COMP 250

INTRODUC TER SCIENCE

Week 9-1 : I

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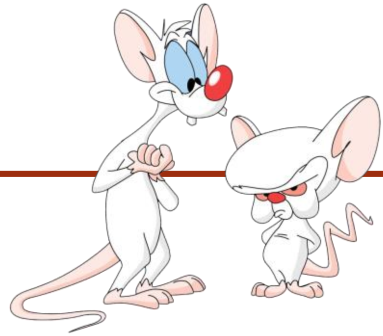
WHAT ARE WE GOING TO DO IN THIS VIDEO?



- Inductive/Recursive definitions
- Inductive/Recursi
  - Mathematical Induction

PROOFS

For all $n \geq 1$,

$$1 + 2 + 3 + \dots + (n-1) + n = \frac{n(n+1)}{2}$$

How can we prove such a statement?

- By "proof", we mean a formal logical argument that convincingly demonstrate the truth of a given proposition.
- Note that "convincingly" is itself not well defined.

EXAMPLE

$$1 + 2 + \ldots + (n - 1) + n$$

Rewrite by considering $n/2$ pairs :

$$1 + 2 \quad \frac{\;}{2} \left(\frac{\;}{2} \right) \quad (- 1) + n$$

If n is even, then adding up the $n/2$ pairs gives

$$n/2 * (n + 1)$$

- What if n is odd?

EXAMPLE

- What if n is odd? Then, n-1 is even. So,

$$1 + 2 + \ldots + (n-1) + n$$

$$= \left(\frac{n-1}{2} + 1\right)$$

$$= \frac{n+1}{2} * n$$

which is the same formula as before.

RECURSIVE (INDUCTIVE) DEFINITION

- Some set of elements can be define recursively/inductively.

- A recursive/inductive definition consists of the following:
  - A *base clause*
    Which one or more basic/initial eleme
  - One or more *inductive clauses*
    Rules on how to generate “new” elements of the set from “old” ones.
  - A *final clause*
    which simply states that no other element is part of the set.

EXAMPLE – NATURAL NUMBERS

The set of natural numbers can be defined as follows:

- *Base clause*:

  0 is a natural number

- *Inductive clause*:

  If n is a natural number, then $n + 1$ is also mber.

- *Final clause*: Nothing else is a natural number.

MATHEMATICAL INDUCTION

Consider a statement of the form:

"For all $n \geq n_0$, $P(n)$ is true"

where n_0 is some co (n) has value true or false for each n.

- If n is an element of an inductively defined set, then the statement above can be proven using a technique called *mathematical induction*.

(WEAK) MATHEMATICAL INDUCTION

To prove a property by mathematical induction, we proceed as follows:

- *Base case*

  Show that the property holds for the basic/initial elements of the set.

- *Induction step*

  Assume the property hold for some elem ction Hypothesis)

  Show that the property also holds for any element generated from n using the inductive clauses.

- *Conclusion*

  The property holds for all elements.

EXAMPLE

“For all $n \geq n_0$, $P(n)$ is true”

For all $n \geq 1$,

$$1 + 2 + 3 + \ldots + (n - \;) \qquad \frac{n(n+1)}{2}$$

This is a property of natural numbers. Since this is a set that can be defined inductively, we can use mathematical induction to prove such property!

PROOF BY MATHEMATICAL INDUCTION

We need to prove the following:

- Base case:
  $P(n_0)$ is true, i.e. the pro

  this case is 1.

- Induction step:

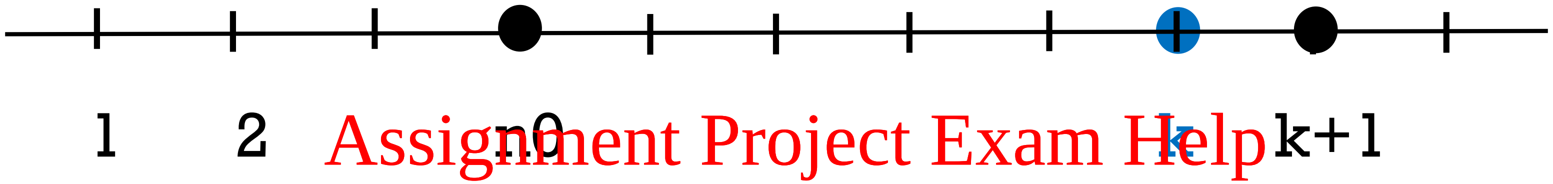
  IH: Assume $P(k)$ is true, i.e. the property element k.
  Prove that $P(k+1)$ is true, i.e. the property holds for $k+1$.

Base case:

P(n_0) is true.

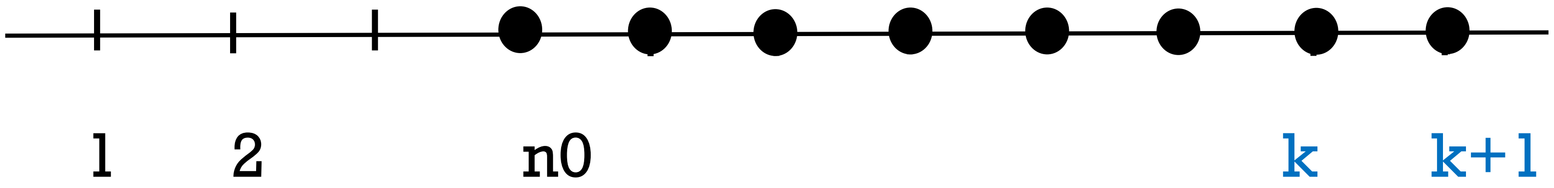
Induction step:

For any k >= n_0, if P(k) is true then P(k+1) is true.



Thus we have proved:

For any n >= n_0, P(n) is true.



BACK TO THE PROOF

For all $n \geq 1$, $1 + 2 + 3 + \; \ldots . + (n-1) + n = \frac{n(n+1)}{2}$

- Base case: $n = 1$, to prov

$$1 = \frac{1 * (1 \quad\quad)}{}$$

$$1 = \frac{2}{2} = 1$$

✓

BACK TO THE PROOF

- Induction step:

  IH: Assume that it holds for k, that is

$$k(k+1)$$

Statement: For all $n \geq 1$, $1 + 2 + 3 + \ \ldots . + (n-1) + n = \frac{n(n+1)}{2}$

BACK TO THE PROOF

- Induction step:

  IH: Assume that it holds for k, that is

$$\frac{k(k+1)}{}$$

Prove it for $k+1$:

$$1 + 2 + \cdots + k \quad (\quad)$$

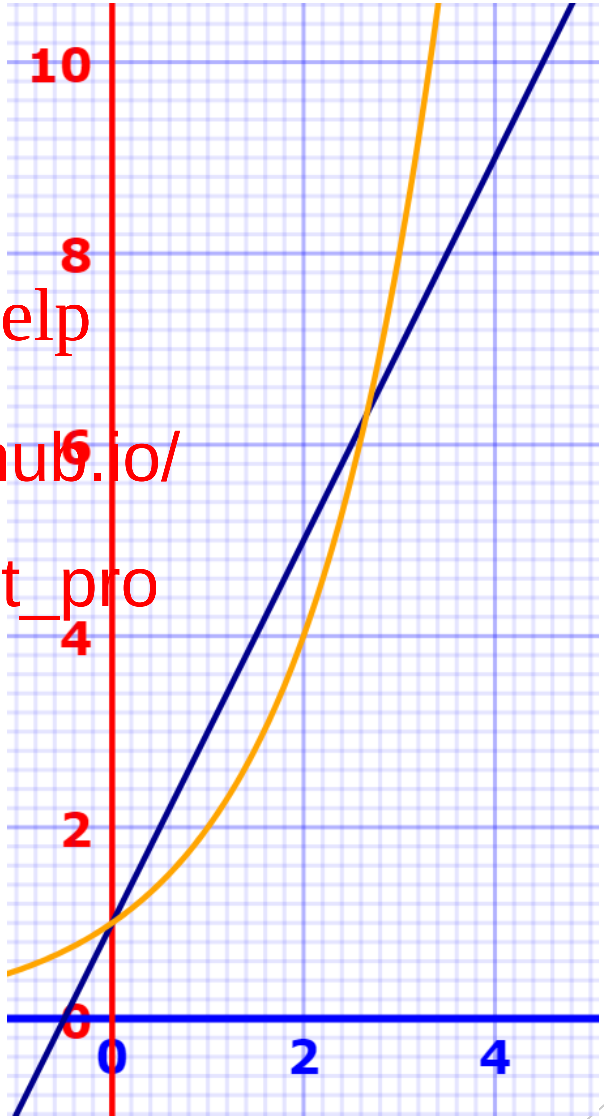
$$= \frac{k(k+1)}{2} + (k+1), \text{ by IH}$$

$$= (k+1) * \left(\frac{k}{2} + 1\right) = \frac{(k+1)(k+2)}{2}$$

EXAMPLE 2

- Prove the following statement:

For all $n \geq 3$, $2n + 1 < 2^n$.



EXAMPLE 2

<u>Statement:</u> For all $n \geq 3$, $2n + 1 < 2^n$.

- Note: $\mathrm{P}(n)$ is false for

But that has nothing to do *e.*

EXAMPLE 2

<u>Statement:</u> For all $n \geq 3$, $2n + 1 < 2^n$.

Proof: (by mathematical in

- Base case ($n = 3$):

$$2 * 3 + 1 = 7$$

Statement: For all $n \geq 3$, $2n + 1 < 2^n$.

EXAMPLE 2

- Induction step:

IH: Assume $2 * k + 1 < 2^k$.

Statement: For all $n \geq 3$, $2n + 1 < 2^n$.

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- Induction step:

IH: Assume $2 * k + 1 < 2^k$.

Prove it for $k + 1$:

$(\quad)$

$= 2 * k +$

Statement: For all $n \geq 3$, $2n + 1 < 2^n$.

EXAMPLE 2

- Induction step:

IH: Assume $2 * k + 1 < 2^k$.

Prove it for $k + 1$:

$2 \quad (\qquad) \qquad 1$

$= 2 * k +$

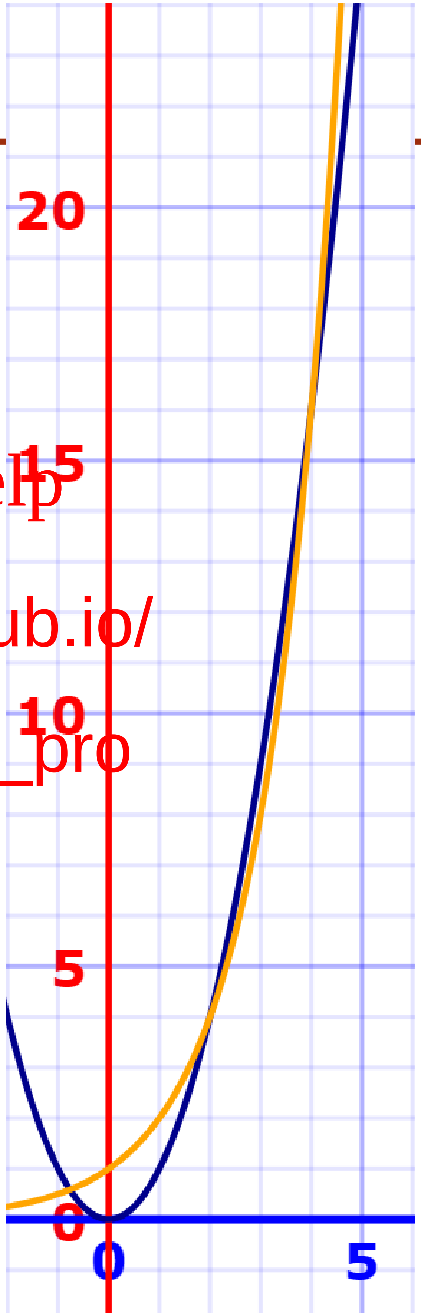
$< 2^k + 2$, by IH

$< 2^k + 2^k$, for $k \geq 3$

$= 2^{k+1}$

EXAMPLE 3

Statement: For all $n \geq 5$, $n^2 < 2^n$.



EXAMPLE 3

Statement: For all $n \geq 5$, $n^2 < 2^n$.

Proof: (by mathematical in

- Base case ($n = 5$):

$$5^2 = 25 <$$

EXAMPLE 3

Statement: For all $n \geq 5$, $n^2 < 2^n$.

- Induction step.

What should we assume?

What do we need to prove?

EXAMPLE 3

Statement: For all $n \geq 5$, $n^2 < 2^n$.

- Induction step.

What should we assume?

What do we need to prove? $(k+1)^2 < 2^{(k+1)}$

EXAMPLE 3

Statement: For all $n \geq 5$, $n^2 < 2^n$.

- Induction step.

  IH: $k^2 < 2^k$ for a $k \geq 5$

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Statement: For all $n \geq 5$, $n^2 < 2^n$.

- Induction step.

  IH: $k^2 < 2^k$ for a $k \geq 5$

$(\quad)$

$< 2^k$ by IH

$< 2^k + 2^k$, by Example 2

$= 2^{k+1}$

(STRONG) MATHEMATICAL INDUCTION

- Sometimes one would like to assume the induction hypothesis not only for the previous element, but also for smaller elements. This leads to a logically equivalent proof method called *strong (or complete) mathematical induction*.

- To prove a property by strong e proceed as follows:
  - Induction step

    Assume the property hold *for all elements* le itrary k. (Induction Hypothesis)
    Show that the property also holds for the k element which was generated using the inductive clauses.
  - Conclusion
    The property holds for all elements.

FIBONACCI NUMBERS

- The Fibonacci sequence is one of the most common example of a recursively-defined set.

- Consider the following s

1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144,

Let f_n denote the nth Fibonacci number. How can we define the sequence above?

FIBONACCI NUMBERS – INDUCTIVE DEFINITION

- Fibonacci sequence: 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, …

- Base clause:
  $f_0 = f_1 = 1$ are Fibonacci

- Inductive clause:
  If f_{n-1} and f_{n-2} are Fibonacci numbers, then $f_n = f_{n-1} + f_{n-2}$ is a Fibonacci number.

EXAMPLE 4

Statement: For all $n \geq 0$, $f_n \leq \left(\frac{7}{4}\right)^n$

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Proof: (by strong mathematical induction)

- Induction step

IH: Let k be ≥ 0, and assume that for any ch that $0 \leq i < k$ then

$$f_i \leq \left(\frac{7}{4}\right)^i$$

EXAMPLE 4

Statement: For all $n \geq 0$, $f_n \leq \left(\frac{7}{4}\right)^n$

Proof: (by strong mathematical induction)

- Induction step

  IH: Let k be ≥ 0, and assume that for any ch that $0 \leq i < k$ then

$$f_i \leq \left(\frac{7}{4}\right)$$

To show: $f_k \leq \left(\frac{7}{4}\right)^k$

EXAMPLE 4

There are 3 possible cases:

1. $k = 0$

   $f_0 = 1$ and $\left(\frac{7}{4}\right)^0 = 1$, so t

2. $k = 1$

   $f_1 = 1$ and $\left(\frac{7}{4}\right)^1 > 1$, so the claim holds.

EXAMPLE 4

There are 3 possible cases:

3. $k > 1$

$$f_k = f_{k-1} + f_{k-2}$$

$$\left(-\right)^{k-1} \quad \left(-\right)^{k-2} \text{, by IH}$$

$$= \left(\frac{7}{4}\right)^{k} \left(\quad - \quad \right) \left(\frac{7}{4}\right)^{k-2} \left(\frac{11}{4}\right)$$

$$= \left(\frac{7}{4}\right)^{k-2} \left(\frac{44}{16}\right)$$

$$< \left(\frac{7}{4}\right)^{k-2} \left(\frac{49}{16}\right) = \left(\frac{7}{4}\right)^{k-2} \left(\frac{7}{4}\right)^{2}$$

$$= \left(\frac{7}{4}\right)^{k}$$

RECOMMENDED EXERCISES

1. Prove that for all $n \geq 0, \sum_{i=0}^{n} 2^i = 2^{n+1} - 1$

2. Prove that for all $n \geq 0, \sum_{i=0}^{n} i^3 = \left(\frac{n(n+1)}{2}\right)^2$

3. Consider the following re n ('+') on natural numbers:

   - Base clause:

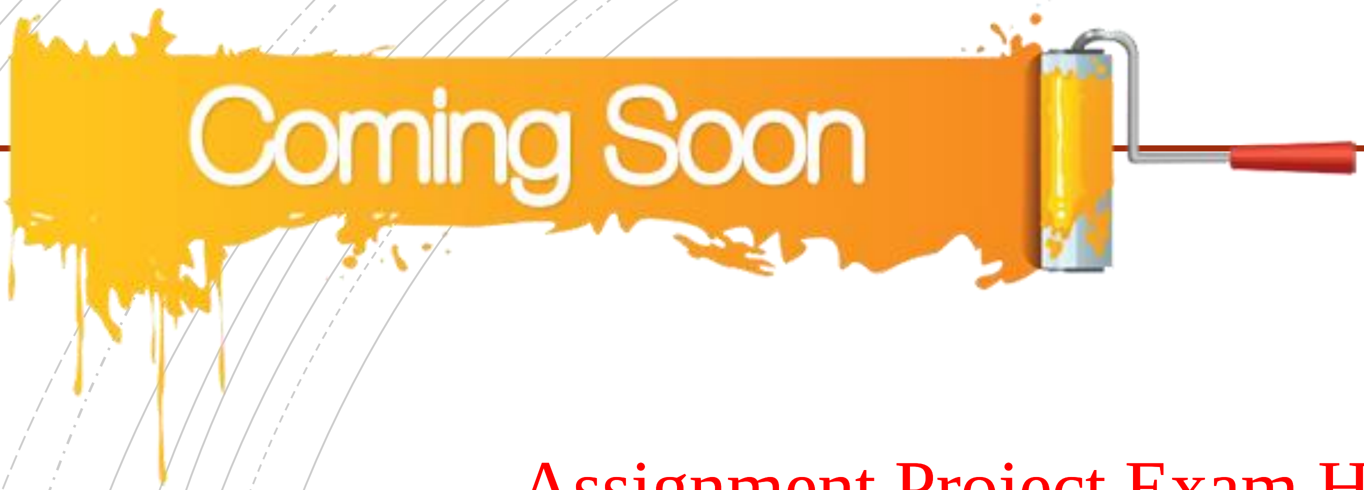
$$0 + m =$$

   - Inductive clause:

$$(n + 1) + m = (n + m) + 1$$

Prove that addition is associative, i.e. for all natural numbers $(a + b) + c = a + (b + c)$

Hint: use mathematical induction on a



In the next

- Recursiv