MAT 311 Abstract Algebra

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Spring 2024

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1 Sets and Relations

Definition: What is Abstract Algebra

- Algebra: procedures for performing operations, i.e. $+, -, \times, \div$, and methods for solving equations. It uses bldspecific operations on **specific** objects.
- Abstract Algebra: discuss **general** structures and the relationships between the elements of these structures.

1.1 Sets

Definition: Set

A set is a collection of objects. These objects are called "elements". A set is typically uppercase, and elements are typically lowercase.

Set Notation

1. List Notation:

$$B = \{\text{John}, \text{Paul}, \text{Ringo}, \text{George}\}$$

$$\mathbb{N} = \{1, 2, 3, \dots\}$$

2. Set-builder Notation:

$$B = \{b : b \text{ is a Beatle}\}\$$

Well-Defined Sets

Sets must be **well-defined**. That is, given set S and any element x, either $x \in S$ or $x \notin S$.

Definition: Subset

A set A is a subset of set B, written as $A \subseteq B$, if every element of A is also in B. Note: every non-empty set has at least two subsets:

- The set itself
- Ø

Definition: Proper Subset

If $A \subseteq B$ but $A \neq B$, then A is a **proper subset** of B, written $A \subset B$ or $A \subsetneq B$. Note: A set B is an *improper subset* of itself.

Definition: Cartesian Product

Let A and B be sets. The set $A \times B = \{(a,b) : a \in A \text{ and } b \in B\}$ is the cartesian product of A and B. Note: $A \times B = B \times A \iff A = B$, or $A \times B = \emptyset$.

Example

Let $A = \{c : c \text{ is a primary color}\}\$ and let $B = \{\epsilon, \delta\}$. Find:

1.
$$B \times B = \{(\epsilon, \epsilon), (\epsilon, \delta), (\delta, \epsilon), (\delta, \delta)\}$$

2.
$$A \times \emptyset = \emptyset$$

1.2 Relations

Definition: Relation

A **relation** between sets A and B is a subset \mathcal{R} of $A \times B$. It is a collection of ordered pairs. Note: $(a,b) \in \mathcal{R} \equiv a\mathcal{R}b$ means "a is related to b".

Definition: Function

A function is a relation in which no two of the ordered pairs have the same first term. Note: if $f : \mathbb{R} \to \mathbb{R}$ is a function, then is passes the vertical-line test.

Definition: One-to-One

A function is **one-to-one**, or **injective**, if no two ordered pairs have the same <u>second</u> term.

To prove f is one-to-one, first assume that $f(x_1) = f(x_2)$, then show that $x_1 = x_2$.

Definition: Onto

A function $f: X \to Y$ is **onto**, or **surjective**, if the codomain is equal to the range, meaning every element $y \in Y$ has some $x \in X$ such that f(x) = y.

Definition: One-to-One Correspondence

A function $f: X \to Y$ is a **one-to-one correspondence**, or a **bijection**, if it is both one-to-one and onto.

1.3 Partitions and Equivalence Relations

Definition: Partition

A partition of a set S is a collection of non-empty subsets of S such that:

- 1. The union of these subsets is S.
- 2. These subsets are pairwise disjoint.

Note: these subsets are called **cells** of the partition.

Definition: Equivalence Relation

An equivalence relation \mathcal{R} on a set S must be:

- 1. Reflexive, meaning $x\mathcal{R}x \quad \forall x \in S$.
- 2. Symmetric, meaning if $x\mathcal{R}y$, then $y\mathcal{R}x$.
- 3. Transitive, meaning if $x\mathcal{R}y$ and $y\mathcal{R}z$, then $x\mathcal{R}z$.

Definition: Equivalence Class

 $\overline{x} = \{y \in S : x \mathcal{R} y\}$ is the equivalence class of x

Example

Let $S = \mathbb{R}$. Define $x\mathcal{R}y$ iff $x \geq y$. Is \mathcal{R} an equivalence relation on S?

- 1. Is \mathcal{R} reflexive? $\forall x \in S, x\mathcal{R}x$, so YES.
- 2. Is \mathcal{R} symmetric? Consider 5 and 1: $5 \ge 1$ but $1 \not\ge 5$, so NO.
- 3. Is \mathcal{R} transitive? If $x \geq y$ and $y \geq z$ then $x \geq z$, so YES.

Since \mathcal{R} is not symmetric, it is not an equivalence relation on S.

Note on Partition Cells and Equivalence Classes

Partitions give rise to equivalence relations and vice versa. The *cells* of the partition are analogous to the *equivalence classes* of the equivalence relation.

2 Binary Operations

Definition: Binary Operation

A binary operation * on a set S is a function from $S \times S$ into $S, *: S \times S \to S$. That is, * is a rule which assigns to each ordered pair $(a,b) \in S \times S$ exactly one element $a*b \in S$.

Condition 1: Uniquely Defined

For all $a, b \in S \times S$, a * b must be **uniquely defined**. This means that * cannot be undefined for any a * b, and each a * b must have exactly one result, not two or more.

Condition 2: Closed under *

S must be **closed** under *. That is,

 $\forall a, b \in S, \qquad a * b \in S.$

Definition: Commutative

A binary operation * on a set S is commutative if

 $\forall a, b \in S, \qquad a * b = b * a.$

Definition: Associative

A binary operation * on a set S is associative if

 $\forall a, b, c \in S, \qquad a * (b * c) = (a * b) * c.$

2.1 Finite Sets

Example

Let $S = \{a, b, c, d\}$. Define a binary operation * on S using the following table. Complete the table so that * is commutative.

Note: * is commutative iff the table is symmetric along the main diagonal.

Is * associative? Why or why not? **No**,

$$a * (b * c) = a * c = a$$

 $(a * b) * c) = d * c = b$

Example

Suppose that * is associative and commutative operation on a set S. Show that $H = \{a \in S : a * a = a\}$ is closed under *. Note that the elements of H are called **idenmptents** of the binary operation *.

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Proof. Let $a, b \in H$. Show $a * b \in H$.

We know a * a = a and b * b = b. Show (a * b) * (a * b) = a * b.

$$LHS = (a*b)*(a*b)$$

$$= a*(b*a)*b$$

$$= a*(a*b)*b$$

$$= (a*a)*(b*b)$$

$$= a*b$$

$$= RHS$$
since * is associative
since * is associative

Thus, H is closed under *.

3 Isomorphic Binary Structures

Definition: Binary Algebraic Structure

A binary algebraic structure $\langle S, * \rangle$ is a set S together with a binary operation *.

Definition: Isomorphism

Let < S, *> and < S', *' be binary structures. An **isomorphism** of S with S' is a *one-to-one* function $\phi: S \mapsto S'$ such that

$$\forall x, y \in S, \qquad \phi(x * y) = \phi(x) *' \phi(y).$$

Notation: $\langle S, * \rangle \simeq \langle S', *' \rangle$