

# Van Rossum Distance with Optimal Lag

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## 1 Introduction

A spike train is a sequence of times representing neural spikes recorded over a short time interval. It can be important to determine whether two recorded trains are “the same”, perhaps meaning responses to the same stimulus. A number of algorithms have been proposed for this problem; the most biologically motivated is the van Rossum distance (van Rossum 2001), which might represent the difference in a synapse’s response to two spike trains.

The standard van Rossum distance is defined as follows. Let  $S = \{s_i\}, i = 1, \dots, M$  and  $T = \{t_j\}, j = 1, \dots, N$  be the increasing times of two spike trains and let  $\tau$  be the characteristic time for the response to a spike. Assume the responses sum in time. Let  $\mathbf{1}()$  be the indicator function (0 or 1 as the argument is false or true), then the responses are idealized as follows. The van Rossum transform of the spike train  $S = \{s_i\}$  is a real function  $\mathfrak{R}_{\tau,S}$  defined by Equation 1 below, where  $\tau$  is a positive parameter.

$$\mathfrak{R}_{\tau,S}(u) = \sum_{i=1}^M e^{-(u-s_i)/\tau} \mathbf{1}(u > s_i) \quad (1)$$

This expression decays exponentially in  $u$ , so is in  $L^p$  for any  $p$ . The van Rossum distance,  $D$ , is the  $L^2$  distance between  $\mathfrak{R}_S$  and  $\mathfrak{R}_T$ .

$$D^2 = \int_{-\infty}^{\infty} (\mathfrak{R}_{\tau,S}(u) - \mathfrak{R}_{\tau,T}(u))^2 du \quad (2)$$

The distance may be computed with linear complexity (Houghman and Kreuz, 2012).

Often the trains  $S$  and  $T$  are not directly comparable because of different time bases, so the lagged distance is more useful.<sup>1</sup>

$$D^2(c) = \int (\mathfrak{R}_S(u) - \mathfrak{R}_T(u - c))^2 du \quad (3)$$

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<sup>1</sup>In this document the parameter  $\tau$  is usually omitted and left to context. Unless specified, the limits of integration are always  $-\infty$  and  $\infty$ .

Then the optimal lag,  $c_{opt}$ , minimizes  $D(c)$  and the optimal distance is  $D(c_{opt})$ . This document develops a new algorithm for finding the exact value of  $c_{opt}$  and for computing  $D(c_{opt})$ .

## 1.1 Why use van Rossum at all?

Several other spike train distances have been proposed and the van Rossum distance is not always the recommended alternative. For example, a recent article (Satuvuori and Kreuz (2018)) looks at five distance algorithms: Victor-Purpura, van Rossum, ISI, SPIKE and RI-SPIKE. The authors distinguish between rate coding and temporal coding, and find that the first two (“spike-resolved”) perform poorly at comparing temporal codes except at very low spike rates. For the Van Rossum distance, “low spike rates” means values of  $\tau$  much less than the mean inter-spike time, so the van Rossum transform is essentially a sum of delta functions - not so attractive. Using the same implementation of van Rossum, Cutts and Eglen (2014) come to a similar conclusion.

On the positive side, the van Rossum transform is exactly the Single Exponential Model for post-synaptic current (Yamazaki et. al., 2022), so is directly relevant to simulations using that model. Also, the ability to calculate an exact lag value has independent interest. The lag based on any metric can of course be approximated, but perhaps only with significant implementation complexity, performance cost, or loss of accuracy.

The deficiencies noted by Satuvuori and Kreuz can in part be overcome by normalization. If one considers two spike trains, S and T, very similar except that T has twice the spike rate of S, then the van Rossum distance of Equation 2 is insensitive to the detailed shapes of the spike trains simply because  $\mathfrak{R}_T \approx 2\mathfrak{R}_S$ . If one wishes a distance insensitive to rate but sensitive to time information, one can normalize the van Rossum distance.

$$D_{norm}^2 = \int (\mathfrak{R}_S/M - \mathfrak{R}_T/N)^2$$

## 2 The Optimal Lag

This section proves that there are only finitely many possible choices for  $c_{opt}$  and develops an efficient algorithm for finding it. Without loss of generality,  $\tau = 1$  is assumed.

### 2.1 The Theorem

The main theorem asserts that there are only a finite number of possibilities for  $c_{opt}$ , obtained by matching a particular  $s_i$  with a particular  $t_j$ . Some formulas developed during the proof of the theorem form the basis for the algorithms.

#### 2.1.1 Lemma 1

Let  $\mathfrak{R}_S()$  and  $\mathfrak{R}_T()$  be as above, then

$$\int \Re_S(u) \Re_T(u) du = 1/2 \sum_{i,j} e^{-|s_i - t_j|} \quad (4)$$

Proof

Multiply the product of sums term by term:

$$\Re_S(u) \Re_T(u) = \sum_{i,j} e^{-(2u - s_i - t_j)} \mathbf{1}(u > \max(s_i, t_j))$$

Take the integral inside the sum and use:

$$\int_{t=a}^{\infty} e^{-2(u-b)} du = 1/2 e^{-2(a-b)}$$

Substitute  $b = (s_i + t_j)/2$  and  $a = \max(s_i, t_j)$ , and use  $\max(s, t) - (s+t)/2 = |s - t|/2$ :

$$\int \Re_S(u) \Re_T(u) du = 1/2 \sum_{i,j} e^{-|s_i - t_j|} \square$$

### 2.1.2 Corollary

$$\int \Re_S^2 = \sum_{1 \leq i \leq j \leq N} e^{s_i - s_j} - N/2 \quad (5)$$

Proof

Substitute  $\Re_S$  for  $\Re_T$  in Equation 4 and rearrange terms.  $\square$

### 2.1.3 Lemma 2

Let  $X$  be a finite set of reals, then  $\sum_{x \in X} e^{-|x-c|}$  is maximized over all reals by some  $c \in X$ .

Proof

Consider  $F(c) = \sum_{i=1}^N e^{-|x_i - c|}$ , where  $\{x_1, \dots, x_N\} = X$ .

$F()$  is positive, continuous and tends to 0 at  $\pm\infty$ , so it must have a maximum.

For any  $c$  not in  $X$ ,  $F''(c) = F(c) > 0$ . So no such  $c$  can be a maximum.

The result follows.  $\square$

### 2.1.4 Theorem

The minimum for  $D(c)$  is taken on by some  $c$  of the form  $s_i - t_j$ .

Proof:

We can write

$$D^2(c_{opt}) = \int \Re_S^2 + \int \Re_T^2 - 2 \max(c) \int \Re_S(u) \Re_T(u - c) du \quad (6)$$

so  $c_{opt} = \operatorname{argmax}(c) \int \Re_S(u) \Re_T(u - c) du$ , which by Lemma 1 is  $\operatorname{argmax}(c) \sum e^{-|s_i - t_j - c|}$ , so Lemma 2 shows  $c_{opt} = s_i - t_j$  for some  $i$  and  $j$ .  $\square$

## 2.2 The Algorithms

This section develops recursive algorithms to calculate the terms in Equation 6.

### 2.2.1 Algorithm 1 - VRDnormSQ

The goal is to calculate  $\int \mathfrak{R}_S^2$  using the representation in Equation 5. In order to calculate the sum, create a recursion for a numerically accurate calculation.

*Algorithm VRDnormSQ*  
Input:  $S$  - vector of increasing times  
Output:  $\int \mathfrak{R}_S(u)^2 du$   
Step 1. Compute  $A_k = \sum_{i|i \leq k} e^{s_i - s_k}$  for  $k = 1, \dots, N$  by the recursion:  
 $A_1 = 1, A_{k+1} = 1 + A_k e^{s_k - s_{k+1}}$   
Step 2. Return  $\sum_{k=1}^N A_k - N/2$

### 2.2.2 Algorithm 2 - VRDcorr

The goal is to calculate  $\max(c) \int \mathfrak{R}_S(u) \mathfrak{R}_T(u - c) du$ , together with the maximizing  $c$ , using the representation in Equation 4. First we sort the set of all differences between  $s$  and  $t$ . This defines the nondecreasing sequence  $X = x_1, \dots, x_{MN} = \text{sort}(\{s_i - t_j\})$ . Equation 4 becomes

$$2 \int \mathfrak{R}_S(u) \mathfrak{R}_T(u - c) dt = \sum_{x \in X} e^{-|x - c|}$$

Since  $c \in X$ , this can be written

$$2 \int \mathfrak{R}_S(u) \mathfrak{R}_T(u - c) du = \sum_{x|x < c} e^{x - c} + 1 + \sum_{x|x > c} e^{c - x} \quad (7)$$

If we define

$$A_k = \sum_{i|i \leq k} e^{x_i - x_k}, B_k = \sum_{i|i \geq k} e^{x_k - x_i}$$

it is sufficient to maximize  $A_k + B_k$ , then  $c_{opt} = x_{k_{max}}$ .

This is essentially a quadratic algorithm (really  $O(MN \log(M + N))$  because of the sort.) So it is more expensive than the fast algorithm without a lag (linear), but if you need the lag, it is less expensive than trying out all  $MN$  possible lags (cubic.)

*Algorithm VRDcorr*  
Inputs: Two increasing sequences,  $S$  and  $T$ .  
Outputs: The optimal lag,  $C$ , and the correlation  $\text{Corr} = \int \mathfrak{R}_S(u) \mathfrak{R}_T(u - C)$   
Step 1: Compute  $\{x_k\} = \text{sort}(\{s_i - t_j | i \leq M, j \leq N\})$   
Step 2: Compute  $A_k = \sum_{i=1}^k e^{x_i - x_k}$  for  $k=1, \dots, MN$  by the recursion:  
 $A_1 = 1, A_{k+1} = 1 + e^{x_k - x_{k+1}} A_k$   
Step 3: Compute  $B_k = \sum_{i=k}^{MN} e^{x_k - x_i}$  for  $k = MN, \dots, 1$  by the recursion:  
 $B_{MN} = 1, B_{k-1} = 1 + e^{x_{k-1} - x_k} B_k$   
Step 4: Compute  $k_{max} = \text{argmax}(k) A_k + B_k$   
 $C = x_{k_{max}}$   
 $\text{Corr} = (A_{k_{max}} + B_{k_{max}} - 1)/2$

### 2.2.3 Algorithm 3 - VRDfastCorr

The algorithm VRDcorr automatically applies the optimal lag, but it can be useful to compute a correlation without lag, especially if that computation is much faster. Algorithm VRDfastCorr uses essentially the same logic as Houghman and Kreuz (2012) to compute  $\int \mathfrak{R}_S(u)\mathfrak{R}_T(u)du$  in time  $O(M+N)$ .

The main idea is to precompute partial sums of exponentials in the  $t$  direction. With care to account for  $s_i == t_j$  only once, write Equation 4 as

$$\int \mathfrak{R}_S \mathfrak{R}_T = 1/2 \left( \sum_{i,j | s_i < t_j} e^{s_i - t_j} + \sum_{i,j | s_i \geq t_j} e^{t_j - s_i} \right) \triangleq 1/2(U + V) \quad (8)$$

Calculate  $U_i = \sum_{j | t_j > s_i} e^{s_i - t_j}$  and  $V_i = \sum_{j | t_j \leq s_i} e^{t_j - s_i}$  by the following recurrences.

*Algorithm VRDfastCorr*

Input: Two increasing sequences  $S$  and  $T$ .

Output: The correlation  $\int \mathfrak{R}_S \mathfrak{R}_T$

Step 1: Calculate  $U_k$  by the recurrence:

$$U_N = \sum_{j | t_j > s_N} e^{s_N - t_j}; U_{k-1} = e^{s_{k-1} - s_k} U_k + \sum_{j | s_{k-1} < t_j \leq s_k} e^{s_{k-1} - t_j}$$

Step 2: Similarly, calculate  $V_k$  by the recurrence:

$$V_1 = \sum_{j | t_j \leq s_1} e^{t_j - s_1}; V_{k+1} = e^{s_k - s_{k+1}} V_k + \sum_{j | s_k \leq t_j < s_{k+1}} e^{t_j - s_{k+1}}$$

Step 3: Return  $1/2 \sum (U_k + V_k)$

## 2.3 Implementation

A C++ implementation of these algorithms is available at <https://github.com/edmundbutler/vanrossumlag.git>.

The main entry point is `laggedVRD()` with the arguments:

```
bool laggedVRD(sIn, tIn, tau,
               sNorm, tNorm,
               corr, lag)
```

Inputs: `sIn`, `tIn` – the spike trains  
`tau` – the scale

Outputs: `sNorm`, `tNorm` – L2 norms of the transforms  
`corr` – the maximized correlation  
`lag` – the maximizing lag

The outputs may be used to calculate the lagged distance

$$D^2 = sNorm^2 + tNorm^2 - 2 * corr$$

or the correlation coefficient

$$CC = corr / (sNorm * tNorm)$$

Also, the normalized VR distance,  $D_{norm}^2 \triangleq \int (\Re_S(u)/M - \Re_T(u)/N)^2 du$  is given by

$$D_{norm}^2 = sNorm^2/M^2 + tNorm^2/N^2 - 2corr/(MN) \quad (9)$$

### 3 Simulation Experiments

Simulated spike trains were created to answer three questions:

1. Do the two correlation calculations get the same answer?
2. What is the CPU performance of the implementations?
3. Does the calculated value of  $c_{opt}$  depend strongly on the scale parameter  $\tau$ ?

The third question could be significant for the following reason. It turns out that the function VRDfastCorr() is much faster than VRDcorr(), so one would like to use VRDcorr() with a single value of  $\tau$  to calculate  $c_{opt}$  and then use VRDfastCorr() to explore the effect of changes to  $\tau$ ; this strategy requires a stable value for  $c_{opt}$ .

#### 3.1 The Simulation

The simulation creates base trains and noised trains. A base train consists of the times of a Poisson process with rate  $1/\gamma$ . A noised train takes a base train and modifies it in three ways, depending on small parameters  $\alpha$  and  $\beta$ .

1. Dropping out a proportion  $\alpha$  of the times.
2. Combining with the times of an independent Poisson process with rate  $\alpha/\gamma$ .
3. Adding noise to each time. The noises are IID centered Uniform random variables with width  $\beta\gamma$ .

Step 2 adds about the same number of noise spikes as is dropped in Step 1. The net effect is that both base and noised trains have average inter-spike time  $\tau$ . (So this simulation avoids rate-discrimination effects, to be addressed in a follow-on study.)

The values  $\alpha = 0.1$ , and  $\beta = 0.03$  are always used in this document. Figure 1 shows 50 noised trains using the same base train with  $\gamma = 0.025$ . (The value  $\gamma = 1$  is used for the rest of this document.)

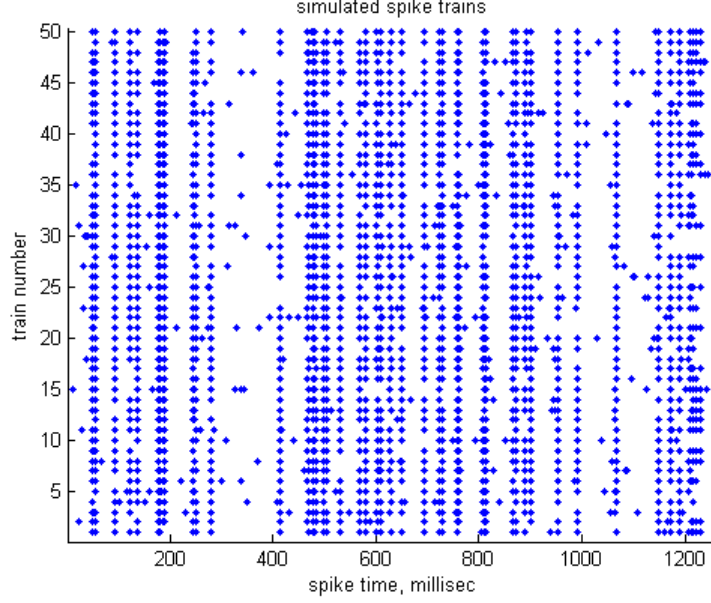


Figure 1 - Simulated spike trains, based on a 40 Hz original ( $\gamma = 0.025$ ).

### 3.2 VRDcorr() and VRDfastCorr() Compute the Same Correlation

The two algorithms VRDcorr() and VRDfastCorr() compute correlations using different algorithms. The outputs may be compared as follows.

Let  $S$  and  $T$  be spike trains. VRDcorr() computes  $c_{opt}$  and  $Corr = VRDcorr(S, T) = \int \Re_S(u) \Re_T(u - c_{opt}) du$ . Then translate the spike train  $T$ ,  $t_{opt}[i] = t[i] - c_{opt}$ . A second version of the correlation is given by  $Corr2 = VRDfastCorr(S, T_{opt})$ .

The above procedure is done 10,000 times each for train lengths between 10 and 1,000. At each train length a Root Mean Square error is computed, based on  $Err = Corr - Corr2$ . The results are plotted in Figure 2. The increase in error is partly due to the larger number of terms in Equation 1, but in any case the error is negligible ( $\leq 2 \times 10^{-12}$ ).

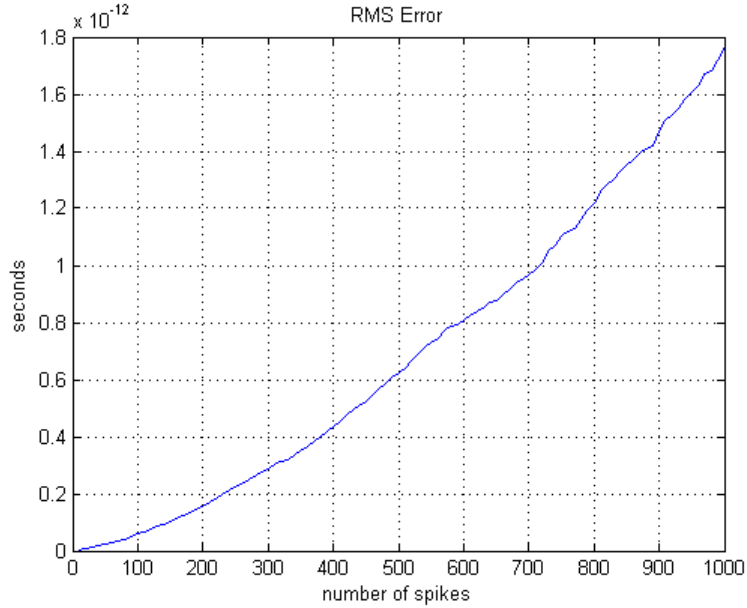


Figure 2. Errors in computation of correlations by `VRDcorr()` and `VRDfastCorr()`

### 3.3 CPU performance

Simulations were run under Windows 11 using an Intel 11th generation I9 processor. Executable were built using Microsoft Visual C++ 2022. (These performance times cannot be directly compared with older measurements because the support is so much better now.)

The simulation described in section 3.2 was instrumented to determine CPU timing. Predicted performance was that `VRDcorr()`'s CPU usage would grow approximately as  $N^2$ , while `VRDfastCorr()` should grow as  $N$ , where  $N$  is the number of spikes in a train. Figures 3 and 4 show the expected behavior.

More significantly, `VRDfastCorr()` is very fast, requiring only tens of microseconds for quite long trains. `VRDcorr()` is much slower; Figure 5 shows the ratio. `VRDcorr()` requires tens of milliseconds; as a practical matter, that speed would be adequate for many uses.

How expensive is the sort in `VRDcorr()`? The current implementation uses the sort from the C++ standard library. As displayed in Figure 6, this sort uses rather more than half of the CPU time. Most likely that could be improved by taking advantage of the special structure of the data.



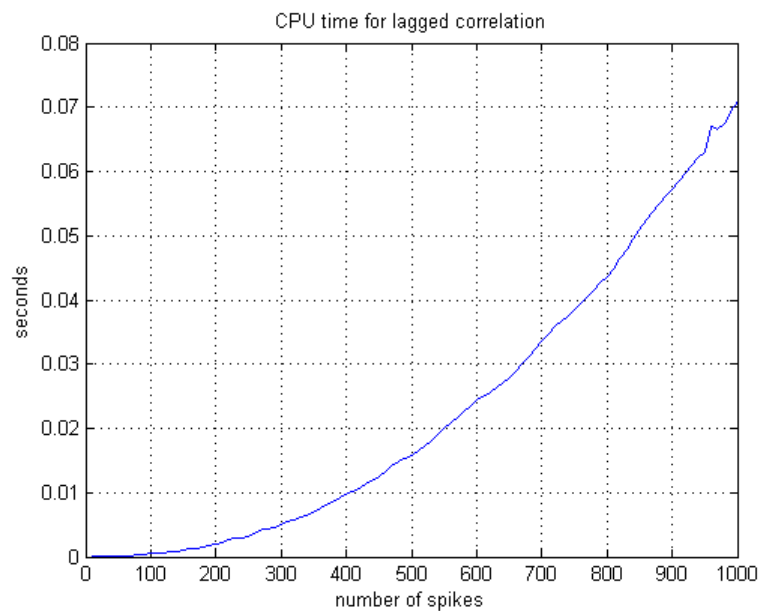


Figure 3. CPU time for a call to `VRDcorr()` as a function of the number of spikes

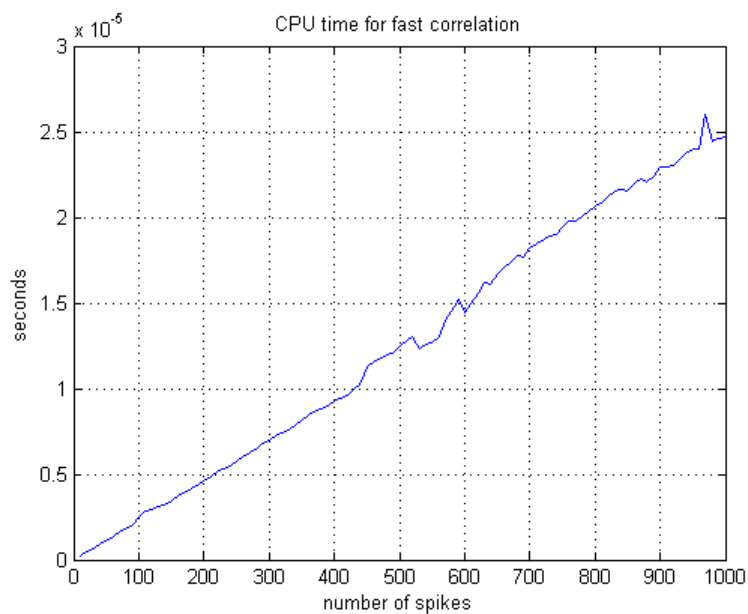


Figure 4. CPU time for a call to `VRDfastCorr()` as a function of the number of spikes

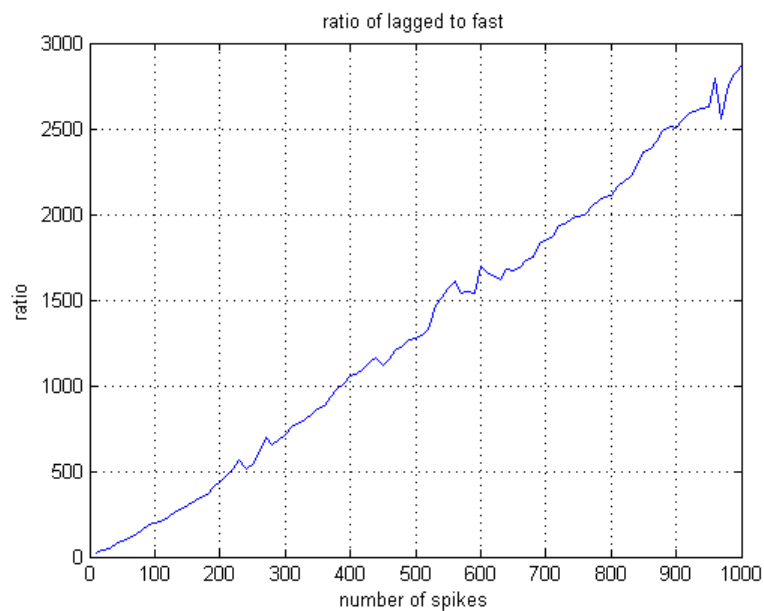


Figure 5. Ratio of CPU times: VRDcorr/VRDfastCorr

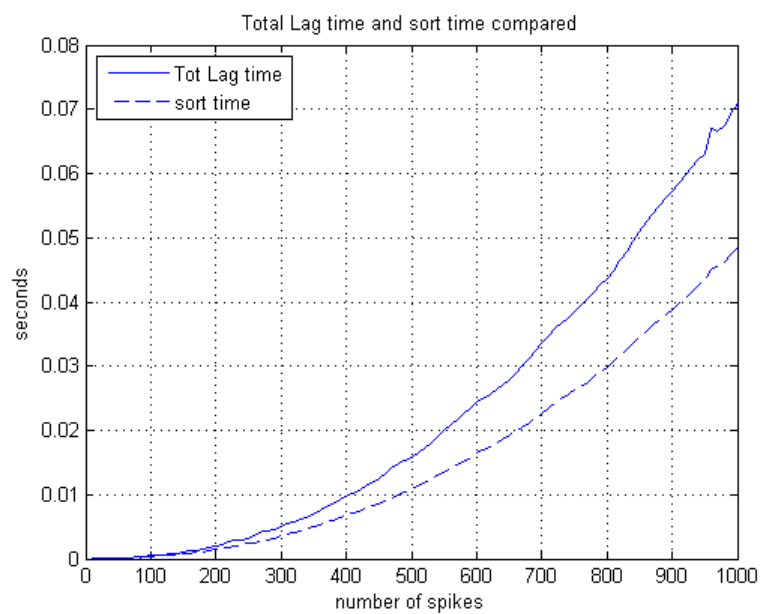


Figure 6. Comparing sort CPU time with total CPU time for VRDcorr

### 3.4 Dependence of Lag on Scale

How meaningful is the lag  $c_{opt}$ ? When there is a known correct value for the lag,  $c_{opt}$  should well approximate that value. Also, the calculation of  $c_{opt}$  should not depend strongly on the value of  $\tau$  used in the computation.

A simulation's results suggests the algorithm VRDcorr meets these goals rather well. 100,000 simulation scenarios were generated, each containing a base train and a noised train with 100 spikes on the average. The VRDcorr algorithm calculated  $c_{opt}$  for the based/noised pair using 41 values for  $\tau$ :  $1.1^i$  for  $i = -20, \dots, 20$ , so  $\tau$  varied between 0.1486 and 6.7275. In each scenario, the set of 41  $c_{opt}$  values had mean and standard deviation computed.

The predicted value for  $c_{opt}$  is 0. The values for  $\tau$  were chosen to cover a reasonable range about the expected interspike time of 1. Note: For  $\tau$  much bigger than 6, the van Rossum transform approximates a sum of delta functions and the correlation degenerates due to noise in the individual spikes. Nothing interesting happens when  $\tau$  is less than 0.14.

Figures 7 and 8 below histogram these sets of means and standard deviations.

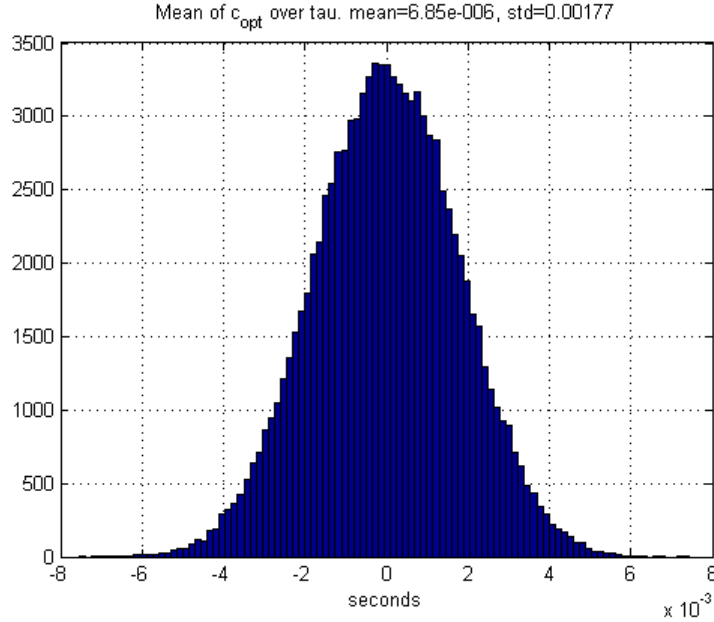


Figure 7. Histogram of calculated  $c_{opt}$  averaged over  $\tau$ .

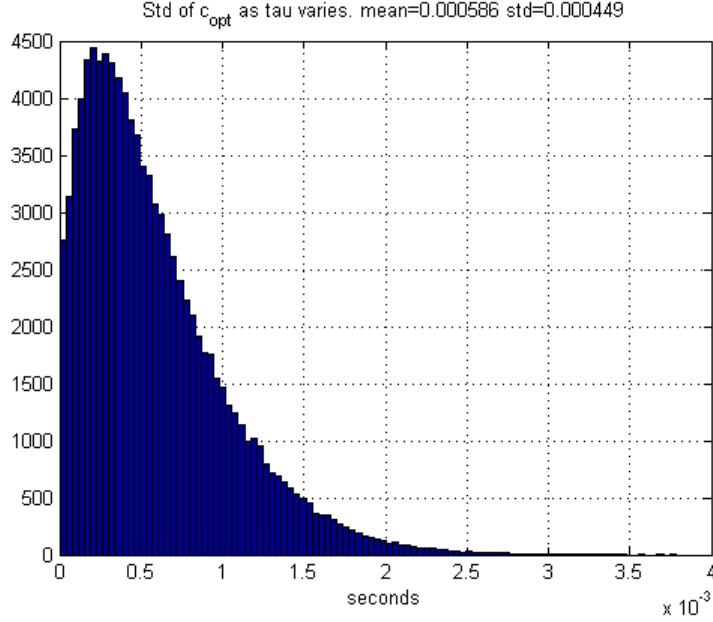


Figure 8. Histogram of variation in  $c_{opt}$  in  $\tau$

In short,  $c_{opt}$  is calculated correctly within about a millisecond. For any one scenario, the variation of  $c_{opt}$  with  $\tau$  tended to be less than half a millisecond. The errors are small compared to the average interspike time of one second, and probably are produced by step 3 in the noising process, which adds a random value of standard deviation 8.7 milliseconds to each spike time. The errors in  $c_{opt}$  are significantly less than the noise in individual spike times.

So, for this simulation varying  $\tau$  within a factor of 6 of the average interspike time, the  $c_{opt}$  calculation has mean error about  $10^{-3}$  of an average interspike time. The calculation bias (overall mean displayed in Figure 7) is about a  $10^{-5}$  of the interspike time. The error is not sensitive to  $\tau$ .

## 4 Conclusion and Further Work

This document develops a new algorithm for computing the optimal lag for von Rossum's distance based on a minor but surprising theorem. The computation is efficient, accurate, and natural. The algorithm selects a pair of spike times to compare; over a broad range of the parameter  $\tau$  it chooses well (in a simulation.)

But it remains to be shown that the algorithm can assist with the analysis of laboratory data. The next step would be the analysis of openly available datasets. Do the algorithms work? Can they lead to useful insights? Can they do anything not easy to accomplish with openly available tools?

## 4.1 Generalizations

This document develops only the standard von Rossum distance for simplicity of exposition, but the analysis can be extended in three useful ways.

Equation 1 defines the von Rossum transform as the convolution of a finite set with a kernel of the form  $k(u) = e^{-(u-s)/\tau} \mathbf{1}(s > u)$ , but any appropriate kernel  $k(u)$  would work. The main theorem also holds for weighted sums, so the more general transform is the following.

$$\mathfrak{S}_{S,W}(u) = \sum_{i=1}^M w_i k(u - s_i)$$

Under the assumptions that  $k(\cdot)$  supports a version of Lemma 1 and Lemma 2 and that  $w_i$  and  $v_j$  are nonnegative, then

$$c_{\text{opt}} = \operatorname{argmax}(c) \int (\mathfrak{S}_{S,W}(u) - \mathfrak{S}_{T,V}(u - c))^2 du$$

is taken on at one of a finite number of points.<sup>2</sup>

It seems possible to use these ideas to create a new distance similar to van Rossum’s but sensitive to spike intervals rather than just spike times (the distinction between  $D(\text{spike})$  and  $D(\text{interval})$  in Victor and Purpura (1998).)

Also,  $c_{\text{opt}} = \operatorname{argmin}(c) \int \mathfrak{S}_{S,W}(u) \mathfrak{S}_{T,V}(u - c) du$ , so variations on  $D$  are possible without affecting  $c_{\text{opt}}$ . For example, a correlation coefficient and a normalized distance can be useful.

$$D_{\text{norm}}(c) = \int (\mathfrak{R}_S(u)/M - \mathfrak{R}_T(u - c)/N)^2 du$$

It might be interesting to repeat some of the literature’s metric comparisons replacing the standard van Rossum metric with  $D_{\text{norm}}$  or a new interval-sensitive version.

## References

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<sup>2</sup>One example is  $k(u) = e^{-|u|/\tau}$ .

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