

Computer Networks

TCP Slow Start (§6.5.10)



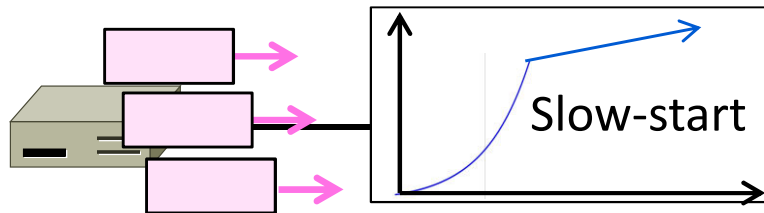
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Topic

- How TCP implements AIMD, part 1
 - “Slow start” is a component of the AI portion of AIMD



Recall

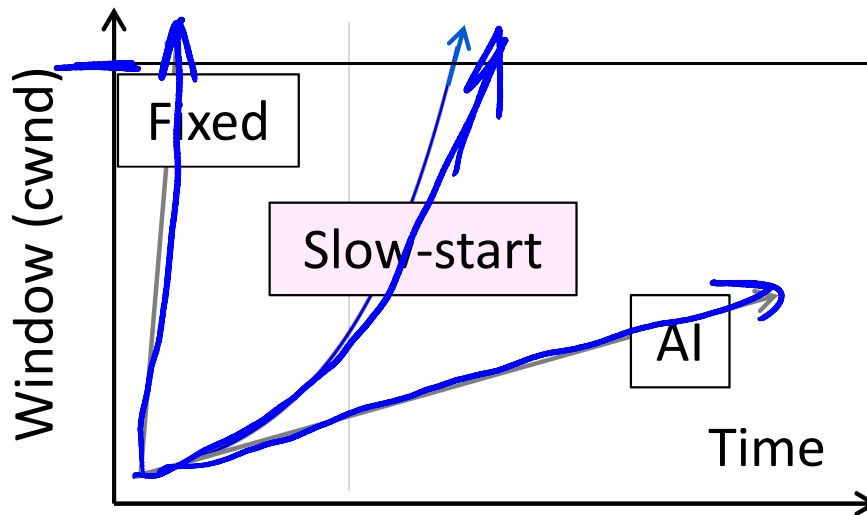
- We want TCP to follow an AIMD control law for a good allocation
- Sender uses a congestion window or cwnd to set its rate ($\approx \text{cwnd}/\text{RTT}$)
- Sender uses packet loss as the network congestion signal
- Need TCP to work across a very large range of rates and RTTs

TCP Startup Problem

- We want to quickly near the right rate, $cwnd_{IDEAL}$, but it varies greatly
 - Fixed sliding window doesn't adapt and is rough on the network (loss!)
 - AI with small bursts adapts $cwnd$ gently to the network, but might take a long time to become efficient

Slow-Start Solution

- Start by doubling cwnd every RTT
 - Exponential growth (1, 2, 4, 8, 16, ...)
 - Start slow, quickly reach large values

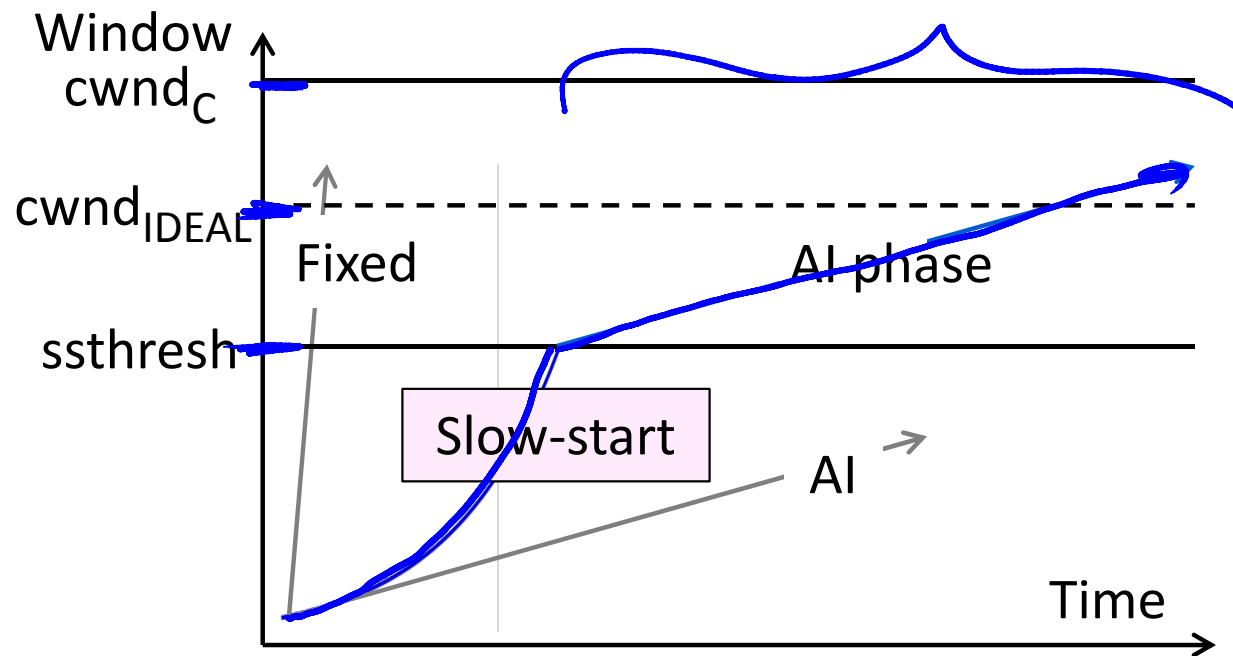


Slow-Start Solution (2)

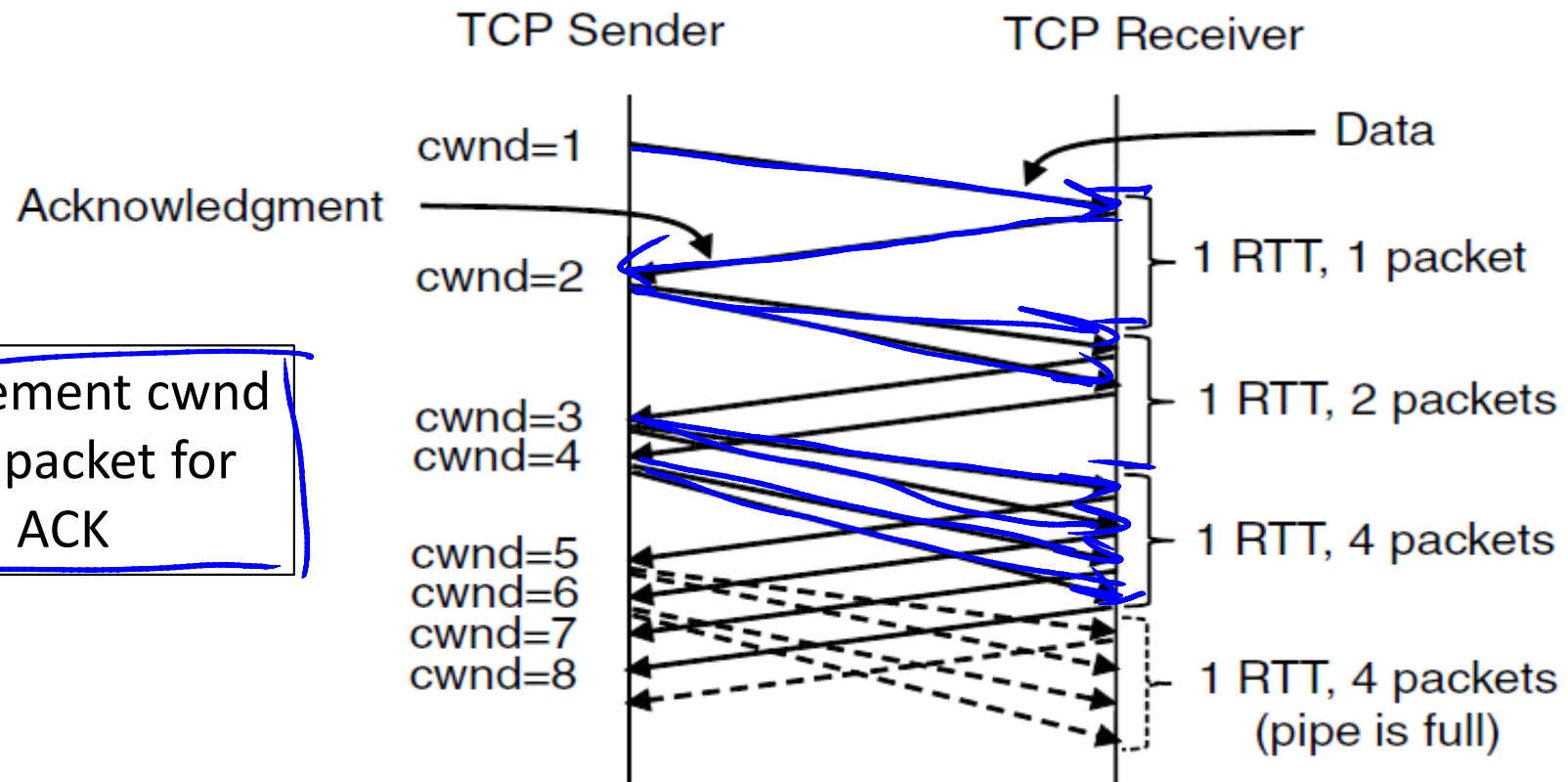
- Eventually packet loss will occur when the network is congested
 - Loss timeout tells us cwnd is too large
 - Next time, switch to AI beforehand
 - Slowly adapt cwnd near right value
- In terms of cwnd:
 - Expect loss for $cwnd_c \approx 2BD + \text{queue}$
 - Use $ssthresh = cwnd_c / 2$ to switch to AI

Slow-Start Solution (3)

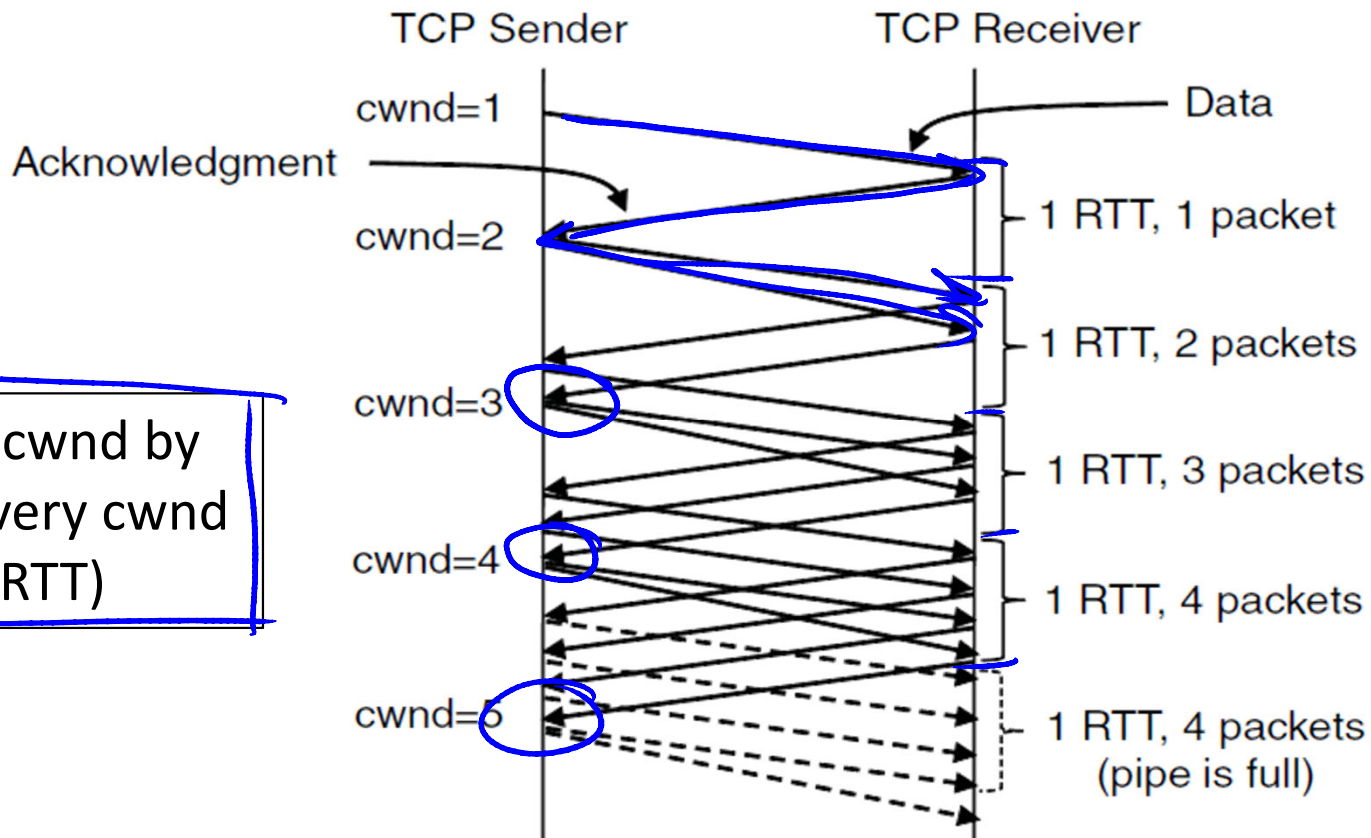
- Combined behavior, after first time
 - Most time spend near right value



Slow-Start (Doubling) Timeline



Additive Increase Timeline



Increment cwnd by 1 packet every cwnd ACKs (or 1 RTT)

TCP Tahoe (Implementation)

- Initial slow-start (doubling) phase
 - Start with $\text{cwnd} = 1$ (or small value)
 - $\text{cwnd} += 1$ packet per ACK
- Later Additive Increase phase
 - $\text{cwnd} += \frac{1}{\text{cwnd}}$ packets per ACK
 - Roughly adds 1 packet per RTT
- Switching threshold (initially infinity)
 - Switch to AI when $\text{cwnd} > \text{ssthresh}$
 - Set $\text{ssthresh} = \text{cwnd}/2$ after loss
 - Begin with slow-start after timeout

Timeout Misfortunes

- Why do a slow-start after timeout?
 - Instead of MD cwnd (for AIMD)
- Timeouts are sufficiently long that the ACK clock will have run down
 - Slow-start ramps up the ACK clock
- We need to detect loss before a timeout to get to full AIMD
 - Done in TCP Reno (next time)

END

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