(m+ax-by)} \leq 3^{\max\{P(1), P(2), \dots, P(d)\}}$$

which contradicts (1).

Comment. We present another variant of the solution above.

Denote the degree of P by k and its leading coefficient by p . Consider any positive integer n and let $a = Q(n)$. Again, denote by b the multiplicative order of 3 modulo $2^a - 1$. Since $2^a - 1 \mid 3^{P(n)} - 1$, we have $b \mid P(n)$. Moreover, since $2^{Q(n+at)} - 1 \mid 3^{P(n+at)} - 1$ and $a = Q(n) \mid Q(n+at)$ for each positive integer t , we have $2^a - 1 \mid 3^{P(n+at)} - 1$, hence $b \mid P(n+at)$ as well.

Therefore, b divides $\gcd\{P(n+at) : t \geq 0\}$; hence it also divides the number

$$\sum_{i=0}^k (-1)^{k-i} \binom{k}{i} P(n+ai) = p \cdot k! \cdot a^k.$$

Finally, we get $b \mid \gcd(P(n), k! \cdot p \cdot Q(n)^k)$, which is bounded by the same arguments as in the beginning of the solution. So $3^b - 1$ is bounded, and hence $2^{Q(n)} - 1$ is bounded as well.