Our flour concept
At the level of the assembly language, 'orunflave is a situation / condition rubick expresses the fact that the result of
returned condition rubicly expresses the fact that he rubicly
in I am I am man in a
for it ( Son imrigned). The flags CF and OF rule is
for it (Son insigned). The flags crowd or well be set jor specific cases of overflow.
Addition the overlow is signaled
Addition  in the unsigned interpretation, the overflow is signaled by setting CF = 1 ruhenerus the result didn't fit the reserved.  spaa for it (on a byte (o, 255]) Otherwise CF:0.
by petting cf = 1 ruhenerus the result duming for
ma for it ( on a befte (0, 255]) Otherwise CF.
01/1/01 1
Ex: mov al, 100 1 100 \( \int \( \text{0.7255} \) but 100 +200 = 300 \( \delta \) (0,255)
mov $bl,200$ , $260 \in (0,255)$ =) $CF=1$ , add al, $bl$
and the transfer of the second of the
- in signed interpretation, the ownflow is signature by
restring OF=1 ruhenerus the base 2 addition reflects an incar reset mathematic result (in the signed interpretation). Otherwise OF=0. On a byte: (-126, 127).  There are only a situations that can issue overflaw  Ex: Anoval
ruled maturnate the first ( son the significant of
uise 01 =0. On a sign. C 120, (21).
Ex: Amou at
$A_{\infty} = O((A_{\infty} \cup I))$ :
1. $\frac{1}{0}$ and $\frac{1}{2}$
Ex: Imau al, 100   100 € [-128, 127)
Ex: $1 \text{max}$ al, $100 \mid 100 \in [-128, 127)$ $1 \text{max}$ bl, $100 \mid 100 \neq 100 = 200 \notin (-128, 127) = )$
add as be I so we bring the value in the domain by
add al, bl /=) we bring the value in the domain by outtracting (256)=) 200-256=-56=>0p

Scanned with CamScanner

Subtraction:

- in the unsigned interpretation, the deenflow is signaled by detting CP = L nuhenerus there exists a borrace from a mon-existent position or in athor woods, the result didn't fit the reserved space for it. Otherwise CF = 0.

Ex: mar al, 100 | mar bl, 101 |  $100 - 101 = -1 \notin [0, 255]^2) CF = 1$ .
Sub al, bl

- in the signed interpretation, the creefler is signaled by setting DF=1 ruhenerus the base a subtraction reflects an incorred result (in Signed interpretation). Otherwise

Ex: Amou are only a retrections that can issue orienflace for subtration:

 $\frac{1}{0 \dots } \quad \text{and} \quad \frac{1}{1 \dots }$ 

Ex: max al,  $|00| | |00 \in \{124, 124\}$ max bl,  $-|00| | -|00| \in \{-24, 124\}$ . Bub al, bl  $|00-(-|00|) = 200 \in \{-124, 124\} = 1$ 

200-256=-56No +-(-)=- which is incorrect =) 0f=1 -

Multiplication:

The multiplication operation does not produce our flow,

the reserved apace bedy enough for both nontripre tations. The

Chaision was taken to set Cf=0 F=0 whenever the result

is the same are of the operations. And CF=0.F=1 for the opposite- for the rases

byte # byte = byte

word \* mord = mad = ) Of = Cf = 0 ( ma meelt plication

amend \* duard = dward

And for the races:

byte \* byte = word word \* or cf = 1.

word \* word = dword = ) Of = Cf = 1.

Ex: mov al, 200 | 200  $\in$  (0,255) mov bel, 200 | 200 \* 200 = 40000  $\in$  (0,65536) =) mul be | =) on a word

Division

- in the case of deluision, if it happens, then it well result in the program crashing (fatal Error) = division overflow. The values of CF = OF are irrelevant. The function of didn't fet the reserved space

ex: mov ax,4096

mar bl, 10; hog6: [0 - 409, 12=6

div bl luit hog & [0,256] =)

-) fATAL ERROR/enach.

There are methods to deal with the overflow conept. and the assembles gives us 2 specific instructions: AD ( (ac ruth carry) and SBB ( subtraction with borrow).

Ex: new meed to compute the value of Dx: Ax + Bx: Bx. add ax 1 bx | -> to abtain a correct result, rue make me that the transport dight is not lost.

There is mo "iadd" or "hub" because freen if they existed they reveall work exactly the same as "ADD" and "SUB". Thus is because in base a addition and subtraction are performed the same insepted DENTLY of the INTER PRETARION. There exists, inju" and "iDiu" because this rule does not apply to meetiplication and division who work differently in bath interpretations.

The programmer can avoid orienflave setections by.

sirroy odota typer (ex: hock instead of byte), cheking the
imput range or wing instructions like: (jo-jimp if OF=1,
jmo--1-OF=0) jc--1-CF=+; jmc--1-CF=0).

Ex:

add all bl fine avoid ; instructions avoid:

- if the addition of the 2 regorters touriets in an overflow by that fitting the reserved space of the jump rul be performed.