

# Signals and Communication Theory

## *Discrete Fourier Transform*

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# Introduction

In wireless communication, we use the **Discrete Fourier Transform (DFT)** for equalization strategies and OFDM. The DFT equations are given by

$$\begin{aligned}\textbf{Analysis:} \quad X[k] &= \sum_{n=0}^{N-1} x[n] e^{-j \frac{2\pi}{N} nk}, \\ \textbf{Synthesis:} \quad x[n] &= \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j \frac{2\pi}{N} nk}.\end{aligned}$$

In this case, note that the frequency variable  $f$  is discrete, i.e.,  $f = \frac{k}{N}$ . Similarly, the time variable  $n$  is discrete. Additionally, the signal  $x[n]$  and its corresponding DFT  $X[k]$  are periodic signals with period  $N$ .

# Modulo operation

Since the signals  $x[n]$  and  $X[k]$  are periodic with period  $N$ , shifts are given in terms of a modulo by  $N$  operation denoted as  $(\cdot)_N$ . This ensures that the argument falls in  $[0, N - 1]$  as shown below

|         |    |    |    |    |    |   |   |   |   |   |   |
|---------|----|----|----|----|----|---|---|---|---|---|---|
| $n$     | -5 | -4 | -3 | -2 | -1 | 0 | 1 | 2 | 3 | 4 | 5 |
| $(n)_4$ | 3  | 0  | 1  | 2  | 3  | 0 | 1 | 2 | 3 | 0 | 1 |

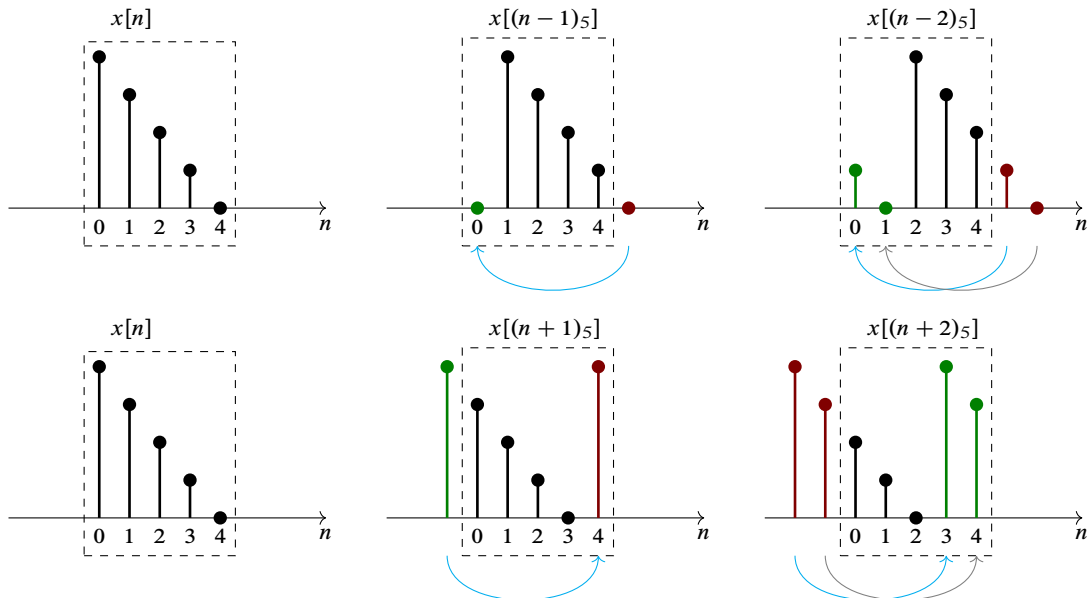
The modulo operation  $(n)_N$  can be computed as

$$(n)_N = r,$$

where  $r$  is a positive integer  $r < N$  such that  $n = qN + r$  and  $q$  is an integer.

# Circular Shift

Using modulo operation, the circular shift is denoted as  $x[(n - k)_N]$ .



# Circular Convolution

The convolution between two discrete-time periodic signal,  $x_1[n]$  and  $x_2[n]$ , is called **Circular convolution** and is defined as

$$y[n] = x_1[n] \circledast x_2[n] = \sum_{\ell=0}^{N-1} x_1[\ell]x_2[(n - \ell)_N].$$

The resulting signal  $y[n]$  is periodic with period  $N$ .

## Example 1

*Calculate the circular convolution between the following signals:*

$$x_1[n] = \begin{cases} 1, & \text{if } n = 0 \\ -1, & \text{if } n = 1 \\ 2, & \text{if } n = 2. \end{cases} \quad x_2[n] = \begin{cases} 0, & \text{if } n = 0 \\ 1, & \text{if } n = 1 \\ 2, & \text{if } n = 2. \end{cases}$$



## Example 1. Answer

**Answer:** Substituting  $x_1[n]$  and  $x_2[n]$  with  $N = 3$  into the definition of circular convolution, we obtain

$$y[n] = \sum_{\ell=0}^{N-1} x_1[\ell]x_2[(n-\ell)_N]$$

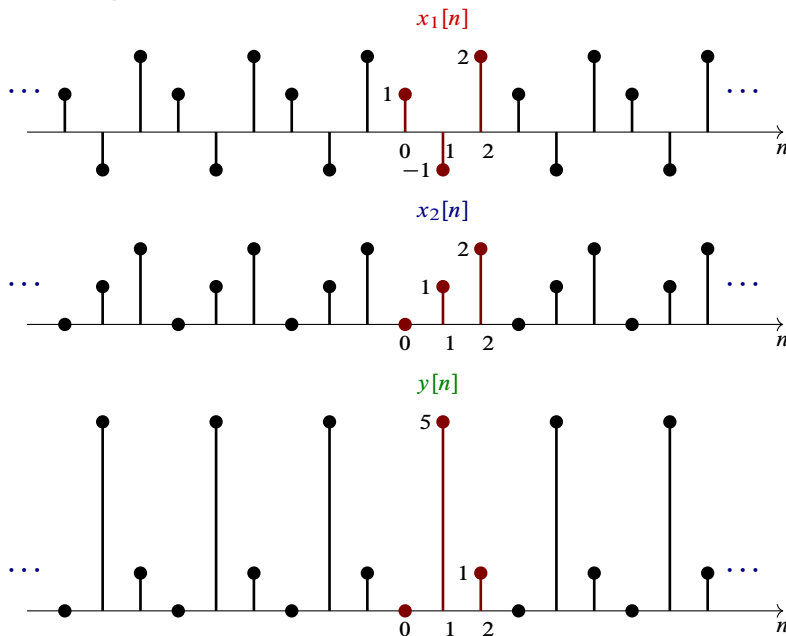
$$y[0] = \sum_{\ell=0}^2 x_1[\ell]x_2[(-\ell)_N] = x_1[0]x_2[\mathbf{0}] + x_1[1]x_2[\mathbf{2}] + x_1[2]x_2[\mathbf{1}] = 0$$

$$y[1] = \sum_{\ell=0}^2 x_1[\ell]x_2[(1-\ell)_N] = x_1[0]x_2[\mathbf{1}] + x_1[1]x_2[\mathbf{0}] + x_1[2]x_2[\mathbf{2}] = 5$$

$$y[2] = \sum_{\ell=0}^2 x_1[\ell]x_2[(2-\ell)_N] = x_1[0]x_2[\mathbf{2}] + x_1[1]x_2[\mathbf{1}] + x_1[2]x_2[\mathbf{0}] = 1$$

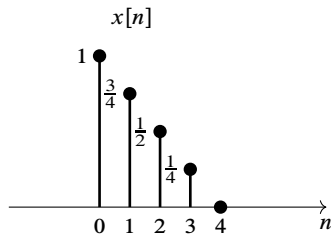
$$y[n] = \begin{cases} 0, & \text{if } n = 0 \\ 5, & \text{if } n = 1 \\ 1, & \text{if } n = 2. \end{cases}$$

## Example 1. Graphs



## Example 2

*Find the DFT of the following signal:*



## Example 2. Answer

Substituting  $x[n]$  into the analysis equation, we have

$$\begin{aligned} X[k] &= \sum_{n=0}^{N-1} x[n] e^{-j \frac{2\pi}{N} kn} \\ &= \sum_{n=0}^4 x[n] e^{-j \frac{2\pi}{5} kn} \\ &= 1 + \frac{3}{4} e^{-j \frac{2\pi}{5} k} + \frac{1}{2} e^{-j \frac{4\pi}{5} k} + \frac{1}{4} e^{-j \frac{6\pi}{5} k} \end{aligned}$$

| $k$ | $X[k]$              | $ X[k] $ | $\angle X[k]$ |
|-----|---------------------|----------|---------------|
| 0   | 2.5                 | 2.5      | 0             |
| 1   | $0.625 - 0.860239j$ | 1.063314 | $-0.3\pi$     |
| 2   | $0.625 - 0.203075j$ | 0.657164 | $-0.1\pi$     |
| 3   | $0.625 + 0.203075j$ | 0.657164 | $0.1\pi$      |
| 4   | $0.625 + 0.860239j$ | 1.063314 | $0.3\pi$      |

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# Linearity

The DFT satisfies the linearity property

$$\alpha x_1[n] + \beta x_2[n] \leftrightarrow \alpha X_1[k] + \beta X_2[k]$$

## Proof

$$\begin{aligned}\sum_{n=0}^{N-1} (\alpha x_1[n] + \beta x_2[n]) e^{-j \frac{2\pi}{N} kn} &= \sum_{n=0}^{N-1} \alpha x_1[n] e^{-j \frac{2\pi}{N} kn} + \sum_{n=0}^{N-1} \beta x_2[n] e^{-j \frac{2\pi}{N} kn} \\ &= \alpha \sum_{n=0}^{N-1} x_1[n] e^{-j \frac{2\pi}{N} kn} + \beta \sum_{n=0}^{N-1} x_2[n] e^{-j \frac{2\pi}{N} kn} \\ &= \alpha X_1[k] + \beta X_2[k]\end{aligned}$$

# Periodicity

The Discrete Fourier Transform and inverse Discrete Fourier Transform are periodic sequences with period  $N$ .

$$X[k] = X[k + N]$$

$$x[n] = x[n + N]$$

## Proof

$$\begin{aligned}x[n + N] &= \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{-j \frac{2\pi}{N} k(n+N)} \\&= \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{-j \frac{2\pi}{N} k n} e^{-j 2\pi k} \\&= \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{-j \frac{2\pi}{N} k n}, \quad \text{Since } e^{-j 2\pi k} = 1 \\&= x[n]\end{aligned}$$

# Time shifting

The time shifting property is expressed by

$$x[(n - n_0)_N] \leftrightarrow e^{-j\frac{2\pi}{N}kn_0} X[k]$$

**Proof.** Considering  $n - n_0 = qN + m$ , we obtain  $(n - n_0)_N = m$ . Thus

$$\begin{aligned}\sum_{n=0}^{N-1} x[(n - n_0)_N] e^{-j\frac{2\pi}{N}kn} &= e^{-j\frac{2\pi}{N}kn_0} \sum_{n=0}^{N-1} x[(n - n_0)_N] e^{-j\frac{2\pi}{N}k(n-n_0)} \\ &= e^{-j\frac{2\pi}{N}kn_0} \sum_{m=0}^{N-1} x[m] e^{-j\frac{2\pi}{N}km} \\ &= e^{-j\frac{2\pi}{N}kn_0} X[k]\end{aligned}$$



# Frequency shifting

The Modulation property is expressed by

$$e^{j\frac{2\pi}{N}k_0n}x[n] \leftrightarrow X[(k - k_0)_N]$$

**Proof.** By considering  $k - k_0 = qN + m$ , we have  $(k - k_0)_N = m$ . Thus

$$\begin{aligned}\sum_{n=0}^{N-1} e^{j\frac{2\pi}{N}k_0n}x[n]e^{-j\frac{2\pi}{N}kn} &= \sum_{n=0}^{N-1} x[n]e^{-j\frac{2\pi}{N}(k-k_0)n} \\ &= X[(k - k_0)_N]\end{aligned}$$

# Time reversal

This property satisfies

$$x[(-n)_N] \leftrightarrow X[(-k)_N]$$

**Proof.** Using  $-n = qN + m$ , we have  $(-n)_N = m$ . Therefore

$$\begin{aligned} \sum_{n=0}^{N-1} x[(-n)_N] e^{-j \frac{2\pi}{N} kn} &= \sum_m x[m] e^{-j \frac{2\pi}{N} k(-qN-m)} \\ &= \sum_{m=0}^{N-1} x[m] e^{-j \frac{2\pi}{N} (-k)m} \\ &= \sum_{m=0}^{N-1} x[m] e^{-j \frac{2\pi}{N} (-k)_N m} \\ &= X[(-k)_N] \end{aligned}$$

# Complex conjugate

This property satisfies

$$x^*[n] \leftrightarrow X^*[(-k)_N]$$

## Proof

$$\begin{aligned}\sum_{n=0}^{N-1} x^*[n] e^{-j \frac{2\pi}{N} k n} &= \left( \sum_{n=0}^{N-1} x[n] e^{j \frac{2\pi}{N} k n} \right)^* \\ &= \left( \sum_{n=0}^{N-1} x[n] e^{-j \frac{2\pi}{N} (-k) n} \right)^* \\ &= \left( \sum_{n=0}^{N-1} x[n] e^{-j \frac{2\pi}{N} (-k)_N n} \right)^* \\ &= X^*[(-k)_N]\end{aligned}$$

# Duality

The duality property is expressed as

$$X[n] \leftrightarrow Nx[(-k)_N]$$

## Proof

$$\begin{aligned}\sum_{n=0}^{N-1} X[n] e^{-j \frac{2\pi}{N} kn} &= \sum_{n=0}^{N-1} X[n] e^{j \frac{2\pi}{N} (-k)n} \\ &= \sum_{n=0}^{N-1} X[n] e^{j \frac{2\pi}{N} (-k)_N n} \\ &= Nx[(-k)_N]\end{aligned}$$

# Circular Convolution

The Discrete Fourier transform of the circular convolution is related as

$$x[n] \circledast h[n] \leftrightarrow X[k]H[k]$$

## Proof

$$\begin{aligned}\sum_{n=0}^{N-1} x[n] \circledast h[n] e^{-j \frac{2\pi}{N} kn} &= \sum_{n=0}^{N-1} \sum_{m=0}^{N-1} x[m] h[(n-m)_N] e^{-j \frac{2\pi}{N} kn} \\ &= \sum_{m=0}^{N-1} x[m] \sum_{n=0}^{N-1} h[(n-m)_N] e^{-j \frac{2\pi}{N} kn} \\ &= \sum_{m=0}^{N-1} x[m] e^{-j \frac{2\pi}{N} km} H[k] \\ &= X[k] H[k]\end{aligned}$$

time shifting prop.

# Multiplication

The Discrete Fourier transform of the multiplication is related as

$$x[n]h[n] \leftrightarrow \frac{1}{N}X[k] \circledast H[k]$$

## Proof

$$\begin{aligned}\sum_{n=0}^{N-1} x[n]h[n]e^{-j\frac{2\pi}{N}kn} &= \sum_{n=0}^{N-1} x[n]\frac{1}{N}\sum_{m=0}^{N-1} H[m]e^{j\frac{2\pi}{N}mn}e^{-j\frac{2\pi}{N}kn} \\&= \frac{1}{N}\sum_{m=0}^{N-1} H[m]\sum_{n=0}^{N-1} x[n]e^{j\frac{2\pi}{N}mn}e^{-j\frac{2\pi}{N}kn} \\&= \frac{1}{N}\sum_{m=0}^{N-1} H[m]X[(n-m)_N] \\&= \frac{1}{N}H[k] \circledast X[k]\end{aligned}$$

# Parseval's relation

Parseval's relation is defined as

$$\sum_{n=0}^{N-1} x_1[n]x_2^*[n] = \frac{1}{N} \sum_{k=0}^{N-1} X_1[k]X_2^*[k]$$

## Proof

$$\begin{aligned} \sum_{n=0}^{N-1} x_1[n]x_2^*[n] &= \sum_{n=0}^{N-1} \frac{1}{N} \sum_{k=0}^{N-1} X_1[k]e^{j\frac{2\pi}{N}kn} x_2^*[n] \\ &= \frac{1}{N} \sum_{k=0}^{N-1} X_1[k] \sum_{n=0}^{N-1} e^{j\frac{2\pi}{N}kn} x_2^*[n] \\ &= \frac{1}{N} \sum_{k=0}^{N-1} X_1[k]X_2^*[k] \end{aligned}$$

# Magnitude spectrum for real-valued signals

Using the fact that  $x[n] = x^*[n]$  and the complex conjugate property, we have

$$X[k] = X^*[(-k)_N].$$

Thus

$$|X[k]| = |X[(-k)_N]|.$$

This means, the magnitude spectrum is an **even function**.



# Phase spectrum for real-valued signals

Using  $X[k] = X^*[(-k)_N]$  gives

$$\angle X[k] = -\angle X[(-k)_N].$$

This means, the phase spectrum is an **odd function**.

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# Multiplications and Additions in DFT Computation

Consider the following example with  $N = 8$ :

$$\begin{aligned} X[k] &= \sum_{n=0}^7 x[n] e^{-j \frac{\pi}{4} kn} \\ &= x[0] \cdot 1 + x[1] e^{-j \frac{\pi}{4} k} + x[2] e^{-j \frac{\pi}{2} k} + x[3] e^{-j \frac{3\pi}{4} k} \\ &\quad + x[4] e^{-j \pi k} + x[5] e^{-j \frac{5\pi}{4} k} + x[6] e^{-j 3\pi k} + x[7] e^{-j \frac{7\pi}{4} k}. \end{aligned}$$

In this particular example, each value of  $X[k]$  requires 8 complex multiplications and 7 complex additions. Thus, the total computation of the DFT requires 64 complex multiplications and 56 complex additions.

**In general, the computation of the DFT performs  $N^2$  complex multiplications and  $N(N - 1)$  complex additions.**

# FFT Algorithm

We assume that  $N$  can be expressed as  $N = 2^\ell$ . Now in order to obtain the  $N$ -point DFT, we split the signal into even and odd-indexing signals as follows

$$\begin{aligned} X[k] &= \sum_{n=0}^{N-1} x[n] e^{-j \frac{2\pi}{N} kn} \\ &= \sum_{n=0}^{N/2-1} x[2n] e^{-j \frac{2\pi}{N} k 2n} + \sum_{n=0}^{N/2-1} x[2n+1] e^{-j \frac{2\pi}{N} k (2n+1)} \\ &= \underbrace{\sum_{n=0}^{N/2-1} x[2n] e^{-j \frac{2\pi}{N/2} kn}}_{N/2\text{-point DFT}} + e^{-j \frac{2\pi}{N} k} \underbrace{\sum_{n=0}^{N/2-1} x[2n+1] e^{-j \frac{2\pi}{N/2} kn}}_{N/2\text{-point DFT}} \end{aligned}$$

**In this way, neglecting the multiplications by  $e^{-j \frac{2\pi}{N} k}$ , the evaluation of the DFT requires  $\left(\frac{N}{2}\right)^2 + \left(\frac{N}{2}\right)^2 = \frac{N^2}{2}$  complex multiplications. That is, a reduction by 2.**

## FFT Algorithm, cont.

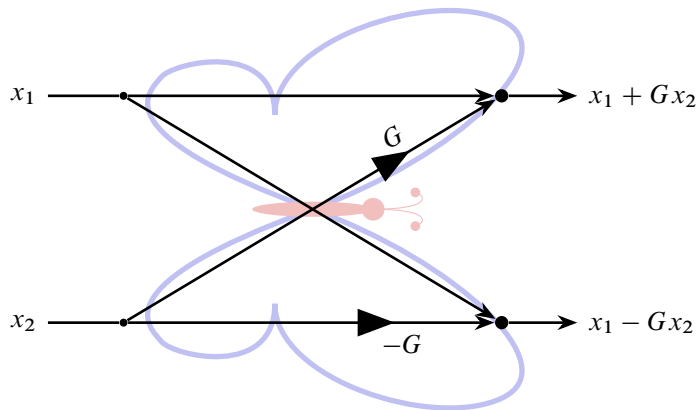
The splitting operation can continue until only pairs of the original signal occur. In that case the 2-point DFT is given by

$$X[k] = \sum_{n=0}^1 x[n]e^{-j\pi kn},$$
$$X[0] = x[0] + x[1],$$
$$X[1] = x[0] - x[1],$$

which does not require multiplications.

**Evaluating the DFT with the FFT algorithm requires  $N \log_2 N$  complex multiplications.**

# Butterfly Pattern

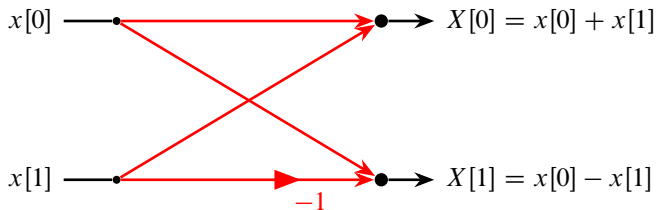


## 2-FFT Butterfly Structure

$$X[k] = \sum_{n=0}^1 x[n]e^{-j\pi kn},$$

$$X[0] = x[0] + x[1],$$

$$X[1] = x[0] - x[1].$$



## 4-FFT Algorithm

$$\begin{aligned}X[k] &= \sum_{n=0}^3 x[n]e^{-j\frac{\pi}{2}kn}, \\&= \underbrace{\sum_{n=0}^1 x[2n]e^{-j\pi kn}}_{\text{2-DFT}} + e^{-j\frac{\pi}{2}k} \underbrace{\sum_{n=0}^1 x[2n+1]e^{-j\pi kn}}_{\text{2-DFT}} \\&= X_e[k] + e^{-j\frac{\pi}{2}k} X_o[k]\end{aligned}$$

$$X[0] = X_e[0] + X_o[0]$$

$$X_e[0] = x[0] + x[2]$$

$$X[1] = X_e[1] - jX_o[1]$$

$$X_e[1] = x[0] - x[2]$$

$$X[2] = X_e[0] - X_o[0]$$

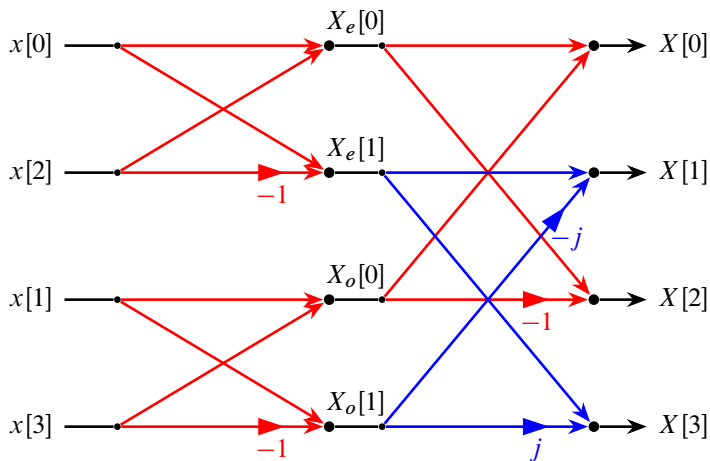
$$X_o[0] = x[1] + x[3]$$

$$X[3] = X_e[1] + jX_o[1]$$

$$X_o[1] = x[1] - x[3]$$



## 4-FFT Butterfly Structure



## 8-FFT Algorithm

$$\begin{aligned} X[k] &= \sum_{n=0}^7 x[n] e^{-j \frac{\pi}{4} kn} = \underbrace{\sum_{n=0}^3 x[2n] e^{-j \frac{\pi}{2} kn}}_{\text{4-DFT}} + e^{-j \frac{\pi}{4} k} \underbrace{\sum_{n=0}^3 x[2n+1] e^{-j \frac{\pi}{2} kn}}_{\text{4-DFT}} \\ &= X_e[k] + e^{-j \frac{\pi}{4} k} X_o[k] \end{aligned}$$

$$X[0] = X_e[0] + X_o[0]$$

$$X[1] = X_e[1] + e^{-j \frac{\pi}{4}} X_o[1]$$

$$X[2] = X_e[2] - j X_o[2]$$

$$X[3] = X_e[3] - e^{j \frac{\pi}{4}} X_o[3]$$

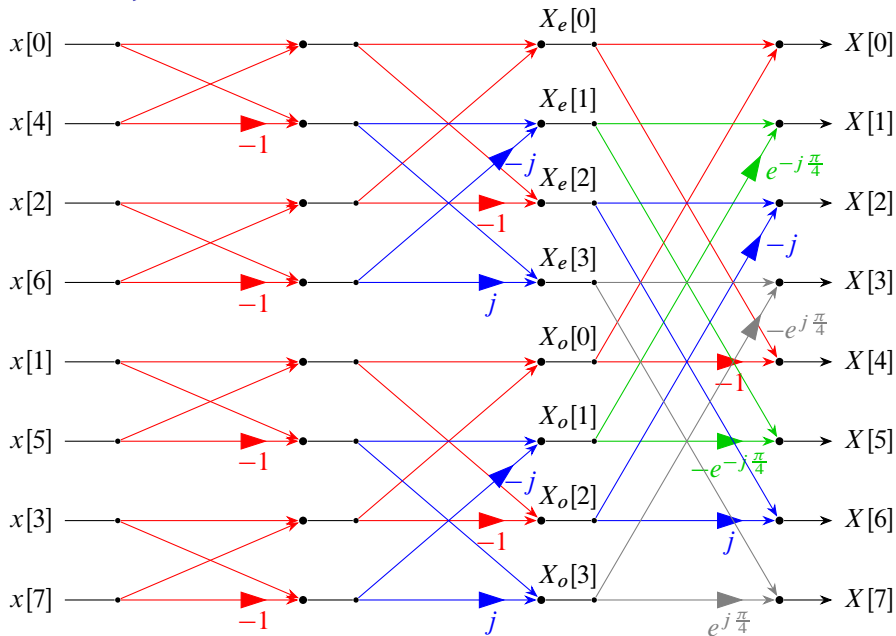
$$X[4] = X_e[0] - X_o[0]$$

$$X[5] = X_e[1] - e^{-j \frac{\pi}{4}} X_o[1]$$

$$X[6] = X_e[2] + j X_o[2]$$

$$X[7] = X_e[3] + e^{j \frac{\pi}{4}} X_o[3]$$

# 8-FFT Butterfly Structure



# Comparison of Complex Multiplications

| $N$  | $\text{DFT}(N^2)$ | $\text{FFT}(N \log_2 N)$ |
|------|-------------------|--------------------------|
| 32   | 1024              | 160                      |
| 64   | 4096              | 384                      |
| 128  | 16384             | 896                      |
| 256  | 65536             | 2024                     |
| 512  | 262144            | 4608                     |
| 1024 | 1048576           | 10240                    |

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# Summary of Fourier Analysis/Synthesis

## Continuous-Time

## Discrete-Time

### Fourier Transform (FT)

$$X(f) = \int_{-\infty}^{\infty} x(t) e^{-j2\pi f t} dt$$
$$x(t) = \int_{-\infty}^{\infty} X(f) e^{j2\pi f t} df$$

Continuous-  
Frequency

### Discrete-Time Fourier Transform (DTFT)

$$X(e^{j2\pi f}) = \sum_{n=-\infty}^{\infty} x[n] e^{-j2\pi f n}$$
$$x[n] = \int_{-\frac{1}{2}}^{\frac{1}{2}} X(e^{j2\pi f}) e^{j2\pi f n} df$$

### Fourier Series (FS)

$$c[n] = f_0 \int_{-\frac{1}{2f_0}}^{\frac{1}{2f_0}} x(t) e^{-j2\pi n f_0 t} dt$$
$$x(t) = \sum_{n=-\infty}^{\infty} c[n] e^{j2\pi n f_0 t}$$

Discrete-  
Frequency

### Discrete Fourier Transform (DFT)

$$X[k] = \sum_{n=0}^{N-1} x[n] e^{-j \frac{2\pi}{N} nk}$$
$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j \frac{2\pi}{N} nk}$$

# Homework

- Problems 6.79–6.82 from [3]

[3] Hwei Hsu, *Schaum's Outline of Signals and Systems*, Second Edition, 2010, McGraw Hill