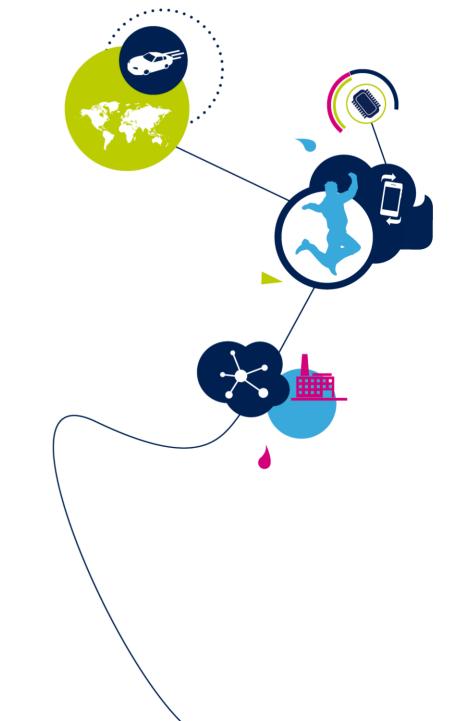
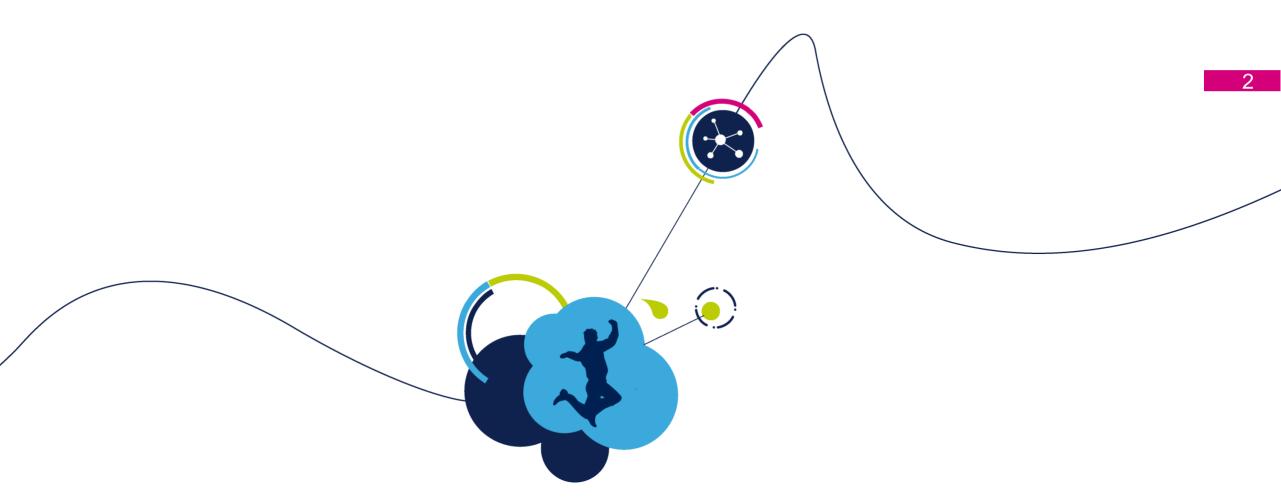
STM32MP1软件架构







STM32MP1 Flash内存映射



Flash partitions (minimal)

Size	Component	Comment
Remaining area	userfs	包括用户数据以及STM32实例应用
768MB	rootfs	Linux根文件系统,包括所有用户空间的程序(可执行文件,软件库,)以及内核模块
16MB	vendorfs	存放第三方专有二进制文件,使其与开源License例如GPL V3分开
64MB	bootfs	启动文件系统包括: - (可选) 初始RAM文件系统。该文件系统可以被复制到外部DDR,在还没有挂装大的根文件系统供linux使用 - Linux kernel device tree (可以做成统一的Flattened Image Tree - FIT) - Linux kernel U-Boot image (可以做成统一的Flattened Image Tree - FIT) - 除NOR之外的所有Flash: 由U-Boot显示的启动splash screen image - U-Boot distro 配置文件extlinux.conf (可以做成统一的Flattened Image Tree – FIT)
2MB	ssbl	存放第二级启动加载器Second Stage Boot Loader (SSBL)。这里是U-Boot, 包括附在文件尾的device tree blob (dtb)
256kB to 512kB (*)	fsbl	存放第一级启动加载器First Stage Boot Loader 。可以是ARM Trusted Firmware (TF-A) 或者U-Boot Secondary Program Loader (SPL)包括文件尾的device tree blob (dtb)。至少需要两份。注意: 由于ROM代码的RAM要求,FSBL大小限制在247KB。



(*): the partition size depends on the flash technology, to be aligned on block erase size for NOR (256kB) / NAND (512kB)

Flash partitions (可选) 4

Size	Component	Comment
256kB (*)	logo	供NOR flash存放boot loader splash screen image (对于其他的Flash, 该 image存放在bootfs分区)
256kB to 512kB (*)	teeh	OP-TEE header
256kB to 512kB (*)	teed	OP-TEE pageable code and data
256kB to 512kB (*)	teex	OP-TEE pager

(*): the partition size depends on the flash technology, to be aligned on block erase size for NOR (256kB) / NAND (512kB)



SD card memory mapping

ROM代码首 先查找在SD card头部的 GPT 分区表

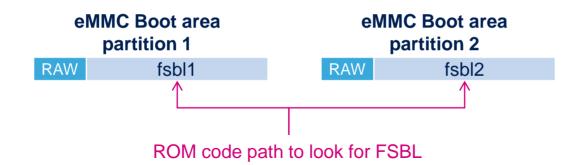
SD card				
Se	condary GPT Header			
EXT4	userfs			
EXT4	rootfs			
EXT4	vendorfs			
EXT4	bootfs			
RAW	teex			
RAW	teed			
RAW	teeh			
RAW	ssbl			
RAW	fsbl2	\leftarrow		
RAW	fsbl1	\leftarrow		
Р	rimary GPT Header			
•	Protective MBR			

- 两个fsbl分区时为了在更新 时可以保证一个有效分区
- 存在ssbl和tee分区是因为 fsbl不支持文件系统
- bootfs是个 EXT4分区。U-Boot使用该分区。
- rootfs和userfs是EXT4文 件系统。Linux使用该分区。

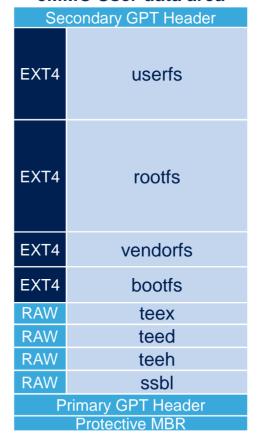


eMMC memory mapping

eMMC与SD card类似。唯一特别的 是它存在两个引导区。所以,我们 讲fsbl1和fsbl2放到这两个引导区。



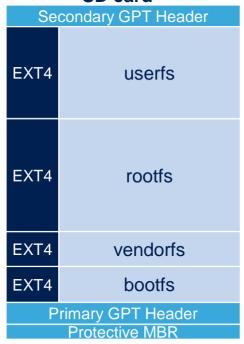
eMMC User data area





NOR memory mapping _____

SD card



注意: 用户基于成本考虑, NOR Flash一般不大, 需要第二级存储空间,例如SD卡。也可以使用 eMMC or NAND.

QSPI NOR

RAW	teex
RAW	teed
RAW	teeh
RAW	logo
RAW	ssbl
RAW	fsbl2
RAW	fsbl1

Offset 256kB Offset 0



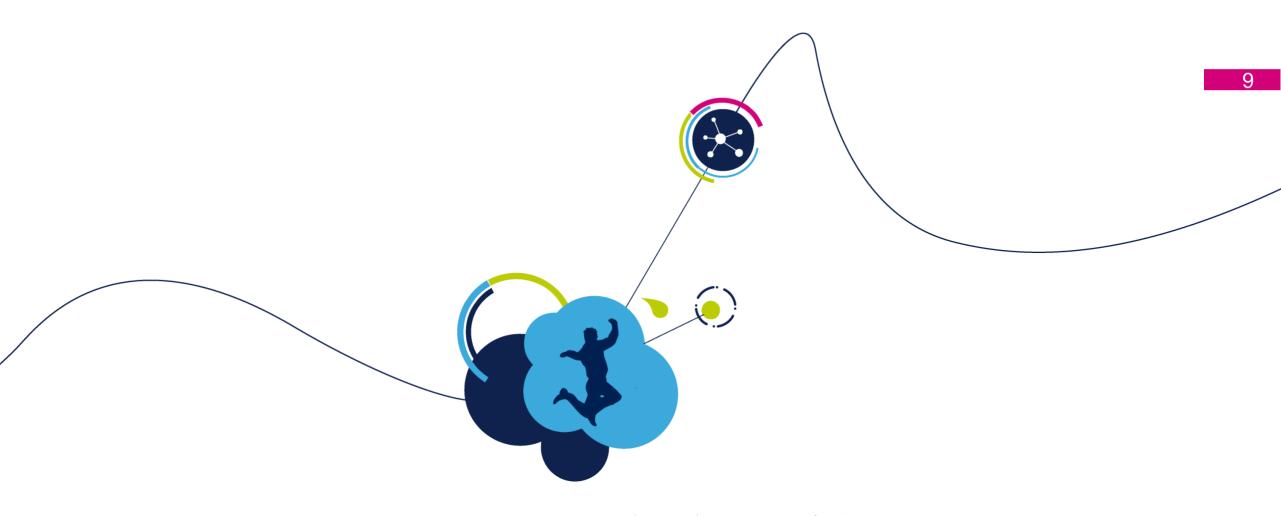
NAND memory mapping 8

NAND

Bad Block Table (BBT)				
		UBIFS	userfs	
MTD	UBI	UBIFS	rootfs	
		UBIFS	vendorfs	
		UBIFS	bootfs	
MTD			teexN	
MTD			teex1	
	MTD		teedN	
	MTD		teed1	
	MTD		teehN	
MTD			teeh1	
MTD			ssblN	
MTD			ssbl1	
Skip Bad Block		Block	fsbIN	
		DIUCK	fsbl1	
			M	

- NAND成本低但需要注意坏 块管理
- NAND由其物理性质本身容 易引起位翻转,故ECC是必 须的

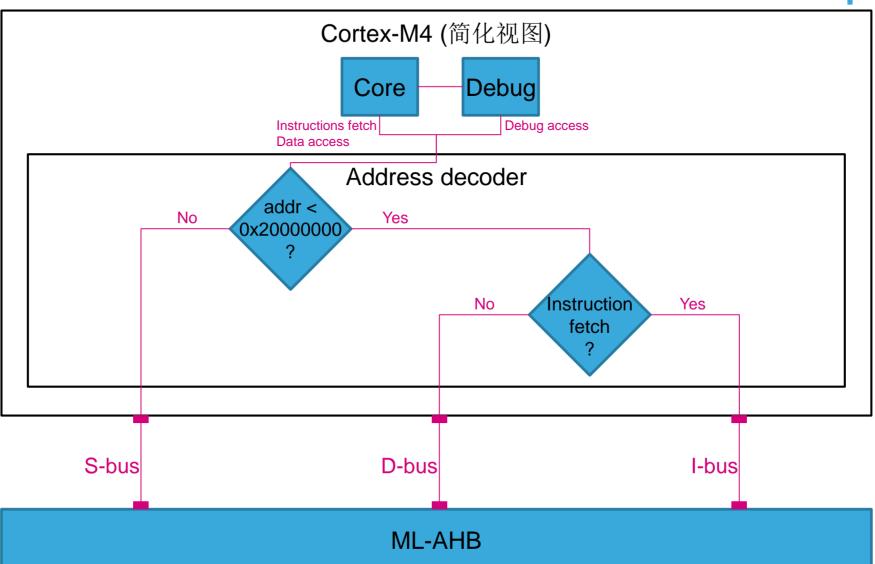




STM32MP1 RAM内存映射



Cortex-M4 ports 10

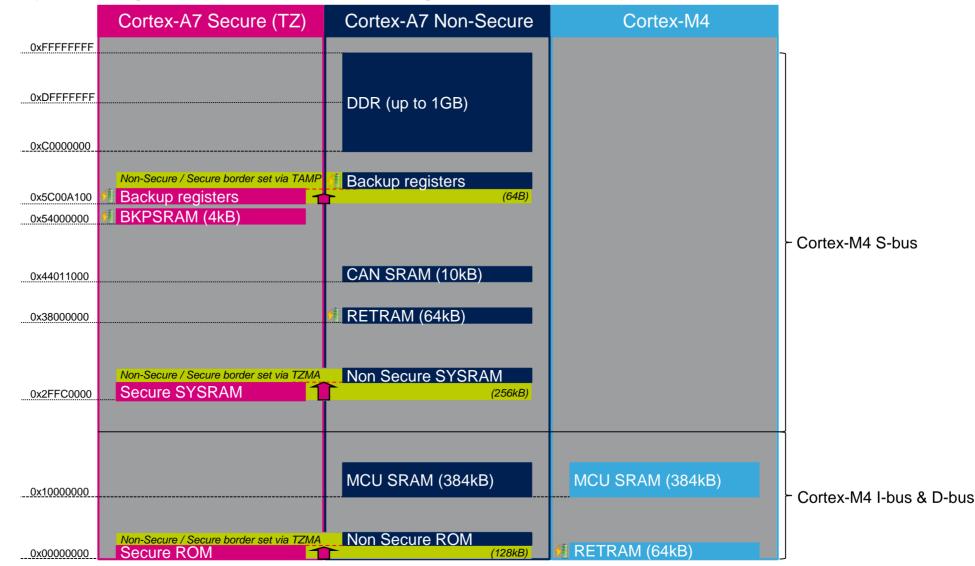




egend

Software memory mapping

• The memory mapping below is a subset of all regions that are really exposed at hardware level.





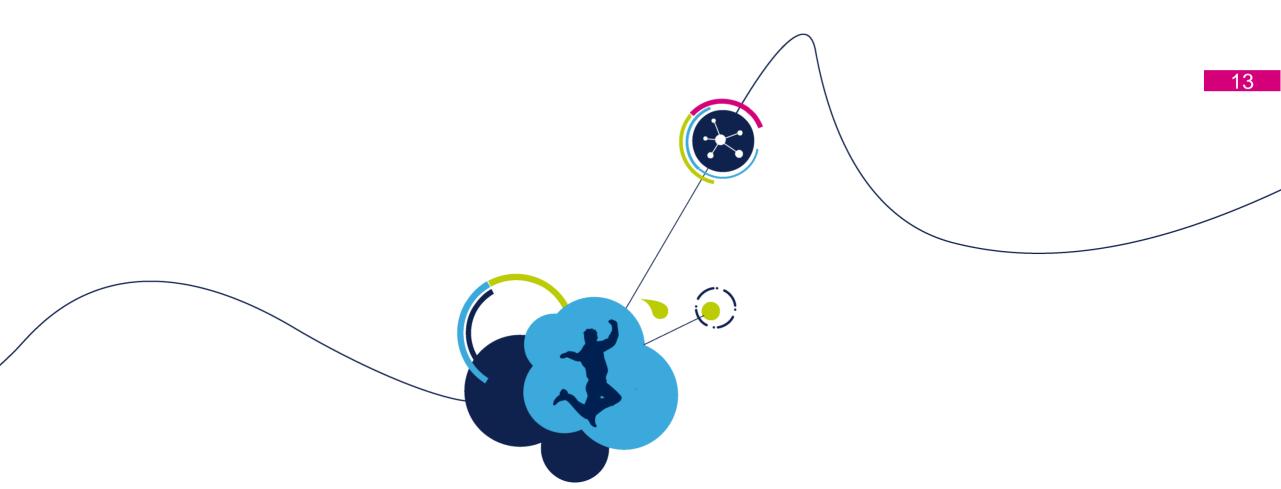
Shared RAM memory mapping 12

Notice that each core may not see the same regions at the same address, as already explained on previous slide

	Cortex-A7 Non-Secure	Cortex-M4
MCU SRAM4 (64kB)	DMA buffers	
MCU SRAM3 (64kB)	Inter Process Commu	unication (IPC) buffers
MCU SRAM2 (128kB)		Data
MCU SRAM1 (128kB)		Code
RETRAM (64kB)		Code & Data Vector table



Each customer can of course tune this mapping (regions location and sizes) to fit with his product needs

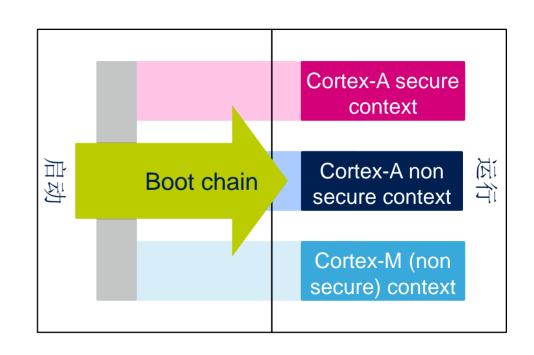


STM32MP1软件架构



STM32MP1多核软件架构

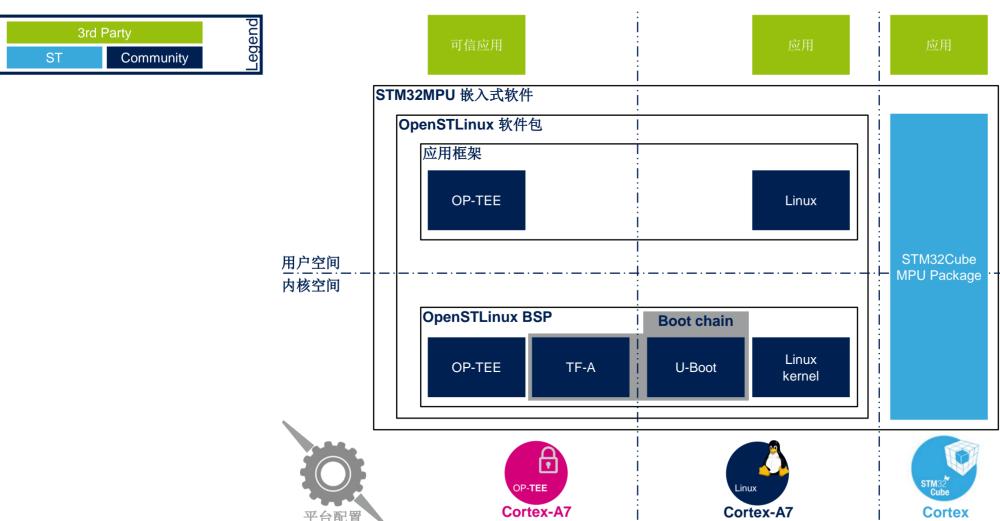
- 系统软件运行情境划分
 - 按照内核
 - 按照安全(同一个内核的两种不同状态)
- 系统中同时存在多个软件运行情境
 - Arm Cortex-A secure (Trustzone) 运行OP-TEE
 - Arm Cortex-A non secure运行Linux
 - Arm Cortex-M (non secure)运行STM32Cube
- 运行时外设分配
 - 指定或者共享





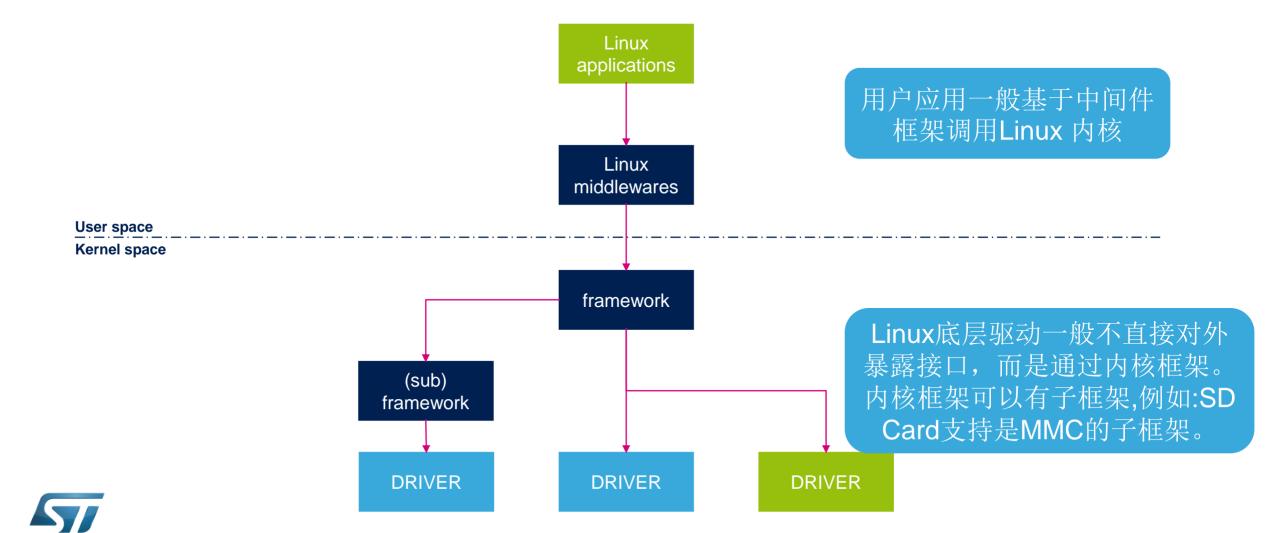
STM32MP1嵌入式软件组成 15

Non-Secure I





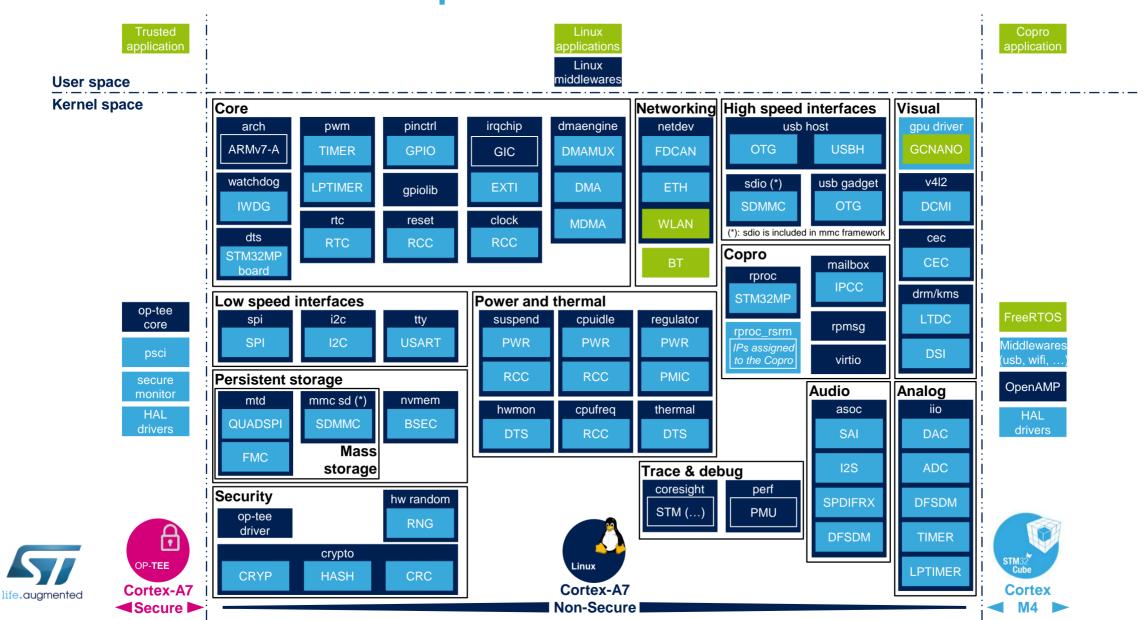
Linux framework & driver 16



egend

- UPPERCASE = 外设驱动

OpenSTLinux + STM32Cube 17



Open-Embedded 用户空间应用 18

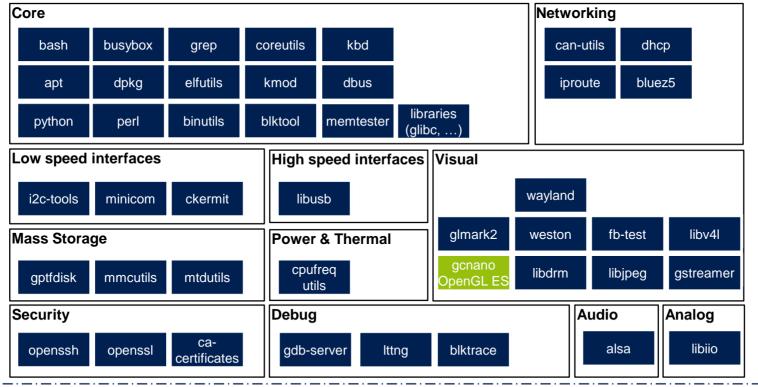
这里仅列出了一部分应用

egend.

• 用户可以根据需要进行定制







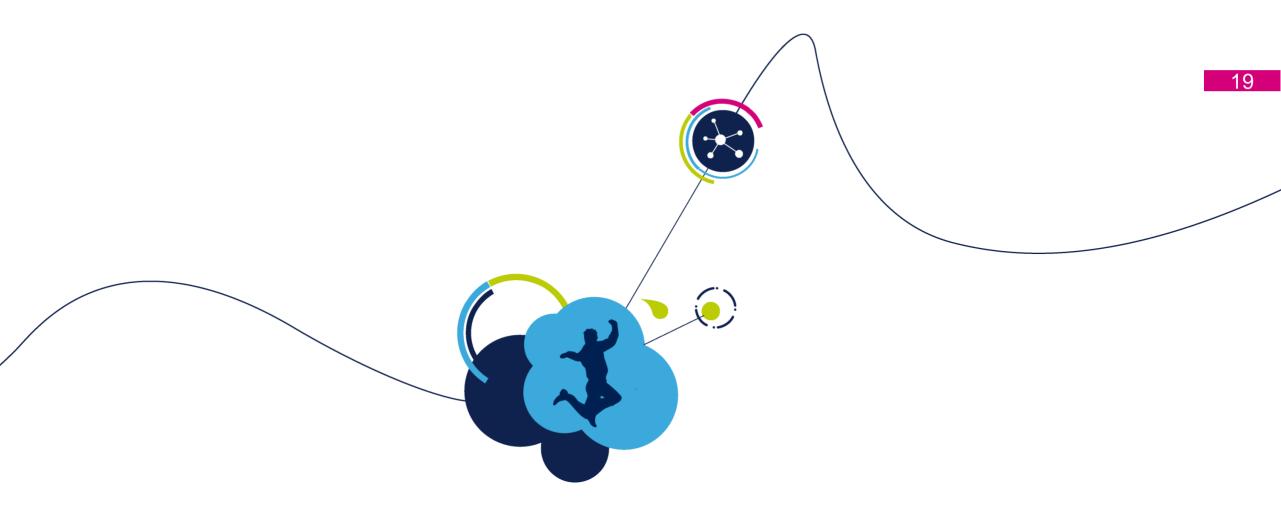


User space

Kernel space



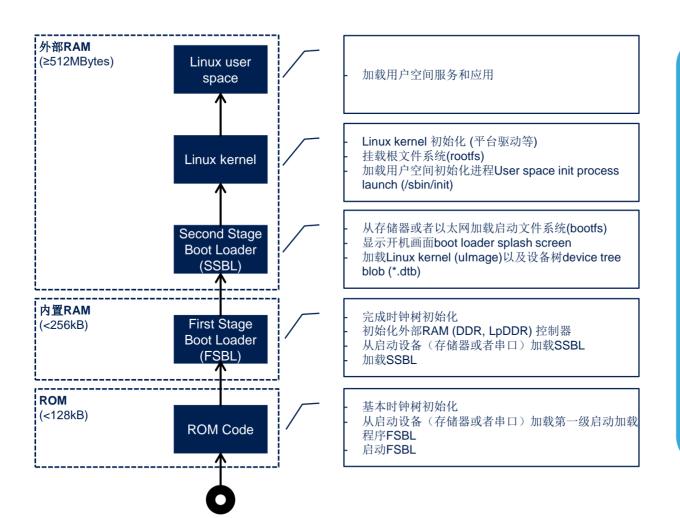




STM32MP1 启动流程



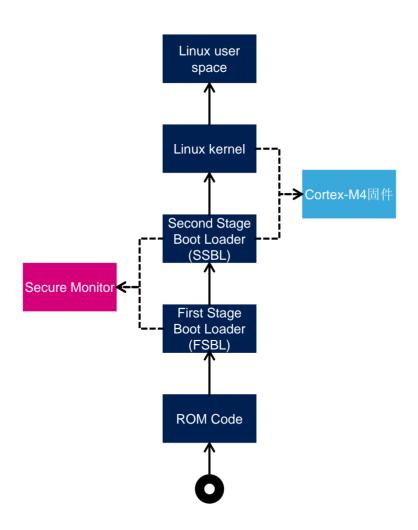
标准Linux启动过程



Linux启动如 同发射火箭, 是一个多级 启动流程。



STM32MP1 启动流程 21



- Secure Monitor可以由 FSBL或者SSBL启动
- Cortex-M4固件可以由 SSBL或者Linux启动



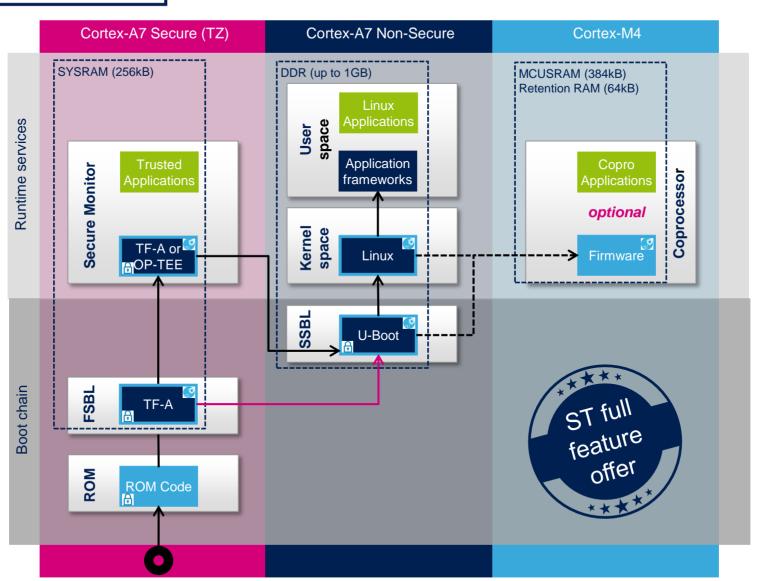
STM32MP1安全启动流程。

TF-A

- BSD许可,用户 可根据需要定制
- ARM为信任执行 而开发,适合安全 敏感场景
- 功能已被ARMv8 平台所验证

U-Boot

• GPL v2许可



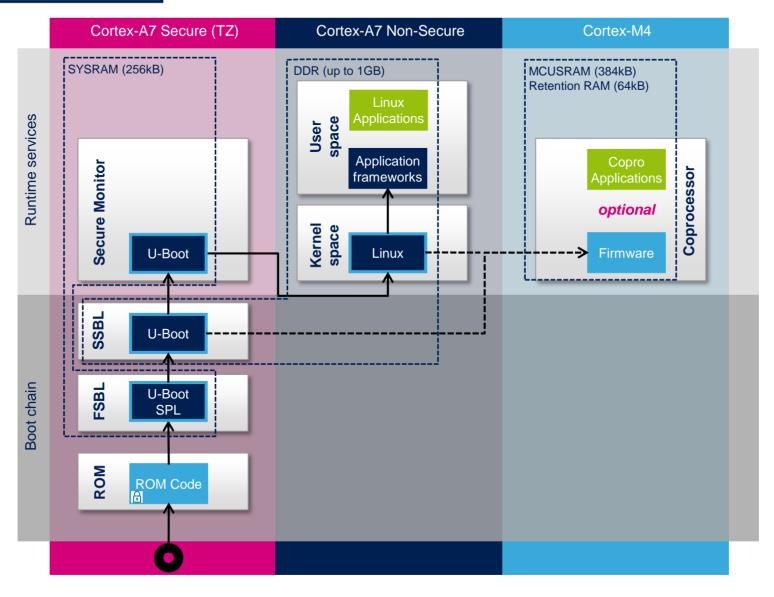
• 安全启动流程是默认方案

• Secure Monitor: 如果没有Secure OS (OP-TEE is optional),则使用TF-A

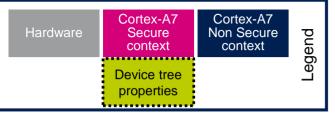


STM32MP1非安全启动流程

U-Boot SPL来自 同一U-Boot源代 码,但是不同的配 置。较小的U-Boot SPL用来加 载较大的U-Boot。







启动过程中的外设初始化 24

剩下的外		剩下的外设	<u>.</u>		(*)			
		Pinctrl	GPIO (boot devices, co	nsole, LEDs)	GPIO (*)			
		Serial console	JART4					
		Power supply	PMIC, I2C4	MIC, I2C4				
外设初始化	<u>'</u>	System time	STGEN	STGEN				
71 24 147 141		Clock tree	PWR, RCC (PLLs setup	0)		PWR, RCC		
		DDR	DDR CTRL / PHY					
		Boot devices	SDMMC, FMC, QSPI, UART, USART, OTG		如左所列举 + ETH, USBH			
			ETZPC, TZC					
	安全		TF-A					
启动过程		ROM code	在设备树中去掉不使用的节点 u-boot,dm-spl	U-Boot pre-reloc	U-Boot	Linux		
			u-boot,dm-pre-reloc					
			U-Boot SPL					
	Reset	ROM	FSBL	SSBL pre-reloc	SSBL	Runtime		



(*): U-Boot和Linux可以使用同一设备树,但是U-Boot只需一部分外设

STM32MP1启动模式配置。

BOOT pins	TAMP_REG[20] (Force Serial)	OTP WORD 3 Primary boot source	OTP WORD 3 Secondary boot source	Boot source #1	Boot source #2 if #1 fails	Boot source if #2 fails
b000	x (don't care)	x (don't care)	x (don't care)	Serial	-	-
b001	!= 0xFF	0 (virgin)	0 (virgin)	QSPI NOR	Serial	-
b010	!= 0xFF	0 (virgin)	0 (virgin)	еММС	Serial	-
b011	!= 0xFF	0 (virgin)	0 (virgin)	FMC NAND	Serial	-
b100	x (don't care)	x (don't care)	x (don't care)	NoBoot	-	-
b101	!= 0xFF	0 (virgin)	0 (virgin)	SD-Card	Serial	-
b110	!= 0xFF	0 (virgin)	0 (virgin)	Serial	-	-
b111	!= 0xFF	0 (virgin)	0 (virgin)	QSPI NAND	Serial	-
!= b100	!= 0xFF	Primary ¹	0 (virgin)	Primary ¹	Serial	-
!= b100	!= 0xFF	0 (virgin)	Secondary ¹	Secondary ¹	Serial	-
!= b100	!= 0xFF	Primary ¹	Secondary ¹	Primary ¹	Secondary ¹	Serial
!= b100	0xFF	x (don't care)	x (don't care)	Serial	-	-

0	No secondary boot source is defined
1	FMC NAND
2	QSPI NOR
3	еММС
4	SD
5	QSPI NAND

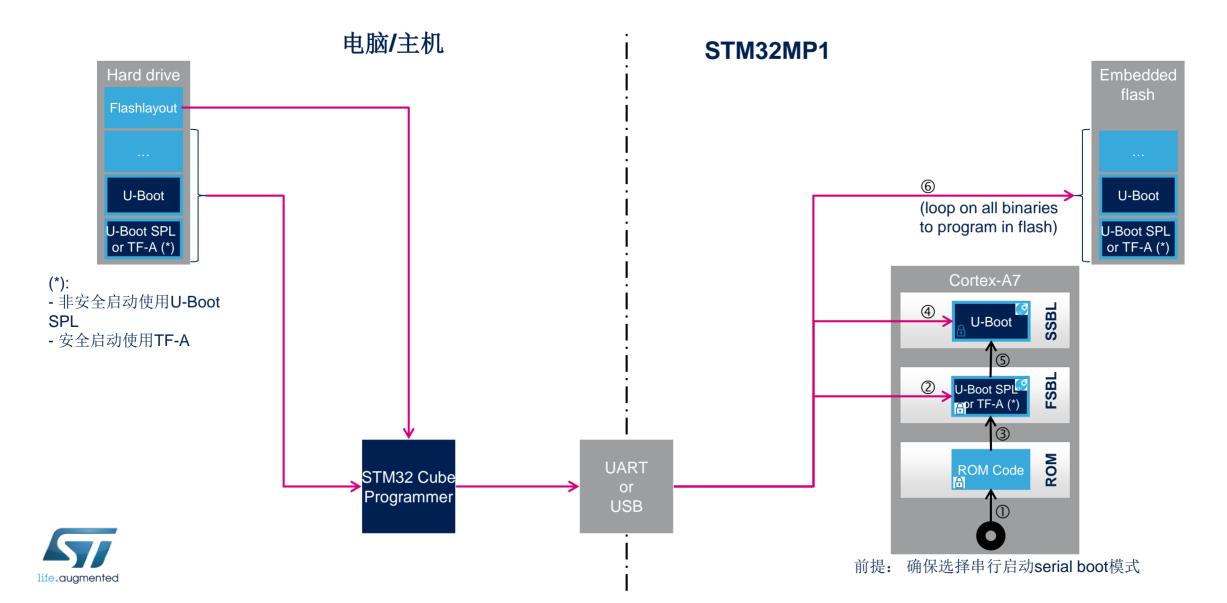
0	No primary boot source is defined
1	FMC NAND
2	QSPI NOR
3	еММС
4	SD
5	QSPI NAND

¹Primary and Secondary are fields of OTP WORD3.



例:若强制系统总是SD card启动,则可将OTP WORD 3的Primary boot source设为4

使用STM32CubeProgrammer进行烧录 型







STM32MP1多核通信



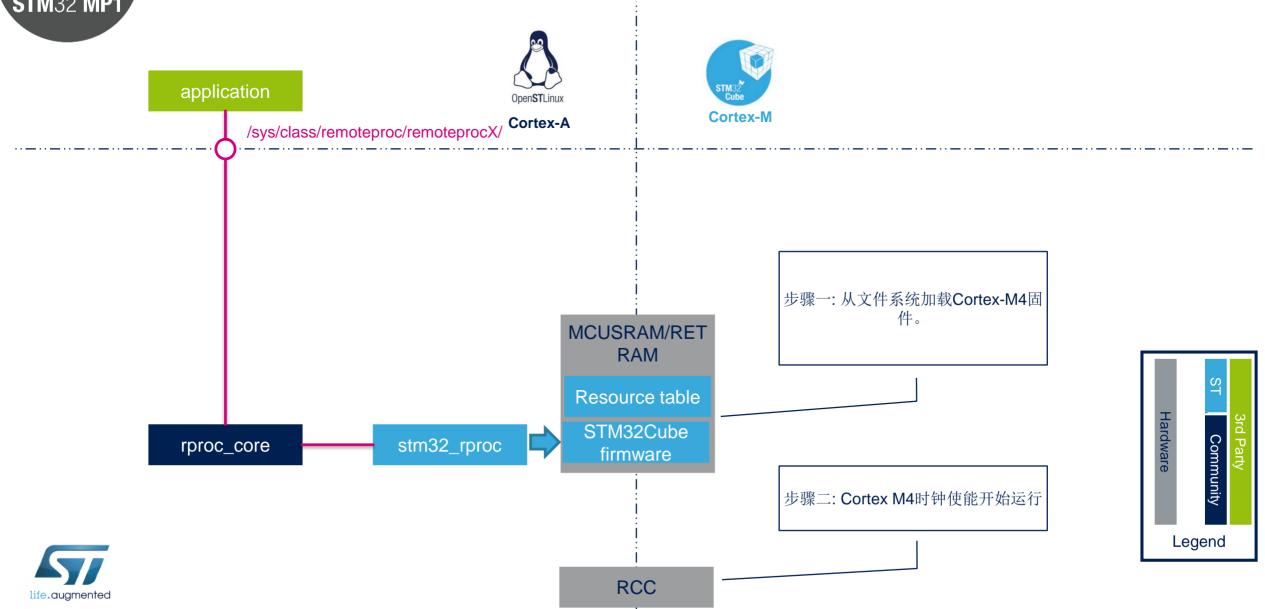
多核通信的目标 28

- · Cortex-M4的生命周期管理
- Cortex-A7与Cortex-M4的消息通讯





STM32MP1加载Cortex-M4固件



加载Cortex-M4固件

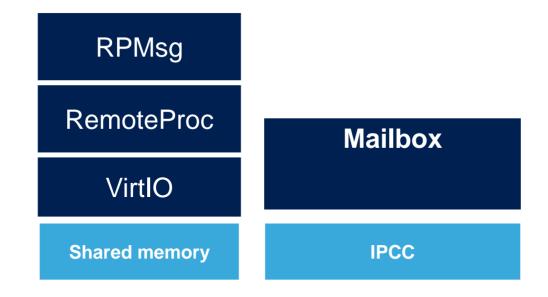
- Cortex-M4固件由Cortex-A使用Linux RemoteProc框架进行加载
- · Cortex-M4固件在文件系统中的格式ELF
- 固件中包含Cortex-A7和Cortex-M4都可以访问的资源表".resource_table"
 - Remote processor trace buffer.
 - Vring 队列和关联的Buffer
- 加载固件的两种方式
 - 自动: 在Linux启动时, lib/firmware存在设备树定义的名字
 - 手工: 使用sysfs接口 (推荐,因为文件系统可以后加载)

注意:根据需要可以在U-Boot中加载Cortex-M4固件



多核消息通信通用框架

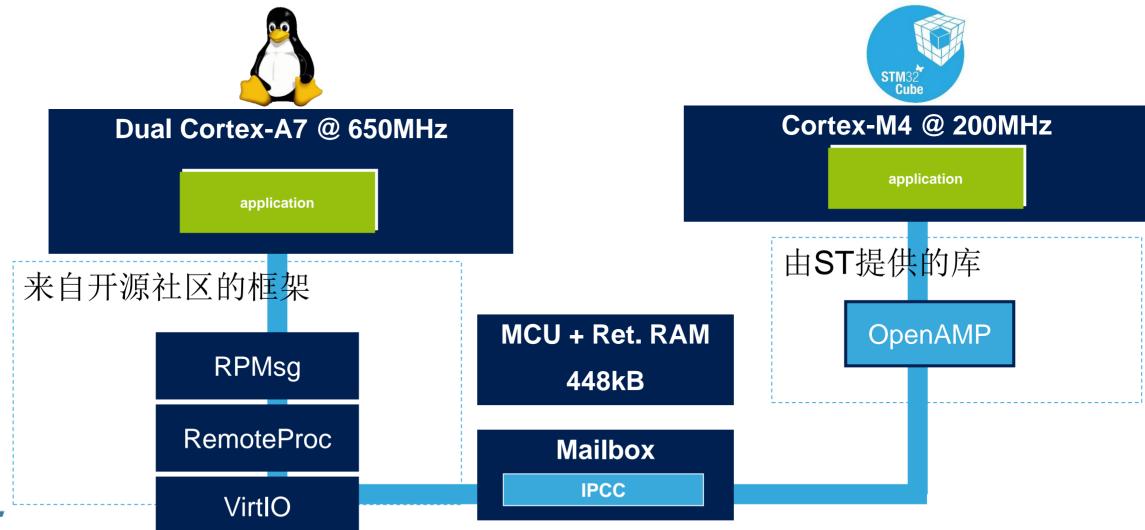
- 消息服务基于共享内存,使用RPMsg和Virtio框架
- 信号通知/mailbox服务则基于内部IPCC外设





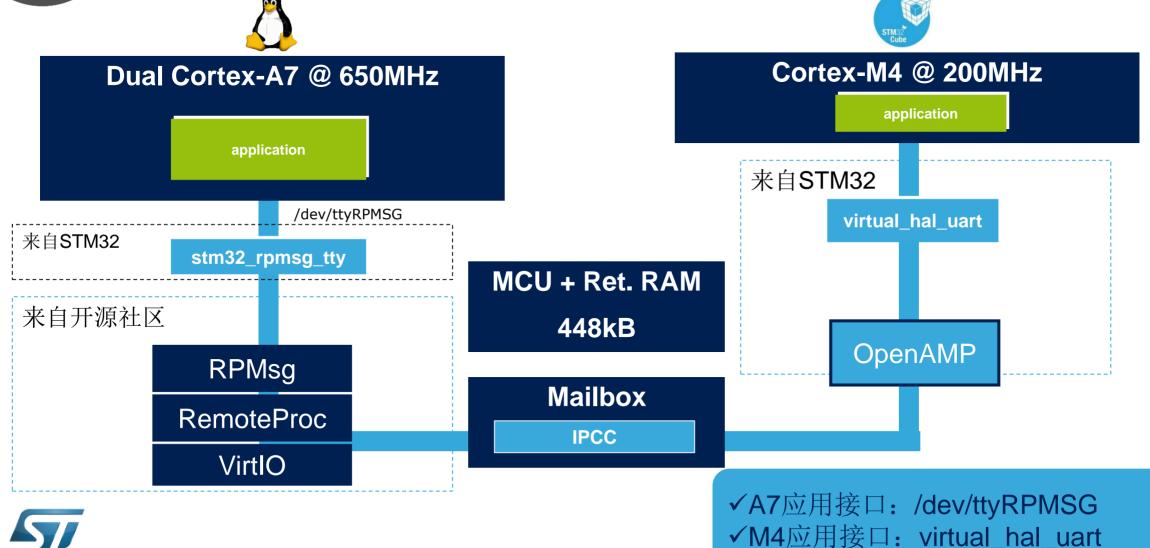


STM32MP1多核通信框架





STM32MP1多核消息通信应用接口



STM32MP1软件架构总结 34

- STM32MP1 Flash内存布局
- STM32MP1 RAM内存布局
- STM32MP1 多核软件架构
 - Cortex-A7
 - Secure
 - Non-secure
 - Bootloader
 - Kernel space
 - User space
 - Cortex-M4
- STM32MP1启动流程 life.augment STM32MP1多核通信