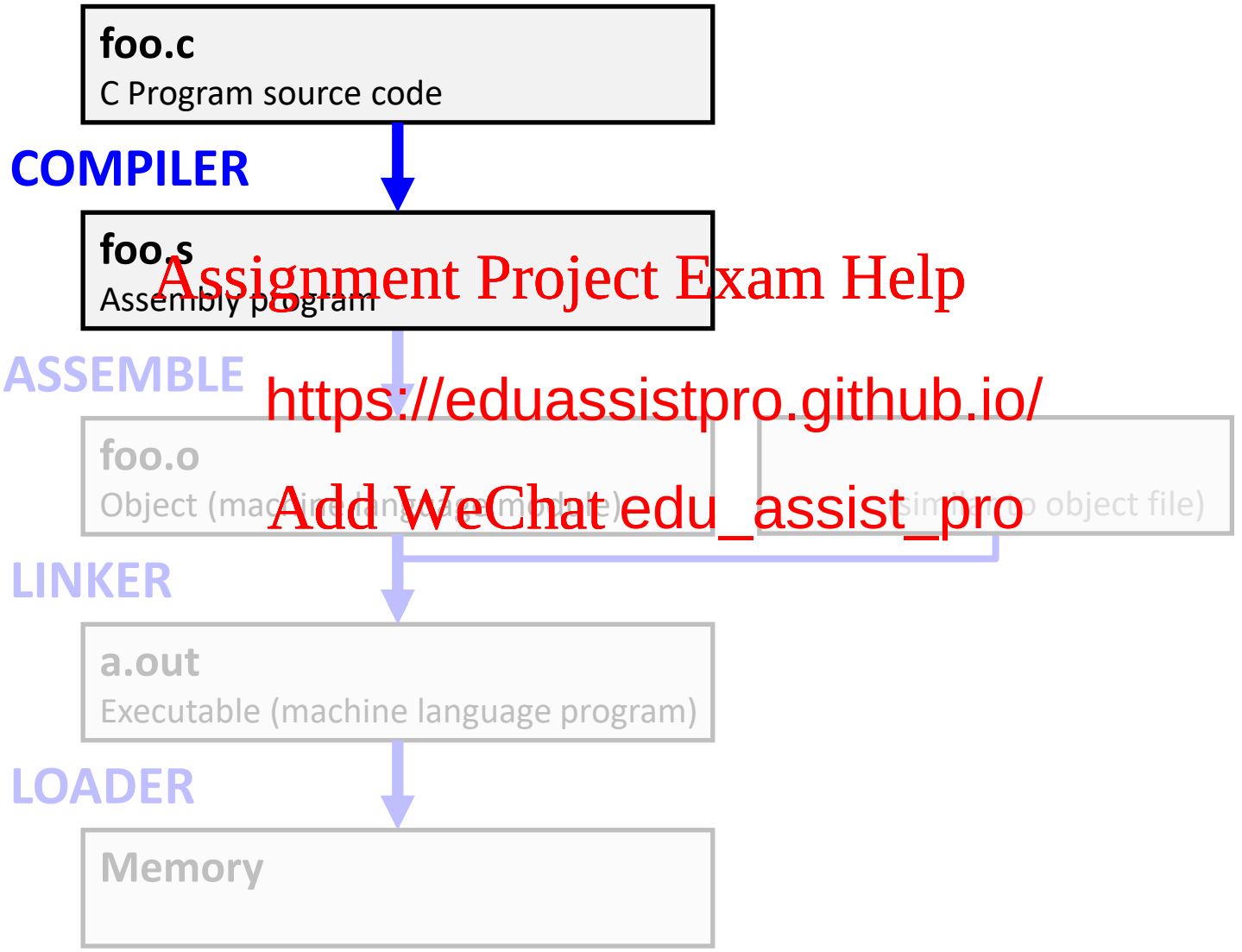
Outline

- Compiler
- Assembler
- Linker
- Loader
- Example

Steps to Starting a Program



Compiler

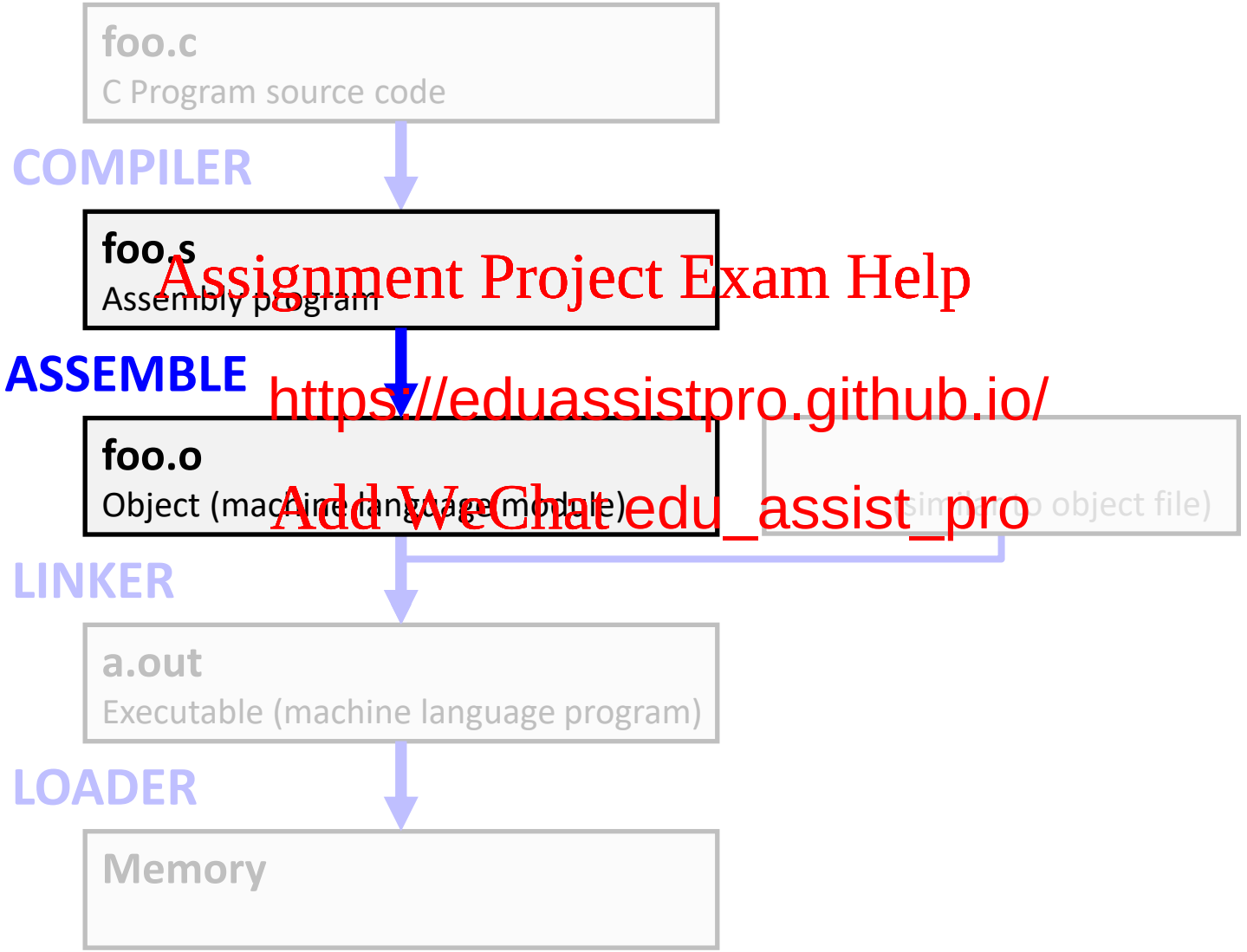
- Input: High-Level Language Code (e.g., C)
- Output: Assembly Language Code (e.g., MIPS)
  - In MARS we use **.as** ut **.s** is a common file extension for assem
- Note: Output *may* contain pseu ti ns
  - Assembler understands these intsructions, but not the machine

Compiler and Standards

- Compiler generates assembly code and directives that respect conventions
  - For example, functio tions
  - There are many mor ata representation and function linkage which are *beyond of this course*

Fine print… Application Binary Interface (ABI) standard formalizes these aspects, in addition to the encoding of object files and executable files (e.g., ELF).

Steps to Starting a Program



Assembler

- Reads and Uses Directives
- Replace Pseudoinstructions
- Produce Machine La
- Creates Object File

Assembler Directives (B.2, B.9, B.10)

- Directives provide directions to assembler, but do not produce machine instructions
  - `.text`: Subsequent it nstructions) segment
  - `.data`: Subsequent it segment
  - `.globl sym`: declares `sym` global ference from other files
  - `.asciiz str`: Store string `str` in memory and null-terminate it
  - `.word w1…wn`: Store *n* 32-bit words in successive memory locations

Pseudoinstruction Replacement

Assembler treats convenient variations of machine language instructions as if real (see B.10)

Pseudo (MAL):	Real (TAL):
`subu $sp,$sp,32`	
`sd $a0,32($sp)`	
	`sw $a0,0($sp)`
`addu $t0,$t6,1`	`addi $t0,$t6,1`
`ble $t0,100,loop`	`slti $at,$t0,101`
	`bne $at,$0,loop`
`la $a0,str`	`lui $at,left(str)`
	`ori $a0,$at,right(str)`
`mul $t7,$t6,$t6`	`mult $14, $14`
	`mflo $15`

NOTE: left() and right() to get lower and upper half worlds does not exist in MARS, but this is necessary in this example

Producing Machine Language (1/2)

- Simple instructions for Assembler
  - Arithmetic, Logical, Shifts, and so on
  - All necessa struction already
- What about
  - PC-Relative
  - Once pseudoinstructions are replaced by real ones, we know by how many instructions to branch
- So these 2 cases are handled easily

Producing Machine Language (2/2)

- What about jumps (`j` and `jal`)?
  - Jumps require **absolute address**
- What about references to data?
  - `la` gets broken up int
  - These will require the full 32-bit add data
- These can't be determined yet
  - Must wait to see where this code will appear in final program
- Two tables are used to help assembly and later resolution of addresses

1st Table: Symbol Table

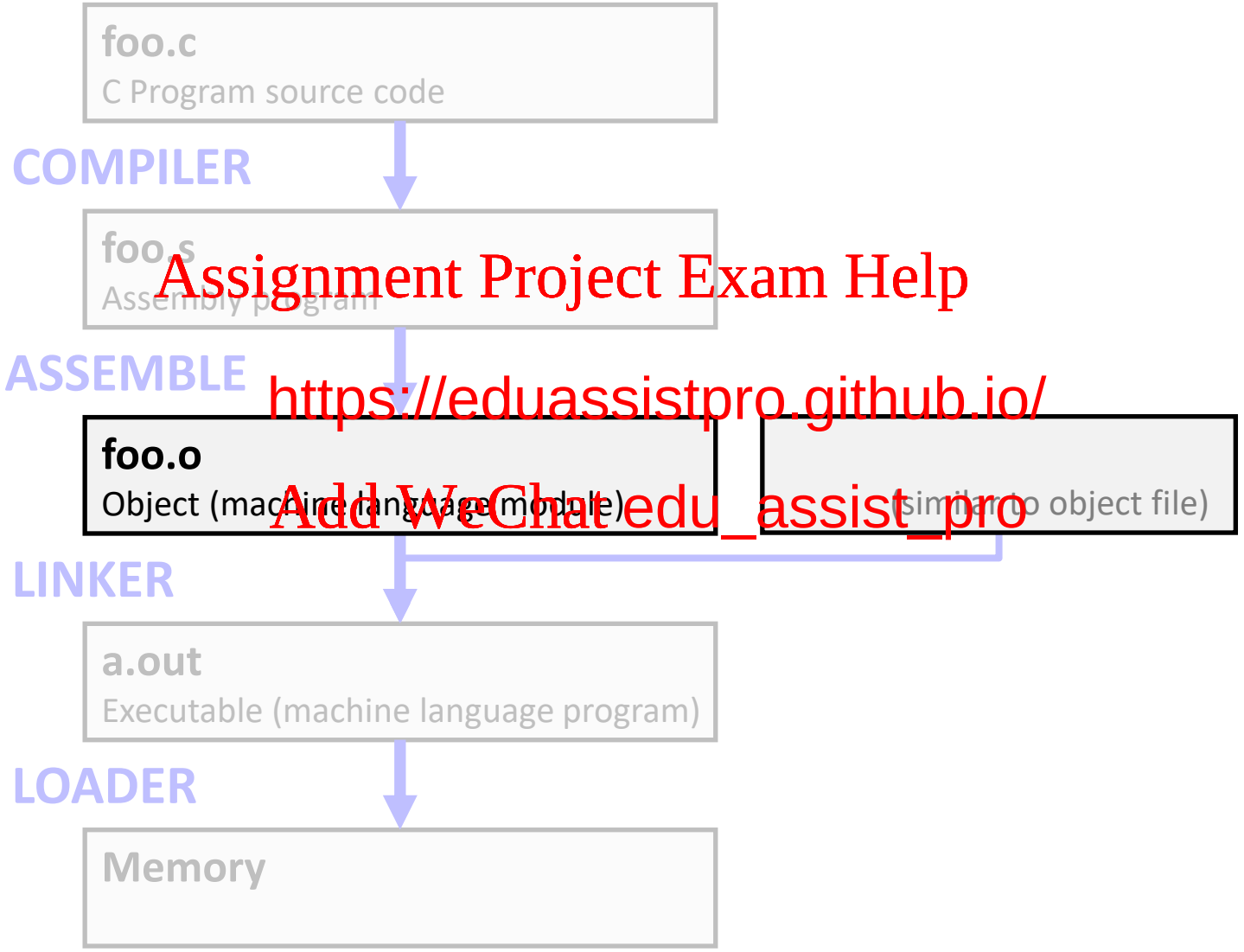
- Symbol table: List of "items" in this file that may be used by this and other files
- What are they?
  - Labels: function calli
  - Data: anything in the `.data` secti les which may be accessed across files
- First Pass: record label-address pairs
- Second Pass: produce machine code
  - Result: can jump to a label later in code without first declaring it

2nd Table: Relocation Table

- Relocation Table: line numbers of "items" in this file which need the address filled in (or fixed up) later.
- What are they
  - Any label ju
    - Internal (i.e. label inside th
    - external (including lib files)
  - Any absolute address of piece of data
    - Such as used by the load address **`la`** pseudo-instruction:

```
la $destination,label
```

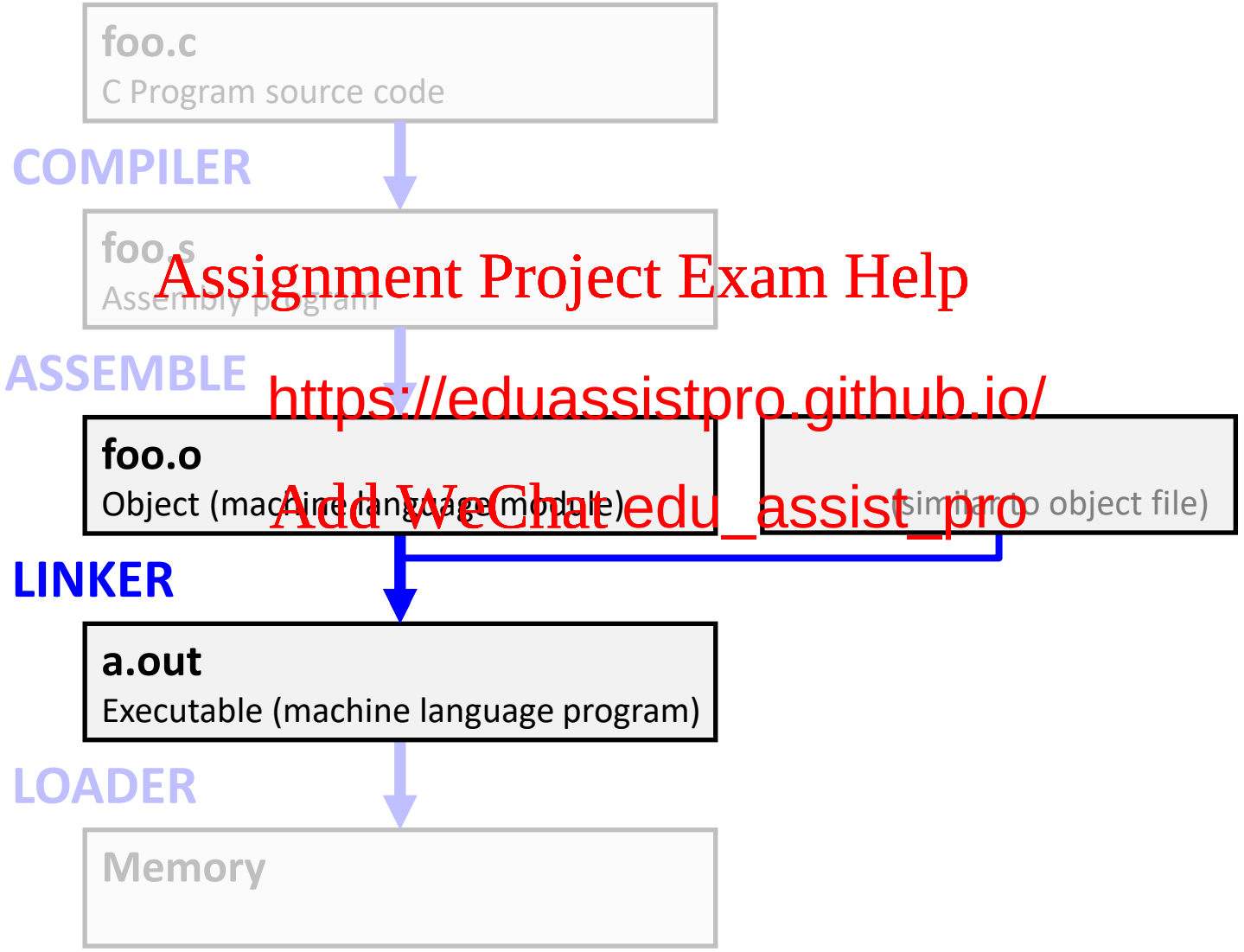
Steps to Starting a Program



Object File Format

- object file header: size and position of the other pieces of the object file
- text segment: the ma
- data segment: binary e data in the source file
- relocation table: identifies lines of need to be "handled"
- symbol table: list of this file's labels and data that can be referenced
- debugging information

Steps to Starting a Program



Link Editor/Linker (1/2)

- What does Link Editor do?
- Combines several object (.o) files into a single executable ("linking")
- Enables *Separate Co*
  - Changes to one file do not require _ tion of whole program
    - Linux kernel source: > 6 M lines of code
    - Windows OS source: > 40 M lines of code
  - Code in file called a module
  - Link Editor name from editing the "links" in jump and link instructions

Link Editor/Linker (2/2)

- Step 1: Combine text segment from each .o file
- Step 2: Combine data segment from each .o file, and concatenate this on ents
- Step 3: Resolve Refe
  - Go through Relocation Table
  - Handle each entry using the Symbol Table
    - That is, fill in all absolute addresses

Four Types of Addresses

- PC-Relative Addressing (**`beq`**, **`bne`**):
  - **never** fix up (never "relocate")
- Absolute Ad
  - always relo
- External Reference (usua
  - always relocate
- Symbolic Data Reference (often **`lui`** and **`ori`**):
  - always relocate

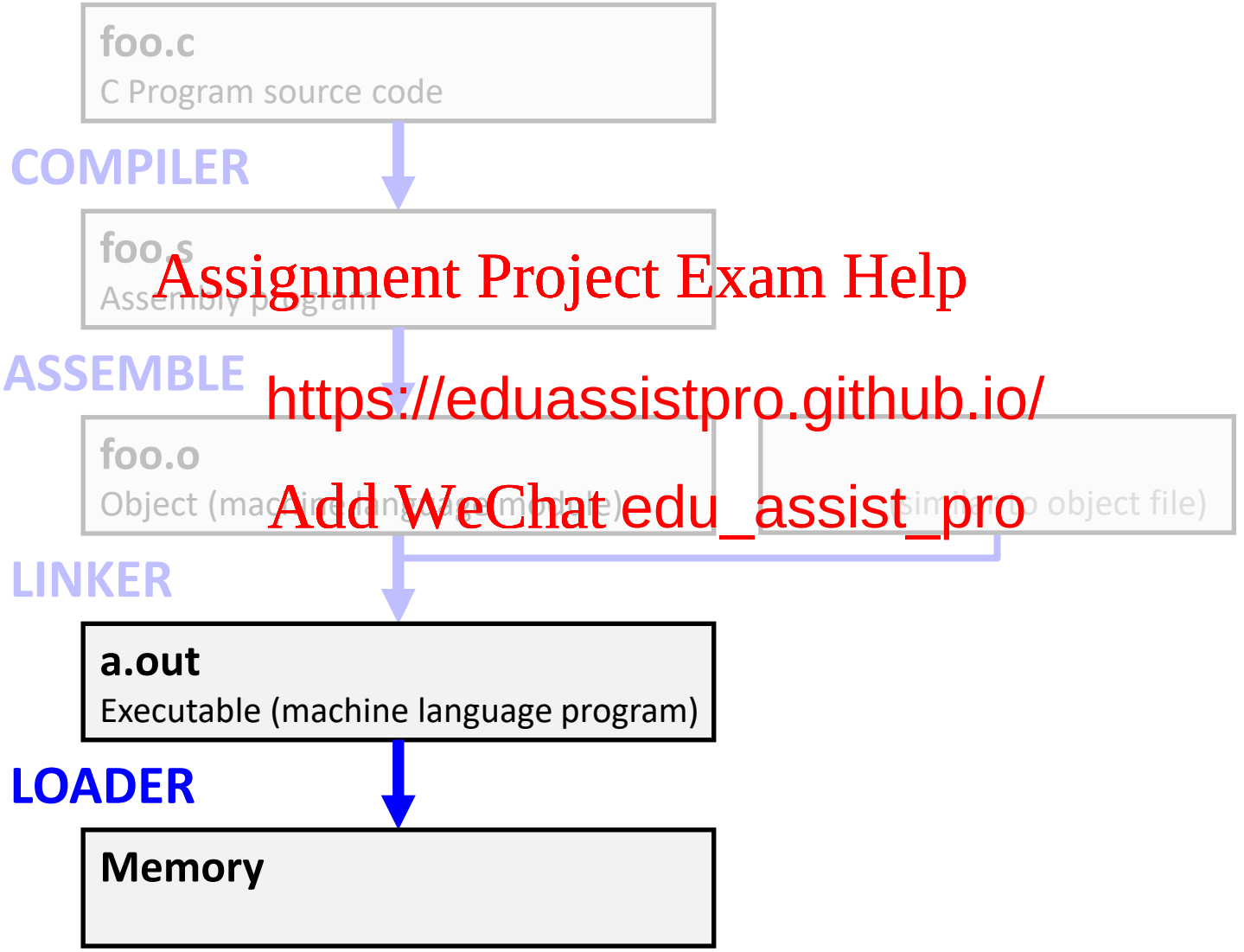
Resolving References (1/2)

- Linker *assumes* first word of first text segment is at address 0x00000000
- Linker knows:
  - Length of each text
  - Ordering of text and data segment
- Linker calculates:
  - Absolute address of each label to be jumped to (internal or external) and each piece of data being referenced

Resolving References (2/2)

- To resolve references:
  - Search for reference (data or label) in all symbol tables
  - If not found, search l le, for `printf`)
  - once absolute addre the machine code appropriately
- Output of linker:
  - Executable file containing text and data (plus a file header)
- May not have library object files resolved if dynamically loaded

Steps to Starting a Program



Loader (1/3)

- Executable files are stored on disk.
- When one is to be run, loader's job is to load it into memory and start it running.
- In reality, loader is t (OS)
  - Loading is one of the OS tasks

Loader (2/3)

- So what does a loader do?
- Reads executable file's header to determine size of text and data segments
- Creates new addres large enough to hold text and data segments, along w k segment
- Copies instructions and data from executable file into the new address space

Loader (3/3)

- Copies arguments passed to the program onto the stack
- Initializes machine registers
  - Most registers clear ust be initialized to top of the stack memory sp
- Jumps to start-up routine that c gram's arguments from stack to registers and sets the PC
  - If main routine returns, start-up routine terminates program with the exit system call

Dynamic Linking

- Some operating systems allow “dynamic linking”
- Both the loader *and* the linker are part of the operating system - so modules can be linke me
- If a module is needed it need not be loaded again
- Called DLLs in Windows, `.so` in Unix (Dynamically Linked Library / Shared Object)

C → Asm → Obj → Exe → Run

Compile C Source

Let us consider compilation of the following code…

```
#include <stdio.h>
int main (int argc, char *argv[]) {
        int i;
        int prod = 0;
        for (i = 0; i <= 100; i = i + 1) {
          prod = prod + i * i;
        }
        printf ("The sum squares from 0 .. 100 is %d\n", prod);
}
```

C → **Asm** → Obj → Exe → Run

Identify Pseudoinstructions

```
.text
   .align     2
   .globl     main
main:
   subu $sp,$sp,32
   sw   $ra, 20($s
   sd   $a0, 32($s
   sw   $0, 24($sp)
   sw   $0, 28($sp)
loop:
   lw   $t6, 28($sp)
   mul  $t7, $t6,$t6
   lw   $t8, 24($sp)
   addu $t9,$t8,$t7
   sw   $t9, 24($sp)
   addu $t0, $t6, 1
```

```
   sw   $t0, 28($sp)
   ble  $t0,100, loop
   la   $a0, str
   lw   $a1, 24($sp)
   jal printf
              , $0
              28($sp)
              $sp,32
.d
.align 0
str:
.asciiz       "The product
   from 0 .. 100 is %d\n"
```


FINE PRINT: The modification of the stack pointer may look strange, but this is ultimately from a *real example of compilation*… a number of the real details are being omitted here (ABI,etc.), some of which we will see later.

C → Asm → **Obj** → Exe → Run

Remove Pseudoinstructions, Assign Addresses

```
00 addiu $29,$29,-32
04 sw  $31,20($29)
08 sw  $4, 32($29)
0c sw  $5, 36($29)
10 sw    $0, 24($29)
14 sw    $0, 28($29)
18 lw   $14, 28($29)
1c mult $14, $14
20 mflo        $15
24 lw   $24, 24($29)
28 addu $25,$24,$15
2c sw   $25, 24($29)
30 addiu $8,$14, 1
34 sw  $8,28($29)
38 slti       $1,$8, 101
3c bne $1,$0, loop
```

```
40 lui $4, l.str
44 ori $4,$4, r.str
48 lw  $5,24($29)
4c jal printf
50 addu       $2, $0, $0
          1,20($29)
          9,$29,32
5c
```

C → Asm → **Obj** → Exe → Run

Symbol Table Entries

- Symbol Table

Label	Address
main:	0x00000000
loop:	0x000
str:	0x10000430
printf:	-

- Relocation Table

Address	Instruction/Type	Dependency
0x0000004c	jal	printf

C → Asm → **Obj** → Exe → Run

Edit Local Addresses

```
00 addiu $29,$29,-32
04 sw    $31,20($29)
08 sw    $4, 32($29)
0c sw    $5, 36($29)
10 sw    $0, 24($29)
14 sw    $0, 28($29)
18 lw    $14, 28($29)
1c multu $14, $14
20 mflo  $15
24 lw    $24, 24($29)
28 addu  $25,$24,$15
2c sw    $25, 24($29)
30 addiu $8,$14, 1
34 sw    $8,28($29)
```

```
38 slti  $1,$8, 101
3c bne   $1,$0, -10
40 lui   $4, 0x1000
44 ori   $4,$4,0x0430
           5,24($29)
5                $0, $0
5                20($29)
5               ,$29,32
5c jr    $31
```


Can fix several of these labels now, while others (0x4c) are left for later

C → Asm → **Obj** → Exe → Run

```
0x000000  00100111101111011111111111100000
0x000004  10101111101111110000000000010100
0x000008  10101111101001000000000000100000
0x00000c  10101111101001010000000000100100
0x000010  10101111101000000000000000011000
0x000014  10101111101000000000000000011100
0x000018  10001111101011100000000000011100
0x0                                   11001
0x0                                   10010
0x0                                   11001
0x0                                   00001
0x00002c  10101111101010                00
0x000030  00100011110010                01
0x000034  10101111101110                00
0x000038  00101001000000010000000001100101
0x00003c  00010100001000001111111111110111
0x000040  00111100000001000001000000000000
0x000044  00110100100001000000010000110000
0x000048  10001111101001010000000000011000
0x00004c  00001100000100000000000011101100
0x000050  00000000000000000001000000100001
0x000054  10001111101111110000000000010100
0x000058  00100111101111010000000000100000
0x00005c  00000011111000000000000000001000
```

C → Asm → Obj → **Exe** → Run

- Combine with object file containing “printf”
- Edit absolute addresses
  - In this case edit **`jal`** actual address of printf
- Output single binary

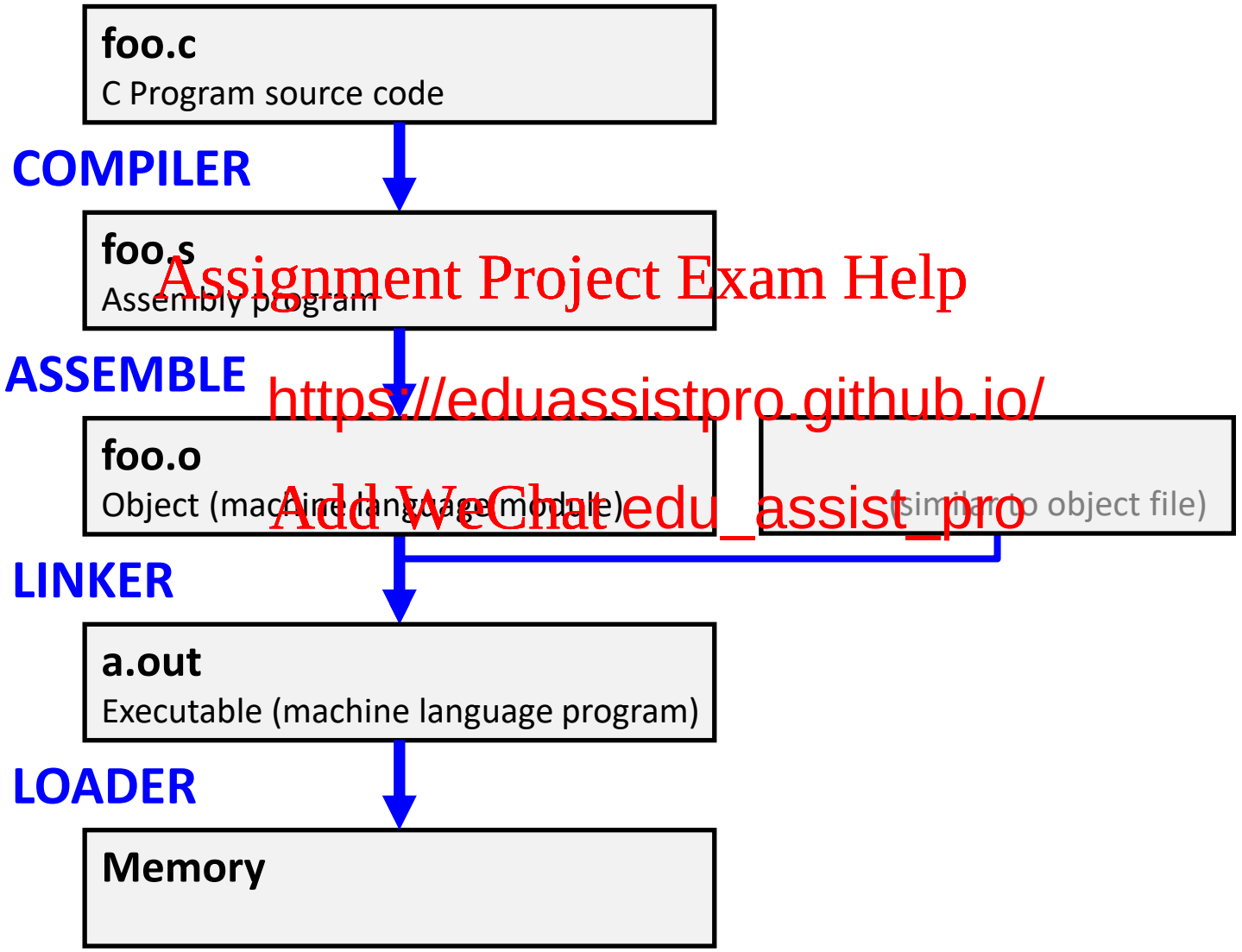
Things to Remember 1/3

- ***Stored Program concept*** means instructions just like data, so we can take data from storage, and keep transforming it until we're ready to load registers and jump to routine to begin execution
- Compiler → Assembler → Linke der)
- Assembler does ***2 passes*** to reso sses, handling internal forward references
- Linker enables ***separate compilation***, libraries that need not be compiled, and resolves remaining addresses

Things to Remember (2/3)

- Compiler converts a single HLL file into a single assembly language file
- Assembler removes , converts what it can to machine languag cklist for the linker (relocation table). This changes le into a .o file
- Linker combines several .o files and resolves absolute addresses
- Loader loads executable into memory and begins execution

Steps to Starting a Program



Review and More Information

- Textbook 5th edition, A.2 and A.3
  - (B2 and B3 of 4th edition)
- Chapter 2 Section 1 arting your program.

hello.c with gcc on Windows 10

```
/* hello.c */
#include <stdio.h>
int main( int ar                    ) {
    printf("Hell
}
```

hello.s X86 Assembly

```
        .file   "hello.c"
        .def    ___main;    .scl    2;      .type   32;     .endef
        .section .rdata,"dr"
LC0:
        .ascii "Hello COMP273\0"
        .text
        .globl  _main
        .def    _main;      .scl    2;      .type   32;     .endef
_main:
LFB13:
        .cfi_startproc
        pushl   %ebp
        .cfi_def_cfa_offset 8
        .cfi_offset 5, -8
        movl    %esp, %ebp
        .cfi_def_cfa_register 5
        andl    $-16, %esp
        subl    $16, %esp
        call    ___main
        movl    $LC0, (%esp)
        call    _printf
        movl    $0, %eax
        leave
        .cfi_restore 5
        .cfi_def_cfa 4, 4
        ret
        .cfi_endproc
LFE13:
        .ident  "GCC: (Rev3, Built by MSYS2 project) 5.2.0"
        .def    _printf;    .scl    2;      .type   32;     .endef
```

hello.o

Portable Executable (PE) Examined with Browser

Running the example