

11. Protection in the OS

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Segments

- It really makes no difference whether a bug that destroys somebody else's data is an accidental mistake or malicious.
 - a faulty process may also try to execute data or perform operations on machine code
- Solution is to ensure that physical memory areas used for different purposes are protected from each other. This is the primary function of the MMU.
- Paged virtual memory implement but:
 - Pages are fixed size
 - Areas we want to be managed by process logic
- Introduce idea of segments which are a part of memory in a process address space of any size determined by process logic:
- In virtual memory systems, a segment must be an integral number of pages.
- When a process is executing, it is not allowed to modify any word in memory that isn't within a segment owned by the process
- MMU *protection* operates at segment level, *physical placement* at page level.

Violations of Memory Protection

- All memory accesses must be checked to see if a **segment violation** has occurred.
 - If so a memory fault is signalled.
 - In a virtual memory system the handler will check if a page fault has occurred, but if there is no page on disk for this access, process is aborted before damage is done.
- Since this must be done every memory cycle it must be done by the MMU.
- The MMU maintains an effective address and a size. If a memory access is made to a segment, the MMU generates a memory fault if the segment is not currently active.
- The MMU is a device with its own registers and is controlled by the CPU.
 - What if a faulty or malicious program accesses the MMU?
 - The MMU must be especially protected...
 - Only the OS should be allowed to modify the MMU.
 - How do we stop other programs from doing so?

User and System Mode

- Hardware support is needed to control access to the MMU
 - At all times, the processor is either in **user mode** or in **system mode**
 - The instructions that access the MMU are called **privileged instructions**
 - Privileged instructions can only be executed in system mode
- To implement this, the CPU contains a flip flop (called the **CPU mode bit**) which indicates the current mode. This is often in the **status register**.
- The hardware executes the instructions in **user mode** until it encounters a privileged instruction. When this happens, the processor switches to **system mode**.
 - If in system mode, the processor executes the instruction.
 - If in user mode, cause an exception which transfers control to the OS.
- It is essential to ensure that the OS can execute privileged instructions in system mode.
- User processes must not be able to change to system mode without passing control back to the OS. (To do otherwise would be a hacker's dream!)
- On every exception the CPU automatically flips to system mode. The OS controls the location and content of all handler routines.
- A privileged instruction is used to set user mode before a user process is run.

More on System Mode

- Every operation that might be dangerous must be executed in system mode
 - Only the Operating System is allowed to perform these operations
 - The user program can only request such an operation to be performed
- Every **I/O operation** is potentially dangerous
 - I/O is complicated, so bugs are likely if users do it themselves
 - Users don't want to write I/O code anyway, since it's so complicated
 - Incorrect I/O could be destroying data
- When your program wants to perform a **request**: a data structure that describes what it really does is build a
 - Then your program executes a TRAP.
 - The OS takes control from your program and examines the request you made.
 - If a legitimate request (e.g. read a line from my own file) OS then executes privileged instructions to perform the action
 - If not legitimate (e.g. read data from somebody else's file) OS refuses the action!

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Levels of Protection

- We have examined the low level basis of protection
 - Implementation requires cooperation between hardware and OS.
 - But there are many higher levels of protection which an OS uses to determine whether a requested operation is legitimate (**protection policy**).
- We want to have both
 - **Security** so that is prevented
 - **Flexibility** so that convenience
- When an OS supports multiple users, each process belongs to a specific user and each user has a level that determines what his/her processes are allowed to do
 - **Administrator** (**root**) can do almost anything
 - Ordinary users can modify only their own data and read data of other users who permit it.

Files

- Data in secondary memory is stored in **files**.
 - A file is a set of logically associated sectors.
- The MMU and system/user modes give protection to information in memory while a program is running
- Data that is stored permanently in files needs **file protection**, which is enforced by the OS
- File access is essential
 - User programs are absolutely prohibited from doing any I/O.
 - To do I/O, a user program sends a request to the OS through a request rating system.
 - The OS checks the request, to determine whether the user is the owner of the file; if the request is OK, the OS performs the actual machine language instructions to do the I/O

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