Project

CMPEN 331 – Computer Organization and Design

Late submission is not accepted and will result in not getting any credits for the project

In the project, you just need to implement what was described in the honor option section in lab 5, with the addition of the implementation and generation of the bit stream without errors.

- 1. Write a report that contains the following:
 - Your Verilog design code. Use:
 - i. Device: Zyboboard (XC7Z010- -1CLG400C)
 - Your Verilog® Test Bench design code. Add "timescale 1ns/1ps" as the first line of your test bench file. ii.
 - The waveforms resulting as requested from item 9 above. iii.
 - The design schematics from the Xilinx synthesis of your design. Do not use any area constraints. iv.
 - Snapshot of the I/O Planning and v.
 - Snapshot of the floor planning vi.
 - vii. The design should be free from errors when synthesized, implemented and generated of the bitstream.

The report format will be as follows: Assignment Project Exam Help

- 2. REPORT FORMAT: Free form, but it must be:
 - a. One report per
 - b. Have a cover s https://eduassistpro.github.yio.de doing in the project.
 - d. You should include an introduction for the proje ms the connection between all the stages and what whall be benefit having the du ass Suputer or rangation field.

 or PDF. If you know LaTex,
 - e. Use Microsoft word and your report should be u you should upload the Tex file in addition to the PDF file.
 - f. Single spaced

The following part is not mandatory but any student will choose to do this part in addition to the previous part, will take 5 points extra to the total grade of the course:

In this extra points project, the students are implementing a pipeline CPU using the Xilinx design package for FPGAs. You can use any information available in previous labs if needed.

3. Pipelining

As described in lab 3

4. Circuits of the Instruction Fetch Stage

As described in lab 3

5. Circuits of the Instruction Decode Stage

As described in lab 3

6. Circuits of the Execution Stage

As described in lab 4

7. Circuits of the Memory Access Stage

As described in lab 4

8. Circuits of the Write Back Stage

As described in lab 5

9. Control Hazards and Delayed Branch

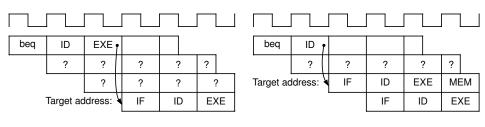
on the two source operards:

The control hazard occurs when a pipelined CPU executes a branch or jump instruction. The jump target address a jump instruction (jr, j, or jal) can be determined in the ID stage and it will be written into PC at the end of the ID stage. But because the pipelined CPU fetches instruction during every clock cycle, the next instruction is being fetched during the ID stage of the jump instruction. The control hazard caused by a conditional branch instruction (beg or bne) becomes more serious than that of a jump instruction because the condition must be evaluated in addition to the calculation of the branch of the target address. Figure 1 shows an example when we calculate the branch target address in the EXE or the ID stage respectively. There are mainly two methods to deal with the instruction(s) next to branch or jump instruction. One method is to cancel it (them). The other is to let it (them) be executed. The second method is called a delayed branch. The position in between the location of a jump or branch instruction and the junicontrancy carretade estrate tatled delay slogs. MIPS (metoprocess or without interlocked pipeline stages) ISA (instruction set architecture) adopts a one delay slot mechanism: the instruction located in delay slot is always executed no matter wither the branch is taken or not as shown in figure 2. In figure where the branch is taken; t is the 2 (a) shows the case w branch target address. https://eduassistpro.githors.gi conditional branch instructions finish the executions in th e no problem for calculating the branch target address within the ID stage. For checking form an exclusive OR (XOR)

 $rsrtequ = (da^db);$ // (da == db)

where the rsrtequ signal indicates where da or db are equal or not. Both da and db should be the state of the art data. Referring to figures 3 and 4, we use the outputs of the multiplexers for internal forwarding as da and db. This is the reason why we put the forwarding to the ID stage instead of to the EXE stage. Because the delayed branch, the return address of the MIPS jal instruction is PC+8. Figure 5 illustrates the execution of the jal instruction. The instruction located in delay slot (PC + 4) was already executed before transferring control to a function (or a subroutine). The return address should be PC+8, which is written into \$31 register in the WB stage by the jal instruction. The return form subroutine can be done by the instruction of jr \$31. The jr rs instruction reads the content of register rs and writes it into the PC.

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(a) Branch is determined in EXE stage

(b) Branch is determined in ID stage

Figure 1 Determining whether a branch is taken or not taken

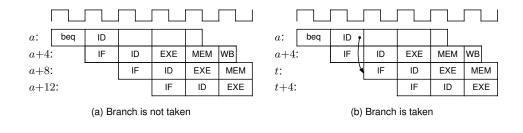
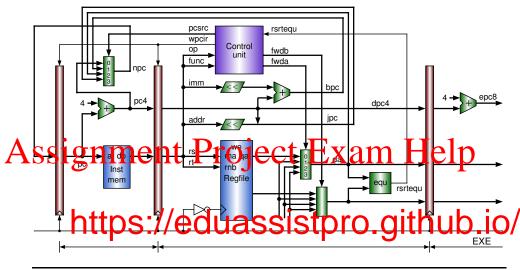


Figure 2 Delayed branch



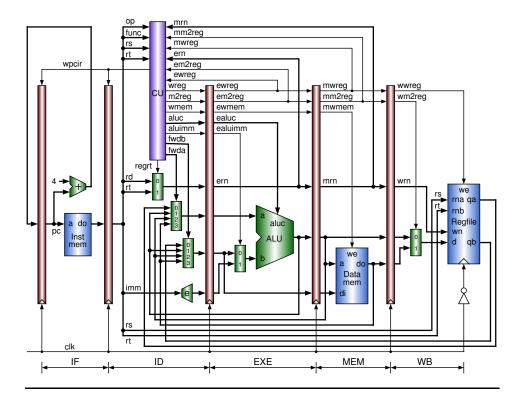


Figure 4 Mechanism of internal forwarding and pipeline stall

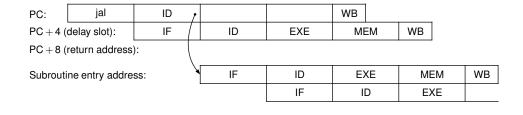


Figure 5 Return address of the function call instruction

For your reference, figure 6 illustrate the detailed circuit of the pipelined CPU, plus instruction memory and data memory. The PC can be considered as the first pipeline located in front of the IF stage, and a register of the register file can be considered as the sixth (last) pipeline register at the end of the WB stage.

In the IF stage, an instruction is fetched from instruction memory, and the PC is incremented by 4 if the instruction in the ID stage is neither a branch nor a jump instruction, and there is no pipeline stall. There are four sources supported Project Exam Help

```
bpc: branch tar
da: target addre
jpc: jump targe https://eduassistpro.github.io/
```

The selection of the next PC (npc) is done by a 32-bit 4-to-1 multiplexer whose selection signal is pesrc (PC source), generated by the control unit in th

In the ID stage, two reaction control is extended and the instruction is not func) by the control unit.

The selection signal of the multiplexer for ALU's input e.g. A is named fwda (forward A) and the other for ALU's input B, is named fwdb (forward B). if there is no data hazard, the multiplexer selects the data read from the register file. The inverse of the stall signal is used as the write enable for the PC and the IF/ID pipeline register (wpcir). The stall signal becomes true when an instruction in the ID stage uses the result of an lw instruction which is in the EXE stage. Thus, the stall signal can be generated by the following Verilog HDL code.

```
stall = ewreg & em2reg & (ern!=0) & (i_rs & (ern== rs) | i_rt & (ern
== rt));
```

where i_rs and i_rt indicate that an instruction uses the contents of the rs register and the rt register respectively.

There is an important thing we must not to forget. The pipeline stall is implemented by prohibiting the updates of the PC and the IF/ID pipeline register. But the instruction that is already in the IF/ID register will be decoded and fed to the next pipeline stage. This will result in an instruction being executed twice. To prvent an instruction from being executed twice, we must cancel the first instruction. Canceling an instruction is easy: prevent it from updating the states of the cpu and memory. All the control signals that will be used in the following stages are saved into the ID/EXE registers. In the EXE stage, in addition to the operation performed by the ALU, the PC+8 operation is carried out by an adder for generating the return address for the jal instruction. The shift amount (sa) for a shift

instruction can be extracted from the immediate field (eimm). If the instruction in the EXE stage is a jal, PC+8 is selected and the destination register number (ern) is set to 31 (done by f component). Otherwise, the ALU output is selected and let ern=ern0 (rd or rt in the EXE stage).

<u>In the MEM stage</u>, if the instruction is an sw, the data mb will be written into the data memory addressed by malu. If the instruction is an lw, the memory data addressed by malu is read out. Other instructions do nothing in this stage.

<u>In the WB stage</u>, an instruction is graduated by writing the result, either the ALU result or memory data, into a register file. The destination register number is wrn (register number in the WB stage). And the write enable signal is wwreg (register write enable in WB stage)

10. Test Program and Simulation Waveform

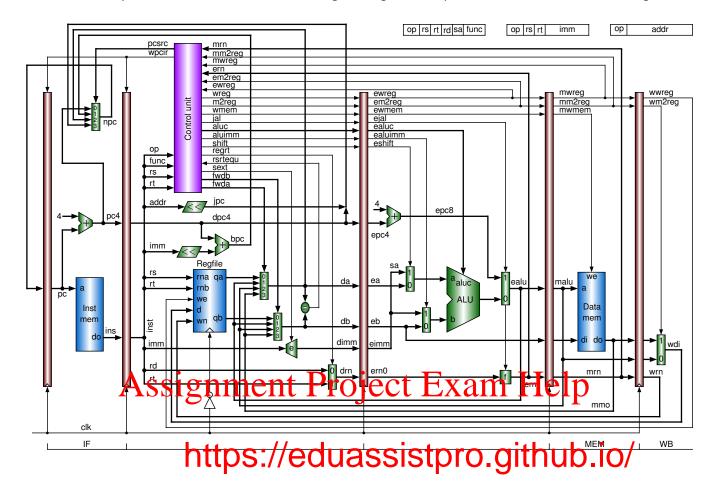
Write a Verilog code that implement the following instructions to verify the correctness of your pipelined CPU design. The code should be used to initialize the instruction memory block. The register file should be all initialized to zeros. In the test program, it is aimed to check the 20 instructions. The main part of the test program is a subroutine in which four 32-bit memory words are summed by a for loop. After returning from the subroutine, the sum is stored in the data memory by a sw instruction. A code pattern that causes pipeline stall is also prepared within the loop. Word address is used to assign the content of each word (a 32 bit instruction). The parenthesized hexaded matching in the center of each line is the byte address (PC).

```
y, rmmdule instruction mem
        input [31:0] a;
        output [31:0] a; https://eduassistpro.github.io/
        wire
               [31:0] rom
        // rom[word addr] = instruction
                                            // (pc)
        assign rom[\overline{6}'h00] = A32h3107070
assign rom[6'h01] = A32h3107070
        assign rom[6'h02] = 32'h0c00001b;
                                            // (08)
        assign rom[6'h03] = 32'h20050004;
                                           // (0c) dslot1: addi $5, $0,
        assign rom[6'h04] = 32'hac820000;
                                           // (10) return: sw
                                                                 $2, 0($4)
        assign rom[6'h05] = 32'h8c890000;
                                           // (14)
                                                                 $9, 0($4)
                                                            lw
        assign rom[6'h06] = 32'h01244022;
                                                                 $8, $9, $4
                                           // (18)
                                                           sub
        assign rom[6'h07] = 32'h20050003;
                                           // (1c)
                                                           addi $5, $0,
                                           // (20) loop2: addi $5, $5, -1
        assign rom[6'h08] = 32'h20a5ffff;
                                           // (24)
                                                            ori $8, $5, 0xffff
        assign rom[6'h09] = 32'h34a8ffff;
                                                            xori $8, $8, 0x5555
        assign rom[6'h0a] = 32'h39085555;
                                           // (28)
        assign rom[6'h0b] = 32'h2009ffff;
                                           // (2c)
                                                            addi $9, $0, -1
        assign rom[6'h0c] = 32'h312affff;
                                           // (30)
                                                            andi $10,$9,0xffff
                                           // (34)
        assign rom[6'h0d] = 32'h01493025;
                                                                 $6, $10, $9
                                                            or
        assign rom[6'h0e] = 32'h01494026;
                                                                 $8, $10, $9
                                           // (38)
                                                            xor
        assign rom[6'h0f] = 32'h01463824;
                                           // (3c)
                                                            and $7, $10, $6
                                           // (40)
                                                                 $5, $0, shift
        assign rom[6'h10] = 32'h10a00003;
                                                            beq
        assign rom[6'h11] = 32'h00000000;
                                           // (44) dslot2: nop
        assign rom[6'h12] = 32'h08000008;
                                           // (48)
                                                                 loop2
        assign rom[6'h13] = 32'h00000000;
                                           // (4c) dslot3: nop
                                           // (50) shift: addi $5, $0, -1
        assign rom[6'h14] = 32'h2005ffff;
        assign rom[6'h15] = 32'h000543c0;
                                           // (54)
                                                                 $8, $5, 15
                                                            sll
        assign rom[6'h16] = 32'h00084400;
                                           // (58)
                                                            sll
                                                                 $8, $8, 16
        assign rom[6'h17] = 32'h00084403;
                                           // (5c)
                                                                 $8, $8, 16
                                                            sra
                                           // (60)
                                                                 $8, $8, 15
        assign rom[6'h18] = 32'h000843c2;
                                                            srl
                                           // (64) finish: j
        assign rom[6'h19] = 32'h08000019;
                                                                 finish
        assign rom[6'h1a] = 32'h00000000; // (68) dslot4: nop
```

```
add $8, $0, $0
    assign rom[6'h1b] = 32'h00004020; // (6c) sum:
    assign rom[6'h1c] = 32'h8c890000; // (70) loop:
                                                               $9, 0($4)
                                                         lw
    assign rom[6'h1d] = 32'h01094020; // (74) stall: add <math>$8, $8, $9
    assign rom[6'h1e] = 32'h20a5ffff; // (78)
assign rom[6'h1f] = 32'h14a0fffc; // (7c)
                                                         addi $5, $5, -1
                                                         bne $5, $0, loop
    assign rom[6'h20] = 32'h20840004; // (80) dslot5: addi $4, $4, 4
    assign rom[6'h21] = 32'h03e00008; // (84)
                                                          jr
                                                               $31
    assign rom[6'h22] = 32'h00081000; // (88) dslot6: sll $2, $8, 0
                                  // use 6-bit word address to read rom
    assign inst = rom[a[7:2]];
endmodule
```

below is the test data that should be stored in the data memory. Four 32-bit words in the memory will be read by lw instructions. The test program will store a word in the location next to the four words.

```
module data memory (clk,dataout,datain,addr,we); // data memory, ram
             input clk;
                                                                                                                                                    // clock
             input [31:0] addr;
                                                                                                                                                    // ram address
             input [31:0] datain;
                                                                                                                                                    // data in (to memory)
                                                                                                                                                    // write enable
             input
                                                             we;
             output [31:0] dataout;
                                                                                                                                                    // data out (from memory)
             reg [31:0] ram [0:31];
                                                                                                                                                    // ram cells: 32 words * 32 bits
             assign data Atssignment;
             always @ (posedge Dk)
                           if (we) ram[addr[6:2]] = datain;
                                                                                                                                                    // write ram
             end
             integer i;
                                                                           https://eduassistpro.github.io/
             initial begin
                           for (i = 0; i
                                         ram[i] = 0;
                           // ram[word_addAldatWeChat edu_assistated assistated as
                           ram[5'h15] = 32'h00000027;
                           ram[5'h16] = 32'h00000079;
                                                                                                                                                    // (58) data[2] ca + 79 = 143
                                                                                                                                                  // (5c) data[3] 143 + 115 = 258
                           ram[5'h17] = 32'h00000115;
                           // ram[5'h18] should be 0x00000258, the sum stored by sw instruction
             end
endmodule
```



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Figure 7 illustrates an example of waveforms when the pipelined CPU executes the jal instruction (PC = 0x00000008). The instruction in the delay slot (PC = 0x00000000c) is executed also. The taget address of the jal instruction is 0x0000006c, the entry of a subroutine (sum). The result at the EXE stage of the jal instruction is 0x00000010, which is the return address (from the subroutine).

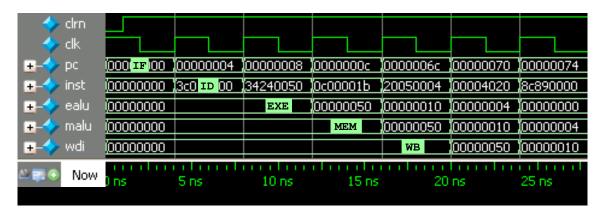


Figure 7 Waveform of the pipelined CPU (call subroutine)

11. Write a Verilog code that implement the instructions shown in item number 8 with the corresponding initialization of data memory using the design shown in **Figure 6**. You need to show your outputs in a similar way as figure 7 with the same signals when the pipelined CPU execute the lw \$9, 0 (\$4) instruction (PC =

0x00000070) and its follow up, the add \$8, \$8, \$9 instruction in the fourth (last round) of the for loop.

- 12. Write a report that contains the following:
 - viii. Your Verilog design code. Use:
 - i. Device: Zyboboard (XC7Z010- -1CLG400C)
 - Your Verilog® Test Bench design code. Add "`timescale 1ns/1ps" as the first line of your test bench file. ix.
 - The waveforms resulting as requested from item 9 above. x.
 - The design schematics from the Xilinx synthesis of your design. Do not use any area constraints. xi.
 - Snapshot of the I/O Planning and xii.
 - Snapshot of the floor planning xiii.
 - The design should be free from errors when synthesized. xiv.
- 13. **REPORT FORMAT:** Free form, but it must be:
 - a. One report per student.
 - b. Have a cover sheet with identification: Title, Class, Your Name, etc.
 - c. You have to write an abstract at the beginning of the project report to describe what you have done in the project.
 - d. Use Microsoft word and it should be uploaded in word format not PDF. If you know LaTex, you should upload the Tex file in addition to the PDF file.
- e. Single spaced ASSIGNMENT Project Exam Help

 14. You have to upload the whole project design file zipped with the word or LaTex with PDF file.

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