completing the parser for Core Language

In Part 1 of the project we did not consider the two productions:

expr -> expr aexpr and expr -> expr1 binop expr2

the problem with the first production is that is is leftrecursive and a naive parser following it would cause an infinite recursion

we transform the production into:

expr -> aexpr1...aexprn, for n>=1 which can be easily mimiked by the parser using the function some

for the application od the binop's we need to have many productions that model the different precedence of the different binop's. The precedences are summarized in the following table:

Precedence	Associativity	Operator
6	Left	Application
5	Right	*
	None	/
4	Right	+
	None	_
3	None	== ~= > >= < <=
2	Right	&
1	Right	1

in order to model the associativity and precedence of the different binop's we follow thw idea that we have studied in chapter 13.8 of the text the new productions for expr are in the next slide

```
expr \rightarrow let defns in expr
           letrec defns in expr
            case expr of alts
            \setminus var_1 \dots var_n \cdot expr
             expr1
expr1 \rightarrow expr2 \mid expr1
   expr2
expr2 \rightarrow expr3 \& expr2
        | expr3
expr3 \ \rightarrow \ expr4 \ relop \ expr4
        expr4
expr4 \rightarrow expr5 + expr4
        | expr5 - expr5 | expr5
expr5 \rightarrow expr6 * expr5
        expr6 / expr6
          expr6
                                       (n \ge 1)
expr6 \rightarrow aexpr_1 \dots aexpr_n
```

Figure 1.3: Grammar expressing operator precedence and associativity

important advice: since aexpr can be simple variables, and, following the last production in the previous table, the parser will search the input for n aexpr, n > 0, it is necessary that the function that recognizes variables, is able to distinguish variables from keywords of the Core Language, such as in, of, let, etc.

Otherwise you risk that such keywords are mixed up with aexpr which may cause the failure of the program.