Compilers

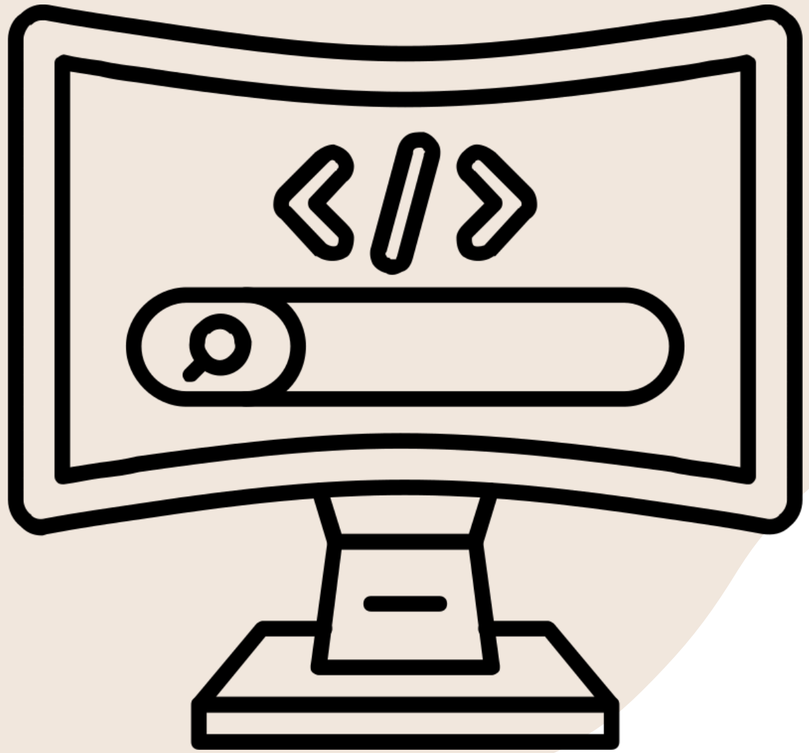
COMPILADOR

Team 04.

INTRODUCTION

Problem solved

This project addresses a fundamental compiler problem.
Given a sequence of tokens we verify if the program is syntactically correct, semantically valid, and then translate it.



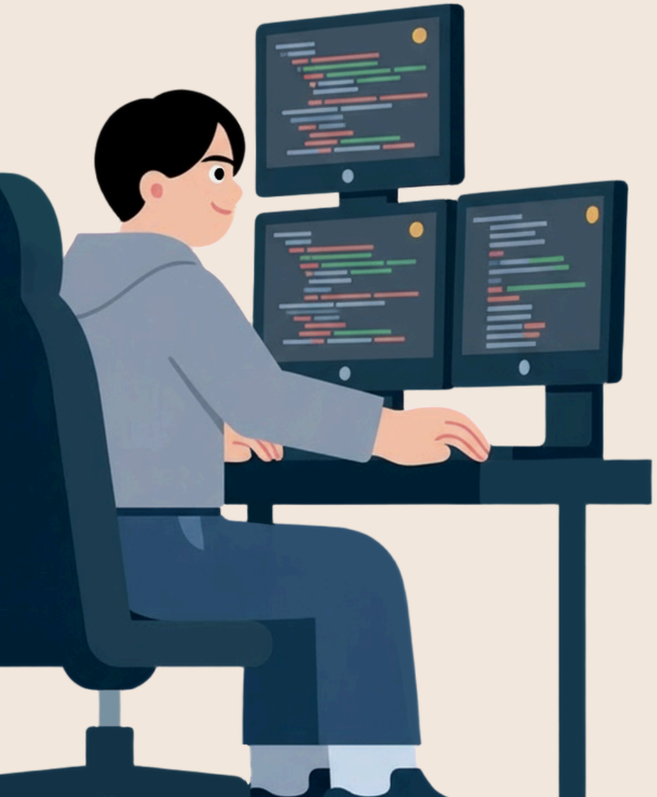
Objective

- Design a parser
- Introduce semantic verification
- Generate Three-Address Code (TAC)
- Compile and execute
- Provide clear feedback

C supported

The project implements a reduced sub-language of C

GENERAL ARCHITECTURE



LEXER
(Tokens)

PARSER + SDT
(Syntax Semantics)

TAC
(Three Address Code)

VM / ASM
(Execution)

L E X E R

01

The lexer is the first phase of the compiler.

It reads the input file character by character and groups them into tokens.

Removes whitespace, newlines, and comments.

02

Recognizes:

- Identifiers
- Numbers
- Keywords
- Operators
- Delimiters
- Produces a linear token stream used by the parser.

03

EXAMPLE

The line int x = 10; translates to:

Keyword Identifier Operator Constant Punctuation.

Tokens produced:

[int] [x] [=] [10] [;]

PARSER

- The parser receives the token stream from the lexer.
- Its job is to check whether the sequence follows the grammar of the C subset.
- We implemented a Bottom-Up parser, based on:
  - An analysis stack
  - The try_reduce() function to apply grammar rules
- Detects syntax errors, such as:
  - Misordered tokens
  - Incomplete structures
  - Invalid declarations

EXAMPLE

Tokens:
int x = 10 ;

Reductions:
int → TYPE
x → ID
TYPE ID = NUM ; → DECLARATION

SEMANTIC ANALYSIS (SDT)

ERRORS

- Parsing Error (Syntax)
  - int x = ; missing item
- SDT Error (Semantic):
  - int x;
  - x = "cadena";   // tipos incompatibles

Types

- int x = 5; OK
- int x = "hola"; error

INTERMEDIATE CODE (THREE ADDRESS CODE – TAC)

Central idea: each operation is broken down into simple instructions with a maximum of three operands (destination = op1 operator op2).

Advantage:

- Machine independence → easier to optimize.
- Clarity in flow control (labels, jumps).
- Basis for translating to ASM.

EXAMPLE

x = a + b * c;

t1 = b * c
x = a + t1

INTERMEDIATE CODE TO MACHINE

Example Code (Left Side TAC vs Right Side ASM):

TAC: t1 = a+ b

ASM: LOAD A, a
ADD A, b
STORE A, t1

THE VIRTUAL MACHINE (EXECUTION)

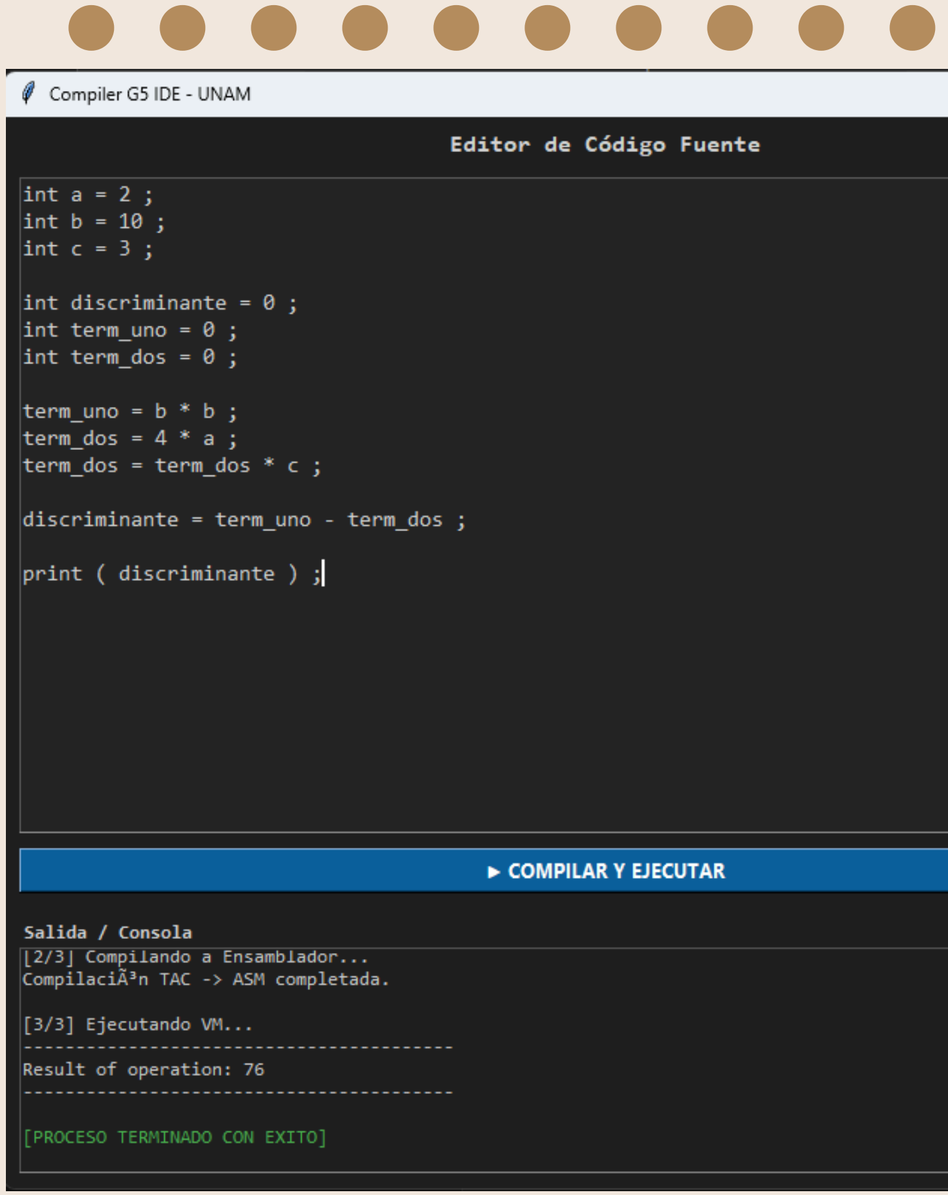
List of VM Features:

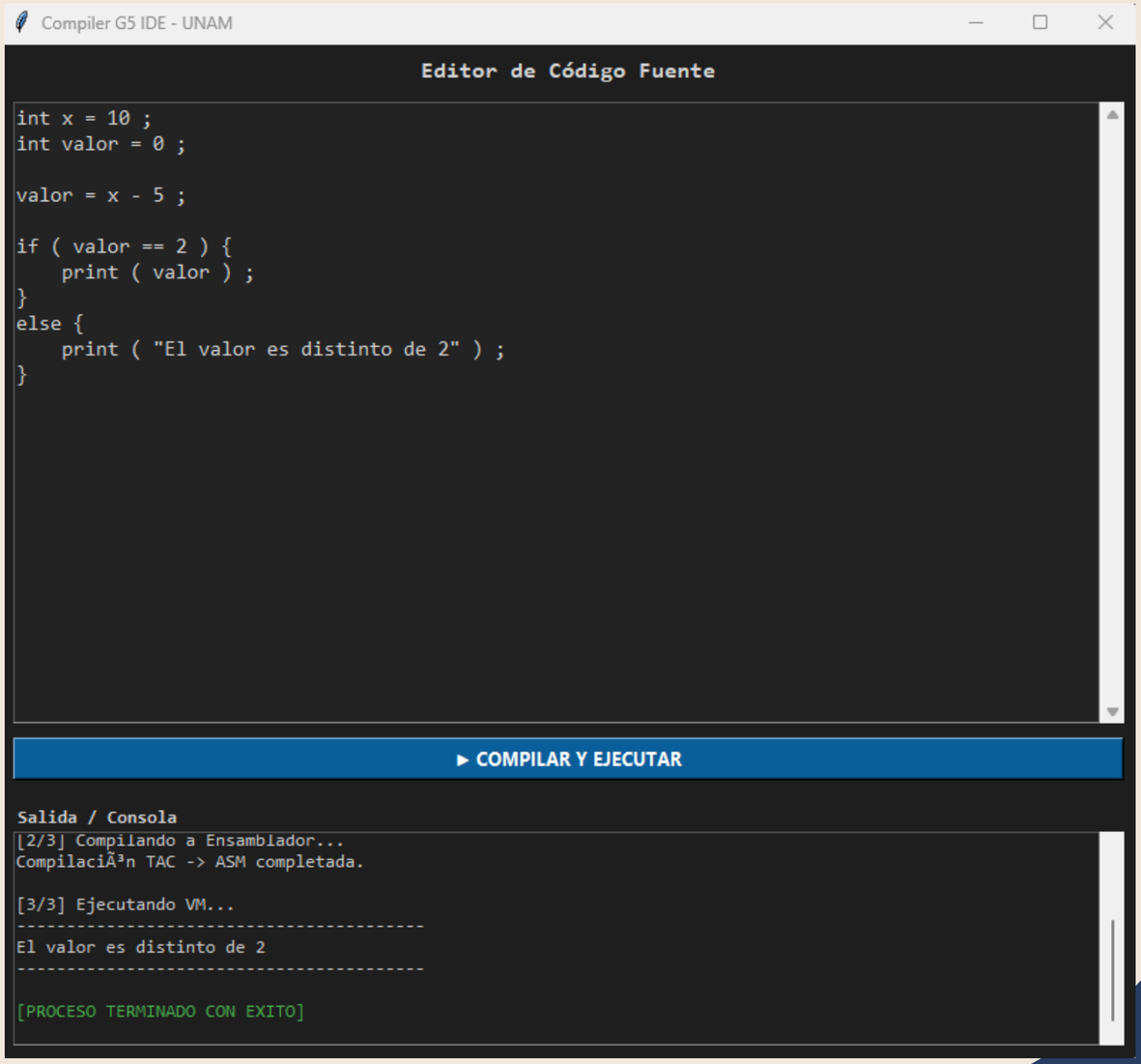
- Accumulator Architecture (Register 'A').
- Data Memory (Dynamic Symbol Table).
- Supported Instructions: Arithmetic (ADD, SUB), Logic (JGT, JLT), I/O (PRINT).

Code snippet from vm_asm.c:

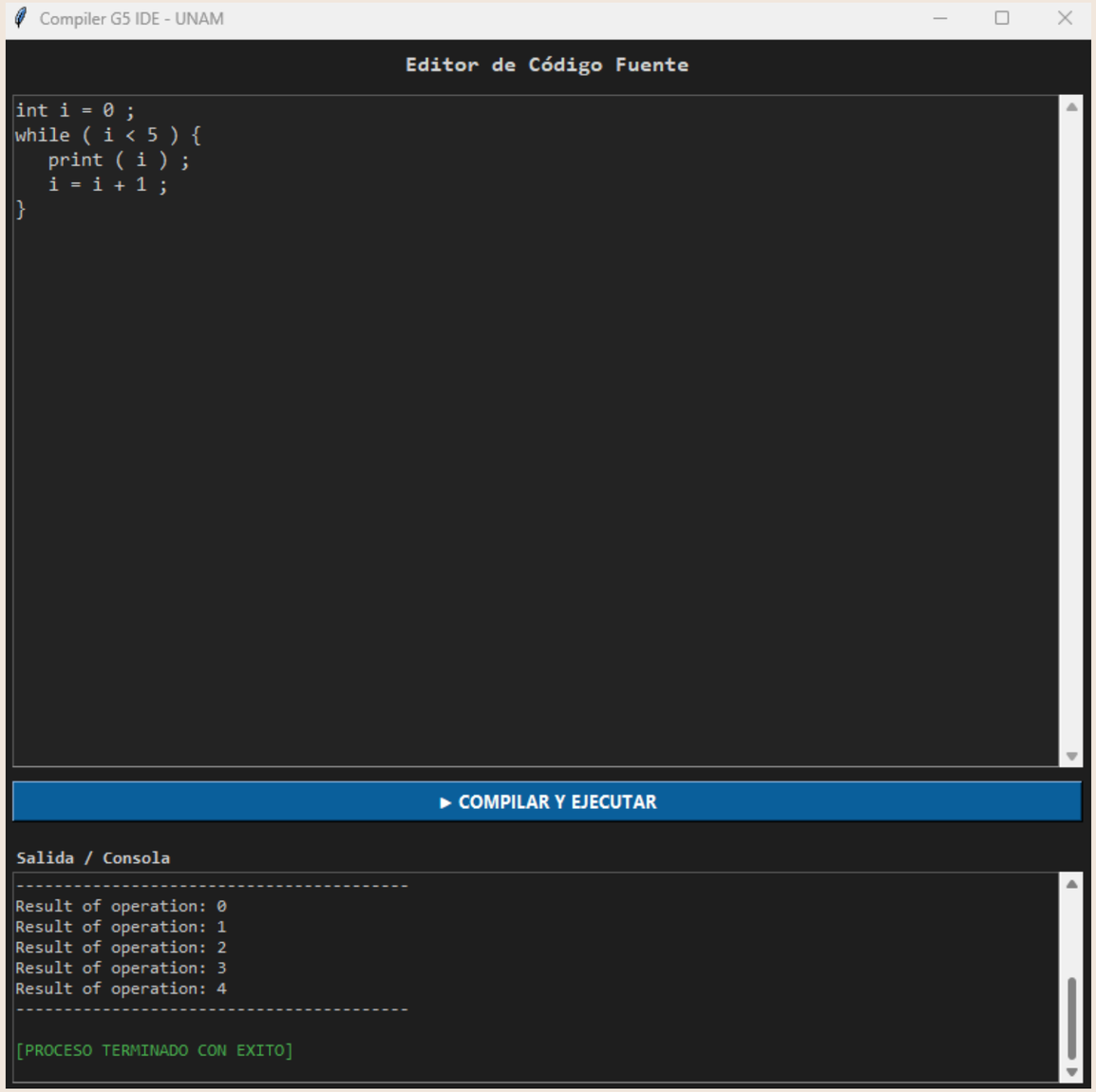
```
case  OP_ADD:
        A += get_value(inst->arg);
        pc++;
```

SUCCESSFULLY COMPLETED





SUCCESSFULLY COMPLETED

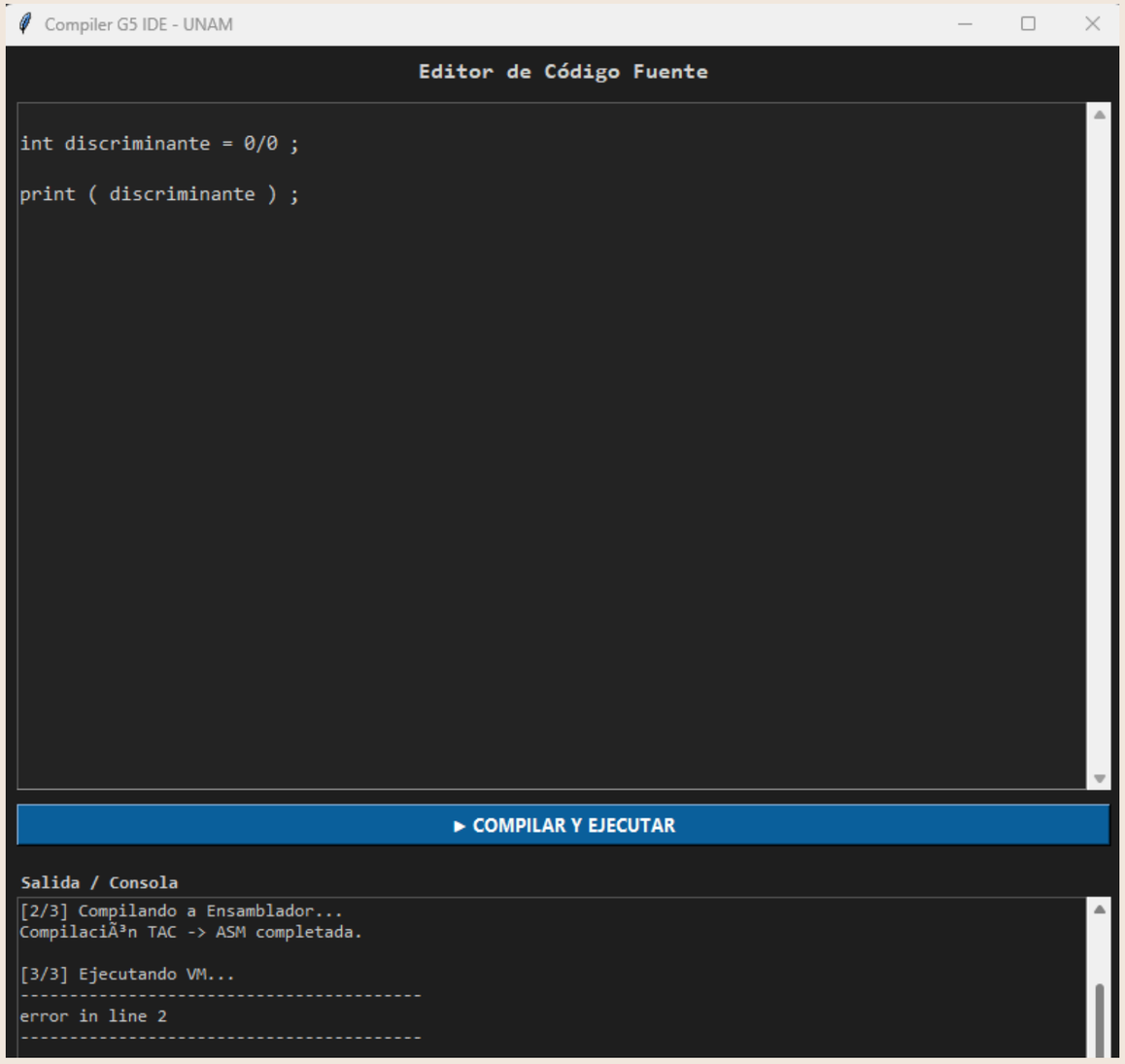


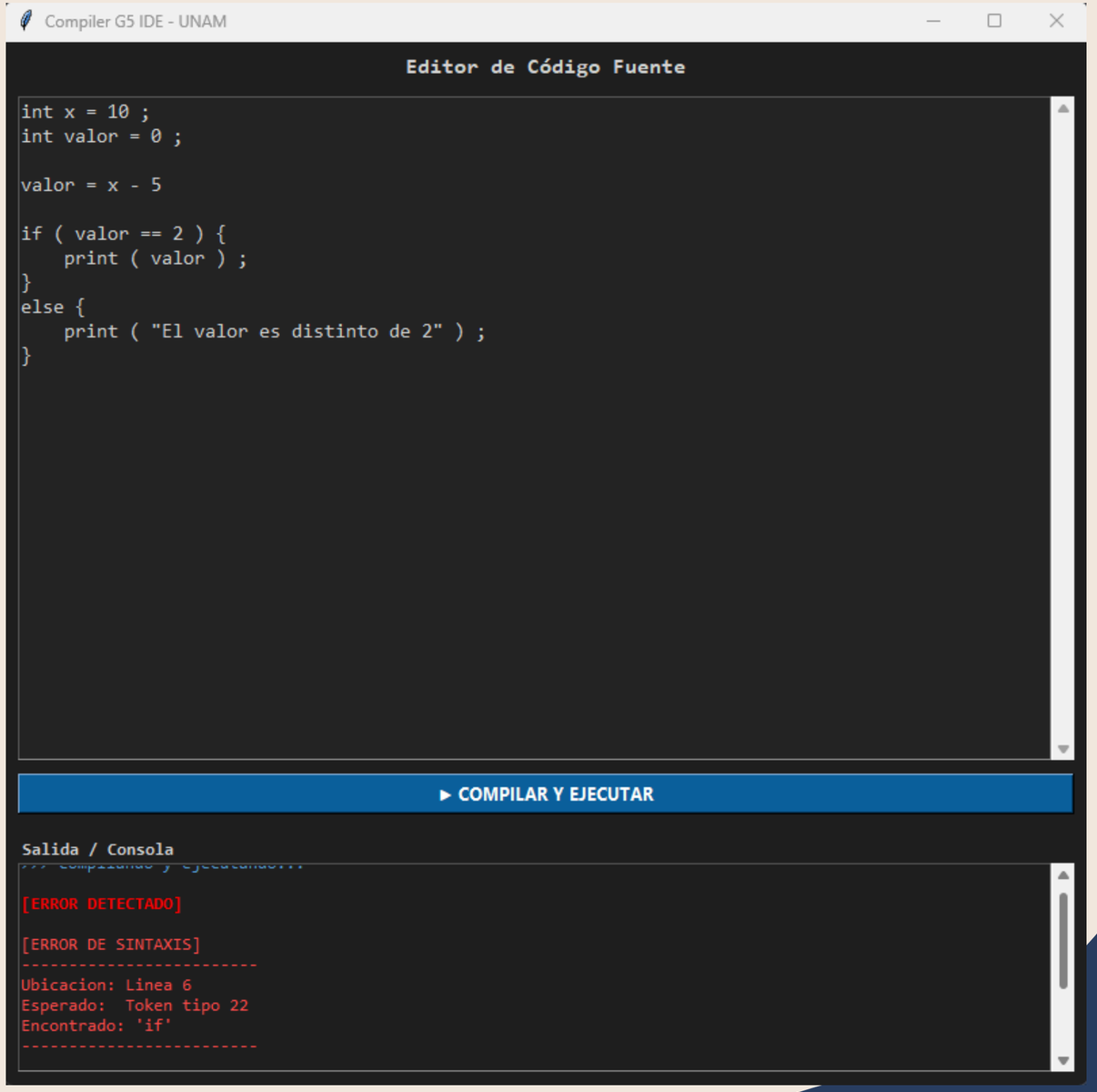

```
PS C:\Users\alons\Documents\Proyecto\unam.fi.compilers.g5.04> ./run.bat examples\temp_gui.txt

[1/3] Lexer - Parser (Analisis)...
----------------------------------
[OK] Sintaxis Correcta. Generando TAC.

[2/3] Compilando a Ensamblador...
Compilaci├│n TAC -> ASM completada.

[3/3] Ejecutando VM...
----------------------------------------
Result of operation: 0
Result of operation: 1
Result of operation: 2
Result of operation: 3
Result of operation: 4
----------------------------------------
PS C:\Users\alons\Documents\Proyecto\unam.fi.compilers.g5.04>
```

Operation failed

```
PS C:\Users\alons\Documents\Proyecto\unam.fi.compilers.g5.04> ./run.bat examples\temp_gui.txt

[1/3] Lexer - Parser (Analisis)...
----------------------------------


[ERROR DE SINTAXIS]
--------------------------
Ubicacion: Linea 2
Esperado:  Token tipo 22
Encontrado: 'while'
--------------------------

[X] SE ENCONTRO UN ERROR. PROCESO DETENIDO.
Revisa el mensaje de arriba para mas detalles.
```

OPERATION FAILED





CONCLUSION

This project demonstrates a reliable, efficient, and fully functional compilation pipeline capable of analyzing, validating, and executing high-level code with precision. By integrating lexical, syntactic, and semantic analysis with TAC generation and virtual-machine execution, we deliver a tool that ensures correctness, prevents critical errors early, and guarantees consistent program behavior. The system's modular design makes it scalable, maintainable, and ready to extend toward larger languages or more complex applications. Overall, this solution provides a solid, professional-grade foundation for any organization seeking to build or enhance a custom compiler or automated code-analysis tool.

THANK YOU

FOR YOUR ATTENTION