COMP9319 Web Data Compression and Search

Semistructured / Tree Data,
XML,
XML Compression

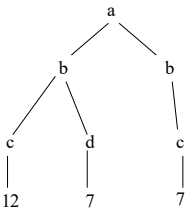
1

1

XPath evaluation

<a><c>12</c><d>7</d><c>7</c>

/ a / b [c = “12”]



2

2

Query evaluation

Top-down

Bottom-up

Hybrid

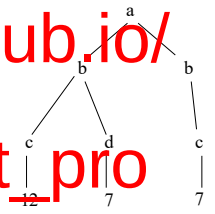
3

3

XPath evaluation

<a><c>12</c><d>7</d><c>7</c>

/ a / b [c = “12”]



4

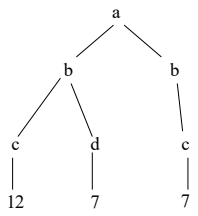
4

XPath evaluation

<a><c>12</c><d>7</d><c>7</c>

/ a / b [c = “12”]

<c>12</c><d>7</d>



5

5

Path indexing

- Traversing graph/tree almost = query processing for semistructured / XML data
- Normally, it requires to traverse the data from the root and return all nodes X reachable by a path matching the given regular path expression
- Motivation: allows the system to answer regular path expressions without traversing the whole graph/tree

6

6

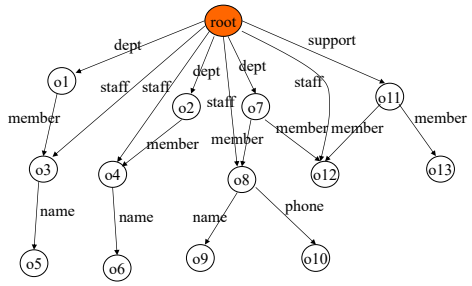
Major Criteria for indexing

- Speed up the search (by cutting the search space down)
- Relatively smaller size than the original data graph/tree
- Easy to maintain (during data loading during updates)

7

7

An Example of DAG Data



8

8

Index graph based on language-equivalence

- a reduced graph tha
  paths from the root i
- The paths from root
  - staff
  - dept/member
  - support/member

9

9

Language-equivalent nodes

a path from the root to x labeled w}

be infinite when there

guage-equivalent (x ≡

y taking the nodes

classes for ≡

10

10

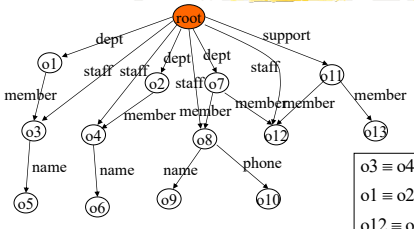
Language-equivalent

- The paths from root to o3
  - staff
  - dept/member
- Paths to o4 happen to be exactly the same 2 sequences
- Same for o8 and o12
- o3 ≡ o4 ≡ o8 ≡ o12

11

11

Equivalence classes

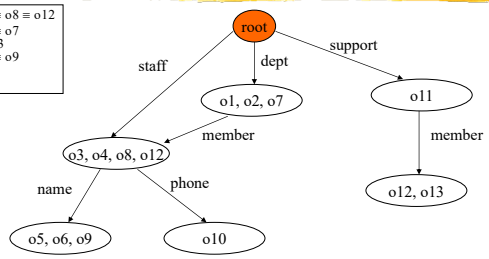


12

12

The index graph

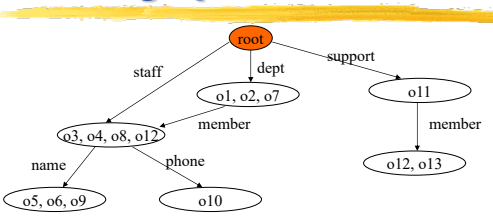
o3 ≡ o4 ≡ o8 ≡ o12
o1 ≡ o2 ≡ o7
o12 ≡ o13
o5 ≡ o6 ≡ o9
o10
o11



13

13

Query processing based on the index graph



dept/member/(name | phone)
-> dept/member/name UNION dept/member/phone
-> {o5, o6, o9} UNION {o10}
-> {o5, o6, o9, o10}

14

14

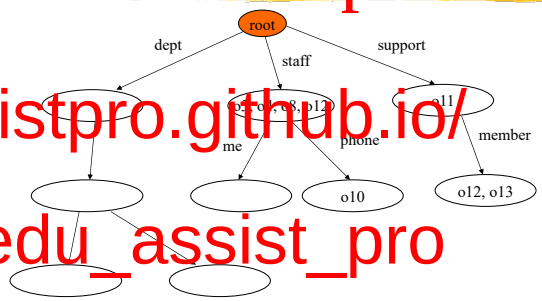
About this indexing scheme

- The index graph is n
- In practice, the inde enough to fit in me
- Construct the index i
  - check two nodes are language-equivalent is very expensive (are PSPACE
  - approximation based on bisimulation exists

15

15

A Data Guide



16

16

About Data Guide

- unique labels at each node
- (hence) extents are no longer disjoint
- query processing proceeds as before
- size of the index may >= data size
- good for data that is regular & has no cycles

17

17

XML-Specific Compressors

- Unqueriable Compression (e.g. XMill):
  - Full-chunked: data commonalities eliminated
  - Very good compression ratio
- Queriable Compression (e.g. XGrind, XPRESS):
  - Fine-grained: data commonalities ignored
  - Inadequate compression ratio and time
  - Support simple path queries with atomic predicate

18

18

XMill

- First specialized compressor for XML data
  - SAX parser (a fast, event-based parser) for parsing XML data
  - Still using gzip as its underlying compressor
  - Clever grouping of data into containers for compression
- Compress XML via three basic techniques
  - Compress the structure separately from the data
  - Group the data values according to their types
  - Apply semantic (specialized) compressors:

19

19

XMill Architecture:

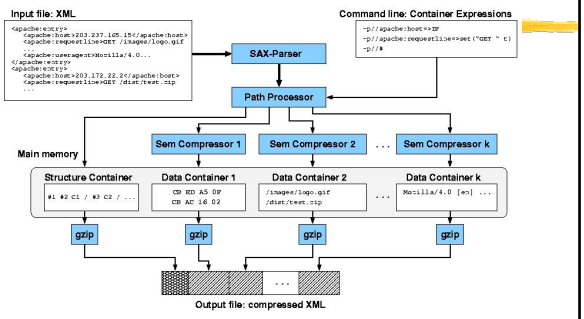


Figure 4: Architecture of the Compressor

20

20

An Example: Web Server Logs

ASCII File 15.9 MB (gzipped 1.6MB):

202.239.238.16|GET / HTTP/1.0|text/html|200|1997/10/01-00:00:02|-|4478|-|-|http://www.net.jp/|Mozilla/3.1[ja](I)

XML-ized apache web log inflates t

```
<apache:entry>
  <apache:host> 202.239.238.16 <
  <apache:requestLine> GET / HTT
  <apache:contentType> text/html </apache:contentType>
  <apache:statusCode> 200</apache:statusCode>
  <apache:date> 1997/10/01-00:00:0</apache:date>
  <apache:byteCount> 4478</apache:byteCount>
  <apache:referer> http://www.net.jp/ </apache:referer>
  <apache:userAgent> Mozilla/3.1$[$ja$]$(I)</apache:userAgent>
</apache:entry>
```

21

21

How Xmill Works: Three Ideas

arately from the data:

gzip Data

=1.75MB

22

22

How Xmill Works: Three Ideas

Group the data values according to their types:

gzip Structure

```
<apache:entry>
. . .
</apache:entry>
```

+

gzip Data1

```
202.23.23.16
224.42.24.55
...
```

+

gzip Data2

```
GET / HTTP/1.0
GET / HTTP/1.1
...
```

=1.33MB

23

23

How Xmill Works: Three Ideas

Apply semantic (specialized) compressors:

gzip Structure + gzip c1(Data1) + gzip c2(Data2) + ... =0.82MB

Examples:
- 8, 16, 32-bit integer encoding (signed/unsigned)
- differential compressing (e.g. 1999, 1995, 2001, 2000, 1995, ...)
- compress lists, records (e.g. 104.32.23.1 → 4 bytes)

Need user input to select the semantic compressor

24

24

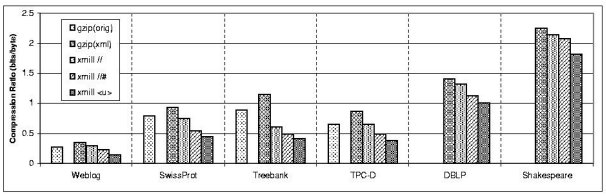
Experiments

Data Source	Original Size	Size in XML	Regular?	Depth	Tags	DataGuide Size
Weblog Data	57.7MB	172MB	yes	1	10	11
SwissProt	98.5MB	158MB	yes	3	92	58
Treebank	39.6MB	53.8MB	no	35	251	339920
TPC-D	34.6MB	119MB	yes	2	43	60
DBLP	-	47.2MB	yes	3	10	145
Shakespeare	-	7.3MB	no	5	21	58

25

25

XML Compression



26

26

Compression Time

27

27

Transfer Time (& Decode)

28

28

XGRIND (Tolani & Haritsa, 2002)

- Encodes elements and attributes using XMill's approach
- **DTD-conscious:** enumerated attributes with k possible values are encoded using a $\log_2 k$-bit scheme
- Data values are encoded using non-adaptive Huffman coding
  - Requires two passes over the input document
  - Separate statistical model for each element/attribute
- Homomorphic compression: compressed document retains original structure

29

XGRIND

Original Fragment:

```
<student name="Alice">
      <a1>78</a1>
      <a2>86</a2>

   <midterm>91</midterm>

   <project>87</project>
</student>
```

Compressed Fragment:

```
T0 A0 nahuff(Alice)
   T1 nahuff(78) /
   T2 nahuff(86) /
   T3 nahuff(91) /
   T4 nahuff(87) /
/
```


30

XGRIND

- Many queries can be carried out entirely in compressed domain
  - Exact-match, prefix-match
- Some others require only decompression of relevant values
  - Range, substring
- Queryability comes at the expense of achievable compression ratio: typically within 65-75% that of XMill