

# COMP2610/COMP6261 - Information Theory

## Tutorial 6: Source Coding

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### 1. Probabilistic inequalities

Suppose a coin is tossed  $n$  times. The coin is known to land “heads” with probability  $p$ . The number of observed “heads” is recorded as a random variable  $X$ .

- (a) What is the *exact* probability of  $X$  being  $n - 1$  or more?
- (b) Using Markov’s inequality, compute a *bound* on the same probability as the previous part.
- (c) Suppose  $n = 2$ . For what values of  $p$  will the bound from Markov’s inequality be within 1% of the exact probability?

### 2. AEP and source coding (cf. Cover & Thomas, Problem 3.7)

A sequence of bits its g

Sequences are code

blocks with more in

$p_0 = 0.995$  and  $p_1 = 0.005$ .

word. Those

- (a) What is the minimum required length of the assigned codeword
- (b) Calculate the probability of observing a 100-bit block that has no a
- (c) (*Harder*) Use Chebyshev’s inequality to bound the probability that no codeword has been assigned. Compare the bound to the proba

### 3. Typical Sets and Smallest $\delta$ -Sufficient Subsets (cf. Cover & Thomas, Problem 3.13)

Let  $X^N$  be an extended ensemble for  $X$  with  $\mathcal{A}_X = \{0, 1\}$  and  $\mathcal{P}_X = \{0.4, 0.6\}$ .

- (a) Calculate the entropy  $H(X)$ .
- (b) Let  $N = 25$  and  $\beta = 0.1$ .
  - i. Which sequences in  $X^N$  fall in the typical set  $T_{N\beta}$ ? (You may find it helpful to refer to Table 1 below.)
  - ii. Compute  $P(\mathbf{x} \in T_{N\beta})$ , the probability of a sequence from  $X^N$  falling in the typical set.
  - iii. With reference to Table 1 below, how many elements are there in  $T_{N\beta}$ ?
  - iv. How many elements are in the smallest  $\delta$ -sufficient subset  $S_\delta$  for  $\delta = 0.9$ ?
  - v. What is the essential bit content  $H_\delta(X^N)$  for  $\delta = 0.9$ ?

#### 4. Source Coding Theorem

Recall that the source coding theorem (for uniform codes) says that for any ensemble  $X$ :

$$(\forall \epsilon > 0) (\forall \delta \in (0, 1)) (\exists N_0) (\forall N > N_0) \left| \frac{1}{N} H_\delta(X^N) - H(X) \right| \leq \epsilon.$$

- (a) Near, an enthusiastic software developer, has just learned about the source coding theorem. He exclaims: “The theorem allows us to pick any  $\epsilon > 0$ . So, if I pick  $\epsilon = H(X) - \epsilon'$ , I get that for sufficiently large  $N$ ,

$$\frac{1}{N} H_\delta(X^N) \geq \epsilon'.$$

*This means that by making  $\epsilon'$  tiny, I can get away with using virtually zero bits per outcome. Great!”*

Is Near’s reasoning correct? Explain why or why not.

- (b) Mello, a skeptical econometrician, has also just learned about the source coding theorem. He complains: “The theorem is not really relevant to me. I am interested in coding blocks of outcomes where each outcome is dependent on the previous outcome, rather than them all being independent of each other. The source coding theorem is not useful in this case.”

Is Mello’s reasoning correct? Explain why or why not.

#### 5. Prefix Codes

Consider the codes  $C_1 = \{0, 01, 1101, 10101\}$ ,  $C_2 = \{00, 01, 100, 101\}$ , and  $C_3 = \{0, 1, 00, 11\}$

- (a) Are  $C_1$ ,  $C_2$ , and  $C_3$  prefix codes? Are they uniquely decodable?  
(b) Construct new prefix codes  $C'_1$ ,  $C'_2$ , and  $C'_3$  that have the same lengths as  $C_1$ ,  $C_2$ , and  $C_3$ , respectively. If this is not possible, explain why.  
(c) Is it possible to construct a prefix code with lengths  $\{4, 4, 4\}$ ? If so, construct one.  
(d) Is it possible to construct a prefix code with lengths  $\{4, 4, 4, 4\}$ ? If so, construct one.

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$k$	$\binom{N}{k}$	$\binom{N}{k} p_1^k p_0^{N-k}$	$-\frac{1}{N} \log_2 p(\mathbf{x})$
0	1	0.000000	1.321928
1	25	0.000000	1.298530
2	300	0.000000	1.275131
3	2300	0.000001	1.251733
4	12650	0.000007	1.228334
5	53130	0.000045	1.204936
6	177100	0.000227	1.181537
7	480700	0.000925	1.158139
8	1081575	0.003121	1.134740
9	2042975	0.008843	1.111342
10	3268760	0.021222	1.087943
11	4457400	0.043410	1.064545
12	5200300	0.079986	1.041146
13	5200300	0.113950	1.017748

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18	480700	0.079986	0.900755
19	177100	0.044203	0.877357
20	56130	0.01989	0.853958
21	12650	0.007104	0.830560
22	2300	0.001937	0.807161
23	300	0.000379	0.783763
24	25	0.000047	0.760364
25	1	0.000003	0.736966

Table 1: Table for Question 3. Column 1 shows  $k$ , the number of 1s in a block of length  $N = 25$ . Column 2 shows the number of such blocks. Column 3 shows the probability  $p(\mathbf{x}) = \binom{N}{k} p_1^k p_0^{N-k}$  of drawing such a block  $\mathbf{x}$ . Column 4 shows the Shannon information per symbol in  $\mathbf{x}$ .