## **ANSWER**

Sequence 1 executes 2+1+2=5 instructions. Sequence 2 executes 4+1+1=6 instructions. Therefore, sequence 1 executes fewer instructions.

We can use the equation for CPU clock cycles based on instruction count and CPI to find the total number of clock cycles for each sequence:

CPU clock cycles = 
$$\sum_{i=1}^{n} (CPI_i \times C_i)$$

This yields

CPU clock cycles<sub>1</sub> = 
$$(2 \times 1) + (1 \times 2) + (2 \times 3) = 2 + 2 + 6 = 10$$
 cycles

CPU clock cycles<sub>2</sub> = 
$$(4 \times 1) + (1 \times 2) + (1 \times 3) = 4 + 2 + 3 = 9$$
 cycles

So code sequence 2 is faster, even though it executes one extra instruction. Since code sequence 2 takes fewer overall clock cycles but has more instructions, it must have a lower CPI. The CPI values can be computed by

$$CPI = \frac{CPU \text{ clock cycles}}{Instruction count}$$

$$CPI_1 = \frac{CPU \text{ clock cycles}_1}{Instruction count_1} = \frac{10}{5} = 2.0$$

$$CPI_2 = \frac{CPU \ clock \ cycles_2}{Instruction \ count_2} = \frac{9}{6} = 1.5$$

Figure 1.14 shows the basic measurements at different levels in the computer and what is being measured in each case. We can see how these factors are combined to yield execution time measured in seconds per program:

 $Time = Seconds/Program = \frac{Instructions}{Program} \times \frac{Clock\ cycles}{Instruction} \times \frac{Seconds}{Clock\ cycle}$ 

Always bear in mind that the only complete and reliable measure of computer performance is time. For example, changing the instruction set to lower the instruction count may lead to an organization with a slower clock cycle time or higher CPI that offsets the improvement in instruction count. Similarly, because CPI depends on type of instructions executed, the code that executes the fewest number of instructions may not be the fastest.