

### Q. 3

Using a sparse representation, each column of the matrix is a `TreeSet<Integer>` instance.

#### ★ Cost of the subtraction operation:

To subtract two columns, you subtract the two `TreeSet` by traversing the `TreeSet` to subtract, and, for each of its `Integer`, either the `Integer` is in the `TreeSet`, and therefore you remove it, either it's not, and you add it.

Looking, inserting or removing in a `TreeSet` is a  $O(\log n)$  cost operation.

But if we assume that both columns are sparse, then it is in  $O(1)$  for each insertion and there are  $O(1)$  insertion, so in the end:  $O(1)$  is the cost of subtraction

#### ★ How many subtractions?

There can be at most  $m$  subtractions to do:  $O(m)$

#### ★ How many lookups? (cost of a lookup)

Instead of looping over the columns, we use an `HashMap` from the lowest non null indices to the `TreeSet`.

Every lookup therefore costs  $O(1)$

#### ★ How many columns on which to perform these steps?

There are  $m$  of them:  $O(m)$

All in all the complexity is:  $O(m^2)$