C:\cygwin\home\zahorjan\cse461\12au\05-retransmission.cp3

CSE 461 – Retransmission

Reliability

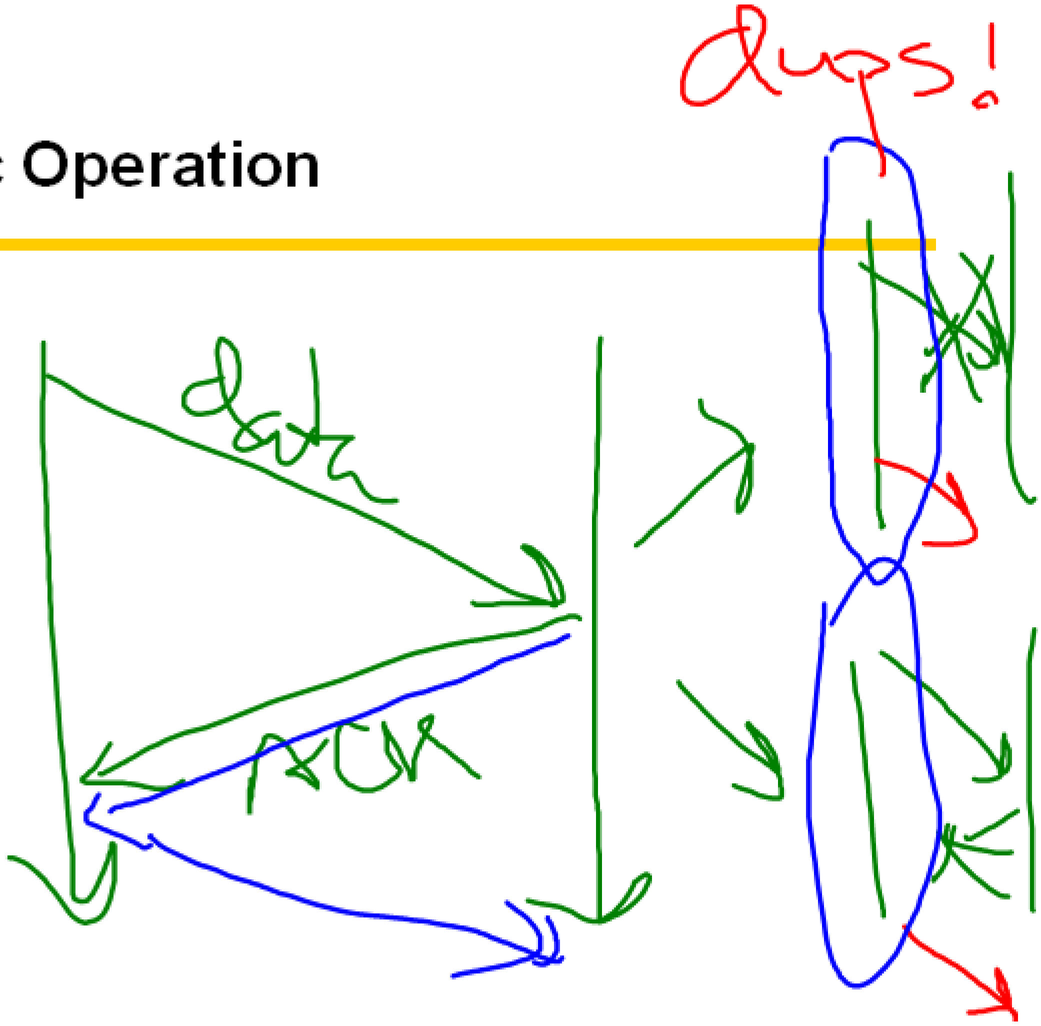
- Two strategies to handle errors:
  - Detect and retransmit, or Automatic Repeat reQuest. (ARQ)
  - Error correcting codes, or *Forward Error Correction* (FEC)

- We're seen error detecting/correcting codes, now look at retransmissions

Operational Assumptions

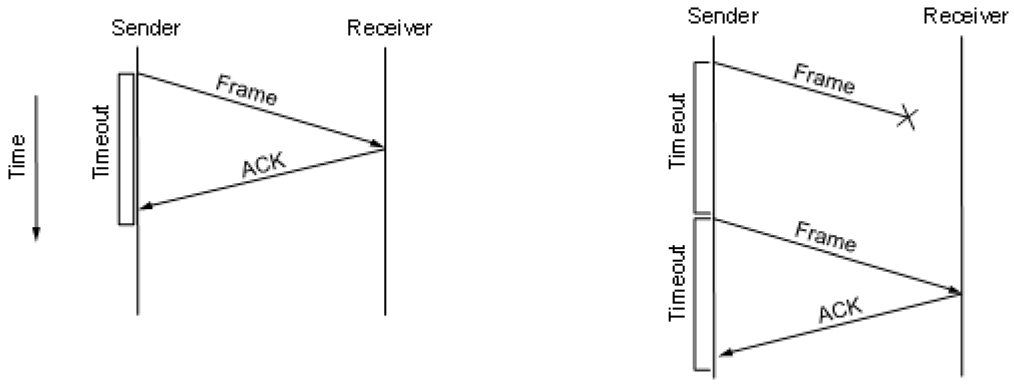
- Frames can simply be undetected by receiver
  - It hears nothing at all
  - Therefore the receiver doesn't do anything in this situation
- Bits in frames can be mangled
  - Receiver detects, corrects when it can, but there are still errors
  - What should receiver do?
- Frames that are received (haven't gone undetected) are received in the order they were sent
  - No overtaking

Basic Operation



4

Retransmissions, or more formally Automatic Repeat Request (ARQ)



- Sender automatically resends after a timeout until a positive acknowledgment (ACK) is obtained from the receiver
- Receiver automatically acknowledges frames (packets) that are not corrupted or lost in the network
- ARQ is generic name for protocols based on this strategy

Two issues

1. How long to set the timeout?
   - Only easy on a direct link, otherwise timing variability
     - Way too long lowers performance
     - Short implies sometimes timeout will be early
2. How to avoid accepting duplicate frames as new
   - Given retransmissions, frame loss, and imprecise timeouts

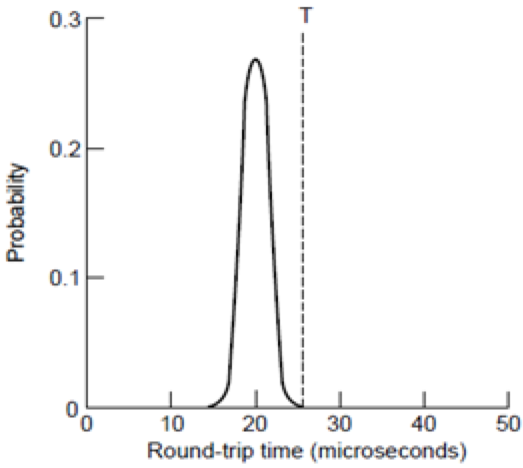
Why might the receiver see duplicates of a single frame?

6

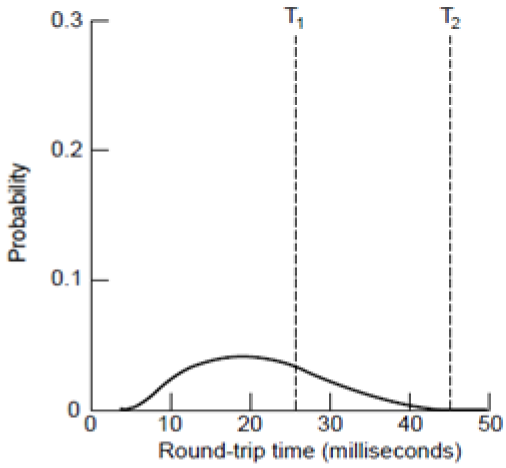
Timeouts

Retransmission timeout depends on round-trip time

- To send frame and receive an acknowledgement
- In general, need to account for variance on complex paths

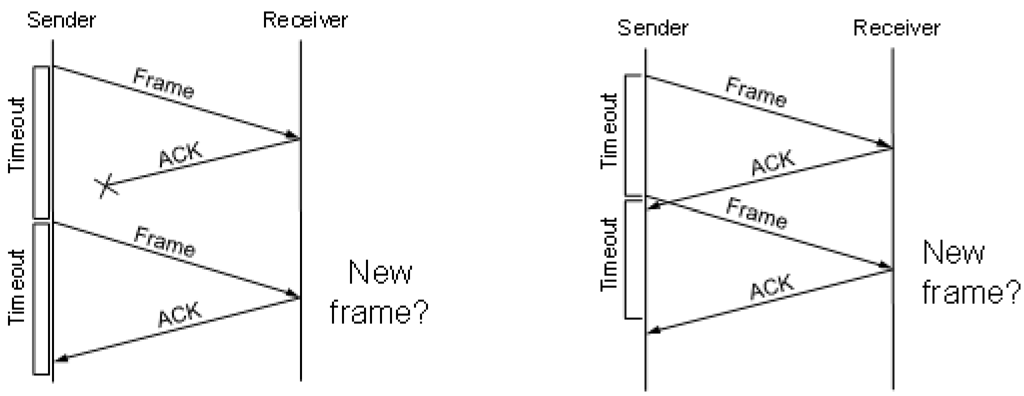


LAN case – small, regular RTT



Internet case – large, varied RTT

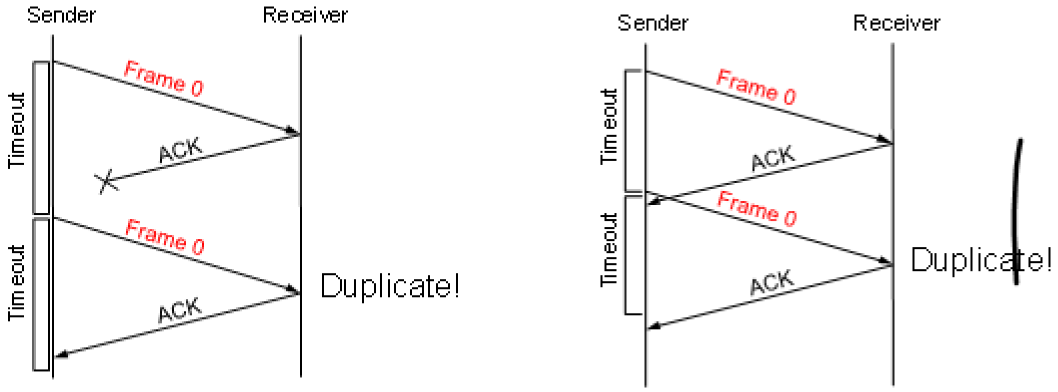
Problem cases (due to loss, timeouts)



- In the case of ACK loss (or poor choice of timeout) the receiver can't distinguish current message from next

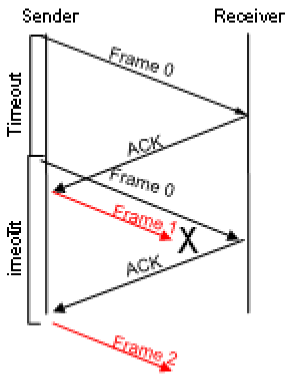
8

The Need for Sequence Numbers

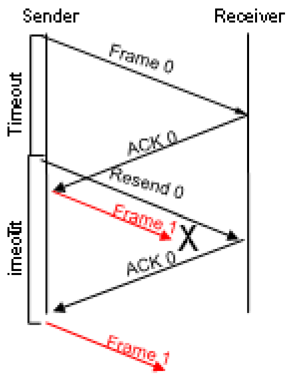


- Frame sequence numbers let receiver tell next frame from duplicate transmission

ACKs need sequence numbers too



The Problem Scenario

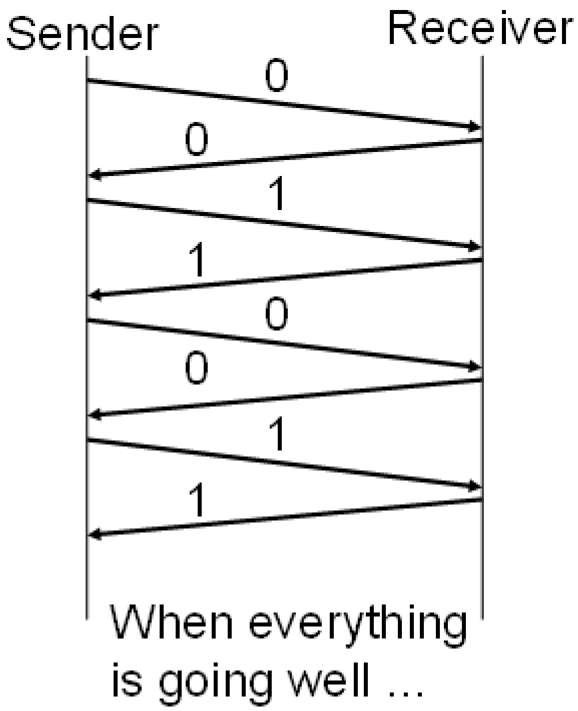


The Solution

- Hm, these things can be tricky!

Stop-and-Wait

- Only one outstanding frame at a time, 0 or 1.
- Retransmissions re-sent with same number
- Number only needs to distinguish between current and next frame
  - A single bit will do

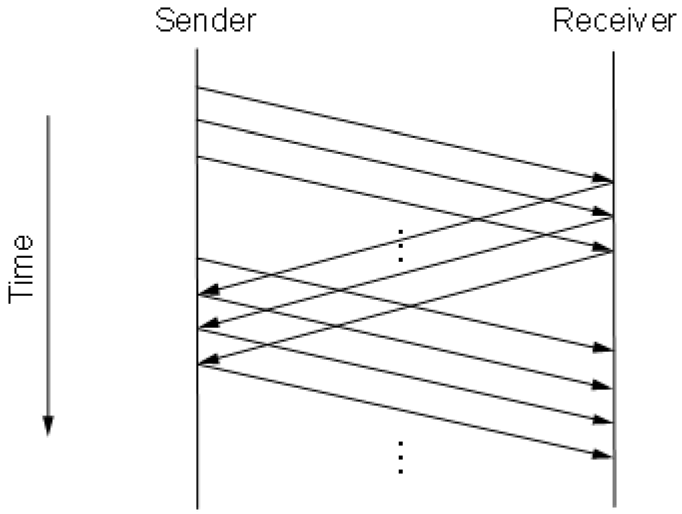


Limitation of Stop-and-Wait

- Lousy performance if transmission time << prop. delay
  - How bad? You do the math
- Want to utilize all available bandwidth
  - Need to keep more data "in flight"
  - How much? The "bandwidth-delay product":
    bits/sec * seconds = bits
- Leads to *Sliding Window Protocol*

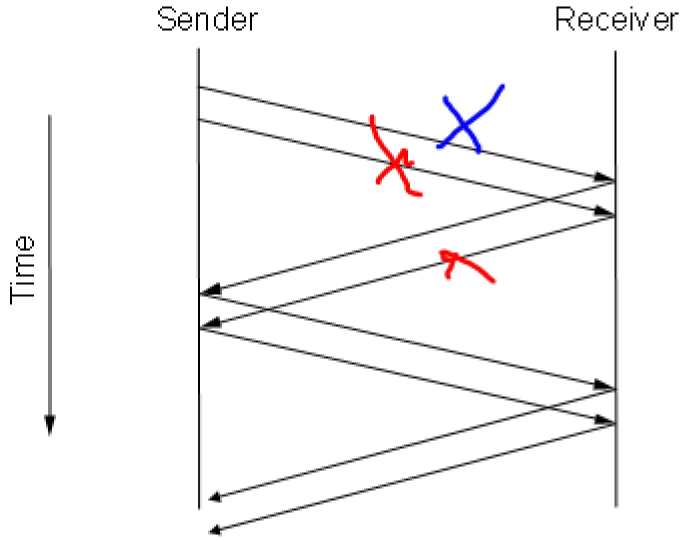
Solution: Allow Multiple Frames in Flight

- This is a form of pipelining



Sliding Window Protocol

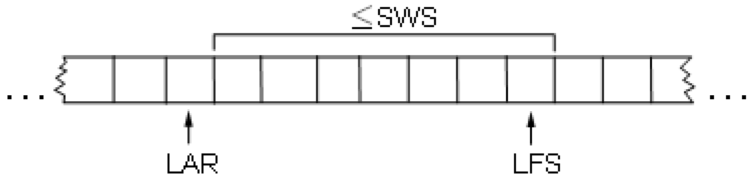
- There is some maximum number of un-ACK'ed frames the sender is allowed to have in flight
  - We call this "the window size"
  - Example: window size = 2



Once the window is full, each ACK'ed frame allows the sender to send one more frame

Sliding Window: Sender

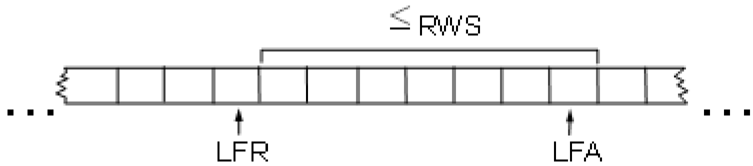
- Assign sequence number to each frame (`SeqNum`)
- Maintain three state variables:
  - send window size (`SWS`)
  - last acknowledgment received (`LAR`)
  - last frame sent (`LFS`)
- Maintain invariant: `LFS - LAR <= SWS`



- Advance `LAR` when ACK arrives
- Buffer up to `SWS` frames

Sliding Window: Receiver

- Maintain three state variables
  - receive window size (**RWS**)
  - largest frame acceptable (**LFA**)
  - last frame received (**LFR**)
- Maintain invariant: **LFA - LFR <= RWS**



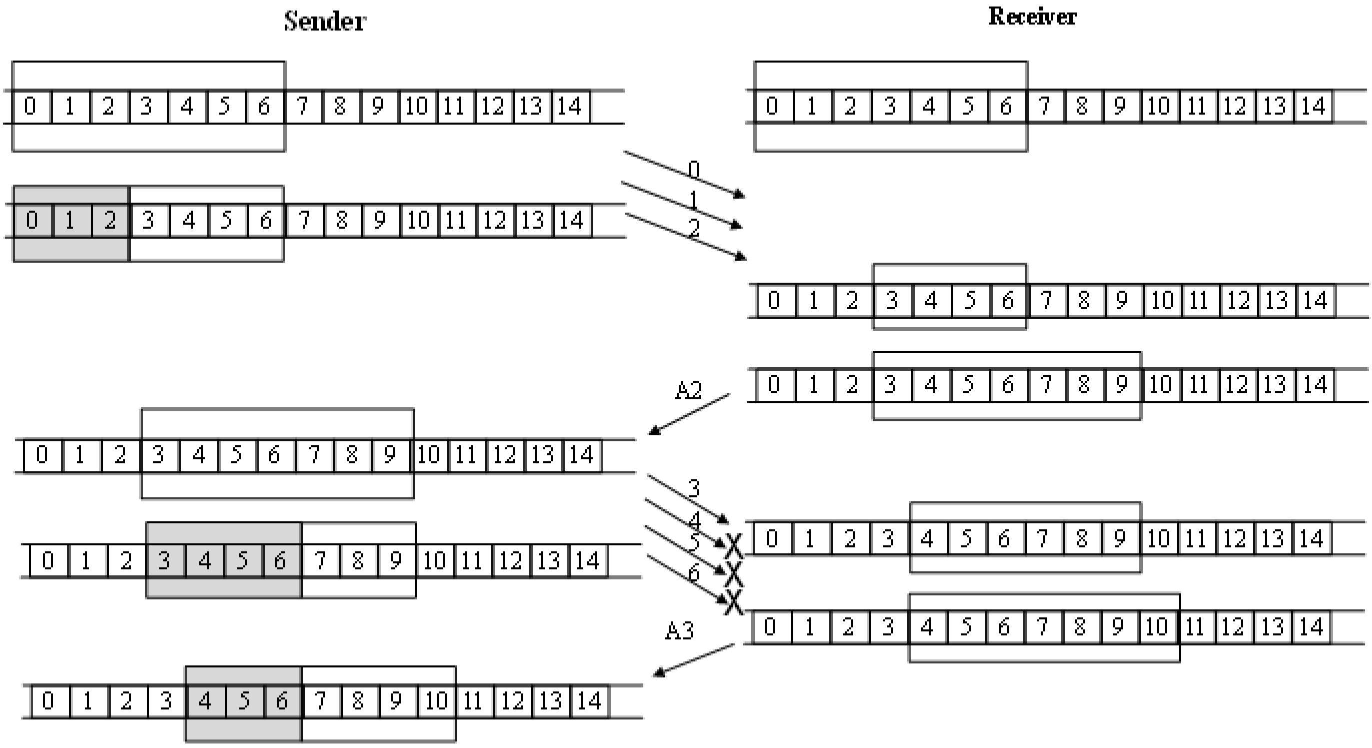
- Frame **SeqNum** arrives:
  - if **LFR** < **SeqNum** ≤ **LFA** ⇒ accept else discard
  - send ACK to tell sender what has arrived (new or repeat)
- Advance **LFR** (and pass to application) as in-order frames arrive
- Need to buffer up to **RWS** frames

Acknowledgement options

- Different options are possible:
- Send cumulative ACKs – send ACK for largest frame such that all frames less than this have been received
  - Robust to ACK loss but not packet loss
- Send individual ACKs
  - Robust to packet loss but not ACK loss!
- Can combine:
  - Idea is to tell the sender what frames the receiver already has
  - Usually have cumulative ACK plus hints

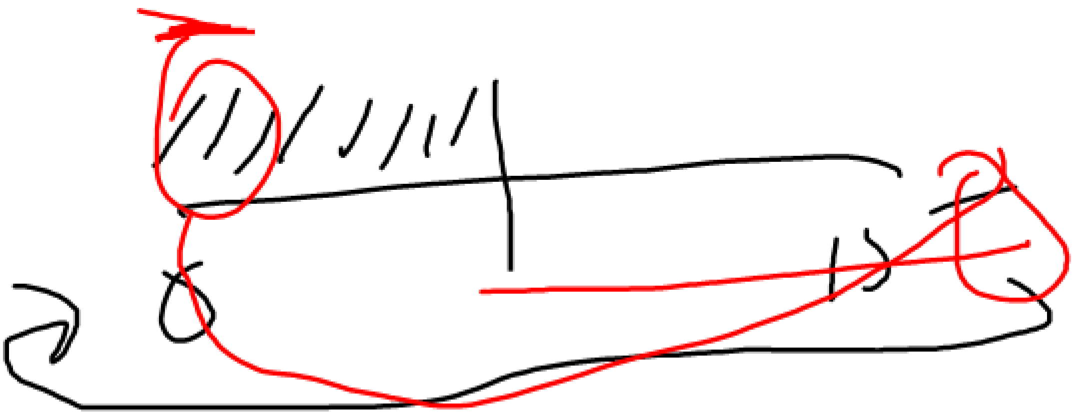
17

Sliding Window Example



Sequence Number Space

- `SeqNum` field is finite; sequence numbers wrap around
- Sequence number space must be larger then number of outstanding frames
- `SWS <= MaxSeqNum-1` is not sufficient
- `SWS < (MaxSeqNum+1)/2` is correct rule
- Intuitively, `SeqNum` “slides” between two halves of sequence number space



Sliding Window Summary

- It is perhaps the best known algorithm in networking
- First role is to enable reliable delivery of packets
  - Timeouts and acknowledgements
  - This has been our focus
- Second role is to enable in order delivery of packets
  - Receiver doesn't pass data up to app until it has packets in order
- Third role is to enable pipelined transmission
  - Crucial for high latency transmissions
- Fourth role is to enable flow control
  - Prevents fast sender from overflowing slow receiver's buffer
  - We will see this when we get to TCP

When to use ARQ or FEC?

- Will depend on the kind of errors and cost of recovery
- Example: Message with 1000 bits, Prob(bit error) 0.001
  - Case 1: random errors
  - Case 2: bursts of 1000 errors
- Q: What to use in Case 1 and 2?

ARQ vs. FEC

- FEC used at low-level to lower residual error rate
- ARQ often used to fix large errors, e.g., packet collision, and with detection to protect against residual errors

- FEC sometimes used at high level too:
  - Real time applications (no time to retransmit!)
  - Nice interaction with broadcast (different receiver errors!)

Example: 802.11

- The standard scheme is:

- PHY: FEC on data via interleaving and a binary convolutional code or LDPC
  - rates from ½ to 5/6.
- PHY header has 16 bit CRC
- Link: 32 bit CRC on frame and retransmission