

Com S 311, Final Exam

Problem	Max Points	Score
1 (Basics)	20	
2 (Hash and Heaps)	20	
3 (Recurrence)	10	
4 (Graph)	15	
5 (Dynamic Programming)	20	
6 (Dynamic Programming)	20	
(Extra Credit-1: Graphs)	20	
(Extra Credit-2: NP)	20	
Total	100 + 40 (EC)	

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 Assignment Project Exam Help

- If you write "do not grade" and nothing else, you will receive 15% credit.
- For all problems that involve writing an algorithm, please write clear pseudocode, not Java or C.

- In general you may use the following algorithms (that is, unless in the problem statement it says otherwise): sorting, breadth-first search, Dijkstra's algorithm, topological sort. You can assume that hash tables are efficient. However, if you modify any of these methods, then you must write a comment explaining the modification. Merely stating the modification does not suffice.

- Level of points for solutions related to design and analysis algorithm

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if designed algorithm is correct and efficient as expected then
  if proof of correctness is unambiguous and justifiable then
    if run-time analysis shows derivation steps and is correct then
      points = 100%
    else
      points = 75%
  else if run-time analysis shows derivation steps and is correct then
    points = 75%
  else
    points = 50%
else if designed algorithm is correct and brute force then
  points = 30%
else if designed algorithm is incorrect then
  points = 0--20% (at the discretion of grader)
else if answer is "DO NOT GRADE" (YOU NEED TO EXPLICITLY WRITE DO NOT GRADE)
  points = 15%
else
  points = 0%
  
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1. Short Answer questions. Do not write explanations. There is *Do not grade* option for this question.
 - (a) True or False: $2^{3n} \in O(2^{2n})$.
 - (b) True or False: $2^{3^n} \in O(2^{2^n})$.
 - (c) True or False: Every divide and conquer algorithm runs in $O(n \log n)$ time.
 - (d) True or False: If G has a topological ordering, then G must be a DAG.
 - (e) True or False: Every undirected graph with more than n edges has a cycle.
 - (f) True or False: The runtime for solving any decision problem in NP is not polynomial with respect to the size of the problem..
 - (g) What is the solution to the recurrence $T(n) = T(n/2), T(1) = O(1)$.
 - (h) What is the runtime for removing the k th smallest element from a binary heap?
 - (i) Assuming that the comparison operation takes $\Theta(1)$ time, what is the runtime of the selection sort algorithm to output the smallest element of an array of size n ?
 - (j) What is the runtime of Prim's Algorithm?
2. *Hashing and Heapsort*
 - (a) Let A_1, A_2 be two arrays each consisting of n integers and T be an integer. Give an algorithm using hashing that checks whether there exist integers x and y such that $x + y$ equals T . State the run-time of your algorithm.
 - (b) Given an array A of n integers and an integer k , give an algorithm that finds the k th smallest element. Express runtime as a function of n and k .
3. Solve the recurrence relation $T(n) = 2T(n/2) + n$.
4. You are given a directed graph G with n vertices and m edges. The graph has a vertex v with the following properties:
 - Number of vertices from which there is a path to v is k_1 .
 - Number of vertices u that have a path from v to u is k_2 .
 - There is no vertex $u \neq v$ such that there is a path from v to u and a path from u to v .

Write the verification algorithm.

5. You are given n jobs numbered $1, 2, \dots, n$ to complete and each job i comes with a difficulty d_i . Each job takes exactly one week to complete irrespective of its difficulty. If you complete job i during week j ($j \leq n$), then you earn a profit of $d_i(n - j)$. You plan to complete all the jobs in n weeks, since the goal is to maximize the profit. Consider greedy algorithm that completes that jobs in the based on the *decreasing order* of difficulty. Use exchange argument to prove that the greedy algorithm produces optimal solution. You may assume that all d_i 's are distinct.
6. Let M be a matrix of integers with n rows (rows numbered 1 to n) and m columns (columns numbered 1 to m). Let $M[i, j]$ denote the entry in i th row and j th column. A *horizontal cut* in M is a sequence $[c_1, c_2, \dots, c_m]$ such that
 - For every $i, 1 \leq c_i \leq n$
 - For $1 \leq i \leq m - 1, c_{i+1} \in \{c_i - 1, c_i, c_i + 1\}$

Given a horizontal cut $[c_1, c_2, \dots, c_m]$, its cost is $M[c_1, 1] + M[c_2, 2] + \dots + M[c_m, m]$. A horizontal cut is a *max-cost cut* if its cost is at least the cost of any other horizontal cut.

Give a dynamic programming algorithm, that takes a matrix M as input, and outputs the *cost of the max-cost cut*. The following needs to be presented as part of your answer.

- (a) Recurrence relation (recursive characterization or recursive definition) describing the value of max-cost cut.
 - (b) Iterative algorithm for computing the value of max-cost cut.
 - (c) Runtime of the iterative algorithm.
- * **Extra Credit-1.** Let G be an undirected graph where every edges has exactly the same weight and let s be a vertex of the graph. Write an algorithm that takes s as input and outputs for every vertex v , the number of shortest path from s to v .
- * **Extra Credit-2.** Let $L = \{S_1, S_2, \dots, S_n\}$ be a collection of sets. A set H is a *blanket* for L if $\forall i \in [1, n] : H \cap S_i \neq \emptyset$. Show that the following decision problem is in NP:
- Input: $L = \{S_1, S_2, \dots$
- Decision: Does L have a b

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