

COMP547A Homework set #4

Due Thursday December 1st, 2022, 23:59

Exercises (from Katz and Lindell's book)

[5%]

4.7 Let F be a pseudorandom function. Show that the following MAC for messages of length $2n$ is insecure: Gen outputs a uniform $k \in \{0, 1\}^n$. To authenticate a message $m_1 \| m_2$ with $|m_1| = |m_2| = n$, compute the tag $F_k(m_1) \| F_k(F_k(m_2))$.

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4.27 Define an appropriate notion of a ϵ -secure *two-time* MAC, and give a construction that meets your definition.

HOMEMADE Question: Achieving Rivest's private-key encryption from a Mac

[5%]

Provide a security definition of a **Mac** that makes the (bit-by-bit) private-key encryption scheme that Rivest described secure in the sense of indistinguishability in the presence of an eavesdropper.

More on back...

12.

[1%]

[4%]

Hint: Prove that if "not CPA-secure" then "DDH problem is efficiently solved".

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13.1 Show that Construction 4.7 for constructing a variable-length MAC from any fixed-length MAC can also be used (with appropriate modifications) to construct a signature scheme for arbitrary-length messages from any signature scheme for messages of fixed length $\ell(n) \geq n$.

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Alice and Bob are a bit confused. They are using a digital signature scheme as a **Mac**. Let $\Pi = (\text{Gen}, \text{Sig}, \text{Vrfy})$ be a digital signature scheme (such as hashed RSA) where $\text{Gen}(1^n)$ to obtain (p_k, s_k) but only share and use s_k as the private-key of a **Mac**.

[5%]

(A) Let $\Pi' = (\text{Gen}', \text{Mac}', \text{Vrfy}')$ be the **Mac** resulting from this idea. Used as a **Mac** they simply set $t := \text{Mac}'_{s_k}(m) := \text{Sig}_{s_k}(m)$. However, since they only use s_k , how will the receiver verify the message-tag pair (m, t) ? In other words, what is $\text{Vrfy}'_{s_k}(m, t)$? Why did I underlined the word "deterministic" above?

[5%]

(B) Show that if Π is a digital signature scheme existentially unforgeable under an adaptive chosen-message attack then Π' is a **Mac** existentially unforgeable under an adaptive chosen-message attack (whether p_k is made public or not).

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(C) Imagine that Alice and Bob use Π' as above, and that p_k is disclosed publicly. Explain how this defeats Rivest's argument seen in class that private-key authentication implies private-key encryption.

More on back...