

1. INTRODUCTION

For the purpose of the report we work over \mathbb{C} , this generalises to any algebraically closed field of arbitrary characteristic. All varieties we consider are normal and projective. Here we give an algorithm to classify Log Del Pezzos admitting a \mathbb{C}^\times action with only log terminal singularities. A variety X of dimension n which admits a torus action of dimension $n - k$ is referred to as complexity k . Here complexity 0 is the study of purely toric varieties, and complexity n is the study of varieties with no possible torus action. This provides essentially a way of grading the difficulty of your problem. Significant progress has been made on this problem before: Süss [?] he classifies log del Pezzo surfaces admitting said action with Picard rank one and index less than 3. Huggenberger [?] she classifies the anticanonical complex of the Cox ring of log del Pezzo surfaces with index 1, this classification was later finished by Ilten, Mishna and Trainor [?] with a view towards higher dimension. This was achieved by looking at polarised complexity one log del Pezzo surfaces. We will show their work fits into our algorithm.

2. POLYHEDRAL DIVISORS

Recall that a toric variety is a normal variety of dimension n containing a dense torus $(\mathbb{C}^*)^n$ with the natural action extending to the variety, there is a one to one correspondence between these varieties and fans inside a lattice $N \cong \mathbb{Z}^n$ upto $\mathrm{GL}_2(\mathbb{Z})$, [?]. Altman et.al [?] establish a similar correspondence for varieties with $T = (\mathbb{C}^*)^{n-k}$ actions where $k \leq n$. We say that this is a torus action of complexity k . They introduce the notion of a polyhedral divisor to recover some of the geometry that a fan encodes in the toric case. In general this applies for any complexity, however the behaviour is easiest to describe in the toric case, then complexity one until you reach the full general case.

Given X , a variety with dimension n admitting of the torus $T = (\mathbb{C}^*)^{n-1}$ action, we can take a Chow quotient Y of X by T , essentially a GIT quotient followed by normalisation. We see that Y will be a variety of dimension k , we can resolve this map to \tilde{X} getting the following diagram

$$\begin{array}{ccc} X & \longleftarrow & \tilde{X} \\ & \searrow & \downarrow \\ & & Y \end{array}$$

Here $Y \cong C$ is a normal curve. In this thesis we will primarily be interested in the case where $C \cong \mathbb{P}^1$. We start by introducing the notion of a tail cone of a given polyhedral cone. This is given F a cone the tail cone δ is the set of $v \in N$ such that for all $u \in F_i$, $\forall v \in \delta$ then $u + v \in F_i$.

Definition 2.1. Let C be a non singular curve then we define a polyhedral divisor to be the pair $(\mathcal{D} = \sum_{i=1}^k F_i \otimes P_i, \delta)$ where

- $P_i \in C$ are divisors on C
- F_i is a cone in $N_{\mathbb{Q}} \cong \mathbb{Q}^{n-1}$ and all F_i have tail cone $\delta \subset N$. We allow the cone F_i to be \emptyset .

Given an element $v \in M$, the dual lattice of N , and the polyhedral divisor \mathcal{D} we define

$$\mathcal{D}(v) = \sum \min_{u \in F_i} \langle u, v \rangle P_i$$

This is defined as a divisor on the following curve

$$Y_{\mathcal{D}} = C - \{P_j\}_j \text{ where } F_j = \emptyset$$

This defines a divisor on a subset of C . We insist that \mathcal{D} satisfies the following conditions:

- $\mathcal{D}(u)$ is Cartier for all $u \in \delta^{\vee}$
- $\mathcal{D}(u)$ is semiample for all $u \in \delta^{\vee}$
- $\mathcal{D}(u)$ is big for all u in the relative interior of δ^{\vee}

This is to ensure that it gives an n -dimensional variety, and to ensure that it is separated [?].

We can now calculate the associated affine variety in both X and \tilde{X} by taking respectively $\text{Spec}/\text{RelSpec}_C$ of the graded ring

$$\bigoplus_{v \in \delta^{\vee}} \mathcal{O}_{Y_{\mathcal{D}}}(\mathcal{D}(v))$$

This gives us an affine variety with $T = \text{Spec } \mathbb{C}[M]$ acting by torus action. Analogous to the toric case, if $F_i \subset F_j$ is a face then we have

$$\bigoplus_{v \in \delta^{\vee}} \mathcal{O}_{Y_{\mathcal{D}}}(\mathcal{D}_{F_j}(v)) \subset \bigoplus_{v \in \delta^{\vee}} \mathcal{O}_{Y_{\mathcal{D}}}(\mathcal{D}_{F_i}(v))$$

This corresponds to an inclusion of schemes. We make the following comment that taking a divisor

$$\mathcal{D} = \sum_{i=1}^n F_i \otimes P_i + \emptyset \otimes P_{n+1}$$

is the same as taking the divisors

$$\mathcal{D}_j = F_i \otimes P_i + \sum_{j=1, j \neq i} \emptyset \otimes P_j$$

and then glueing these affine varieties together along the affine patch defined by $C - P_i - P_j$ for all P_i, P_j .

In the case of surfaces of complexity one we often use the notation of fancy divisors as set out in [?]. This follows the key notion that in the case of $n = 2$ and $k = 1$ we have that every tail fan is either 0 , $\mathbb{Z}_{\geq 0}$ or $\mathbb{Z}_{\leq 0}$. We have n subdivisions of $N \cong \mathbb{Z}$, these should be viewed as the polyhedral divisors over these n points. Note that if we have a closed interval in any of subdivisions this will have tail fan zero and these give rise to a cyclic quotient singularity, with a nice torus quotient, i.e the map to \tilde{X} is a contraction to a point. It is the intervals $[a_1, \infty)$ which provide difficulty, if as polyhedral divisors these are all of the form

$$\mathcal{D}_i = [a_i, \infty) \otimes P_i + \sum_{\substack{j=1 \\ j \neq i}}^n \emptyset \otimes P_j$$

Then this gives rise to a nice quotient map down base curve with respect to the torus action, i.e the map to \tilde{X} is a local isomorphism. If this is not the case however, then we are left with a bad quotient. These are the only two cases that can occur, in the surface case. In the language of fancy divisors we say if we mean the latter case we denote it with \mathbb{Q}^+ , if we mean the other the earlier case, we do not denote it at all. In this way fancy divisor uniquely specify polyhedral fans.

Definition 2.2. *A fancy divisor a smooth curve C is a collection of n subdivisions of \mathbb{Q} with markings \mathbb{Q}^+ , \mathbb{Q}^- , \mathbb{Q}^\pm or no markings at all.*

This is equivalent to a polyhedral divisor provided the the assumptions of ?? are satisfied.

Example 2.3. *One A_1 two A_3 .*

This defines a complexity one surface. In toric varieties full dimensional cones give rise to torus fixed point. Analogously, the same way for varieties of higher complexity every full dimensional subdivision of the plane gives rise to a toric fixed point. In the case of surfaces these fixed points can be classified giving rise to three cases

- **Elliptic** - Around the fixed point in local coordinates, the torus behaves on all coordinates with positive or negative degree. These points are isolated.
- **Parabolic** - These always arise as blowups of elliptic points, these occur when in local coordinates, one of the coordinates is acted trivially upon by the torus. These points lie on a section of the map to Y
- **Hyperbolic** - These are where the the local coordinates are acted in positive and negative degree.

It is easy to see that Hyperbolic points correspond to a subdivision with $\delta = 0$, Parabolic correspond to an unmarked edge going to infinity and Elliptic to a marked point going to infinity.

3. DIVISORS IN COMPLEXITY ONE

We now limit ourselves strictly to complexity one, and the chow quotient Y will now be \mathbb{P}^1 . In the torus setting we know that divisors correspond to rays of the associated fan. Almost exactly the same is true in complexity one: divisors occur as torus invariant divisor, these correspond the codimension 1 polyhedral divisors or they are preimages of the \mathbb{P}^1 . These correspond to a polyhedral divisor \mathcal{D} going off to infinity in a direction, with $\dim(\delta) = \infty$ which for all $P \in \mathbb{P}^1$ we do not have $\mathcal{D}|_P = \emptyset$. Note that this also holds for higher dimensions, with a little bit of extra work. From this it is easy to derive the following theorem

Theorem 3.1 ([?]). *The Picard rank of a complexity one surface defined by a polyhedral fan \mathcal{S} is*

$$\rho_X = \# \text{ Number of parabolic lines} + \sum_{P \in Y} (\#\mathcal{S}_P^{(0)} - 1)$$

where n is the dimension and $\#\mathcal{S}_P^{(0)}$ is the number of points on this slice of the fan. Similar statements can be made in dimension n where the parabolic lines are replaced by x-rays. In a similar style to this we can classify Cartier divisors, we here make no pretense at proof or justification.

Definition 3.2. *A divisorial support function h on a divisorial fan \mathcal{S} is a piecewise linear function on each component of the fan such that*

- *On every polyhedron $\Delta \in \mathcal{S}_{P_i}$ it is a linear function*
- *h is continuous*
- *at all points h has integer slope and integer translation*
- *if \mathcal{D}_1 and \mathcal{D}_2 have the same tail cone, then the linear part of h restricted to them is equal*

We call a support function principal if it is of the form $h(v) = \langle u, v \rangle + D$, this corresponds to a principal Cartier divisor. We call a support function Cartier, if on every component with complete locus the support function is principal. In the case of Fansy divisors, this just correspond to the edge with a marking. We denote h restricted to a component by h_P . We refer to a piecewise linear function with rational slope and rational translation as a \mathbb{Q} support function.

Theorem 3.3 ([?]). *Let X be the variety associated with the divisorial fan \mathcal{S} . There exists a one to correspondence between support functions support function quotiented by principal support functions and Cartier divisors on the complexity one variety. In addition there exists a one to correspondence between \mathbb{Q} support functions support function quotiented by principal support functions and \mathbb{Q} Cartier divisors on the complexity one variety*

Using the above languages we represent the canonical divisor as a Weil divisor, it has the following form

Theorem 3.4 ([?]). *The canonical divisor of a complexity one surface can be represented in the following form*

$$K_X = \sum_{(P,v)} (\mu(v)K_Y(P) + \mu(v) - 1) \cdot D_{(P,v)} - \sum_{\rho} D_{\rho}$$

Here $K_Y(P)$ is the degree of K_Y at P , and $\mu(v)$ is the smallest value k such that $k \cdot v \in \mathbb{N}$. While I have not stated the conditions for linear equivalence these can be seen in [?], and using these you can show that it does not depend on the choice of representative of K_Y . Note that given the singularities and varieties we are working with we know that our K_X will be \mathbb{Q} -Cartier. The fano index is clear and easy to derive from the singularities we have, so all that remains is to check on the conditions for a complexity one divisor to be ample.

Theorem 3.5 ([?]). *A support function h is ample iff for all P we have h_P is strictly concave, and for all polyhedral divisors \mathcal{D} defined on an affine curve we have*

$$- \sum_{P \in \mathbb{P}^1} h_P|_{\mathcal{D}}(0) \in \text{Weil}_{\mathbb{Q}}(Y)$$

is an ample \mathbb{Q} Cartier divisor.

Note that in reality $h_P|_{\mathcal{D}}$ may not be defined at 0 but we can extend the affine function to 0. We finish this recap on divisors by describing the Weil divisor corresponding to a Cartier divisor

Theorem 3.6 ([?]). *Let $h = \sum_P h_P$ be a Cartier divisor on \mathcal{S} then the corresponding Weil divisor is*

$$- \sum_{\rho} h_{\rho}(n_{\rho})D_{\rho} - \sum_{(P,v)} \mu(v)h_P(v)D_{(P,v)}$$

Here n_{ρ} is the generator of the ray inside the tail fan and $\mu(v)$ is as before. Note that is easy to see why we need this μ function. If you start with a closed subinterval $[a, b]$ and try to work out what the corresponding affine variety is, we see that it just the toric variety defined by the cone $(a, 1)$, $(b, 1)$, and then all your calculations can be done in the realm of toric varieties, however there you use the generator of your rays in the lattice, so you need the μ function.

We use the above note to easily calculate the minimal resolution of a complexity one surface. Note that we can split this across affine charts, in the first case if we have the affine chart corresponding to the polyhedral divisor $[a, b]$ then using the above point we can calculate this by the toric methods. In case two where we have a non marked edge going to infinity, we can split this into affine charts $[a_i, \infty)$ this is also a toric chart corresponding to

the cone $(a, 1)$, $(1, 0)$, so once again the resolution is toric. The final case is with a marked edge, however we can take a weighted blowup to resolve the elliptic point, then resolve the resulting singularities by the above methods. To calculate the intersection numbers on the resolution you can either use [Tim], [?] or you can note that the only part that is not toric is the parabolic line, this is defined by glueing together charts coming from $[a'_i, \infty)$, here by smoothness $a'_i \in \mathbb{Z}$, this is isomorphic to the charts defined by $[\sum(a'_i), \infty)$ at P_1 and $[0, \infty)$ for all other P_i . Hence we see that the parabolic line is defined torically as the fan $(\sum(a'_i), 1)$, $(1, 0)$, $(0, -1)$ from this an easy derivation of the intersection number follows.

We can also draw out the graph of divisors on the minimal resolution. For example considering the following log del Pezzo from [?]

$$\{-2, 0\} \otimes P_0 + \left\{-\frac{1}{2}\right\} \otimes P_1 + \left\{-\frac{1}{2}\right\} \otimes P_2$$

gives us the following resolution:

4. ALGORITHMS

We propose two different algorithms for the classification of complexity one log del Pezzo surfaces. These both rely on several key facts

Lemma 4.1. [?] *Let S be a non cyclic complexity one log terminal surface singularity. Then S has, upto isomorphism, a fan over \mathbb{P}^1 with coefficients*

$$\left[\frac{p_1}{q_1}, \infty\right) \otimes P_1 + \left[\frac{p_2}{q_2}, \infty\right) \otimes P_2 + \left[\frac{p_3}{q_3}, \infty\right) \otimes P_3$$

with (q_1, q_2, q_3) satisfying $\sum(1 - \frac{1}{q_i}) < 2$.

Proof. See [?] □

We now use the following crucial lemma

Lemma 4.2. *Let S be a log terminal surface singularity of Gorenstein index l . Let E be an exceptional curve in the minimal resolution. Then $E^2 \geq -2l$ if it is not a trivalent curve and $E^2 \geq -3l$ if it is trivalent.*

Proof. Via the classification of log terminal singularities [?] we have that E intersects at most three other exceptional curves. Denote the discrepancies of these curves d_1, d_2, d_3 , note that any d_i could be equal to zero. Also note that $0 \geq d_i \geq -1$. Denote the discrepancy of E by d . Then we have the formula $dE^2 + \sum d_i = 0$. This rearranges to $d = \frac{(\sum d_i)}{E^2} \leq \frac{-3}{E^2}$ as $d \in \frac{1}{l}\mathbb{Z}$ we get $E^2 \geq -3l$. In the case of a non trivalent curve, we can assume $d_3 = 0$ and we see that $E^2 \geq -2l$. □

Lemma 4.3. *Given a complexity one log del Pezzo surface of index l then there cannot be more than $6l$ points where the polyhedral fan is not the tail fan*

Proof. Taking the minimal resolution of our log del Pezzo, this admits a map to a Hirzebruch surface \mathbb{F}_n . As we are contracting -1 curves our map is invariant under the torus action. Hence this is a torus action on the Hirzebruch surface. Any series of complexity one non toric blowups on a toric surface correspond to blowing up points on a line of invariant points. We note that by the above lemma we cannot get a map to \mathbb{F}_n when $n > 3l$. Hence we the largest possible self intersection of a torus invariant curve on our Hirzebruch surface is $3l$ and the smallest possible intersection on our minimal resolution is $-3l$ so there can only be $6l$ blowups on the curve. \square

Remark. *In the case of index one, we know DuVal singularities only have -2 curves in the resolution hence this bound can be refined to four non general fibers.*

We now, for the sake of classifying abandon the condition on index, replace it with this condition on self intersections. It is easy to check the index post classification. We can now specify how the first algorithm works

PUT ALGORITHM HERE

As an example we illustrate how this can classify Gorenstein log del Pezzo surfaces which have complexity one.

Example 4.4. *As a Gorenstein surface singularity are of DuVal type we see that the minimal resolution can only admit maps to \mathbb{F}_0 , \mathbb{F}_1 or \mathbb{F}_2 . We now ask what possible torus actions can there be on these Hirzebruch surfaces such that we will only have DuVal singularities. For convenience we write the fan of \mathbb{F}_n having rays generated by $(1,0)$, $(0,1)$, $(-1,n)$, $(0,-1)$. We can then write a subtorus as a vector $v \in N$.*

We make the observation that we have classified all log del Pezzo surfaces which do only admit maps to \mathbb{F}_0 or \mathbb{F}_1 ???. This results in only one surface.

For the case of \mathbb{F}_2 either our torus action coincides with a ray of the fan or it does not. There are only four cases where it coincides. However via symmetry we only need to consider the rays $(0,1)$ and $(1,0)$. In the first case your surface has two parabolic lines and hence admits a well defined map to \mathbb{P}^1 . As this a morphism we see that the our torus invariant fibers are 0 curves. The only ways to blow up a zero curve to get values greater than -2 are:

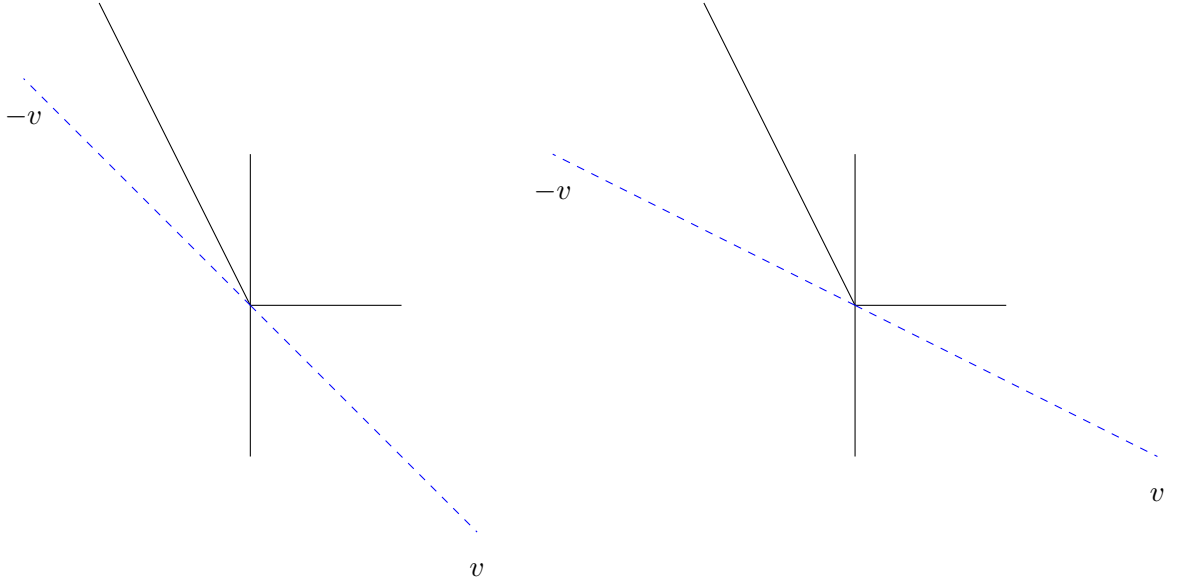
- $-2, -1, -2$
- $-1, -2, \dots, -2, -1$
- $-1, -1$
- 0

We note that with the exception of the second case nothing there are finitely many options. It is clear that in the second case a chain of k -2 curves can only arise on a fiber by blowing

up parabolic lines a total of $k+1$ times. However earlier on we showed that there could only be a total of 4 blowups on the parabolic line hence this is bounded.

Dealing with the case where the vector v does not correspond to a ray of the fan. To do a non toric blowup we need to do a sequence of blowups that will produce a line of torus invariant points. We note that if we do this twice then there will be a well defined map to \mathbb{P}^1 with the two parabolic lines being sections. This means that there would be a map to a Hirzebruch surface such that one of these lines is being sent to the negative section and the other to the positive. As both of these would be fixed by the torus action these have been classified in the previous paragraphs.

As this blow up needs to be a $(1, n)$ or a $(n, 1)$ weighted blowups the only choices are the following



With the above disclaimer we carry on. We know that our standard admit a canonical map to a Hirzebruch surface, so we instead work backwards, take a Hirzebruch surface, look at a subtorus action on it and consider all the ways we can make a basic surface out of it. From here on out our fan for a Hirzebruch surface \mathbb{F}_n will be $(0, 1)$, $(1, 0)$, $(0, -1)$, $(-1, n)$. There are 4 possible ways the diagram of the Weil divisors on a Hirzebruch surface with a given \mathbb{C}^* action can look

- A - Two parabolic lines, this corresponds with the subtorus $(0, \pm 1)$.
- B - One parabolic line, this corresponds to the subtorus $(\pm 1, 0)$.

C - Two elliptic points, connected by a line, this corresponds the sublattice generated by a point lying in between $(-1, 0)$ and $(-1, n)$.

D - Two elliptic points, not connected by a line, this corresponds to any other point.

In the above pictures we have the n curve on the top and the $(-n)$ curve on the bottom, with the two vertical lines being the 0 fibers. The blue lines are parabolic lines, and the red points represent elliptic points.

If X is a log del Pezzo with Y its minimal resolution. The number of Elliptic points can only decrease in the cascade, hence we see that it would map down to one of B, C, D. Next note that if we consider case C or D there is no way to make it non toric without resolving one of the elliptic points. Because of this we have the following theorem

Theorem 4.5. *Let X be a complexity one log del Pezzo, Y its minimal resolution. Then if Y has two elliptic points, then X is toric.*

We now split things into a case by case analysis. In case A, we have 0 fibers so we just substitute them with the all possible choices of fibers in??, using [?] we know that we can only substitute in at most 4 fibers. Hence this case is finished. Note that actually the 4 fibers comes out in the calculations in this case, although that does not prove the general log del Pezzo case.

For case B, if we resolve the elliptic point on Y in the process of our cascade then it admits a canonical \mathbb{P}^1 fibration, hence we can factor it through case A. Hence we only care about ones that preserve the elliptic point. From the fan we know that it requires n blowups of \mathbb{F}_n to resolve the elliptic point, in these cases the elliptic point lies on the intersection of a (-1) curve and a curve C with $C^2 > 0$. We also see that the zero curve intersecting the elliptic point has to be taken to one of the cases ??, we deal with the three cases, nothing happens at the elliptic point, it becomes an A_n singularity and finally the $[-2, -1, -2]$. If $n \neq 1, 2$ then we cannot have an A_n singularity, as we would have a (-1) curve next to a curve with positive self intersection, i.e the top line, so we would have to keep on blowing up till it has negative self intersection, but this would resolve the elliptic point. Also note if $n = 0$ case B does not occur, it would take zero blowups to resolve the elliptic singularity, i.e we just have two parabolic lines, so it is just case A. We deal with $n = 1, 2$ separately. In the case of nothing happening at the elliptic point. We need to blow up two point on the parabolic curve to get a non toric surface. After this we have two (-1) intersecting (-2) curves, we know by ?? there is nothing more we can do at those points, so we are left with only being able to blow up the intersection of a (-1) curve with a positive curve. We know we cannot blow up a third point on the parabolic line or it will stop being a klt singularity. Moving on to the $[-2, -1, -2]$ case, we have all the same possibilities as before, however it may take slightly less blowups than before we reach an acceptable configuration as we have our elliptic point lying on a (-2) curve. This leaves

with only two non toric options from a given Hirzebruch surface.

In the case of $n = 1$ we cannot blowup the elliptic point, the only other difference is not being able to use the klt argument, however the (-1) curve being next to the 0 curve guarantees that it still cannot be less than (-2) , and if we modify the (-1) curve in any other way it would arise from a different configuration and hence has already been classified. In the case $n = 2$ it also falls under the previous classification as we are only allowed one blowup and in this case the two (-1) curves intersect.

In case C we see there is a symmetry between the two elliptic points, so it does not matter which we resolve. When we resolve it we will have (-1) curve adjacent to at most two sets of negative curves. To make a non toric example, we need to blow up one of the lines connecting the parabolic line to the elliptic point. This curve has self intersection greater than 0, using the same argument as before, we can only blow up one point on the parabolic line. As our curve has positive self intersection we know, that we have to have, by the previous argument again, the following set of curves connecting to our parabolic line $[-2, -1, -2, \dots -2]$. In particular the number of -2 curves is greater than two. Using [Ishii, Brieskorn] classification of singularities, we see that to be klt, you have to have at most 3 negative meeting in a point and one of those has to be just -2 curve. Because of this the possible points generating the torus action, up to symmetry, are

- $(n, -1)$ or $(1, -n)$, here this is a $\mathbb{P}(1, n)$ blowup. One component is smooth.
- $(2n - 1, -2)$ or $(2, -2n + 1)$ this weighted blowup gives a $\frac{1}{2}(1, 1)$ singularity in one chart, this is the singularity whose resolution is a just a (-2) curve.

Note that different values of n give different singularities in the other chart. In the first case, they are clearly just A_n singularities. In the second case, you just get $[-3, -2, \dots -2]$, here the chain has $n - 1$ of the (-2) curves. We also need both case of the these torus actions as the vary which side the singularities appear. Of course the Brieskorn classification is much more comprehensive than this, and once the code is written I will start a more comprehensive check as to which or these give klt log del Pezzos.

In case D, it proceeds almost exactly the same as case C. However you do not have the pleasantness of the symmetry. However it is still straight forward what will happen, the elliptic point next to 0 curve and the n curve will have the same possible choices of torus actions as in Case C. For the one on the intersection on the $(-n)$ curve and the 0 curve, 2 of the expected 4 options will not occur due to the presence of the $(-n)$ curve. The Brieskorn classification in this situation gives us $n \leq 5$.