

# CS 2200 Spring 2011 Final Exam

Name: \_\_\_\_\_ GT Number: \_\_\_\_\_

Problem	Points	Lost	Gained	Running Total	TA
1	1				
2	10				
3	15				
4	15				
5	5				
6	10				
7	5				
8	10				
9	10				
10	10				
11	9				
Total	100				

**You may ask for clarification but you are ultimately responsible for the answer you write on the paper.**

**Please look through the entire test before starting. WE MEAN IT!!!**

**Illegible answers are wrong answers.**

**Show your work in the space provided to get any credit for problem-oriented questions.**

**Good luck!**

1. (1 point, 1 min)

Atlanta Hawks have not made it to their division finals during playoffs since

- a) 1969-70
- b) 2009-2010
- c) 1957-58
- d) They came to Atlanta
- e) Who are Atlanta Hawks?
- f) What is playoff?

# CS 2200 Spring 2011 Final Exam

Name: \_\_\_\_\_ GT Number: \_\_\_\_\_

## Memory management

2. (10 points, 5 mins)

a) (2 points) During the time interval  $t_1$ - $t_2$ , the following virtual page accesses are recorded for the three processes P1, P2, and P3, respectively:

- P1: 0, 25, 1, 0, 1, 2, 10, 2, 1, 1, 0
- P2: 0, 1, 101, 102, 103, 0, 1, 104, 105, 106
- P3: 0, 1, 2, 3, 4, 5, 0, 1, 2, 3, 4, 5

The cumulative memory pressure is (circle one)

- 1. 33
- 2. 11
- 3. 12
- 4. 10
- 5. 19

b) (2 points) To implement paging the minimum additional hardware needed in the CPU data path (circle one)

- 1. One Page table implemented in hardware
- 2. Multiple page table (one per process) implemented in hardware
- 3. One Page Table Base Register (PTBR)
- 4. Multiple PTBR (one per process)
- 5. None of the above

c) (2 points) External fragmentation occurs in the following memory management policies (circle one)

- 1. Fixed size partitions
- 2. Variable size partitions
- 3. Paging
- 4. (1) and (2)
- 5. (2) and (3)
- 6. (1) and (3)
- 7. (1), (2), and (3)

d) (2 points) Virtual address 20 bits; page size 1K bytes. Number of entries in the page table (circle one)

- 1. 2048
- 2. 1024
- 3.  $2^{20}$
- 4. Cannot determine with the given data

e) (2 points) Physical address 24 bits; page size 8K bytes; Maximum number of physical page frames possible (circle one)

- 1. 1024
- 2. 2048
- 3.  $2^{24}$
- 4. 8K
- 5. Cannot determine with the given data

# CS 2200 Spring 2011 Final Exam

Name: \_\_\_\_\_ GT Number: \_\_\_\_\_

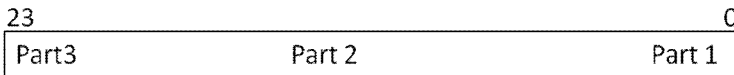
## Cache

3. (15 points, 15 min)

a) (6 points)

A computer has the following characteristics:

- It is **word-addressed** and the address is 24 bits long (little endian)
- It has a **direct-mapped cache** with each block holding 32 words of data
- The direct-mapped cache has **2048 sets**



The above figure shows how the address is divided into three parts to perform a cache access. Name the parts (below) and the number of bits associated with each part.

Part 1: \_\_\_\_\_ (name) \_\_\_\_\_ (number of bits)

Part 2: \_\_\_\_\_ (name) \_\_\_\_\_ (number of bits)

Part 3: \_\_\_\_\_ (name) \_\_\_\_\_ (number of bits)

b) (5 points)

A fully associative cache is initially empty, has only four blocks, and uses the **FIFO replacement policy**. The processor performs a total of eight accesses, to memory blocks A, B, C, D, A, E, A, and B, in that order. For each of these accesses, specify (by filling in the table below) whether it is a cache hit or a cache miss, and the memory block evicted (if any).

Memory Access	Hit/miss	Block evicted from cache
A		
B		
C		
D		
A		
E		
A		
B		

c) (2 points) Average CPI = 1.5; average cache miss per instruction = 3%; miss penalty = 20. The effective CPI is

1. 1.8
2. 2.1
3. 21.5
4. 7.5

# CS 2200 Spring 2011 Final Exam

Name: \_\_\_\_\_ GT Number: \_\_\_\_\_

- d) (2 points) Temporal locality suggest that
1. Once brought into the cache, we should keep the data around as long as possible
  2. On a miss, we should bring in adjacent memory locations into the cache
  3. The memory location being brought in due to a miss is not likely to be referenced in the future
  4. None of the above

## Input/Output

4. (15 points, 10 min)

a) Given the following specifications for a disk drive:

256 bytes per sector  
12 sectors per track  
20 tracks per surface  
3 Platters  
2 surfaces per platter  
Seek time 20 ms  
Rotational speed 3600 RPM

i)(2 points) What is the total capacity of such a drive in bytes assuming normal recording?

ii)(3 points) Assume a zoned bit recording with 3 zones  
Zone 3 (outermost): 8 tracks, 18 sectors per track  
Zone 2: 7 tracks, 14 sectors per track  
Zone 1: 5 tracks, 12 sectors per track

What is the total capacity of this drive with the zoned-bit recording?

# CS 2200 Spring 2011 Final Exam

Name: \_\_\_\_\_ GT Number: \_\_\_\_\_

iii) (6 points) Assuming normal recording, calculate the average time needed to read 2 contiguous sectors from the same track (Note: the head can be anywhere).

b) (2 point) Give an example of a synchronous I/O device and an example of an asynchronous I/O device.

c) (2 points) Programmed transfer

1. Refers to the processor being able to access memory directly
2. Refers to the device controller being able to move data from/to the device to/from the memory directly
3. Refers to the processor moving the data from/to the device to/from the memory using load/store instructions
4. None of the above

# CS 2200 Spring 2011 Final Exam

Name: \_\_\_\_\_ GT Number: \_\_\_\_\_

## Disk Scheduling

5. (5 points, 10 min)

Given the following:

Total number of cylinders in the disk = 1024

Current head position is at cylinder 78, and before that it was 71.

Current requests in order of arrival =

768, 20, 513, 69, 104, 745, 25, 320, 99

a) (1 point) Show the order in which these requests are serviced using FCFS

b) (2 points) Show the order in which these requests are serviced using SSTF

c) (2 points) Show the order in which these requests are serviced using LOOK

# CS 2200 Spring 2011 Final Exam

Name: \_\_\_\_\_ GT Number: \_\_\_\_\_

## File Systems

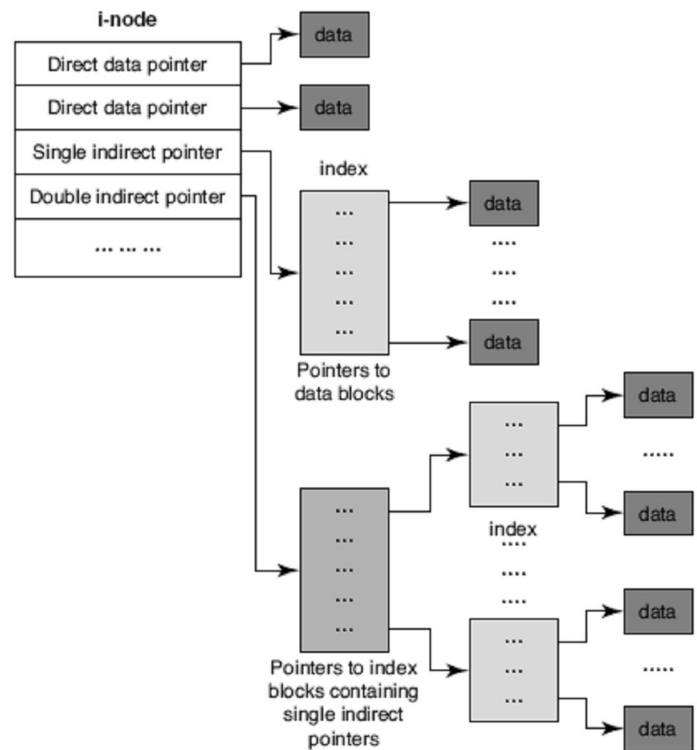
6. (10 points, 15 mins)

Given the following:

Size of index block = 1024 bytes  
Size of Data block = 1024 bytes  
Size of pointer = 8 bytes (to  
index or data blocks)

The i-node consists of  
2 direct data block pointers,  
1 single indirect pointer, and  
1 double indirect pointer.

Note that the index blocks and data  
blocks are allocated on a need basis.  
An index block is used for the i-node  
as well as for the index blocks store  
pointers to other index blocks and data  
blocks (see Figure).



- a) (2 points) How many pointers does each index block contain?
- b) (2 points) How many data blocks are used to store a 2 MB file?
- c) (3 points) How many index blocks (including the i-node for the file) are needed to store a 2 MB file?

# CS 2200 Spring 2011 Final Exam

Name: \_\_\_\_\_ GT Number: \_\_\_\_\_

d) (3 points) What is the largest file size that can be supported in this file system?

7. (5 points, 5 min)

Notes:

- Unix "**touch file1**" command creates a zero byte new file
- Unix "**ln file1 file2**" command creates a hard link
- Unix "**ln -s file1 file2**" command creates a sym link

Fill in the table below. The reference count in the table pertains to the i-node that is affected by the command in that row. If a new i-node is created, show the old reference count for that i-node as 0.

Command	New i-node created (yes/no)	Reference count	
		old	new
<b>touch f1</b>			
<b>ln -s f1 f2</b>			
<b>ln -s f2 f3</b>			
<b>ln f1 f4</b>			
<b>ln f4 f5</b>			

Use this area for rough work for this question.



# CS 2200 Spring 2011 Final Exam

Name: \_\_\_\_\_ GT Number: \_\_\_\_\_

## Parallel Systems

8. (10 points, 10 min)

In a parallel program, it is frequently necessary to divide a range of elements among threads. For example, if four threads are used to zero out all elements of array A which has a 1000 elements numbered 0 through 999, one thread can zero out elements 0 through 249, another zeroes out elements 250 through 499, still another initializes elements 500 through 749, and the remaining threads initializes elements 750 through 999. However, each thread needs to know which part of the array to process. Instead of writing different code for each thread, we can use a **shared counter variable** to assign an index to each thread. The counter starts at value zero. Each thread reads its index value from the counter and then increments the counter. The following code snippet shows the static initialization of the counter and the function called by each thread.

```
int counter = 0; /* shared counter initialized to 0 */
```

```
/* This is the function called by each thread */
```

```
void threadfun(void){  
    int myindex;
```

```
  
    myindex = counter;  
    counter++;
```

```
  
    // Now zero out my part of the array  
    for(i = 0; i < 250; i++)  
        a[myindex*250 + i] = 0;  
}
```

a) (5 points)

What is the problem with the above code?

b) (5 points) How can you fix the above code with minimal loss of concurrency?

# CS 2200 Spring 2011 Final Exam

Name: \_\_\_\_\_ GT Number: \_\_\_\_\_

9. (10 points, 10 min)

- a) (5 points) Assume that the following events happen in the order shown (T1-T5 are threads of the same process):
- T1 executes `thread_mutex_lock(L1);`
  - T2 executes `thread_mutex_lock(L1);`
  - T3 executes `thread_mutex_lock(L2);`
  - T4 executes `thread_mutex_lock(L2);`
  - T5 executes `thread_mutex_lock(L1);`

Assuming that there have been no other calls to the threads library prior to this, show the state of the internal queues in the threads library after the given five calls.

- b) (3 points) Fill in the blanks using a subset of the following phrases exactly once.

<b>Hardware</b>	<b>operating system</b>	<b>page table</b>	<b>cache</b>	<b>registers</b>
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In an SMP that supports hardware cache coherence, the \_\_\_\_\_ is responsible for ensuring that the copies of the same memory location in the different caches are kept consistent; the \_\_\_\_\_ is responsible for ensuring that the TLBs in the different processors are kept consistent; the \_\_\_\_\_ is shared among all the threads of the same process.

- c) (2 points) Deadlock (choose one of the following)
1. Is a condition where threads are not using mutex locks
  2. Is a condition where all the locks variables are in use
  3. A lock variable that is dead
  4. Is a condition where one or more threads are waiting for an event that will never happen
  5. None of the above

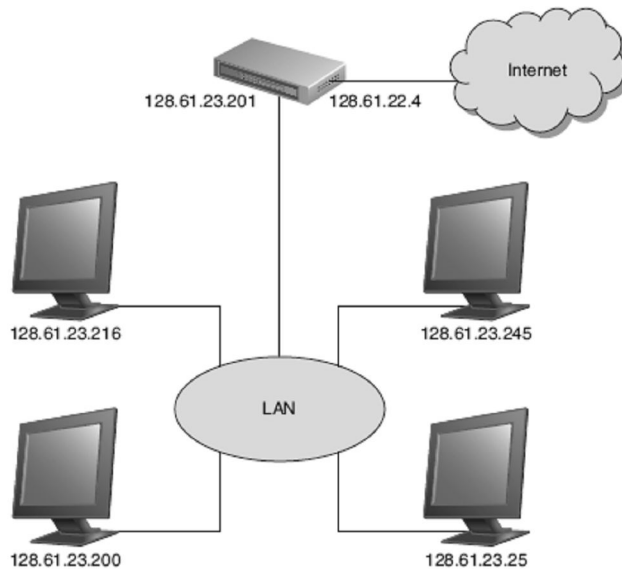
# CS 2200 Spring 2011 Final Exam

Name: \_\_\_\_\_ GT Number: \_\_\_\_\_

## Networking

10. (10 points)

a) (2 points) How many IP Networks are there in the following Figure? Assume that the top 24 bits of the 32-bit address name an IP network.



Your answer:

b) (2 points) (**Answer True/False with justification**) Circuit switching results in inefficient use of network resources.

c) (2 points)(**Answer True/False with justification**) Message switching and packet switching mean exactly the same thing.

- d) (2 points) The sequence number in a packet
1. Gives the destination address
  2. Is needed for message reconstruction at the destination
  3. Assures the integrity of the packet
  4. Is computed using cyclic redundancy check (CRC) algorithm
  5. Is the same for every packet in a given message

# CS 2200 Spring 2011 Final Exam

Name: \_\_\_\_\_ GT Number: \_\_\_\_\_

e) (2 points) Fill in the blanks using a subset of the following phrases exactly once.

<b>Ethernet</b>	<b>Token ring</b>	<b>Stop-and-wait</b>	<b>IP</b>	<b>TCP</b>
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In a LAN, under high load, \_\_\_\_\_ protocol would result in a better performance; under light load, \_\_\_\_\_ protocol would result in a better performance.

11. (10 points, 10 min)

a) (4 points) Given the following:

Message size = 100,000 bytes

Header size per packet = 100 bytes

Packet size = 1100 bytes

How many packets are needed to transmit the message assuming a 10% packet loss? Ignore fractional packet loss. Ignore ACKs. Show your work for partial credit.

b) (5 points)

A message has 10 packets, and the RTT is 2 msec. Assuming that the time to send/receive the packet and the ACK are negligible compared with the propagation time on the medium, and given no packet loss, how much time is required to complete the transmission with the sliding window protocol with a window size of 5?