

# Assignment Project Exam Help

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## Motivation

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We've already seen

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How do we come up with correctness properties in the first place?

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## Structure of a Module

A Haskell program will usually be made up of many modules, each of which exports one or more *data types*.

Typically a module for a data type  $X$  will also provide a set of functions, called *operations*, on

- <https://eduassistpro.github.io/>
- $X$
- to update the data type:

A lot of software can be designed with this structure.

### Example (Data Types)

A dictionary data type, with empty, insert and lookup.

## Data Invariants

One source of properties is *data invariants*.

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### Data Invariants

Data invariants are

Whenever we use  
invariants are mai

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### Example

- That a list of words in a dictionary is always in sorted order.
- That a binary tree satisfies the search tree properties.
- That a date value will never be invalid (e.g. 31/13/2019).

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## Properties for Data Invariants

For a given data type  $X$ , we define a *wellformedness predicate*

$wf :: X \rightarrow \text{Bool}$

For a given value

$x$ .

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### Properties

For each operation, if all input values of type  $X$  satisfy  $wf$ .

then the result  $x$  satisfies

In other words, for each constructor operation  $c$

$wf(c \dots)$ ,

and for each update operation  $u :: X \rightarrow X$  we must show  $wf\ x \implies wf(u\ x)$

**Demo:** Dictionary example, sorted order.

## Stopping External Tampering

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Even with our sorted dictionary example, there's nothing to stop a malicious or clueless programmer from

### Example

The malicious programmer could just add a word directly to the dictionary, unsorted, bypassing our carefully written `insert` function.

We want to prevent this sort of thing from happening.

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## Abstract Data Types

An *abstract* data type (ADT) is a data type where the implementation details of the type and its associated operations are hidden.

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## Abstract Data Types

An *abstract* data type (ADT) is a data type where the implementation details of the type and its associated operations are hidden.

```
newtype Dict
```

```
type Word = String
```

```
type Definition
```

```
emptyDict :: Dict
```

```
insertWord :: Word -> Definition -> Dict -> Dict
```

```
lookup :: Word -> Dict -> Maybe Definition
```

If we don't have access to the implementation of the provided operations, which we know preserve our data invariants cannot be violated if this module is correct.

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## Abstract Data Types

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```
newtype Dict
type Word = String
type Definition
emptyDict :: Dict
insertWord :: Word -> Definition -> Dict -> Dict
lookup :: Word -> Dict -> Maybe Definition
```

If we don't have access to the implementation of the provided operations, which we know preserve our data invariants cannot be violated if this module is correct.

**Demo:** In Haskell, we make ADTs with module headers.

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## Abstract? Data Types

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In general, *abs*

The inverse of *abs*:

Abstract data type

implementation details are hidden, and we no longer have to re

level of implementation.

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## Validation

Suppose we had a `sendEmail` function

```
sendEmail :: String -- email address
```

-

-

It is possible to mix the two, it's possible that the given email address is not valid.

### Question

Suppose that we wanted to make it impossible to call `sendEmail` without first checking that the email address was valid. How would we accomplish this?

## Validation ADTs

We could define a tiny ADT for validated email addresses, where the data invariant is that the contained email address is valid:

```
module EmailADT (Email, checkEmail, sendEmail)
```

```
  newtype E
```

```
  checkEm
```

```
  checkEm
```

```
    | otherwise = Nothing
```

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Then, change the type of sendEmail:

```
  sendEmail :: Email -> String -> IO()
```

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The only way (outside of the EmailADT module) to create a value of type Email is to use checkEmail.

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Then, change the type of sendEmail:

```
  sendEmail :: Email -> String -> IO()
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The only way (outside of the EmailADT module) to create a value of type Email is to use checkEmail.

checkEmail is an example of what we call a *smart constructor*: a constructor that enforces data invariants.

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## Reasoning about ADTs

Consider the following, more traditional example of an ADT interface, the unbounded queue.

```
data Queue
```

```
emptyQueue :: Q
enqueue    :: Int -> Q
front      :: Queue -> Int
dequeue    :: Queue -> Queue -- partial
size       :: Queue -> Int
```

We could try to come up with properties that relate these functions without reference to their implementation, such as:

$$\text{dequeue } (\text{enqueue } x \text{ emptyQueue}) == \text{emptyQueue}$$

However these do not capture functional correctness (usually).

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## Models for ADTs

We could imagine a simple implementation for queues just in terms of lists:

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We could imagine a simple implementation for queues just in terms of lists:

```
emptyQueueL = []
enqueueL a   = (++ [a])
frontL       = head
dequeueL     = tail
sizeL        = length
```

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But this implementation is  $\mathcal{O}(n)$  to enqueue! Unacce

**However!**

This is a dead simple implementation, and trivial to see that it is correct. For a better queue implementation, it should always give the same results as this simple one. Therefore: This implementation serves as a **functional correctness specification** for our Queue type!

## Refinement Relations

The typical approach to connect our model queue to our Queue type is to define a relation, called a *refinement relation*, that relates a Queue to a list and tells us if the two structures represent the same queue conceptually:

```
rel :: Queue -> [In
```

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Then, we show tha

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prop_empty_
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That any query functions for our two types produce equal result  
such as for size:

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prop_size_r fq lq = rel fq lq ==> size fq == size lq
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```
prop_size_r fq lq = rel fq lq ==> size fq == size lq
```


And that each of the queue operations preserves our refinement relation, for example for enqueue:

```
prop_enq_ref fq lq x =  
  rel fq lq ==> rel (enqueue x fq) (enqueueL x lq)
```



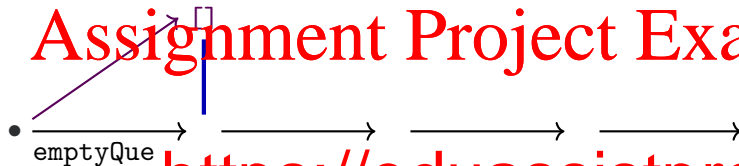
## In Pictures

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## In Pictures

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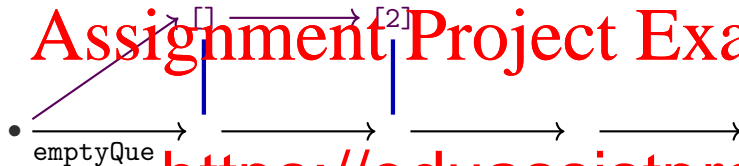


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`prop_empty_r = rel emptyQueue emptyQueue`

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$$\text{prop\_eng\_ref } fq, lq, x =$$

$$\text{rel } fq, lq \Rightarrow \text{ref } (\text{enqueue } x, fq) (\text{enqueue } x, lq)$$

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## In Pictures



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```
prop_enqueue_ref fq lq x =
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```

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## In Pictures

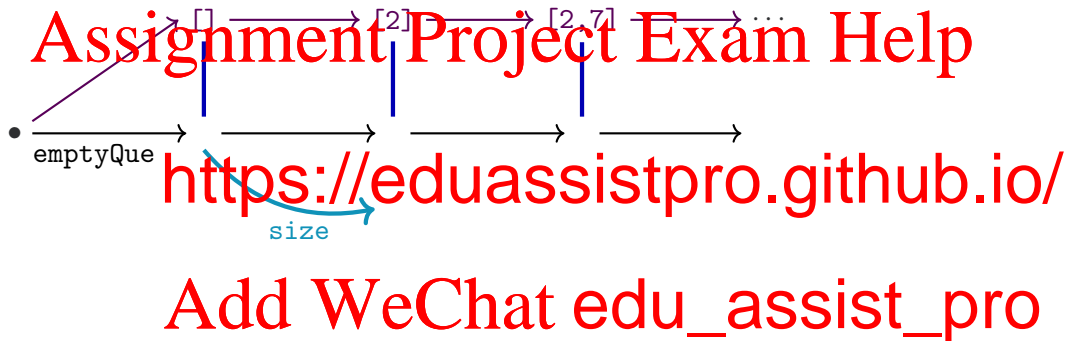


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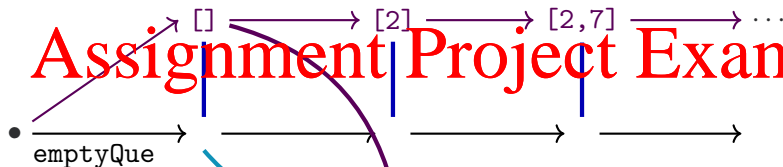
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## In Pictures



## In Pictures



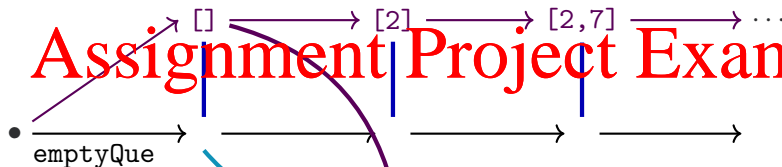
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`prop_size_r fq lq = rel fq lq ==> size fq == sizeL lq`

## In Pictures



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`prop_size_r fq lq = rel fq lq ==> size fq == sizeL lq`

Whenever we use a Queue, we can reason as if it were a list!



## Abstraction Functions

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These refinement relations are very difficult to use with QuickCheck because the  $\text{rel}$   $\text{fq}$   $\text{lq}$  preconditions are very hard to satisfy with randomly generated inputs.

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For this example, it  
the corresponding

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Conceptually, our refinement relation is then just:

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$\lambda \text{fq} \text{lq} \rightarrow \text{absfun } \text{fq} =$

However, we can re-express our properties in a much more QC-friendly format ([Demo](#))

## Fast Queues

Let's use test-driven development! We'll implement a fast Queue with amortised  $\mathcal{O}(1)$  operations.

```
data Queue
```

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```
Int    -- size of the rear
```

We store the rear part of the queue in *reverse order*, to make enqueue

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```
Int -- size of the rear
```

We store the rear part of the queue in *reverse order*, to make enqueue

Thus, converting from our Queue to an abstract list is

```
toAbstract :: Queue -> [Int]
```

```
toAbstract (Q f sf r sr) = f ++ reverse r
```

## Data Refinement

These kinds of properties establish what is known as a *data refinement* from the abstract, slow list model to the fast, concrete Queue implementation.

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### Refinement and Specifications

In general, all *funct*

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There is a limit to the amount of abstraction we can do before they b testing (but not necessarily for proving).

### Warning

While abstraction can simplify proofs, abstraction does not reduce the fundamental complexity of verification, which is provably hard.

## Data Invariants for Queue

In addition to the already-stated refinement properties, we also have some data invariants to maintain for a value  $Q$ :

- 1  $\text{length } f == sf$

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In addition to the already-stated refinement properties, we also have some data invariants to maintain for a value  $Q$  of type  $\text{sf} \times \text{sr}$ :

- 1 `length f == sf`
- 2 `length r == sr`
- 3 `important`

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② `length r == sr`

③ **important**

We will ensure our invariants.

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## Data Invariants for Queue

In addition to the already-stated refinement properties, we also have some data invariants to maintain for a value  $Q$  of type  $\text{Queue}$ :

- 1 `length f == sf`
- 2 `length r == sr`
- 3 **important**

We will ensure our invariants.

Thus, our wellformed predicate is used merely to ensure the outputs of our operations:

```
prop_wf_empty = wellformed (emptyQueue)
prop_wf_enq q = wellformed (enqueue x q)
prop_wf_deq q = size q > 0 ==> wellformed (dequeue q)
```

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## Implementing the Queue

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We will generally implement by:

- Dequeue from the back
- Enqueue to the back
- If necessary, dequeue from the back and appending it to the front.

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## Implementing the Queue

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We will generally implement by:

- Dequeue from the back
- Enqueue to the back
- If necessary, double the array and append it to the front.

This step is slow ( $O(n)$ ) but only happens every  $n$  operations, so the average case amortised complexity of  $O(1)$  time.

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## Amortised Cost

```
enqueue x (Q f sf r sr) = inv3 (Q f sf (x:r) (sr + 1))
```

When we enqueue each of [1..7] to the empty Queue in turn:

Q []

0 []

0

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|   |    |   |    |     |
|---|----|---|----|-----|
| Q | [] | 0 | [] | 0   |
|   |    |   |    | 0 ( |
|   |    |   |    | 1   |

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→ Q [1, 2, 3]      3 [5

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Q []      0 []      0  
          0 (

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→ Q [1, 2, 3]      3 [6

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When we enqueue each of [1..7] to the empty Queue in turn:

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Q [] 0 [] 0
0 (
1
```

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```
→ Q [1,2,3] 3 [5
→ Q [1,2,3] 3 [6
→ Q [1,2,3,4,5,6,7] 7 [] 0 (
```

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Q [] 0 [] 0
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```
→ Q [1, 2, 3] 3 [5
→ Q [1, 2, 3] 3 [6
→ Q [1, 2, 3, 4, 5, 6, 7] 7 [] 0 (
```

Observe that the slow invariant-reestablishing step (\*) happens after 1 step, then 2, then 4...

Extended out, this averages out to  $\mathcal{O}(1)$ .

## Another Example

Consider this ADT interface for a bag of numbers:

```
data Bag
```

```
emptyBag :: Bag
```

```
addToBag :: Int -> Bag
```

```
averageBag :: Bag -> Int
```

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Our conceptual a

```
emptyBagA = []
```

```
addToBagA x xs = x:xs
```

```
averageBagA [] = Nothing
```

```
averageBagA xs = Just (sum xs `div` length xs)
```

But do we need to keep track of all that information in our implementation?

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But do we need to keep track of all that information in our implementation? **No!**

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## Concrete Implementation

Our concrete version will just maintain two integers, the total and the count

```
data Bag = B { total :: Int , count :: Int }
```

```
emptyBag :: Bag
```

```
emptyBag = B 0 0
```

```
addToBag :: Int -> Bag -> Bag
```

```
addToBag x (B t c) = B (x + t) (c + 1)
```

```
averageBag :: Bag -> Maybe Int
```

```
averageBag (B _ 0) = Nothing
```

```
averageBag (B t c) = Just (t `div` c)
```

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## Refinement Functions

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Normally, writing an abstraction function (as we did for Queue) is a good way to express our refine

write such a functi

```
toAbstract :: B -> A
toAbstract (B t c) = ??????
```

Instead, we will go in the other direction, giving us a

```
toConc :: [Int] -> Bag
toConc xs = B (sum xs) (length xs)
```

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## Properties with Refinement Functions

Refinement functions produce properties much like abstraction functions, only with the abstract and concrete layers swapped:

```
prop_ref_em
```

```
  toConc emb
```

```
prop_ref_add x ab =
```

```
  toConc (addToBagA x ab) == addToBag x (toConc ab)
```

```
prop_ref_avg ab =
```

```
  averageBagA ab == averageBag (toConc ab)
```

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## Assignment 1 and Break

# Assignment Project Exam Help

Assignment 1 has b

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## Assignment 1 and Break

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Assignment 1 has b

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### Advice from Alu

The assignments do not involve much coding, but they do invol

*Start early!*

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## Homework

# Assignment Project Exam Help

- 1 Get started o
- 2 Next prog
- 3 Last week's
- 4 This week's quiz is also up, due the following Friday.

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