



COMPILADOR



Team 04.

INTRODUCTION

Problem solved

This project addresses a fundamental compiler problem.

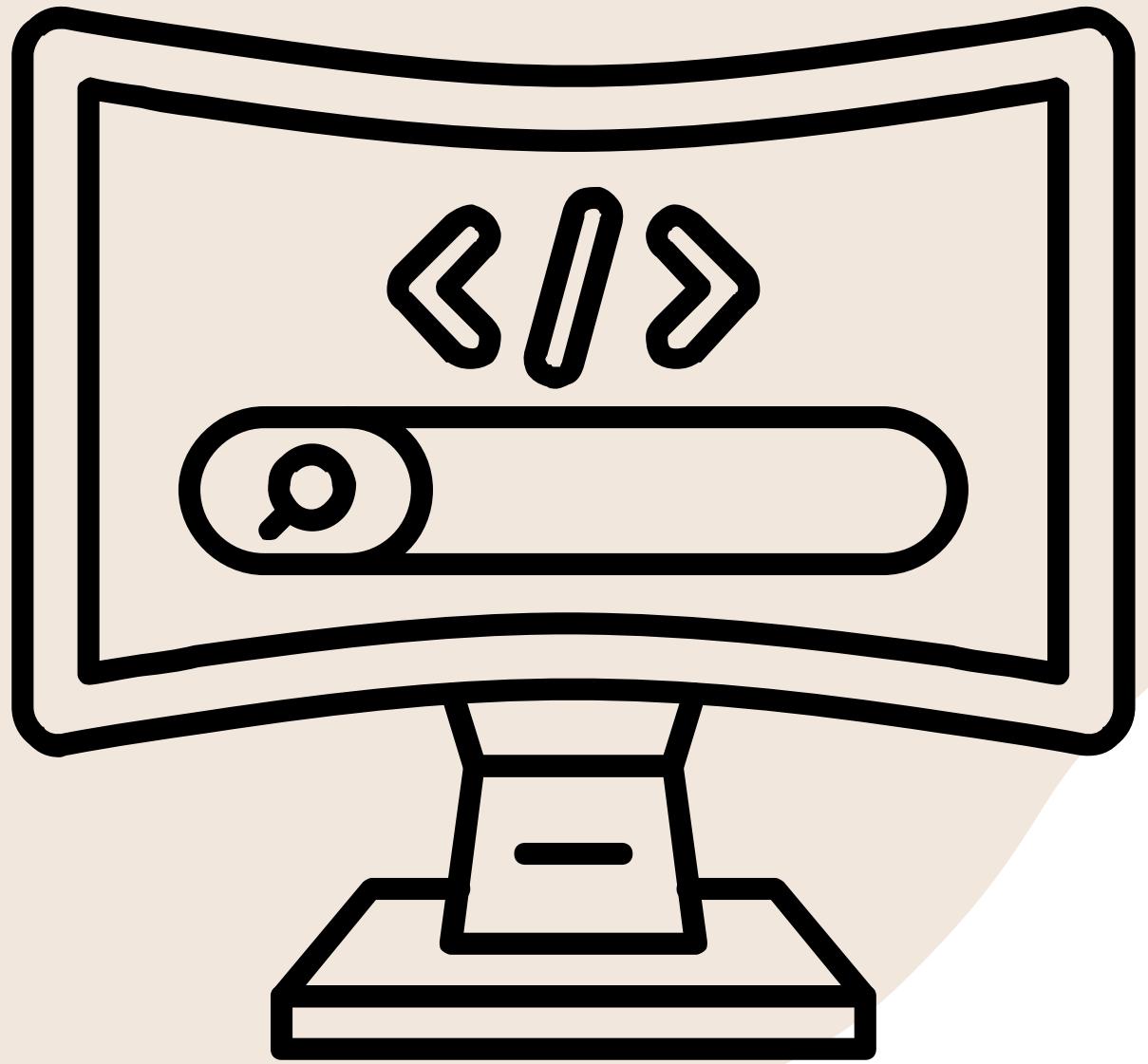
Given a sequence of tokens we verify if the program is syntactically correct, semantically valid, and then translate it.

Objective

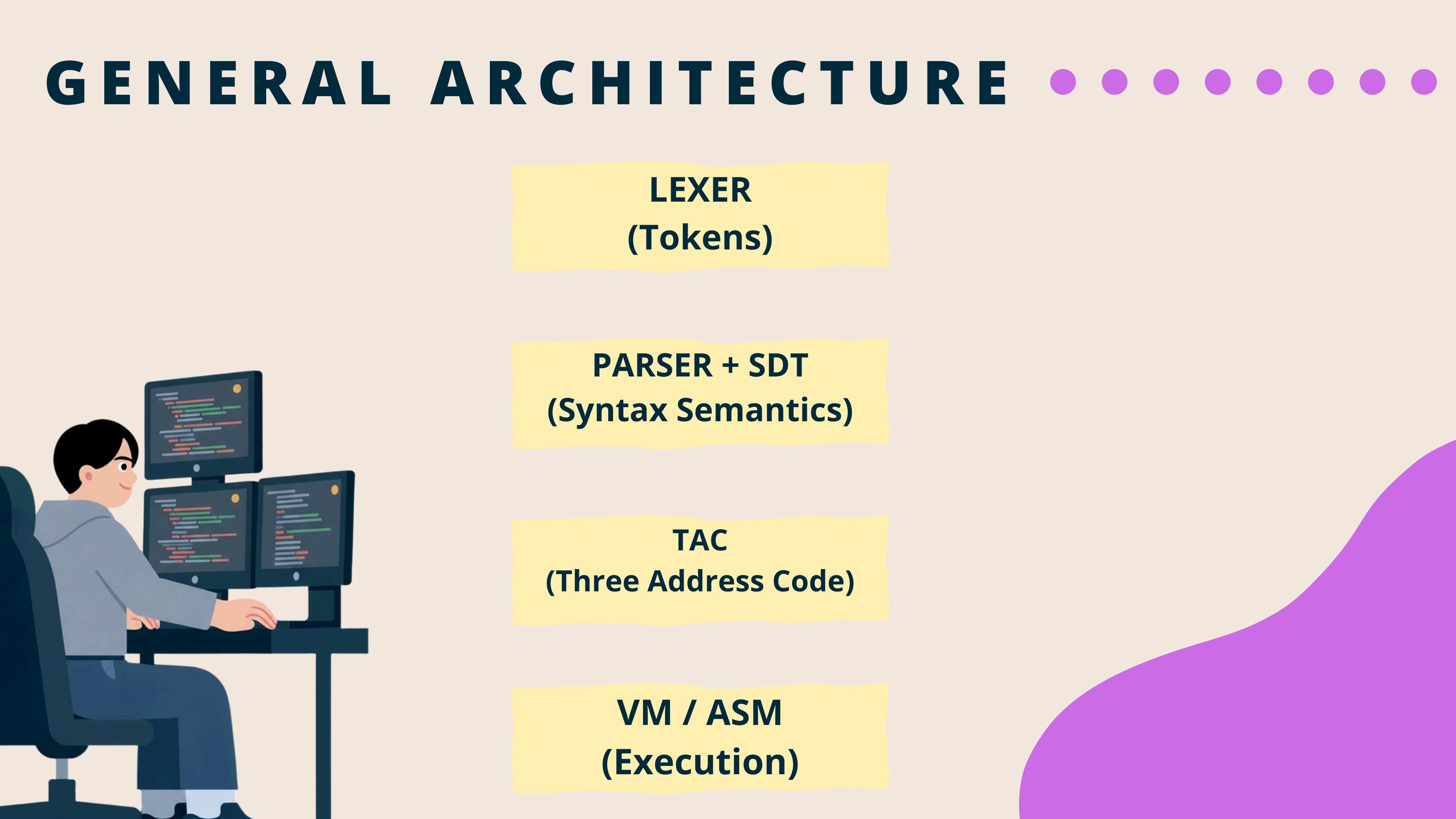
- Design a parser
- Introduce semantic verification
- Generate Three-Address Code (TAC)
- Compile and execute
- Provide clear feedback

C supported

The project implements a reduced sub-language of C



GENERAL ARCHITECTURE



-
-
-
-
-
-
-

LEXER
(Tokens)

PARSER + SDT
(Syntax Semantics)

TAC
(Three Address Code)

VM / ASM
(Execution)



LEXER

01

The lexer is the first phase of the compiler.

It reads the input file character by character and groups them into tokens.

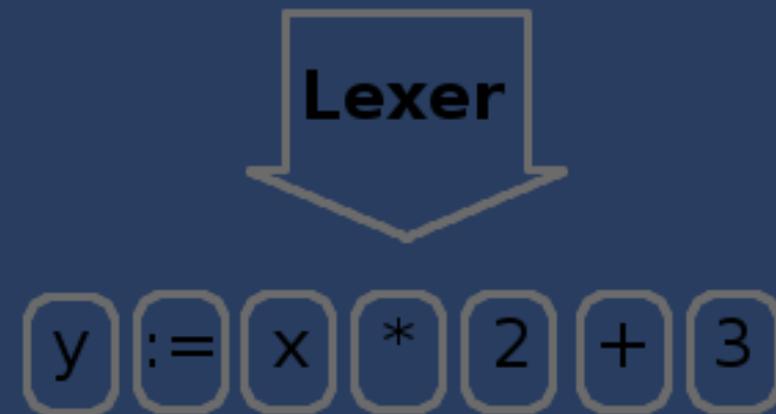
Removes whitespace, newlines, and comments.

02

Recognizes:

- Identifiers
- Numbers
- Keywords
- Operators
- Delimiters
- Produces a linear token stream used by the parser.

$y := x * 2 + 3$



03

EXAMPLE

The line `int x = 10;` translates to:

Keyword Identifier
Operator Constant
Punctuation.

Tokens produced:
[int] [x] [=] [10] [;]

PARSER



- The parser receives the token stream from the lexer.
- Its job is to check whether the sequence follows the grammar of the C subset.
- We implemented a Bottom-Up parser, based on:
 - An analysis stack
 - The `try_reduce()` function to apply grammar rules
- Detects syntax errors, such as:
 - Misordered tokens
 - Incomplete structures
 - Invalid declarations

EXAMPLE

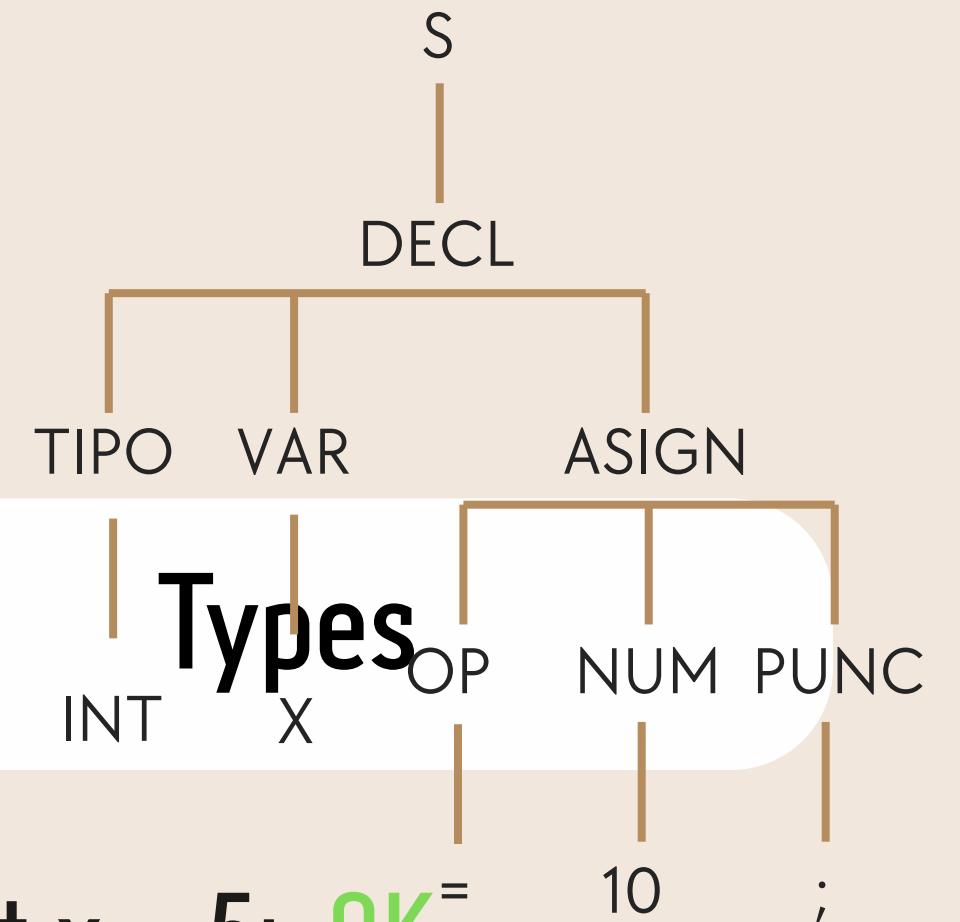
Tokens:
`int x = 10 ;`

Reductions:
 $\text{int} \rightarrow \text{TYPE}$
 $x \rightarrow \text{ID}$

$\text{TYPE ID} = \text{NUM} ; \rightarrow \text{DECLARATION}$



Parser Tree



ERRORS

- Parsing Error (Syntax)
 - int x = ; missing item
- SDT Error (Semantic):
 - int x;
 - x = "cadena"; // tipos incompatibles

Grammar

- [1] S → DECL
- [2] DECL → TIPO VAR ASIGN
- [3] TIPO → KW
- [4] VAR → ID
- [5] ASIGN → OP NUM PUNC
- [5] ASIGN → OP LIT PUNC

INTERMEDIATE CODE (THREE ADDRESS CODE - TAC)



Central idea: each operation is broken down into simple instructions with a maximum of three operands (destination = op1 operator op2).

Advantage:

- Machine independence → easier to optimize.
- Clarity in flow control (labels, jumps).
- Basis for translating to ASM.

EXAMPLE

$x = a + b * c;$

$t1 = b * c$

$x = a + t1$



INTERMEDIATE CODE TO MACHINE



Example Code (Left Side TAC vs Right Side ASM):

TAC: $t1 = a + b$

ASM: LOAD A, a

ADD A, b

STORE A, t1

THE VIRTUAL MACHINE (EXECUTION)



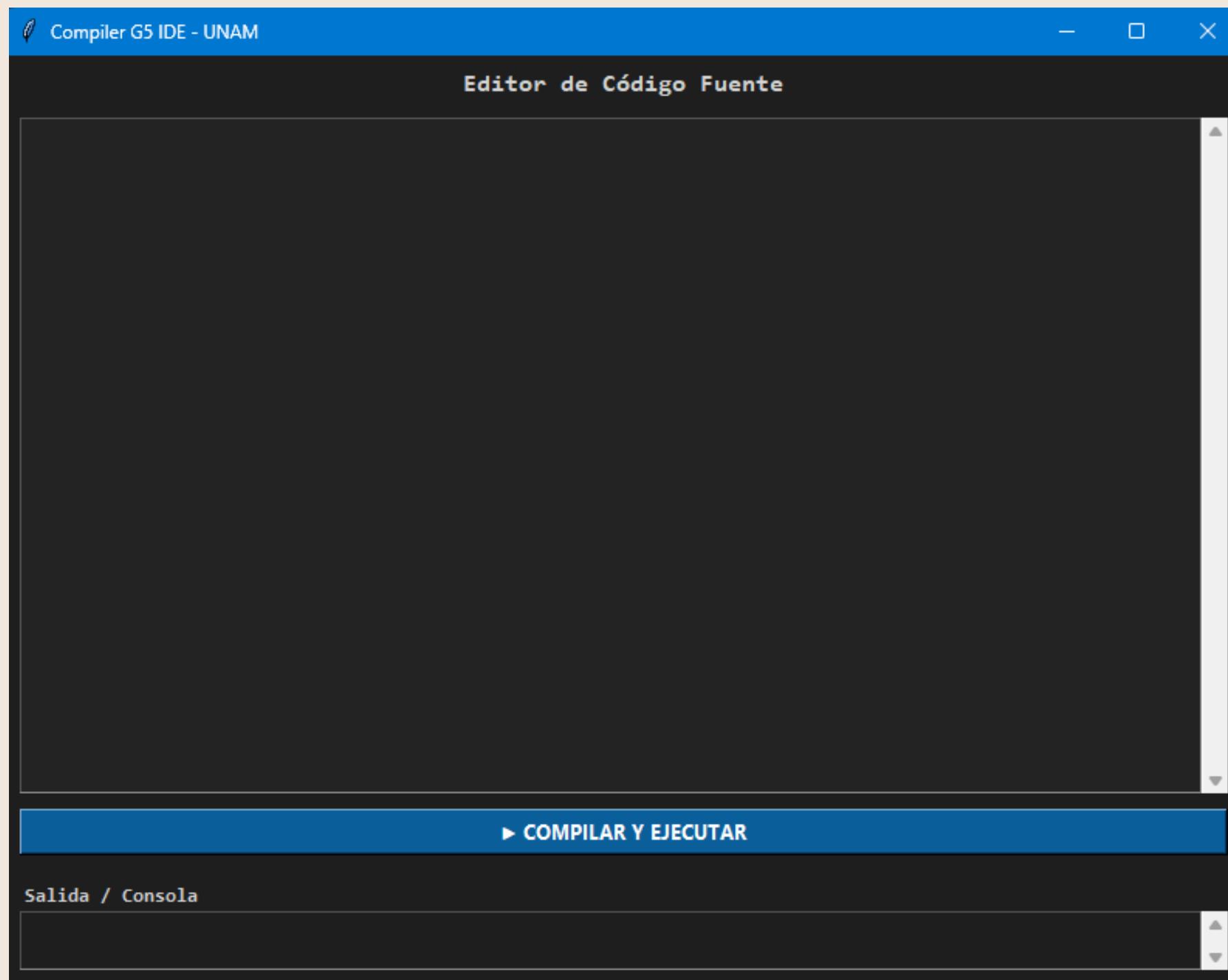
List of VM Features:

- Accumulator Architecture (Register 'A').
- Data Memory (Dynamic Symbol Table).
- Supported Instructions: Arithmetic (ADD, SUB), Logic (JGT, JLT), I/O (PRINT).

Code snippet from
vm_asm.c:

```
case OP_ADD:  
    A += get_value(inst->arg);  
    pc++;
```

GRAPHICAL USER INTERFACE

A screenshot of the Compiler G5 IDE - UNAM. The window title is "Compiler G5 IDE - UNAM". The main area is titled "Editor de Código Fuente". The code in the editor is:

```
int a = 10 ;
int b = 20 ;

if ( a < b ) {
    print ( "A_es_menor" ) ;
    print ( a ) ;
}
```

Below the editor is a blue button labeled "► COMPILEAR Y EJECUTAR".

SUCCESSFULLY COMPLETED (WHILE)



The screenshot shows the Compiler G5 IDE - UNAM interface. The title bar reads "Compiler G5 IDE - UNAM". The main window is titled "Editor de Código Fuente". The code area contains the following Python-like pseudocode:

```
int i = 0 ;
while ( i < 5 ) {
    print ( i ) ;
    i = i + 1 ;
}
```

Below the code area is a blue button labeled "► COMPILAR Y EJECUTAR". The bottom section is titled "Salida / Consola" and displays the output of the executed code:

```
Result of operation: 0
Result of operation: 1
Result of operation: 2
Result of operation: 3
Result of operation: 4
```

At the very bottom of the console output, the message "[PROCESO TERMINADO CON EXITO]" is displayed in green text.

SUCCESSFULLY COMPLETED(IF ELSE)



The screenshot shows the Compiler G5 IDE - UNAM interface. The main window is titled "Editor de Código Fuente" and contains the following C-like code:

```
int a = 10 ;
int b = 20 ;

if ( a < b ) {
    int verdadero = 1 ;
    print ( verdadero ) ;
}
else {
    int falso = 0 ;
    print ( falso ) ;
}

int c = 50 ;
if ( c == 100 ) {
    print ( 1 ) ;
}
else {
    print ( 0 ) ;
}
```

Below the editor is a blue button labeled "► COMPILAR Y EJECUTAR". The bottom section is titled "Salida / Consola" and displays the following output:

```
CompilaciÃ³n TAC -> ASM completada.
[3/3] Ejecutando VM...
-----
Result of operation: 1
Result of operation: 0
-----
[PROCESO TERMINADO CON EXITO]
```

TAC FROM THE LAST EXAMPLE (WHILE)



```
PS C:\Users\alons\Documents\Proyecto\unam.fi.compilers.g5.04> cat .\examples\temp_gui.txt | .\bin\lexer.exe | .\bin\parser.exe
    i = 0
L1:
    if i >= 5 goto L2
    print i
    t1 = i + 1
    i = t1
    goto L1
L2:
```

An other example (IF ELSE)

```
PS C:\Users\alons\Documents\Proyecto\unam.fi.compilers.g5.04> cat .\examples\negativo.txt | .\bin\lexer.exe | .\bin\parser.exe
    a = 10
    b = 20
    if a < b goto L1
    goto L2
L1:
    verdadero = 1
    print verdadero
    goto L3
L2:
    falso = 0
    print falso
L3:
    c = 50
    if c == 100 goto L4
    goto L5
L4:
    print 1
    goto L6
L5:
    print 0
L6:
```

SUCCESSFULLY COMPLETED



Compiler G5 IDE - UNAM

Editor de Código Fuente

```
int a = 2 ;
int b = 10 ;
int c = 3 ;

int discriminante = 0 ;
int term_uno = 0 ;
int term_dos = 0 ;

term_uno = b * b ;
term_dos = 4 * a ;
term_dos = term_dos * c ;

discriminante = term_uno - term_dos ;

print ( discriminante );
```

► COMPILAR Y EJECUTAR

Salida / Consola

```
[2/3] Compilando a Ensamblador...
CompilaciÃ³n TAC -> ASM completada.

[3/3] Ejecutando VM...
-----
Result of operation: 76
-----
[PROCESO TERMINADO CON EXITO]
```

Compiler G5 IDE - UNAM

Editor de Código Fuente

```
int x = 10 ;
int valor = 0 ;

valor = x - 5 ;

if ( valor == 2 ) {
    print ( valor ) ;
}
else {
    print ( "El valor es distinto de 2" ) ;
}
```

► COMPILAR Y EJECUTAR

Salida / Consola

```
[2/3] Compilando a Ensamblador...
CompilaciÃ³n TAC -> ASM completada.

[3/3] Ejecutando VM...
-----
El valor es distinto de 2
-----
[PROCESO TERMINADO CON EXITO]
```

Successfully completed [Console]

```
PS C:\Users\alons\Documents\Proyecto\unam.fi.compilers.g5.04> ./run.bat examples\temp_gui.txt
[1/3] Lexer - Parser (Analisis)...
-----
[OK] Sintaxis Correcta. Generando TAC.

[2/3] Compilando a Ensamblador...
Compilaci|n TAC -> ASM completada.

[3/3] Ejecutando VM...
-----
Result of operation: 0
Result of operation: 1
Result of operation: 2
Result of operation: 3
Result of operation: 4
-----
PS C:\Users\alons\Documents\Proyecto\unam.fi.compilers.g5.04> |
```

Operation failed [Console]

```
PS C:\Users\alons\Documents\Proyecto\unam.fi.compilers.g5.04> ./run.bat examples\temp_gui.txt
[1/3] Lexer - Parser (Analisis)...
-----
[ERROR DE SINTAXIS]
-----
Ubicacion: Linea 2
Esperado: Token tipo 22
Encontrado: 'while'
-----
[X] SE ENCONTRO UN ERROR. PROCESO DETENIDO.
Revisa el mensaje de arriba para mas detalles.
```

OPERATION FAILED



Compiler G5 IDE - UNAM

Editor de Código Fuente

```
int discriminante = 0/0 ;
print ( discriminante ) ;
```

▶ COMPILAR Y EJECUTAR

Salida / Consola

```
[2/3] Compilando a Ensamblador...
CompilaciÃ³n TAC -> ASM completada.

[3/3] Ejecutando VM...
-----
error in line 2
```

Compiler G5 IDE - UNAM

Editor de Código Fuente

```
int x = 10 ;
int valor = 0 ;

valor = x - 5

if ( valor == 2 ) {
    print ( valor ) ;
}
else {
    print ( "El valor es distinto de 2" ) ;
}
```

▶ COMPILAR Y EJECUTAR

Salida / Consola

```
[ERROR DETECTADO]
[ERROR DE SINTAXIS]
-----
Ubicacion: Línea 6
Esperado: Token tipo 22
Encontrado: 'if'
-----
```

CONCLUSION

This project demonstrates a reliable, efficient, and fully functional compilation pipeline capable of analyzing, validating, and executing high-level code with precision. By integrating lexical, syntactic, and semantic analysis with TAC generation and virtual-machine execution, we deliver a tool that ensures correctness, prevents critical errors early, and guarantees consistent program behavior. The system's modular design makes it scalable, maintainable, and ready to extend toward larger languages or more complex applications. Overall, this solution provides a solid, professional-grade foundation for any organization seeking to build or enhance a custom compiler or automated code-analysis tool.

THANK YOU
FOR YOUR
ATTENTION

