



*Islamic University – Gaza*  
*Engineering Faculty*  
*Department of Computer Engineering*  
*ECOM 3010: Computer Architecture Discussion*



## Chapter 2

### Exercises with solutions

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## Discussion exercises

### Exercise 1:

Convert the following C statements to equivalent MIPS assembly language. Assume that the variables f, g, h and j are assigned to registers \$s0, \$s1, \$s2 and \$s3 respectively. Assume that the base address of the array A and B are in registers \$s6 and \$s7 respectively.

a)  $f = g + h + B[4]$

```
lw $t0, 16($s7)
add $s0, $s1, $s2
add $s0, $s0, $t0
```

b)  $f = g - A[B[4]]$

```
lw $t0, 16($s7)
sll $t1, $t0, 2
add $t2, $t1, $s6
lw $t3, 0($t2)
sub $s0, $s1, $t3
```

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### Exercise 2: (2.4 from book)

Why doesn't MIPS have

Since MIPS includes  $\pm 2^{15}$ , add immediate

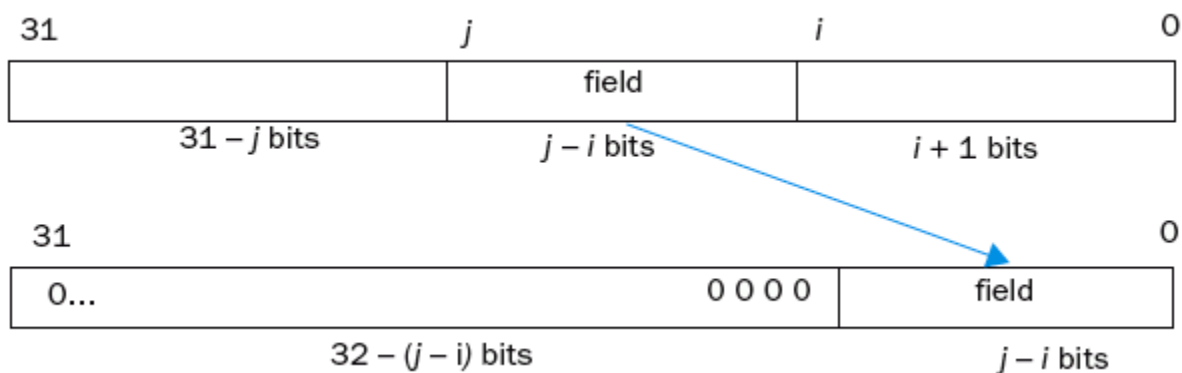
number, so subtract immediate would be redundant

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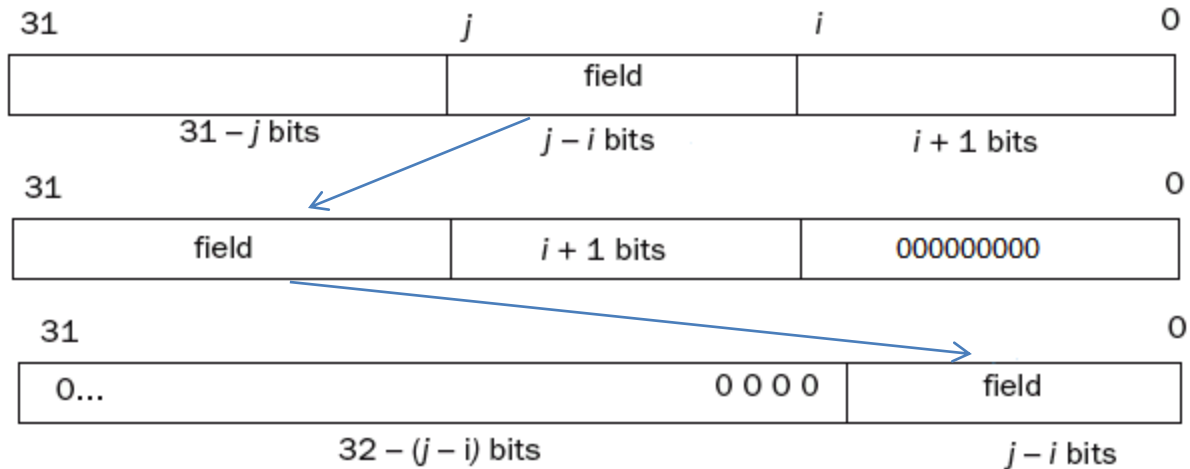
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### Exercise 3: (2.6 from book)

Find the shortest sequence of MIPS instructions that extracts a field for the constant values  $i=5$  and  $j=22$  from register \$t3 and places it in register \$t0.



```
sll $t0, $t3, 9
srl $t0, $t0, 15
```



➤  $31 - j = 31 - 22 = 9$   
`sll $t0, $t3, 9`  
 ➤  $i + 1 + 9 = 5 + 1 + 9 = 15$   
`srl $t0, $t0, 15`

Another solution:

➤  $i + 1 = 6$   
`srl $t0, $t3, 6`  
 ➤  $j - i = 22 - 5 = 17$   
`andi $t0, $t0,`

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#### Exercise 4: (2.32 from book)

Show the single MIPS instruction for this C statement

`b = 25 | a;`  
`ori $t1, $t0, 25`

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#### Exercise 5:

Convert the MIPS instruction to machine language:

`srl $s1, $t2, 3`

`srl` is R-type, opCode is 0 and function is 2

`$s1 = 17` is rd

`$t2 = 10` is rt

rs unused

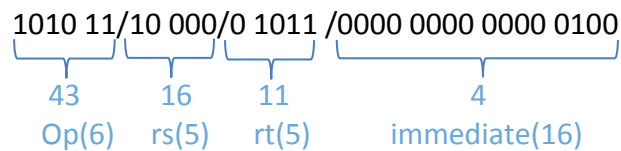
shamt is 3

0000 00/00 000/0 1010 /1000 1/000 11/00 0010 = 0x00A88C2

Op(6) rs(5) rt(5) rd(5) shamt(5) func(6)

**Exercise 6:**

Translate the following machine code to MIPS:



Op = 43 = sw

16 = \$s0

11 = \$t3

sw \$t3, 4(\$s0)

**Exercise 7: (2.37 from book)**

For each pseudoinstruction in the following table, produce a minimal sequence of actual MIPS instructions to accomplish the same thing. In the following table, *big* refers to a specific number that requires 32 bits to represent and *small* to a number that can fit in 16 bits.

Pseudoinstruction	What it accomplishes	Solution
move \$t1, \$t2		\$t1, \$t2, \$zero
clear \$t0		\$t0, \$zero, \$zero
beq \$t1, small		\$t0, \$zero, small 1, \$t0, L
beq \$t2, big, L	if (\$t2 == big) go to L	lui \$t1, upper(big) ori \$t1, \$t1, lower(big) beq \$t2, \$t1, L
li \$t1, small	\$t1 = small	addi \$t1, \$zero, small
li \$t2, big	\$t2 = big	lui \$t2, upper(big) ori \$t2, \$t2, lower(big)
ble \$t3, \$t5, L	if (\$t3 <= \$t5) go to L	slt \$t0, \$t5, \$t3 beq \$t0, \$zero, L
bgt \$t4, \$t5, L	if (\$t4 > \$t5) go to L	slt \$t0, \$t5, \$t4 bne \$t0, \$zero, L
bge \$t5, \$t3, L	if (\$t5 >= \$t3) go to L	slt \$t0, \$t5, \$t3 beq \$t0, \$zero, L
addi \$t0, \$t2, big	\$t0 = \$t2 + big	lui \$t1, upper(big) ori \$t1, \$t1, lower(big) add \$t0, \$t2, \$t1
lw \$t5, big(\$t2)	\$t5 = Memory[\$t2 + big]	lui \$t1, upper(big) ori \$t1, \$t1, lower(big) add \$t1, \$t1, \$t2 lw \$t5, 0(\$t1)

**Exercise 8:**

Convert the following C fragment to equivalent MIPS assembly language. Assume that the variables `a` and `b` are assigned to registers `$s0` and `$s1` respectively. Assume that the base address of the array `D` is in register `$s2`.

```
while(a < 10){
    D[a] = b + a;
    a += 1;
}
Loop: stli $t0, $s0, 10
      beq $t0, $zero, exit
      sll $t1, $s0, 2
      add $t1, $t1, $s2
      add $t2, $s1, $s0
      sw $t2, 0($t1)
      addi $s0, $s0, 1
      j Loop
exit:
```

**Exercise 9:**

Show the effects on memory and registers of the following instructions. Suppose a portion of memory contains the

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0x10000000	0
0x10000004	0

And register `$t0` contains `0x10000000` and `$s0` contains `0x01234567`. Assume each of the following instructions is executed independently of the others, starting with the values given above. *Hint: Don't forget that the MIPS architecture is Big-Endian.*

The memory:

Address	Data
0x10000000	0x12
0x10000001	0x34
0x10000002	0x56
0x10000003	0x78
0x10000004	0x9A
0x10000005	0xBC
0x10000006	0xDE
0x10000007	0xF0

- a) lw \$t1, 0(\$t0)  
\$t1 = 0x12345678
- b) lw \$t2, 4(\$t0)  
\$t2 = 0x9ABCDEF0
- c) lb \$t3, 0(\$t0)  
\$t3 = 0x00000012
- d) lb \$t4, 4(\$t0)  
\$t4 = 0xFFFFF9A → lb is sign extended
- e) lb \$t5, 3(\$t0)  
\$t5 = 0x00000078
- f) lh \$t6, 4(\$t0)  
\$t6 = 0xFFFF9ABC → lh is sign extended
- g) sw \$s0, 0(\$t0)  
at address 0x10000000 will contain 0x01234567

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Address	Data
0x10000000	0x12
0x10000004	
0x10000005	
0x10000006	
0x10000007	0xF0

- h) sb \$s0, 4(\$t0)  
the address 0x10000004 will contain 0x67BCDEF0

Address	Data
0x10000000	0x12
0x10000001	0x34
0x10000002	0x56
0x10000003	0x78
0x10000004	0x67
0x10000005	0xBC
0x10000006	0xDE
0x10000007	0xF0

i) `sb $s0, 7($t0)`

the address 0x10000004 will contain 0x9ABCDE67

Address	Data
0x10000000	0x12
0x10000001	0x34
0x10000002	0x56
0x10000003	0x78
0x10000004	0x9A
0x10000005	0xBC
0x10000006	0xDE
0x10000007	0x67

### Exercise 10:

Convert the following program into machine code.

```

0xFC00000C      start: .....
0xFC000010      loop: addi $t0, $t0, -1
0xFC000001
0xFC000001
0xFC000001
0xFC000001

```

```

addi $t0, $t0, -1
-1 = 0xFFFF

```

0010 00/01 000/0 1000 /1111 1111 1111 1111

8      8      8      -1

Op(6)   rs(5)   rt(5)   immediate(16)

```
sw $t0, 4($t2)
```

1010 11/01 010/0 1000 /0000 0000 0000 0100

43      10      8      4

Op(6)   rs(5)   rt(5)   immediate(16)

```
bne $t0, $t3, loop
```

target address = (immediate \* 4) + address of the following instruction

immediate = (target address – address of the following instruction) / 4

= (FC000010 – FC00001C) / 4

= - C / 4

= -3 → 1111 1111 1111 1101

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Or convert to binary first

= 1111 1100 0000 0000 0000 0000 0001 0000 –  
1111 1100 0000 0000 0000 0000 0001 1100

= 1111 1100 0000 0000 0000 0000 0001 0000 +  
0000 0011 1111 1111 1111 1111 1110 0100

= 1111 1111 1111 1111 1111 1111 1111 0100 / 4 → srl by 2  
= 11 1111 1111 1111 1111 1111 1111 1101

immediate is 16 only so immediate = 1111 1111 1111 1101

0001 01/01 000/0 1011 /1111 1111 1111 1101  
5 8 11 -3  
Op(6) rs(5) rt(5) immediate(16)

j start

target address = last

immediate = first 28

= C00000C / 4 = 3000003

= 1100 0000 0000 0000 0000 0000 1100 / 4 → srl b

= 11 0000 0000 0000 0000 0000 0011

0000 10/11 0000 0000 0000 0000 0000 0011  
2 3000003  
Op(6) immediate (26)

### Exercise 11: (2.38 from book)

Explain why an assembler might have problems directly implementing the branch instruction in the following code sequence:

here: beq \$s0, \$s2, there

...

there: add \$s0, \$s0, \$s0

Show how the assembler might rewrite this code sequence to solve these problems.

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The problem is that we are using PC-relative addressing, so if that address is too far away, we won't be able to use 16 bits to describe where it is relative to the PC.

If there refers to a location further than 128 KB from the PC, the solution would be:

```
here:      bne $s0, $s2, skip
           j  there
skip:
...
there:     add $s0, $s0, $s0
```

If there refers to a location further than 256 MB from the PC, the solution would be:

```
here:      bne $s0, $s2, skip
           lui $ra, there(upper)
           ori $ra, $ra, there(lower)
           jr $ra
skip:
...
there:     add $s0, $s0, $s0
```

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### Exercise 12:

Suppose that you have the following function signature:

```
int sum(int A[], int first, int last)
```

This function calculates the sum of the elements of A from index first and ending with element last. Write a fragment of MIPS assembly language which calls this function and uses it to calculate the average of all values in A. You may assume that the size of the array A is N, the base address of A in \$a0.

#### Average:

```
add $a1, $zero, $zero      # index of first element
addi $a2, $zero, N
addi $a2, $a2, -1           # index of last element is N-1
jal sum
add $t0, $zero, $v0        # Save the return value in $t0
addi $t1, $zero, N         # Load size of array into $t1
div $t2, $t0, $t1          # This form of div is provided as
                           # a pseudoinstruction.
```

**Exercise 13:**

Below is a recursive version of the function BitCount. This function counts the number of bits that are set to 1 in an integer.

Your task is to translate this function into MIPS assembly code. The parameter x is passed to your function in register \$a0. Your function should place the return value in register \$v0.

```
int BitCount(unsigned x) {
    int bit;
    if (x == 0)
        return 0;
    bit = x & 0x1;
    return bit + BitCount(x >> 1);
}
```

BitCount:

```
addi $sp, $sp, -8
sw $s0, 4($sp)
sw $ra, 0($sp)
bne $a0, $0, else
add $v0, $0, $0
addi $sp, $sp, 8
jr $ra
```

else:

```
andi $s0, 1
srl $a0, $a0, 1
jal BitCount
add $v0, $v0, $s0
lw $ra, 0($sp)
lw $s0, 4($sp)
addi $sp, $sp, 8
jr $ra
```

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## Extra exercises

### Exercise 14: (2.29 from book)

Add comments to the following MIPS code describe in one sentence what it computes. Assume that \$a0 and \$a1 are used for the input and both initially contain the integers a and b, respectively. Assume that \$v0 used for the output.

```

                add $t0, $zero, $zero    # initialize running sum $t0 = 0
loop:          beq $a1, $zero, finish    # finished when $a1 is 0
                add $t0, $t0, $a0        # compute running sum of $a0
                sub $a1, $a1, 1          # compute this $a1 times
                j loop
finish:        addi $t0, $t0, 100         # add 100 to a * b
                add $v0, $t0, $zero      # return a * b + 100

```

```

t0=0
while(a1 != 0){
    t0 = t0 + a0;
    a1 = a1 - 1;
}
t0 = t0 + 100;
v0 = t0;

```

The program computes  $a * b + 100$ .

### Exercise 15: (2.34 from book)

The following program tries to copy words from the address in register \$a0 to the address in register \$a1, counting the number of words copied in register \$v0. The program stops copying when it finds a word equal to 0. You do not have to preserve the contents of registers %v1, \$a0 and \$a1. This terminating word should be copied but not counted.

```

                addi $v0, $zero, 0 # Initialize count
loop:lw, $v1, 0($a0) # Read next word from source
                sw $v1, 0($a1) # Write to destination
                addi $a0, $a0, 4 # Advance pointer to next source
                addi $a1, $a1, 4 # Advance pointer to next destination
                beq $v1, $zero, loop # Loop if word copied != zero

```

There are multiple bugs in this MIPS program; fix them and turn in a bug-free version.

Bug 1: Count (\$v0) is initialized to zero, not -1 to avoid counting zero word.

Bug 2: Count (\$v0) is not incremented.

Bug 3: Loops if word copied is equal to zero rather than not equal.

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Bug-free version:

```

addi $v0, $zero, -1 # Initialize to avoid counting zero word
loop:lw, $v1, 0($a0) # Read next word from source
    addi $v0, $v0, 1 # Increment count words copied
    sw $v1, 0($a1) # Write to destination
    addi $a0, $a0, 4 # Advance pointer to next source
    addi $a1, $a1, 4 # Advance pointer to next destination
    bne $v1, $zero, loop # Loop if word copied != zero

```

### Exercise 16:

Convert the following C fragment to equivalent MIPS assembly language. Assume that the variables a, b, c, d, i and x are assigned to registers \$t1, \$t2, \$t3, \$t4, \$s0 and \$s1 respectively. Assume that the base address of the array A and B is in register \$a0 and \$a1 respectively.

a) if ((a<b) && (c==0)) d = 1;

```

    slt $t0, $t1, $t2
    beq $t0, $0, not          #if (a>=b) go to not
    bne $t3, $0, not          #if (c!=0) go to not
not:

```

b) if (a > b) b = a + 10; else b = a - 10;

```

    slt $t0, $0, $t1 # if $0 < $t1 then $t0 = 1 else $t0 = 0
    beq $t0, $0, else # if $t0 == $0 then branch to else
    addi $t2, $t1, 10
    j exit
else: addi $t2, $t1, -10
exit:

```

c) A[x+3] = B[x+2] | 0x10

```

addi $t0, $s1, 2 # $t0 = x+2
sll $t0, $t0, 2 # $t0 = (x+2)*4
add $t1, $a1, $t0 # $t1 = (base address of B + (x + 2))
lw $t2, 0($t1) # $t2 = B[x+2]
ori $t3, $t2, 0x10 # $t3 = B[x+2] | 0x10
addi $t4, $s0, 3 # $t4 = x+3
sll $t4, $t4, 2 # $t4 = (x+3)*4
add $t5, $a0, $t4 # $t5 = (base address of A + (x + 3))
sw $t3, 0($t5) # A[x+3] = $t3

```

```

d) for(int i=0; i<5; i++){
    a += b;
}

    add $s0, $zero, $zero
Loop: stli $t0, $s0, 5
    beq $t0, $zero, exit
    add $t1, $t1, $t2
    addi $s0, $s0, 1
    j Loop
exit:

```

**Exercise 17:**

Convert this high level language code into MIPS code. Do not forget to write MIPS code for the abs(x) procedure. (i saved in \$s0, the address of a in \$s1, y in \$s2)

```

i = 3;
y = y + abs(a[i]);

```

```

    addi $s0, $zero, 3
    add $t0, $s0, $s1
    add $t1, $t0, $s1
    lw $
    jal
    add
    j exitall
abs:
    slt $t0, $s0, $zero
    bne $t0, $zero, else
    add $v0, $a0, $zero
    jr $ra
else:
    sub $v0, $zero, $a0
    jr $ra
exitall:

```

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**Exercise 18:**

For each of the following, write the shortest sequence of MIPS assembly instructions to perform the specified operation.

(Hint:  $12345678 = 188 \times 2^{16} + 24910$   
 $31415924 = 479 \times 2^{16} + 24180$ )

a)  $\$v0 = 12345678$

```

lui $v0, 188
ori $v0, $v0, 24910

```

b) if ( \$t0 < 12345678 ) go to address less

```
lui    $t1, 188
ori    $t1, $t1, 24910
slt    $t2, $t0, $t1
bne    $t2, $zero, less
```

c) t1 = 12345678 + 31415924

```
lui    $t0, 188
ori    $t0, $t0, 24910
lui    $t2, 479
ori    $t2, $t2, 24180
add    $t1, $t0, $t2
```

### Exercise 19:

Given the register values in two's complement representation

\$s1 = 0000 0000 0000 0000 0000 0000 0000 0101 = 5

\$s2 = 0000 0000 0000 0000 0000 0000 0000 0011 = 3

\$s3 = 1111 1111 1111 1111 1111 1111 1111 1100 = -4

What are the values of registers \$s1 through \$s4 after executing the following MIPS instructions:

```
slt    $s1,
slt    $s2,
slt    $s3,
slt    $s4, $s1, $s3      Result
```

slt compares the unsigned values of \$s1 and \$s3, so in the last instruction \$s1 less than \$s3 big unsigned value, so \$s4 = 1.