COMP90050: Advanced Database Systems
Semester 2, 2024
Lab 2b: Intention Locking

In this lab session, we delve into Intention Locks, specifically IS and IX locks.

I (Intention) Lock:

For this section, we will be looking at row level locks for specific S and X locks while the relevant table will be under an IS or IX lock as necessary. Generally, databases manage Intention locks themselves including coordinating the sequence of granting Intention locks at coarse granularities with specific locks at fine grained granularities. To enable us to see some of these functionalities more easily, we are going to explicitly setup Intention locks using SQL.

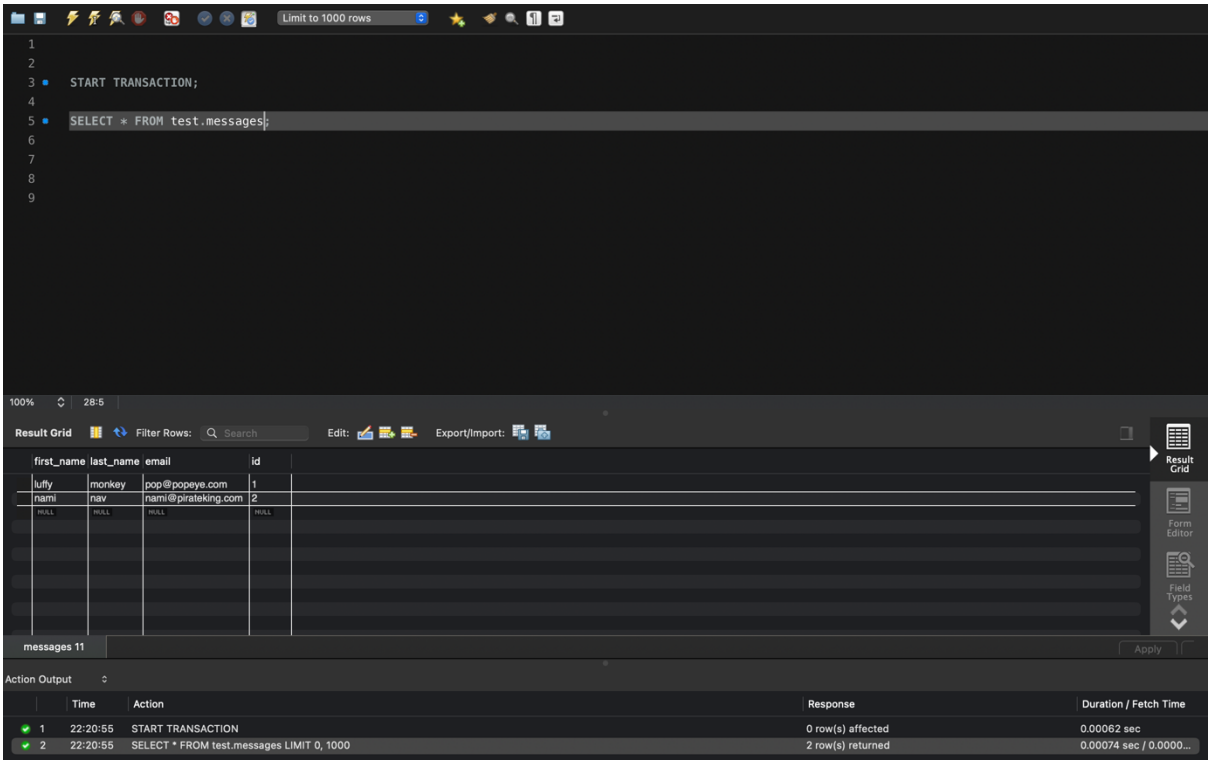
IS Lock:

Note: As part of this Lab, we use the same database table as used in Lab 2a.

The format to explicitly request for an IS lock is as follows:

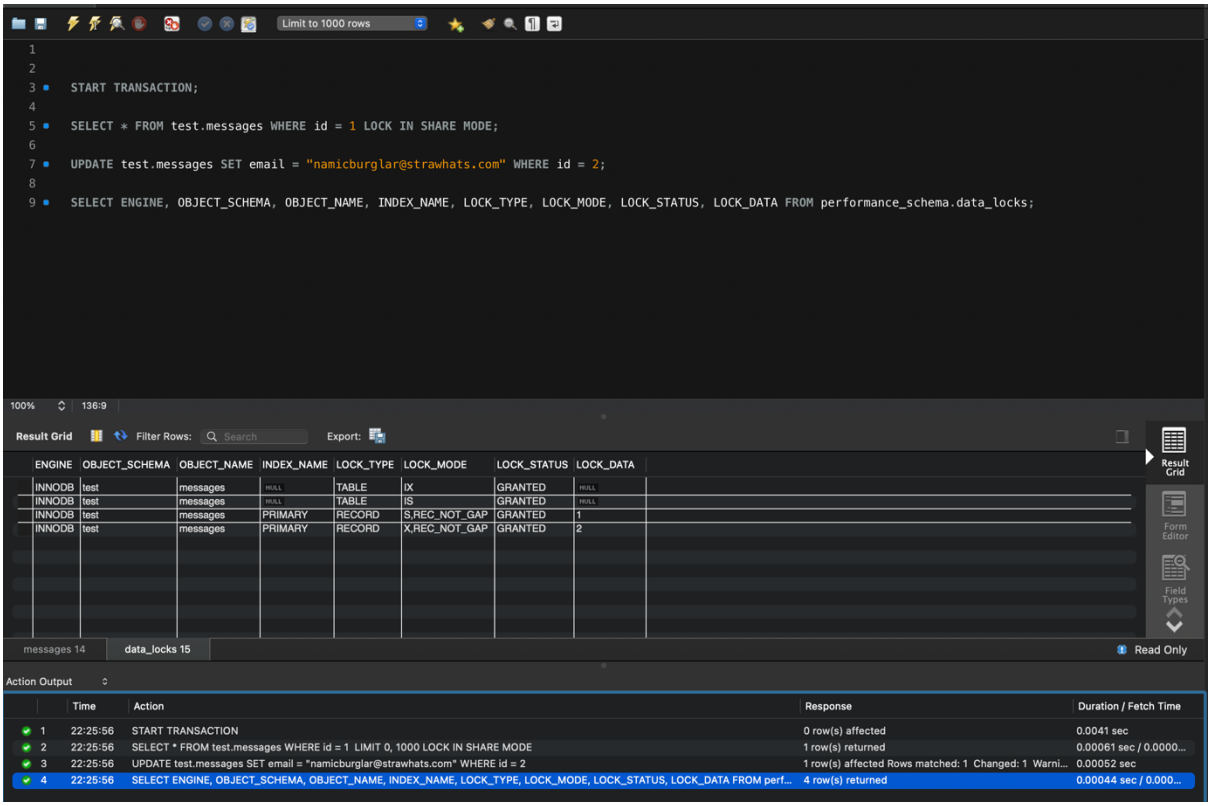
SELECT …. LOCK IN SHARE MODE;

For this process, we are going to start a transaction to ensure that we retain the lock till we complete the transaction, i.e. till we either COMMIT or ROLLBACK.

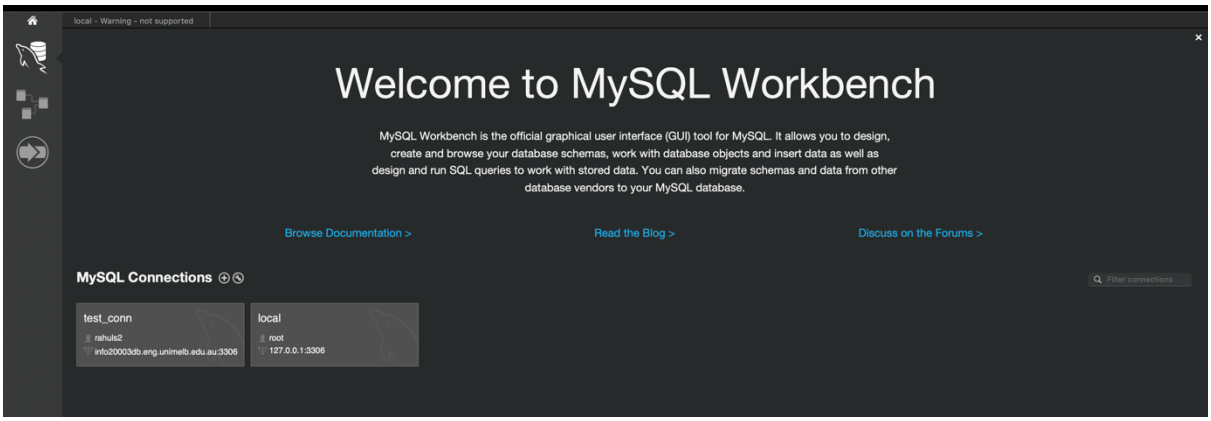


We can observe that there are two rows in the test.messages table at the moment. We get an IS lock on the test.messages table with a specific lock (Shared lock) on a particular row in that table, with row id = 1. Simultaneously, we also try to update one of the columns (email_id) of row with id = 2. When we run these queries in a transaction, the type of locks are held onto by the transaction till we either COMMIT or ROLLBACK. To observe the type of Locks held on specific resources/hierarchy of the database, we can run a SELECT query on the performance_schema.data_locks table, which provides additional context around type of locks currently held on specific resources.

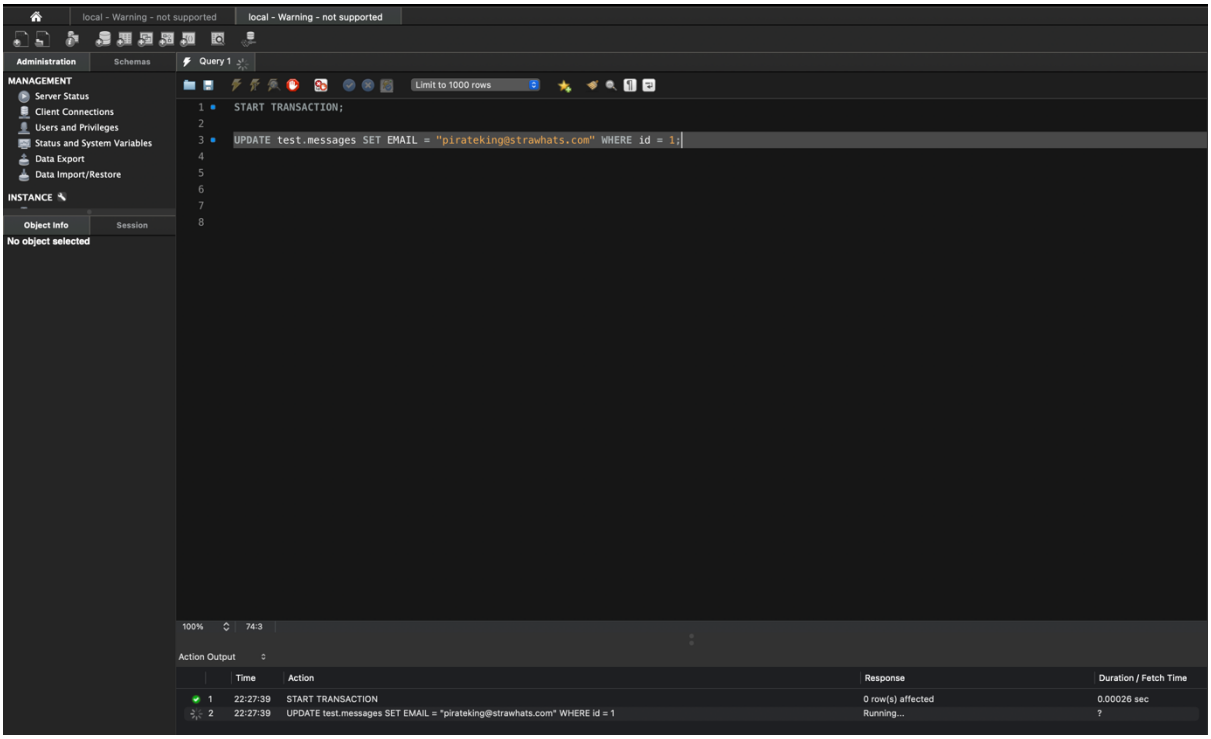
Further information on the performance_schema.data_locks table can be found here: https://dev.mysql.com/doc/refman/8.0/en/performance-schema-data-locks-table.html



We can see that the current session holds an IS and IX lock on the test.messages table, with a specific S lock on row 1 and X lock on row 2. We will now test how these locks persist and work across different transactions. The way we do this is by opening a new Session via the MySQL Workbench home page as follows:

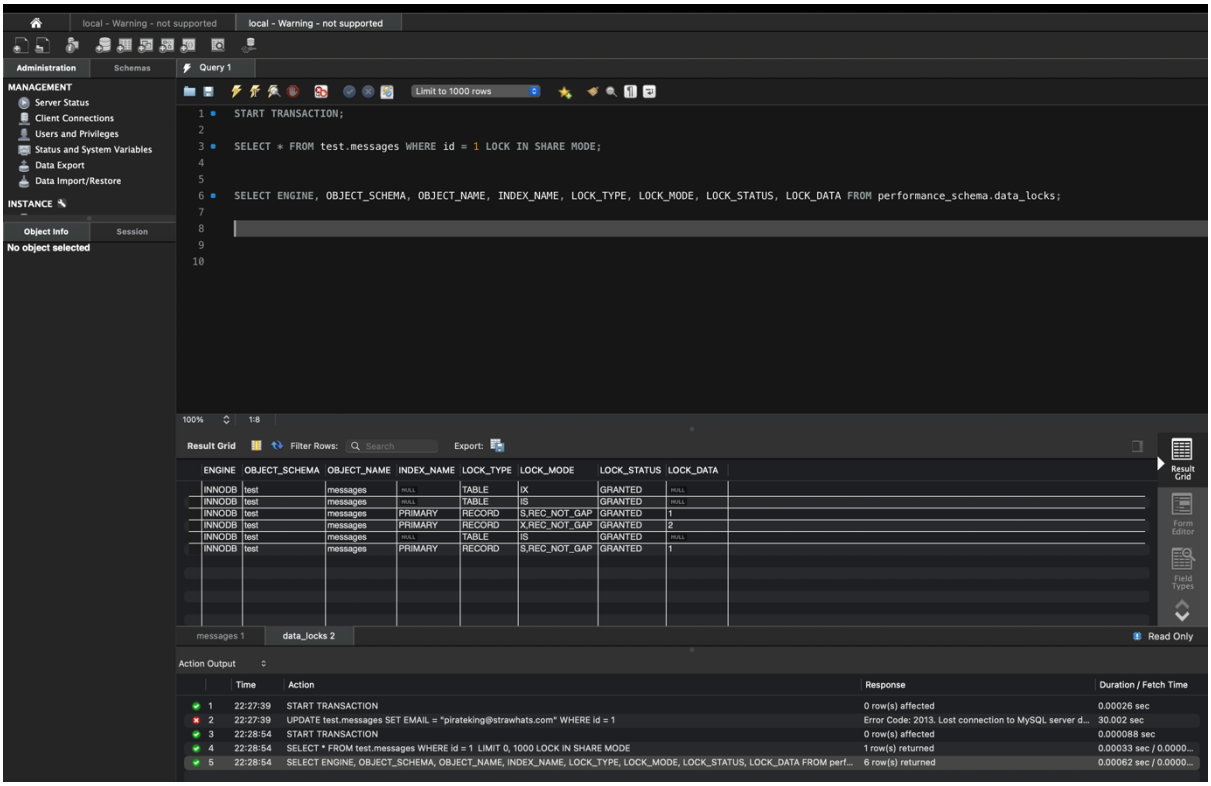


Click on your Database connection to open a new SQL editor tab which will open a new Session with your MySQL server. In your new session, start a transaction and try to update row with id = 1 and execute the queries.



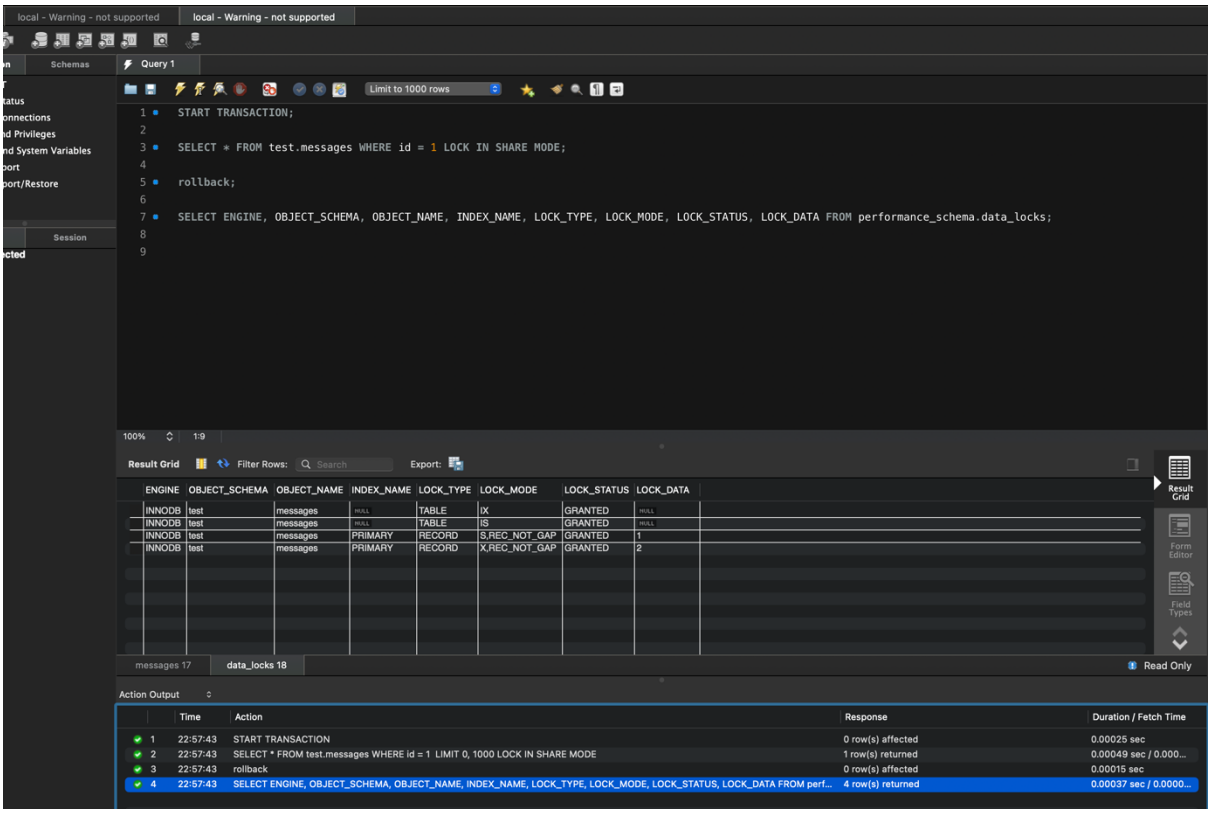
You will observe that the UPDATE query hangs and ends up with an error that the query cannot be executed. This occurs because the previous Session holds a S lock on row with id = 1, which is not compatible with an UPDATE query that tries to get an X lock on it.

However, if we try to get an IS lock on the test.messages table with a specific S lock on row with id = 1, then we will observe that this lock will be granted to the transaction in the new Session as follows:

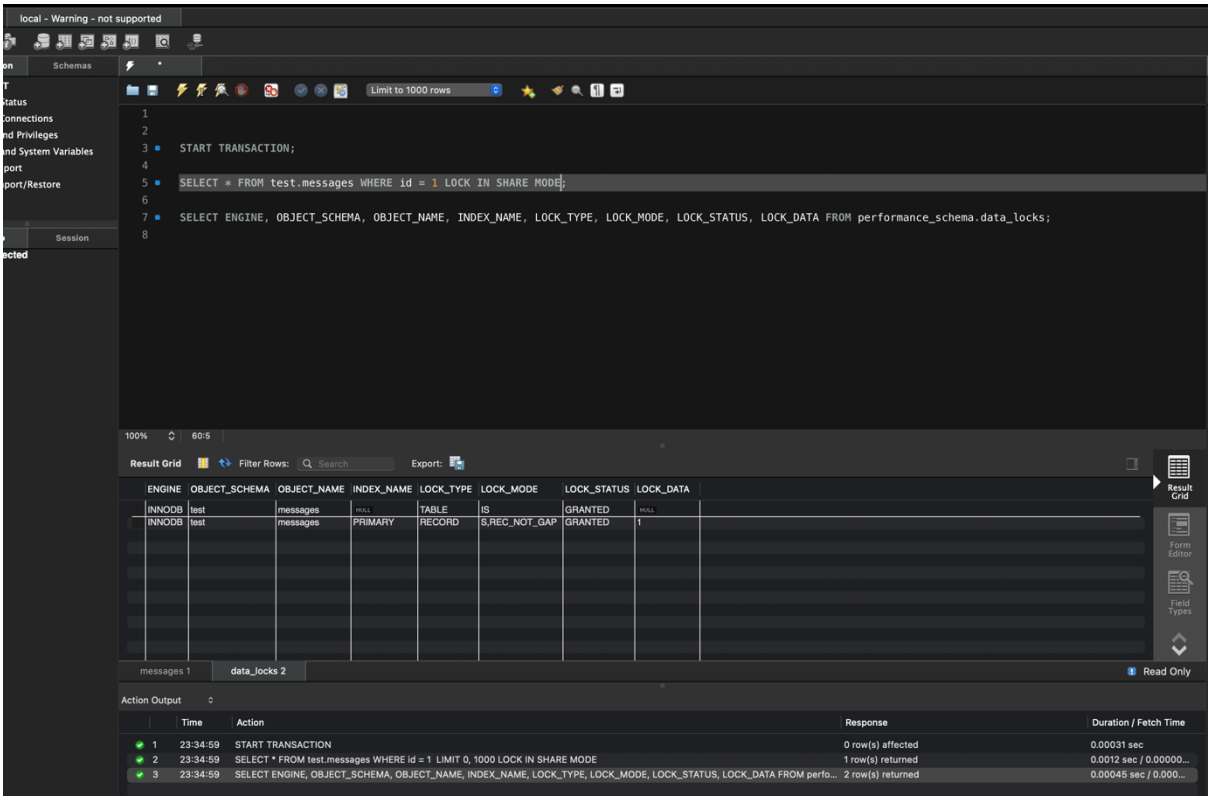


Since the read based Intention and specific locks are compatible, we observe that the new sessions transaction will be granted the IS and S lock.

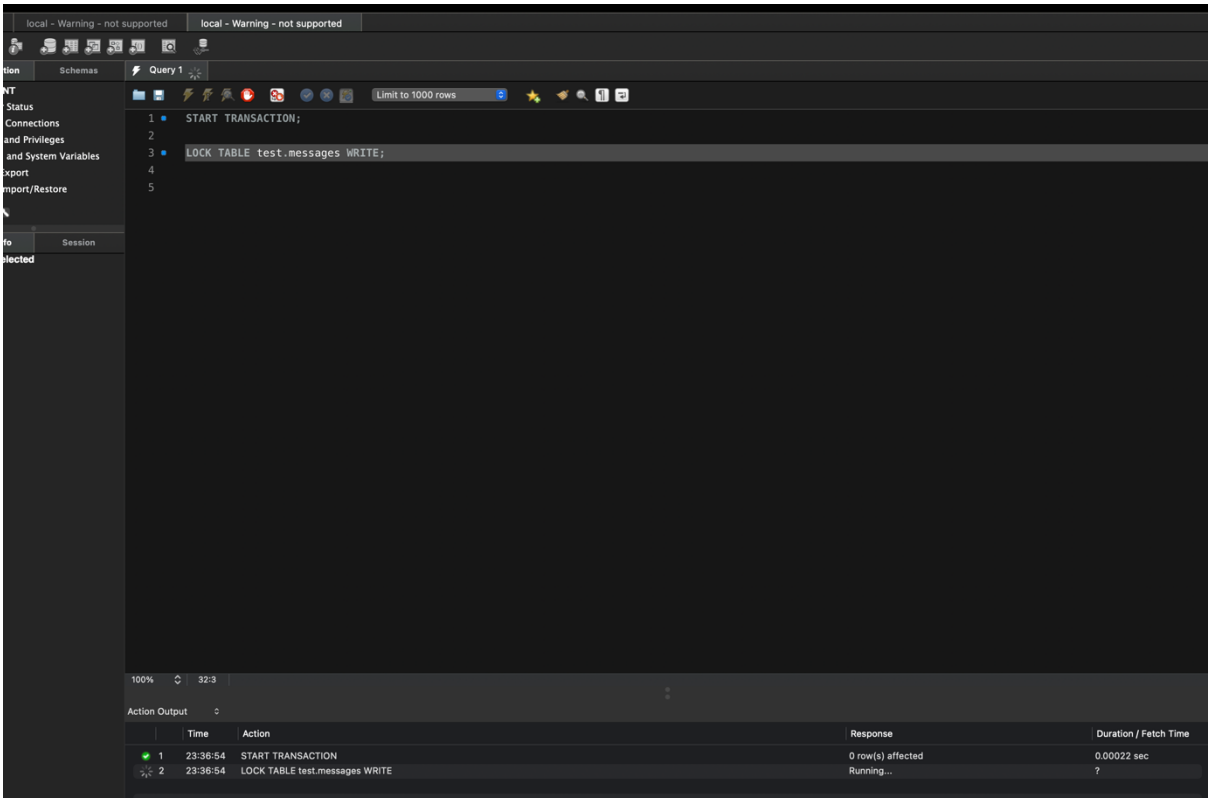
If we rollback the new transaction then we will observe that the IS and S locks held by this transaction will be released. Only the first session will be maintaining its locks now.



We can further understand whether an IS lock is compatible with an X lock by getting an IS lock on the test.messages table in Session 1 as shown below.



Following this, in Session 2 we try to get a table level X lock using the LOCK TABLE... WRITE query as seen in Lab 2a.



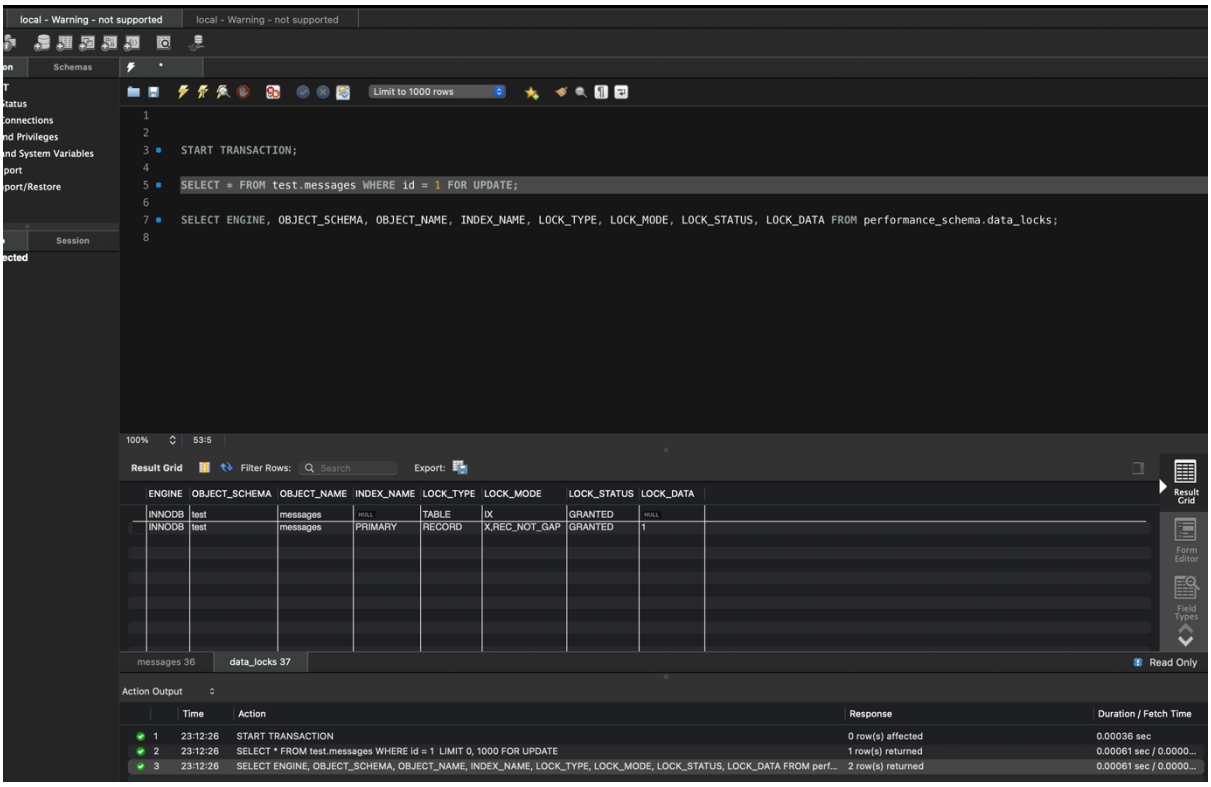
When we try to get the X lock on test.messages (currently under an IS lock by Session 1) from Session 2, we observe that the request for an X lock will be stuck waiting and will

eventually time out. This is because IS locks are not compatible with an X lock. If we try the reverse (trying to get an IS lock on a table with an X lock already), we will observe the same issue.

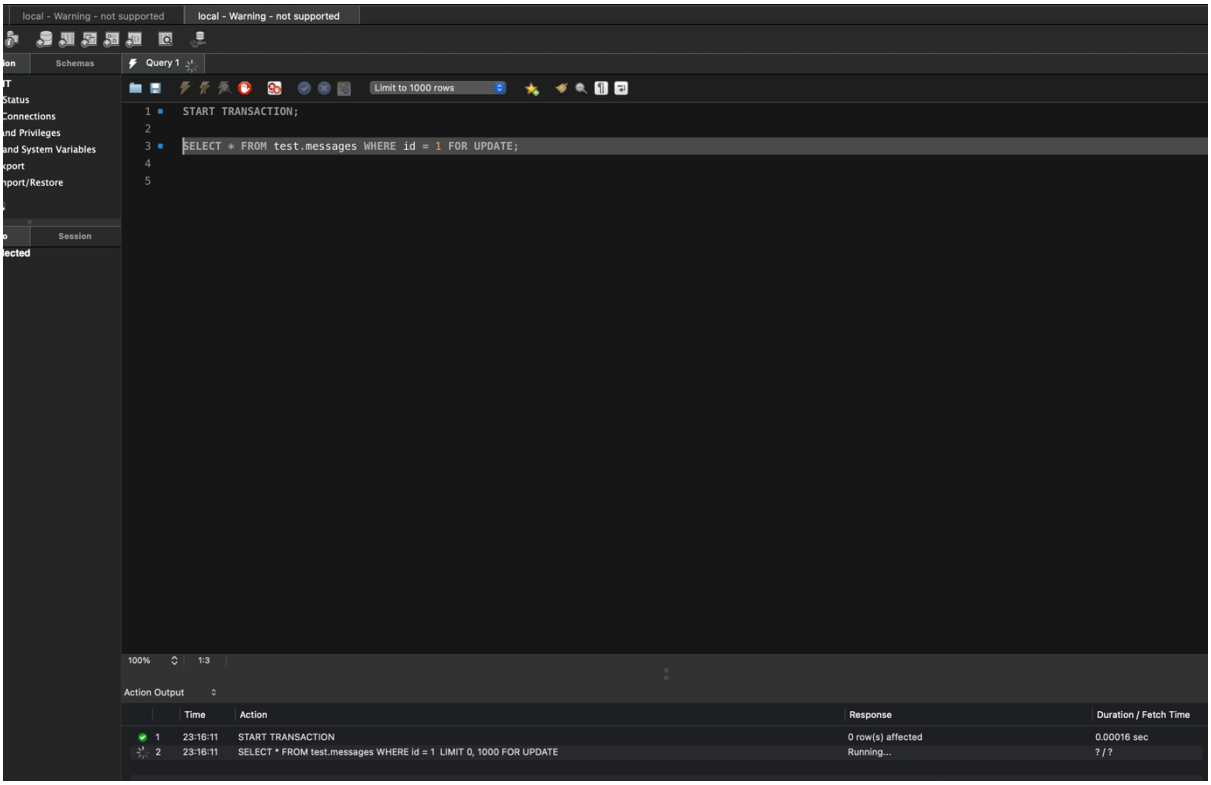
IX Locks:

Similarly, we can retrieve an IX lock explicitly using the format:

SELECT …. FOR UPDATE;

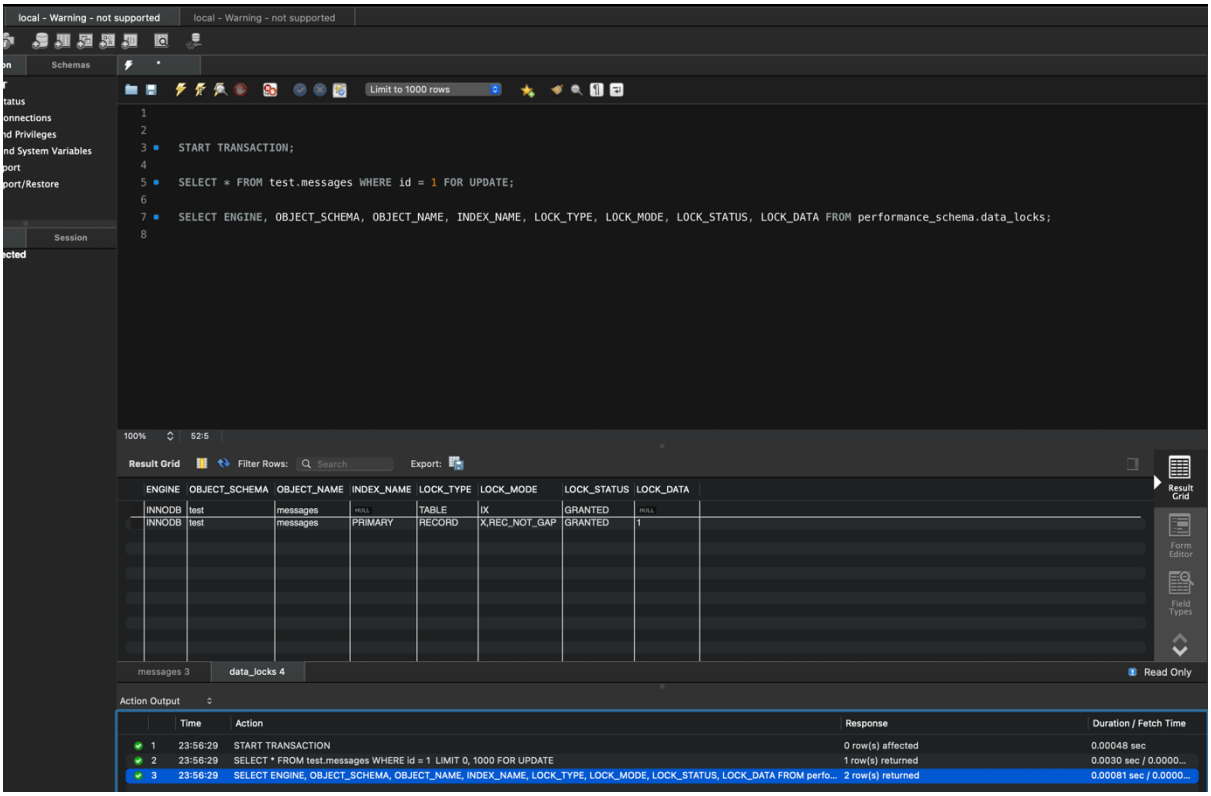


In the above sequence, we have acquired an IX lock on the test.messages table and a specific X lock on the row with id = 1 in the first session. If we open a new session and try to get an X lock on the row with id = 1, the X lock would be in the waiting state and eventually timeout since only 1 session would be allowed to have an X lock on a data object (in our case, a row of data) at one time. This can be seen in the below screenshot:

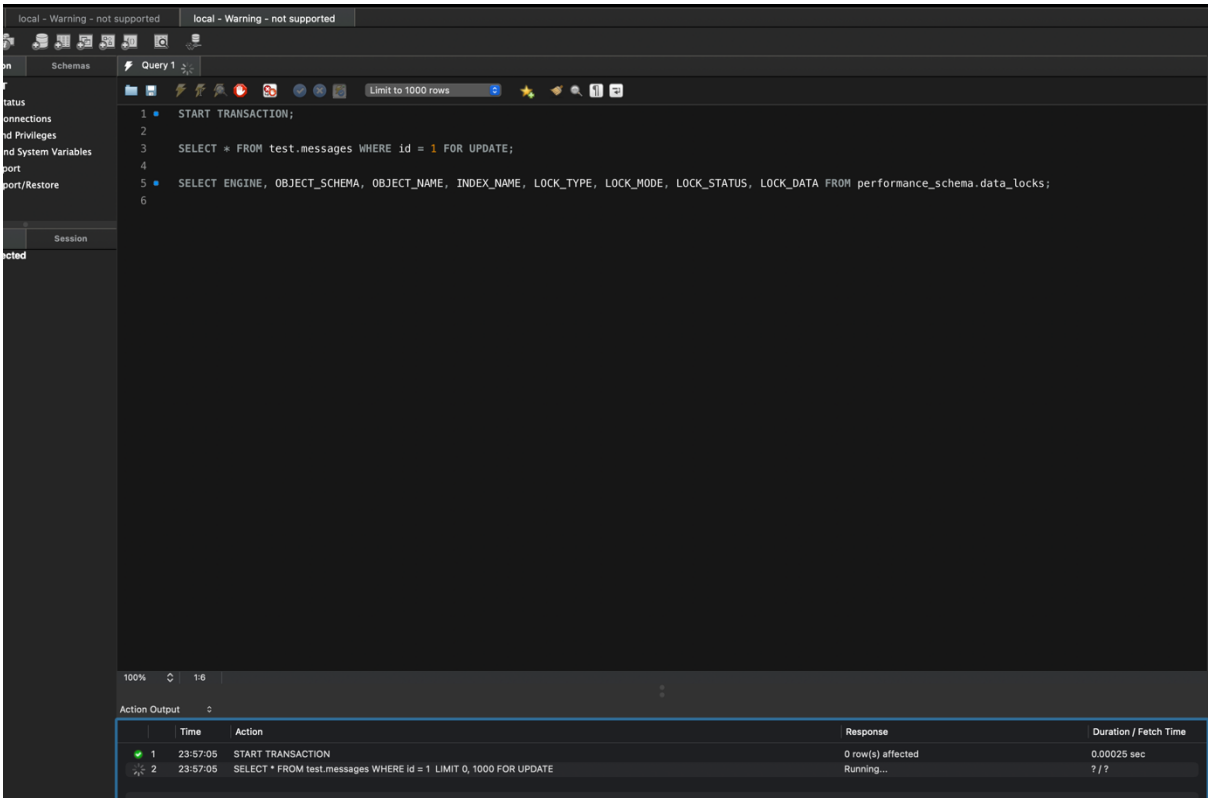


Similarly, if we try to get an IS lock on the same table with a specific S lock on the same row, the query will still timeout despite IX and IS locks being compatible. This is because S locks and X locks on the same resource are not compatible. However, if we try to get a lock on a separate row (id other than 1), the lock would be granted.

We can verify whether an IX lock is compatible with another IX lock by first getting an IX lock on the test.messages table from Session 1.

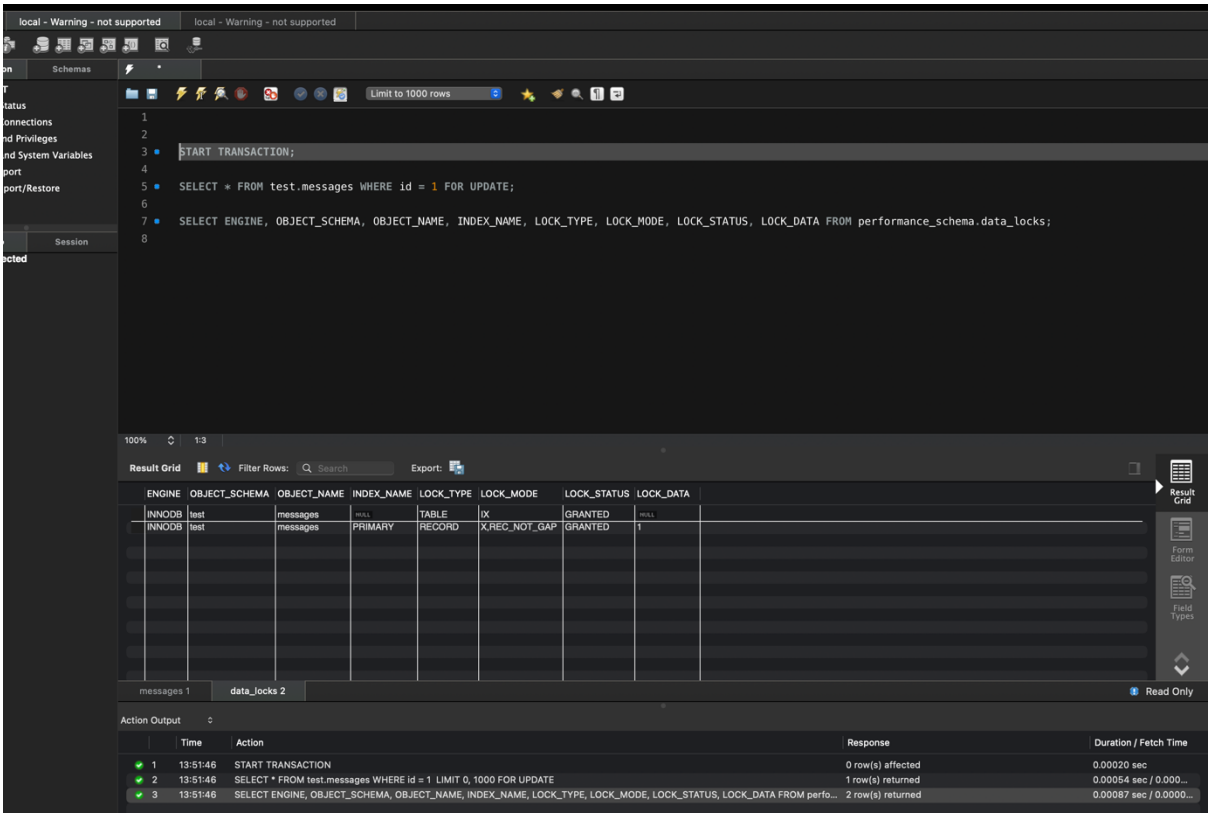


As seen in the output of the performance_schema.data_locks table, we have gotten an IX lock on the test.messages table from Session 1. Next, we try to get an IX lock on the test.messages table from Session 2 as shown below.

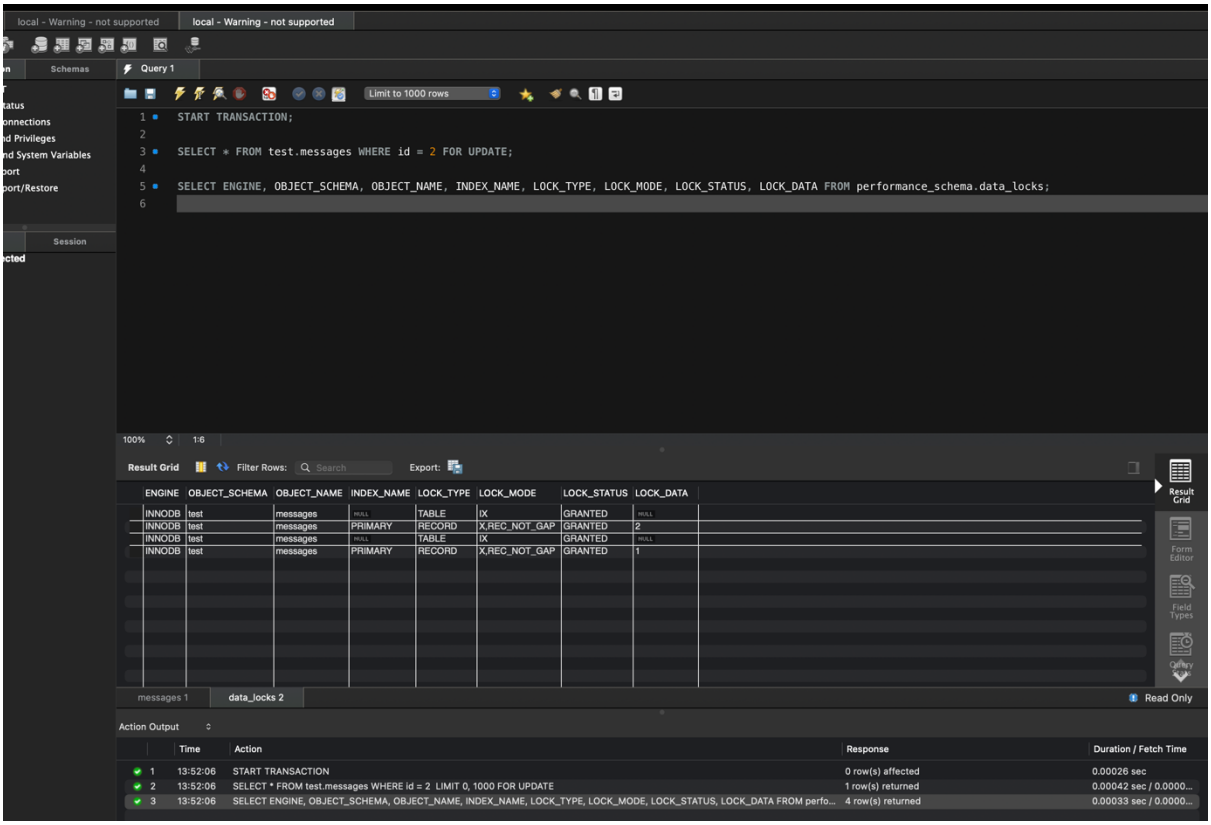


We can observe in the logs for Session 2 that the request for getting a specific X lock on the same row that the transaction in Session 1 holds will be blocked, even if the IX lock in Session 1 is compatible with the IX lock in Session 2 at the table level. This is because an X lock from one transaction is not compatible with an X lock requested by another transaction.

Next we can verify whether an IX of one transaction is compatible with an IX lock from another transaction with specific locks being granted at different database resources at lower level granularities. For this, we first retrieve a specific X lock on the row with id = 1 and an IX lock on the test.messages table as seen below.



In another session, we try to retrieve a specific X lock on a row with id = 2 with an IX lock on the test.messages table as seen below.



We can observe that both transactions are able to get a specific X lock on different rows with IX locks on the test.messages table. This is because IX locks are compatible with requests for an IX lock from other transactions.

There are other types of locks as well in MySQL which can be read about here: https://dev.mysql.com/doc/refman/8.0/en/innodb-locking.html

Try running some of the lock types described in the link above while running READ and WRITE based operations.