

# Computer Networks

## Error Detection (§3.2.2)



David Wetherall (djw@uw.edu)

Professor of Computer Science & Engineering

UNIVERSITY *of* WASHINGTON

# Topic

- Some bits may be received in error due to noise. How do we detect this?
  - ➔ Parity »
  - ➔ Checksums »
  - ➔ CRCs »
- Detection will let us fix the error, for example, by retransmission (later).

# Simple Error Detection – Parity Bit

- Take D data bits, add 1 check bit that is the sum of the D bits
  - Sum is modulo 2 or XOR

10011001 ← parity bit

## Parity Bit (2)

- How well does parity work?
  - What is the distance of the code?  
2
  - How many errors will it detect/correct?  
1 2
- What about larger errors?  
odd # errors

# Checksums

- Idea: sum up data in N-bit words
  - Widely used in, e.g., TCP/IP/UDP



- Stronger protection than parity

# Internet Checksum

- Sum is defined in 1s complement arithmetic (must add back carries)
  - And it's the negative sum
- { *"The checksum field is the 16 bit one's complement of the one's complement sum of all 16 bit words ..."* – RFC 791

1's comp 001 = "1"  
1's comp 110 = "1"      1's comp 111

# Internet Checksum (2)

Sending:

1. Arrange data in 16-bit words
2. Put zero in checksum position, add
3. Add any carryover back to get 16 bits
4. Negate (complement) to get sum

$$\begin{array}{r} 0001f203 \\ f4f5f6f7 \\ + 0000 \\ \hline 2ddf02 \\ + 2 \\ \hline ddf2 \\ 220d \end{array}$$

$1 = 1101$  flip  $0010$

# Internet Checksum (3)

Sending:

1. Arrange data in 16-bit words
2. Put zero in checksum position, add
3. Add any carryover back to get 16 bits
4. Negate (complement) to get sum

```
0001
f203
f4f5
f6f7
+ (0000)
-----
2ddf0
  ↓
  ddf0
+      2
-----
  ddf2
  ↓
  220d
```



# Internet Checksum (4)

Receiving:

1. Arrange data in 16-bit words
2. Checksum will be non-zero, add
3. Add any carryover back to get 16 bits
4. Negate the result and check it is 0

$$\begin{array}{r} 0001 \\ f203 \\ f4f5 \\ f6f7 \\ + 220d \\ \hline 2fffd \\ fffd \\ + 2 \\ \hline ffff \\ 0000 \quad \checkmark \end{array}$$

# Internet Checksum (5)

Receiving:

1. Arrange data in 16-bit words
2. Checksum will be non-zero, add
3. Add any carryover back to get 16 bits
4. Negate the result and check it is 0

```
0001
f203
f4f5
f6f7
+ 220d
-----
2fffd
  ↓
  fffd
+    2
-----
ffff
  ↓
0000
```

# Internet Checksum (6)

- How well does the checksum work?
  - What is the distance of the code? 2
  - How many errors will it detect/correct?

1 0

- What about larger errors?

all burst errors up to 16  
random w/ prob  $1/2^{16}$

# Cyclic Redundancy Check (CRC)

- Even stronger protection
  - Given  $n$  data bits, generate  $k$  check bits such that the  $n+k$  bits are evenly divisible by a generator  $C$
- Example with numbers:
  - $n = 302$ ,  $k = \text{one digit}$ ,  $C = 3$

$$302 \underline{1}$$

$$3020 \div 3 \\ = 2$$

## CRCs (2)

- The catch:
  - It's based on mathematics of finite fields, in which “numbers” represent polynomials
  - e.g, 10011010 is  $x^7 + x^4 + x^3 + x^1$
- What this means:
  - We work with binary values and operate using modulo 2 arithmetic

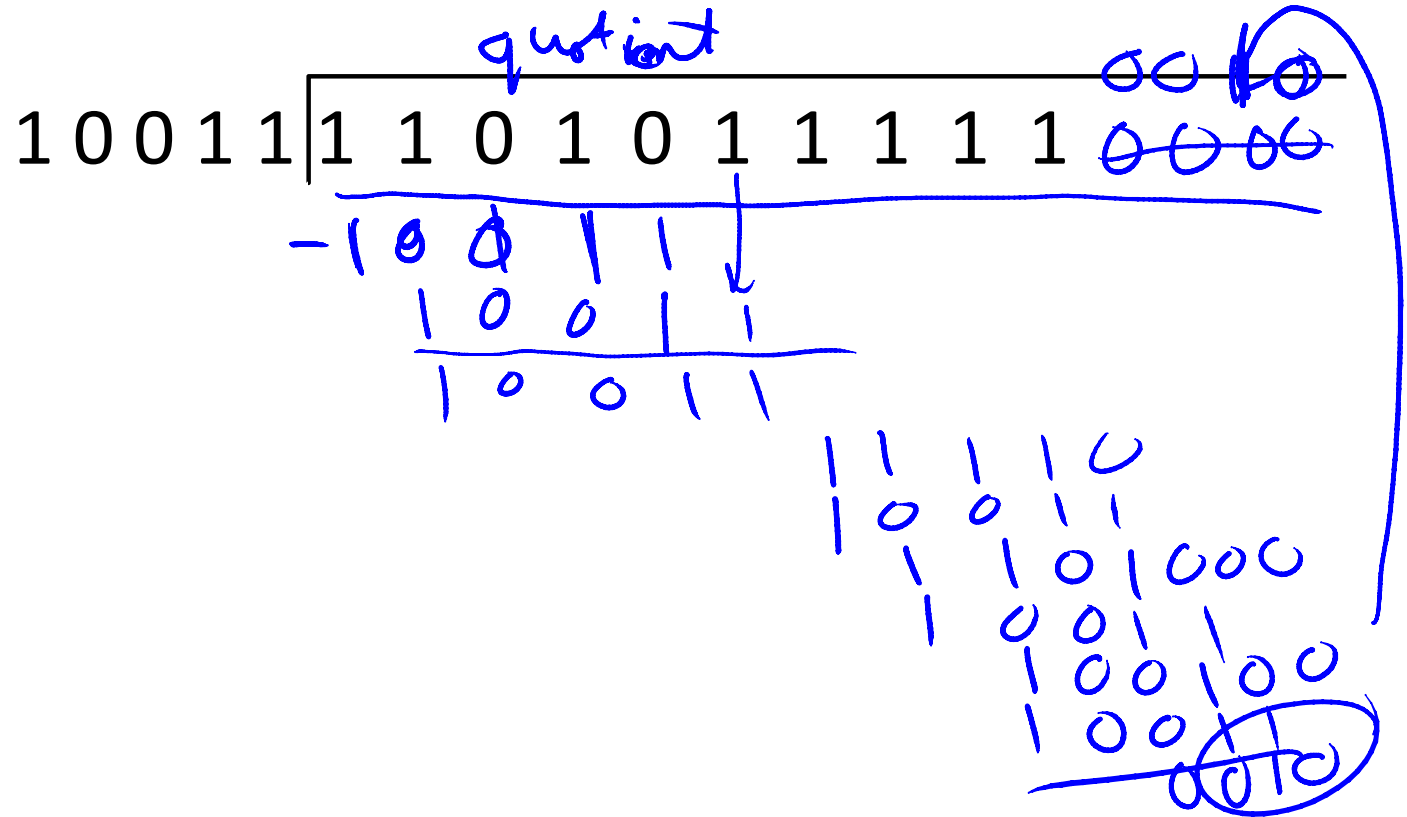
## CRCs (3)

- Send Procedure:
  1. Extend the  $n$  data bits with  $k$  zeros
  2. Divide by the generator value  $C$
  3. Keep remainder, ignore quotient
  4. Adjust  $k$  check bits by remainder
- Receive Procedure:
  1. Divide and check for zero remainder

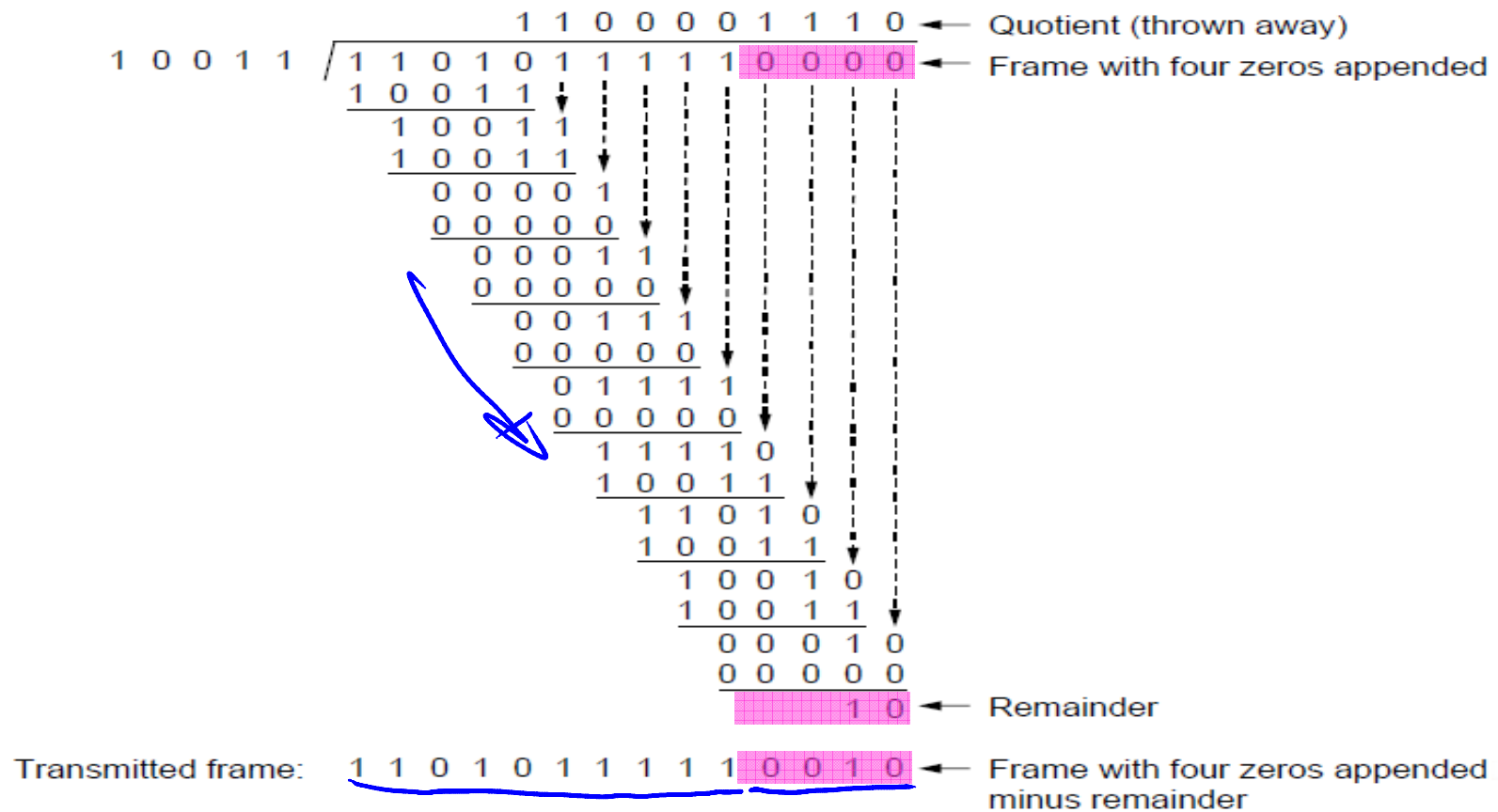
# CRCs (4)

Data bits:  
1101011111

Check bits:  
 $C(x) = x^4 + x^1 + 1$   
 $C = 10011$   
 $k = 4$



# CRCs (5)







## CRCs (6)

- Protection depend on generator

→ Standard CRC-32 is 1 0000 0100  
1100 0001 0001 1101 1011 0111

- Properties:
  - HD=4, detects up to triple bit errors
  - Also odd number of errors
  - And bursts of up to k bits in error
  - Not vulnerable to systematic errors like checksums

# Error Detection in Practice

-  CRCs are widely used on links
  - Ethernet, 802.11, ADSL, Cable ...
-  Checksum used in Internet
  - IP, TCP, UDP ... but it is weak
- Parity
  - Is little used

# END

© 2013 D. Wetherall

Slide material from: TANENBAUM, ANDREW S.; WETHERALL, DAVID J., COMPUTER NETWORKS, 5th Edition, © 2011.  
Electronically reproduced by permission of Pearson Education, Inc., Upper Saddle River, New Jersey