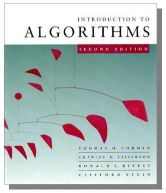
Analysis of Algorithms



LECTURE 23

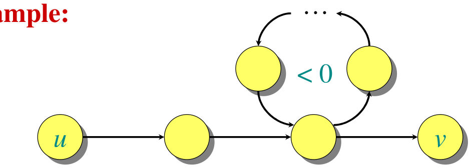
Shortest Paths II

- Bellman-Ford algorithm
- Linear programming and difference constraints
- VLSI layout compaction

Negative-weight cycles

Recall: If a graph $G = (V, E)$ contains a negative-weight cycle, then some shortest paths may not exist.

Example:



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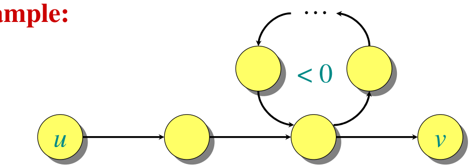


Bellman-Ford algorithm: Finds all shortest-path lengths from a ***source*** $s \in V$ to all $v \in V$ or determines that a negative-weight cycle exists.

Bellman-Ford algorithm

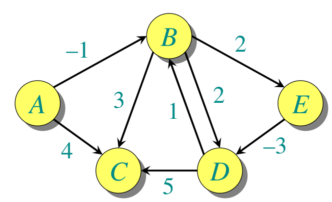
$d[s] \leftarrow 0$
for each $v \in V - \{s\}$
  do $d[v] \leftarrow \infty$
} initialization

for $i \leftarrow 1$ **to** $|V| - 1$
  do for each edge $(u, v) \in E$
    do if $d[v] > d[u] + w(u, v)$
      then $d[v] \leftarrow d[u] + w(u, v)$
} ***relaxation step***

for each edge $(u, v) \in E$
  do if $d[v] > d[u] + w(u, v)$
    then report that a negative-weight cycle exists

At the end, $d[v] = \delta(s, v)$, if no negative-weight cycles. Time $= O(VE)$.

Example of Bellman-Ford



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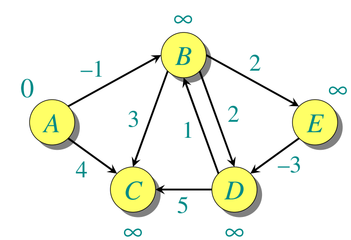
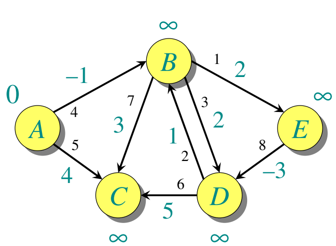


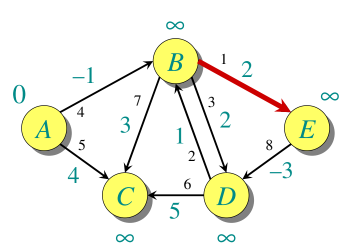
Initialization.

Example of Bellman-Ford



Order of edge relaxation.

Example of Bellman-Ford



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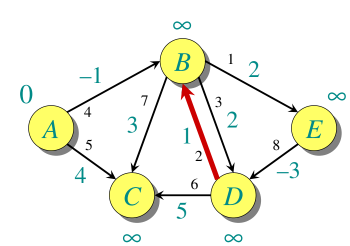

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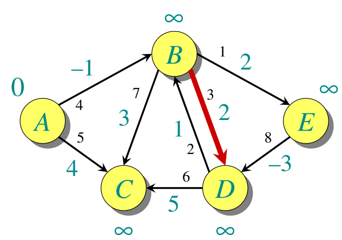

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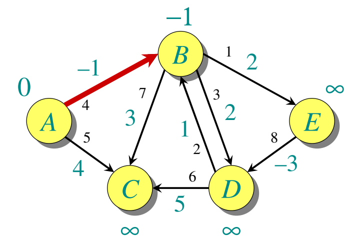

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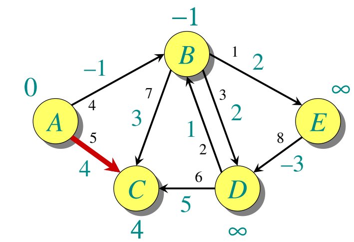

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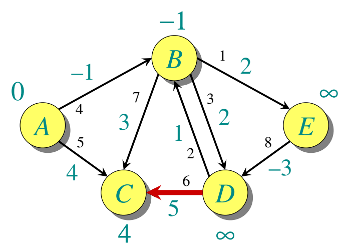

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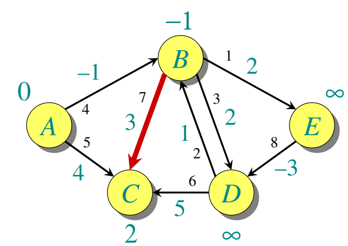

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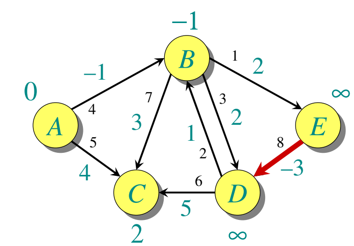

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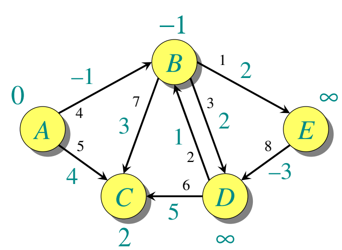
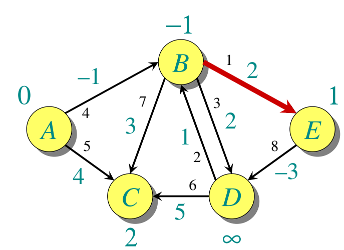


End of pass 1.

Example of Bellman-Ford



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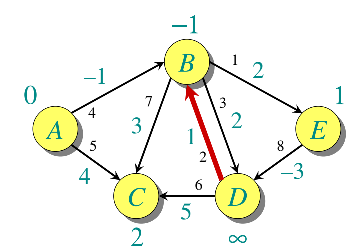

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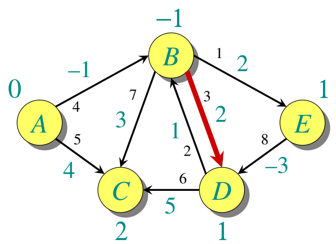

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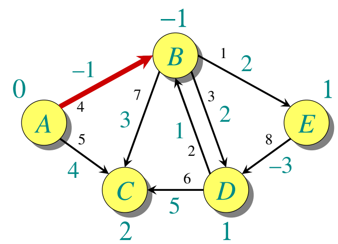

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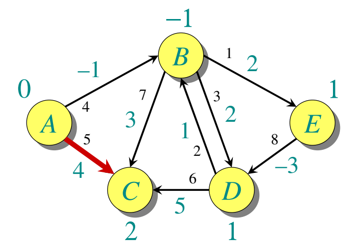

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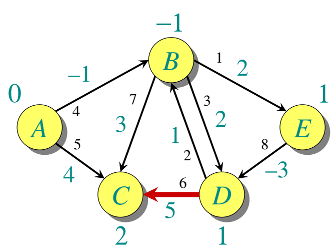

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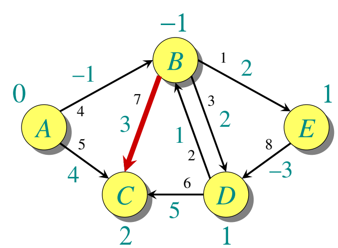

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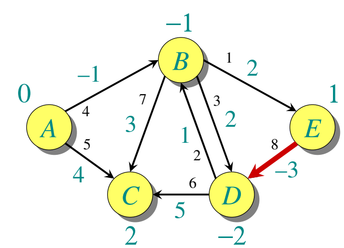

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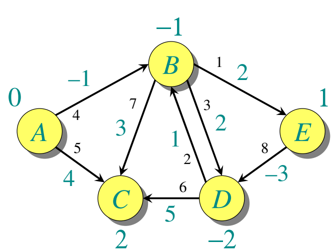

End of pass 2 (and 3 and 4).

Correctness

Theorem. If $G = (V, E)$ contains no negative-weight cycles, then after the Bellman-Ford algorithm executes, $d[v] = \delta(s, v)$ for all $v \in V$.

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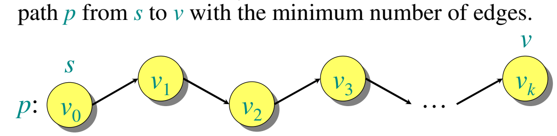
Theorem. If $G = (V, E)$ contains no negative-weight cycles, then after the Bellman-Ford algorithm executes, $d[v] = \delta(s, v)$ for all $v \in V$.

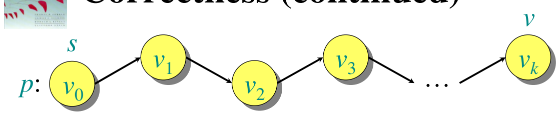
Proof. Let $v \in V$ be any vertex, and consider a shortest path p from s to v with the minimum number of edges.



Since p is a shortest path, we have

$$\delta(s, v_i) = \delta(s, v_{i-1}) + w(v_{i-1}, v_i) .$$

Correctness (continued)



Initially, $d[v_0] = 0 = \delta(s, v_0)$, and $d[v_0]$ is unchanged by subsequent relaxations (because of the lemma from *Shortest Paths I* that $d[v] \ge \delta(s, v)$).

- After 1 pass through E, we have $d[v_1] = \delta(s, v_1)$.
- After 2 passes through E, we have $d[v_2] = \delta(s, v_2)$.

⋮

- After k passes through E, we have $d[v_k] = \delta(s, v_k)$.

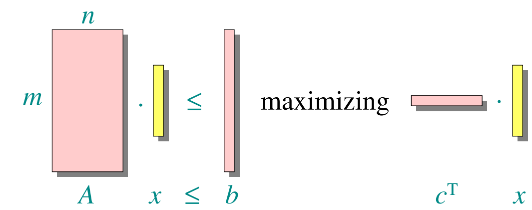
Since G contains no negative-weight cycles, p is simple. Longest simple path has $\le |V| - 1$ edges. □

Detection of negative-weight cycles

Corollary. If a value $d[v]$ fails to converge after $|V| - 1$ passes, there exists a negative-weight cycle in G reachable from s. □

Linear programming

Let A be an $m \times n$ matrix, b be an m-vector, and c be an n-vector. Find an n-vector x that maximizes $c^T x$ subject to $Ax \le b$, or determine that no such solution exists.



Linear-programming algorithms

Algorithms for the general problem

- Simplex methods — practical, but worst-case exponential time.
- Interior-point methods — polynomial time and competes with simplex.

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Feasibility problem: No optimization criterion. Just find x such that $Ax \le b$.

- In general, just as hard as ordinary LP.

Solving a system of difference constraints

Linear programming where each row of A contains exactly one 1, one –1, and the rest 0's.

Example:

$$\left.\begin{aligned} x_1 - x_2 &\le 3 \\ x_2 - x_3 &\le -2 \\ x_1 - x_3 &\le 2 \end{aligned}\right\} \; x_j - x_i \le w_{ij}$$

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Solution:

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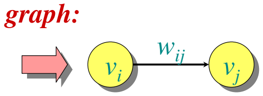
Constraint graph:

$x_j - x_i \le w_{ij}$



(The "A" matrix has dimensions $|E| \times |V|$.)

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Theorem. If the constraint graph contains a negative-weight cycle, then the system of differences is unsatisfiable.

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Proof. Suppose that the negative-weight cycle is $v_1 \to v_2 \to \cdots \to v_k \to v_1$. Then, we have

$$\begin{aligned} x_2 - x_1 &\le w_{12} \\ x_3 - x_2 &\le w_{23} \\ &\vdots \\ x_k - x_{k-1} &\le w_{k-1,\,k} \\ x_1 - x_k &\le w_{k1} \end{aligned}$$

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Therefore, no values for the x_i can satisfy the constraints. ☐

Satisfying the constraints

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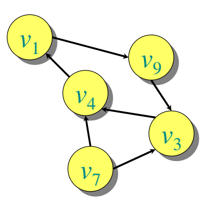
Proof. Add a new vertex s to V with a 0-weight edge to each vertex $v_i \in V$.



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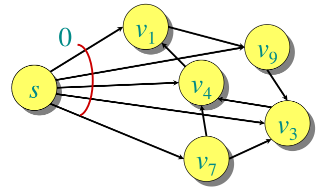
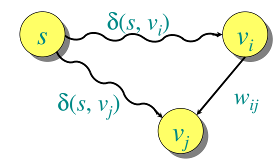


Note: No negative-weight cycles introduced $\Rightarrow$ shortest paths exist.

Proof (continued)

Claim: The assignment $x_i = \delta(s, v_i)$ solves the constraints.

Consider any constraint $x_j - x_i \le w_{ij}$, and consider the shortest paths from s to v_j and v_i:



The triangle inequality gives us $\delta(s,v_j) \le \delta(s, v_i) + w_{ij}$. Since $x_i = \delta(s, v_i)$ and $x_j = \delta(s, v_j)$, the constraint $x_j - x_i \le w_{ij}$ is satisfied. ☐

Bellman-Ford and linear programming

Corollary. The Bellman-Ford algorithm can solve a system of m difference constraints on n variables in $O(mn)$ time. ☐

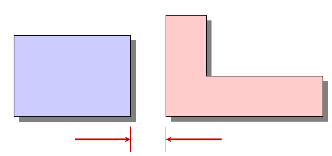
Single-source shortest paths is a simple LP problem.

In fact, Bellman-Ford maximizes $x_1 + x_2 + \cdots + x_n$ subject to the constraints $x_j - x_i \le w_{ij}$ and $x_i \le 0$ (exercise).

Bellman-Ford also minimizes $\max_i\{x_i\} - \min_i\{x_i\}$ (exercise).

Application to VLSI layout compaction

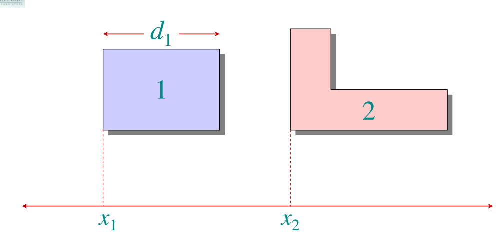
Integrated-circuit features:



minimum separation λ

Problem: Compact (in one dimension) the space between the features of a VLSI layout without bringing any features too close together.

VLSI layout compaction



Constraint: $x_2 - x_1 \ge d_1 + \lambda$

Bellman-Ford minimizes $\max_i\{x_i\} - \min_i\{x_i\}$, which compacts the layout in the x-dimension.