

# CSCI 350 Project 2 Spring 2022

Kernel Threads

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#### 1. Synchronization using locks

Part 1 - lock shared resources & add some functionality

#### 2. Implement kernel level threads

kthread create, kthread id, kthread exit, kthread join

#### 3. Implement mutex API

• kthread\_mutex\_alloc, kthread\_mutex\_dealloc, kthread\_mutex\_lock, kthread mutex unlock



### Part 1



Change the implementation of some existing system calls.

Growproc(), Fork(), Exec(), Exit()

Growproc(), Fork(), and Exit() methods are in proc.c

Exec() function is in **exec.c** 

Note: Some functions might not need changes you have to pick which ones to change.



### Growproc() - example

- Growproc() is responsible for retrieving more memory when the process asks for it
- sz variable alloc()/delloc() checking that
- We have to synchronize sz variable



```
int
growproc(int n)
  uint sz;
 // added for part 1.1
  acquire(&ptable.lock);
  sz = proc->sz;
  if(n > 0){
    if((sz = allocuvm(proc->pgdir, sz, sz + n)) == 0){
      // added for part 1.1
      release(&ptable.lock);
      return -1;
 } else if(n < 0){</pre>
    if((sz = deallocuvm(proc->pgdir, sz, sz + n)) == 0){
      // added for part 1.1
      release(&ptable.lock);
      return -1;
  proc->sz = sz;
  switchuvm(proc);
  // added for part 1.1
  release(&ptable.lock);
  return 0;
```



# Fork()



Fork should duplicate only the calling thread, if other threads exist in the process they will not exist in the new process.

Questions to ask:

Are there any conflicts between shared variables?

Do we need to kill any threads after calling fork?

Is the acquired the lock enough for synchronization or should we put more locks?

```
int
fork(void)
{
  int i, pid;
  struct proc *np;
  struct thread *nt;

acquire(&ptable.lock);

// Allocate process.
  if((np = allocproc()) == 0){
    release(&ptable.lock);
    return -1;
}

nt = np->threads;
```







Should kill the process and all of its threads, remember while a single threads executing *exit*, others threads of the same process might still be running.

### We have to create killall() method:

#### It kills all alive threads:

```
kill_all():
   Create thread pointer *t
   For each thread t:
   Begin
   If ( thread t is not current thread and not running and not unused)
     Make it zombie
   End
   Make current thread zombie
   Kill process
```

```
kill_all();

// Jump into the scheduler, never to return.
thread->state = TINVALID;
proc->state = ZOMBIE;
sched();
panic("zombie exit");
```







The thread performing exec should "tell" other threads of the same process to destroy themselves and only then complete the exec task.

### modify kill\_all() method and create kill\_others():

#### It kills all alive threads but itself:

```
kill_others():
   Create thread pointer *t
   For each thread t:
   Begin
    If ( thread t is not current thread and not running and not unused)
      Make it zombie
   End
```

in exec.c



### Other Hints for Part 1

- Find shared variables and put locks before them
- Always release locks before return statement if it is not released previously
- Use kill\_all and kill\_others in respective functions.
- Important: I did not explain all the steps, you might need to think more about synchronization and find where to put methods, locks etc.
- I will explain how to loop through threads and where to find current thread in later slides.



```
int
growproc(int n)
  uint sz;
  // added for part 1.1
  acquire(&ptable.lock);
  sz = proc->sz;
  if(n > 0){
   if((sz = allocuvm(proc->pgdir, sz, sz + n)) == 0){
     // added for part 1.1
      release(&ptable.lock);
      return -1;
  } else if(n < 0){
    if((sz = deallocuvm(proc->pgdir, sz, sz + n)) == 0){
     // added for part 1.1
      release(&ptable.lock);
      return -1;
  proc->sz = sz;
  switchuvm(proc);
  // added for part 1.1
  release(&ptable.lock);
  return 0;
```



### Part 2



#### Implement thread API for kernel.

```
hkthread_create, kthread_id, kthread_exit, kthread_join
```

You will implement these methods in proc.c and create header file kthread.h

Header file should include:

```
#define NTHREAD 16
int kthread_create(void*(*start_func)(), void* stack, int stack_size);
int kthread_id();
void kthread_exit();
int kthread_join(int thread_id);
```



### Changing thread states / Finding current thread



```
t->state = TZOMBIE;
```

```
t->tid != thread->tid
```

```
enum threadstate { TUNUSED, TEMBRYO, TSLEEPING,
TRUNNABLE, TRUNNING, TZOMBIE, TINVALID, TBLOCKED };
```

```
struct thread {
 int tid;
                               // Thread ID
                               // thread state
 enum threadstate state;
                              // Bottom of kernel stack for this thread
 char *kstack;
 struct proc *parent;
                              // Parent process
 struct trapframe *tf;
                               // Trap frame for current syscall
 struct context *context;
                              // swtch() here to run process
 void *chan;
                              // If non-zero, sleeping on chan
                              // If non-zero, have been killed
 int killed;
};
```

In proc.h





### Changing thread states / Finding current thread

```
t->state = TZOMBIE;
```

```
t->tid != thread->tid
```

```
enum threadstate { TUNUSED, TEMBRYO, TSLEEPING,
TRUNNABLE, TRUNNING, TZOMBIE, TINVALID, TBLOCKED };
```

### Trap Frame of thread

```
*t->tf = *thread->tf;
```

```
struct thread {
  int tid;
                               // Thread ID
  enum threadstate state;
                               // thread state
                               // Bottom of kernel stack for this thread
  char *kstack;
  struct proc *parent;
                               // Parent process
  struct trapframe *tf;
                               // Trap frame for current syscall
  struct context *context;
                               // swtch() here to run process
                               // If non-zero, sleeping on chan
  void *chan;
  int killed;
                               // If non-zero, have been killed
};
```

In proc.h



### How to loop through threads



#### Allocthread method in proc.c

```
struct thread*
43
    allocthread(struct proc * p)
45
                          Create a thread pointer
      struct thread *t;
46
      char *sp;
47
      int found = 0;
48
49
      for(t = p->threads; found != 1 && t < &p->threads[NTHREAD]; t++)
50
51
                                                     Loop threads
        if(t->state == TUNUSED)
52
53
          found = 1;
54
55
          t--;
56
        else if(t->state == TZOMBIE)
57
58
          clearThread(t);
59
          t->state = TUNUSED;
60
          found = 1;
61
62
          t--;
63
64
```

#### Allocproc method in proc.c

```
allocproc(void)
108
       struct proc *p;
       struct thread *t;
110
111
       for(p = ptable.proc; p < &ptable.proc[NPROC]; p++)</pre>
112
113
         II(p->state == uNUSEV)
114
           goto found;
                                            Loop processes
115
       return 0;
116
117
      found:
118
       p->state = USED;
       p->pid = nextpid++;
119
120
121
       t = allocthread(p);
122
123
       if(t == 0)
124
125
         p->state = UNUSED;
126
         return 0;
127
                                         Loop threads
128
       p->threads[0] = *t;
129
       for(t = p->threads; t < &p->threads[NTHREAD]; t++)
130
131
         t->state = TUNUSED:
132
133
       return p;
134 }
```

In proc.c





```
kill_all():
    Create thread pointer *t
For each thread t: -> loop through threads
Begin
    If ( thread t is not current thread and not running and not unused)->
    check its state
    Make it zombie-> change its state
End
Make current thread zombie -> find current thread and change its state
Kill process -> proc->killed = 1
```







- Calling kthread\_create will create a new thread within the context of the calling process. The newly created thread state will be TRUNNABLE. The caller of kthread\_create must allocate a user stack for the new thread to use (it should be enough to allocate a single page i.e., 4K for the thread stack). This does not replace the kernel stack for the thread.
- start\_func is a pointer to the entry function, which the thread will start executing. Upon success, the identifier of the newly created thread is returned. In case of an error, a non-positive value is returned.
- The kernel thread creation system call on real Linux does not receive a user stack pointer. In Linux the kernel allocates the memory for the new thread stack. You will need to create the stack in user mode and send its pointer to the system call in order to be consistent with current memory allocator of xv6.







```
kthread_create(void*(*start_func)(), void* stack, int stack_size)
Create a thread pointer
Allocate the thread using allocthread method
Check if t is 0 -> allocated correctly?
    If not return -1
Else
    Copy current thread's trap frame
    Find stack address of the thread using stack pointer given parameter
    Make stack pointer inside trap frame stack address + stack size
    Update base pointer inside trap frame as stack pointer
                                                                                   t->tf->ebp
    Find address of the start function which is given in parameter
    Make instruction pointer inside trap frame start address_
    return t id
```

This is not the only way to create a thread



### esp, eip, ebp?

### For more information you can refer to xv6 manual chapter 3



#### t->tf->esp

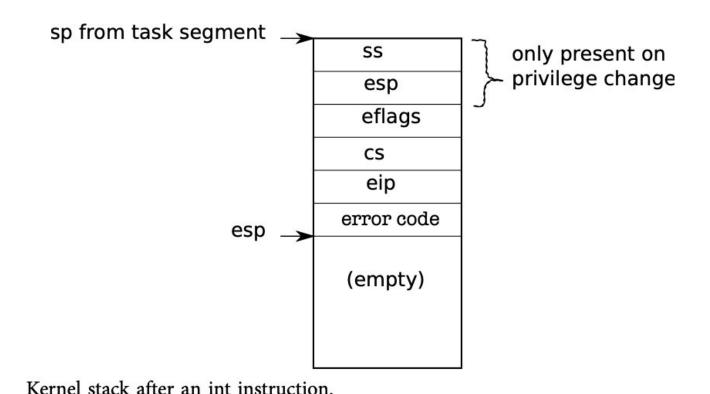
We have to change the stack address of new thread So we use given parameter to make stack address different than current thread

Add given **stack's address** with **stack size** to find where to put stack pointer

#### t->tf->ebp

Initially base pointer and stack pointer point the same place

Instruction pointer points the instructions which will be implemented by thread This pointer has to point stack function at the beginning



xv6 manual: <a href="https://pdos.csail.mit.edu/6.828/2014/xv6/book-rev8.pdf">https://pdos.csail.mit.edu/6.828/2014/xv6/book-rev8.pdf</a>

# kthread\_id



Easiest function to implement in project 2:)

- Upon success, this function returns the caller thread's id. In case of error, a non-positive error identifier is returned.
- Remember, thread id and process id are not identical.



### kthread\_id



Easiest function to implement in project 2:)

#### kthread\_id()

```
If process and thread exists
    return t->t_id
Else
    return -1
```

This is not the only way to return a thread id



### kthread\_exit



- This function terminates the execution of the calling thread.
- If called by a thread (even the main thread) while other threads exist within the same process, it shouldn't terminate the whole process.
- If it is the last running thread, process should terminate. Each thread must explicitly call *kthread exit()* in order to terminate normally.





kthread\_exit()

```
Create a thread pointer
Create a found flag
Loop through all threads to find another thread running
    If t is not current thread (because calling thread is current)
        If t is not Unused, not Zombied and not Invalid
            Make flag true
            Break -> only one running t is enough
If (found)
    Wakeup all waiting using wakeup1()
Else -> not found
    exit()
    wakeup()
Make this thread zombie
Call shed to schedule another thread
```

This is not the only way to exit a thread This is not a complete pseudocode You have to add locks if necessary



# kthread\_join



- This function suspends the execution of the calling thread until the target thread (of the same process), indicated by the argument thread\_id, terminates.
- If the thread has already exited, execution should not be suspended.
- If successful, the function returns zero.
- Otherwise, -1 should be returned to indicate an error.





```
kthread_join(int thread_id)
```

```
Check if thread_id is valid
Create a thread pointer t
Loop through all threads to find target thread id(parameter)
    Make t point target thread with thread_id

If not found
    return -1
While (t->t_id = thread_id and valid)
    Make t sleep using sleep method with a lock

If state of t is zombie
    clearThread(t);

return 0
```

This is not the only way to join threads
This is not a complete pseudocode
You have to add locks if necessary







#### Implement thread API for mutex.

```
kthread_mutex_alloc, kthread_mutex_dealloc, kthread_mutex_lock,
kthread_mutex_unlock
```

You will implement these methods in proc.c and add methods into header file kthread.h

#### Header file should include:

```
#define MAX_MUTEXES 64
int kthread_mutex_alloc();
int kthread_mutex_dealloc(int mutex_id);
int kthread_mutex_lock(int mutex_id);
int kthread_mutex_unlock(int mutex_id);
```







```
enum threadstate { TUNUSED, TEMBRYO, TSLEEPING, TRUNNABLE, TRUNNING, TZOMBIE, TINVALID, TBLOCKED };

Create mutexstate struct similar to this

You need:

Unused, locked, unlocked states
```

```
struct thread {
  int tid;
                             // Thread ID
  enum threadstate state;
                             // thread state
  char *kstack;
                             // Bottom of kernel stack for this thread
  struct proc *parent;
                             // Parent process
  struct trapframe *tf;
                             // Trap frame for current syscall
  struct context *context;
                             // swtch() here to run process
                                                                   You need:
  void *chan;
                             // If non-zero, sleeping on chan
                                                                   mutex id
  int killed;
                             // If non-zero, have been killed
                                                                   Mutex state
};
```

Create kthread mutex struct similar to this

mutex\_id
Mutex state
[other variables depending
on implementation]
Ex: blocked threads array



# Changes in proc.c



```
struct {
   struct spinlock lock;
   struct proc proc[NPROC];
} ptable;
```

You need to have a spin lock on mutex table similar to this

We will discuss other mutex methods in next slides...

```
int nextpid = 1;
int nexttid = 1;
int nextmid = 1;
int nextmid = 1;
```



# kthread\_mutex\_alloc()



- Allocates a mutex object and initializes it; the initial state should be unlocked.
- The function should return the ID of the initialized mutex, or -1 upon failure







```
for (i = 0; i < MAX_MUTEXES; i++)
{
    m = &mtable.mutex[i];
    // if found break
    if(m->mid == mutex_id)
        break;
}
```

```
for(m = mtable.mutex; m < &mtable.mutex[NPROC]; m++);</pre>
```

You can use other ways to do this, i.e. while loop etc.



# kthread\_mutex\_alloc()



```
Create a mutex pointer m (using struct in proc.h)

Loop through all mutex table

If m is unused

m->mid = nextmid++;
This is not a complete pseudocode
You have to add locks if necessary
Initiate all other values

If (&mtable.mutex[NPROC]) -> unused mutex not found
return -1 -> cannot allocate because there is no unused mutex in mutex table

Else
Return m->mid
```

This is not the only way to implement mutex allocation



# kthread\_mutex\_dealloc()



- De-allocates a mutex object which is no longer needed.
- The function should return 0 upon success and -1 upon failure
- For example, if the given mutex is currently locked



### kthread\_mutex\_dealloc()



```
Create a mutex pointer m (using struct in proc.h)

Loop through all mutex table to find given mutex_id

If m is locked -> we can't dealloc
Return -1

This is not a complete pseudocode You have to add locks if necessary

If (&mtable.mutex[NPROC]) -> unused mutex not found
return -1 -> cannot allocate because there is no unused mutex in mutex table

Else -> deallocate all properties of mutex
m->mid = 0
m->state = MUNUSED;
Zero all other values
Return 0

Return -1 -> if no mutex_id is found
```

This is not the only way to implement mutex deallocation



# kthread\_mutex\_lock()



- This function is used by a thread to lock the mutex specified by the argument  $mutex\_id$ .
- If the mutex is already locked by another thread, this call will block the calling thread (change the thread state to *TBLOCKED*) until the mutex is unlocked.

#### in proc.h

```
enum threadstate { TUNUSED, TEMBRYO, TSLEEPING,
TRUNNABLE, TRUNNING, TZOMBIE, TINVALID, TBLOCKED };
```

Add TBLOCKED state if you did not add earlier





```
Create a mutex pointer m (using struct in proc.h)
Loop through all mutex table to find given mutex id
    If m->mid == mutex id
        break;
If (i is MAX_MUTEXES) -> given mutex_id not found
    return -1
while (m->state == M_LOCKED) -
                                              → Spin lock
    sleep
If (m->state != M_UNLOCKED) -> failed
    Return -1
m->state = M LOCKED;
Return 0
```

This is not a complete pseudocode You have to add locks if necessary

YOU HAVE TO IMPLEMENT TBLOCKED!

This is not the only way to implement mutex lock







- This function unlocks the mutex specified by the argument mutex\_id if called by the owning thread, and if there are any blocked threads, one of the threads will acquire the mutex.
- An error will be returned if the mutex was already unlocked.
- The mutex may be owned by one thread and unlocked by another!







This is not a complete pseudocode You have to add locks if necessary

This implementation does not include last part of description as well:

• The mutex may be owned by one thread and unlocked by another!

This is not the only way to implement mutex unlock







```
./sign.pl bootblock
make: execvp: ./sign.pl: Permission denied
Makefile:95: recipe for target 'bootblock' failed
make: *** [bootblock] Error 127
```

```
sudo chmod +x *.pl
```







```
$ threadtest1
3 threadtest1: unknown sys call 23
thread in main -1,process 3
3 threadtest1: unknown sys call 22
3 threadtest1: unknown sys call 22
3 threadtest1: unknown sys call 25
Got id : -1
3 threadtest1: unknown sys call 25
Got id : -1
Finished.
3 threadtest1: unknown sys call 24
```

Make sure to implement system calls for all kthread and mutex methods.



### Common Errors



Most important error in Project 2

```
thread in main 3,process 3
Thread id is: 3
cpu with apicid 0: panic: acquire
80103e81 80103aa8 80104f3c 80104263 80105117 80104ff3 80111e8c 0 0 0
```

Also, panic: release

panic: acquire and panic: release errors mean that program fails to acquire lock because it is already acquired earlier or it cannot released lock because it is already released.



### Common Errors



panic: sched locks

This means process holding multiple locks.

Before calling sched() make sure to release all other locks.

ac(ptable) ac(ttabke) ac(mtable)

rel(mtable)
rel(ttable)
rel(ptable)





# Questions?

